

## CX4640: Homework 2

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### PROBLEM 3.5

“The kind of man who would achieve that end, is at the end of achieving his own ends.” –  
Some Pretentious Moron

#### PART A

The answer here is C. We can find this without computing the residual because a few quick

multiplications reveals that  $\begin{bmatrix} -1 \\ 1 \\ 1 \\ -1 \end{bmatrix}$  is orthogonal to  $\text{span}(A)$ .

### PROBLEM 3.17

$$\vec{v} = \vec{a} - \alpha \vec{e}_1$$
$$\alpha = -\text{sign}(v_1) \|\vec{v}\|$$

$$\vec{v} = \begin{bmatrix} 3 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

This can be verified in MATLAB.

## PROBLEM 3.18

### PART A

3, since there are 3 non-diagonal columns, and Householder makes each column a “diagonal” column by zeroing lower entries.

### PART B

The first column of  $A$  becomes  $\begin{bmatrix} -2 \\ 0 \\ 0 \\ 0 \end{bmatrix}$ .

### PART C

This column is untouched by the transformation since after that row has been zeroed, we work on the  $n - 1 \times n - 1$  submatrix  $A'$ .

### PART D

There are 6 Givens rotations required because each rotation zeros one position in the matrix, and there are 6 sub-diagonal non-zero entries in this matrix.

## PROBLEM 3.28

### PART A

The product  $(I - P_k)(I - P_{k-1})(I - P_{k-2})\dots$  can be rewritten as  $I - P_k - P_{k-1} - P_{k-2} + P_k P_{k-1} + P_{k-1} P_{k-2} + P_k P_{k-1} P_{k-2} \dots$ . Since the matrices  $P_j$  for  $0 < j \leq k$  are all rank-one matrices with each column being orthogonal to each other,  $P^m P^n = 0$  for  $m \neq n$ , reducing the expression to  $(I - P_k - \dots - P_1)$ .

### PART B

This is equivalent to the classical Gram-Schmidt procedure because it can be rewritten as:

$$a_k - P_1 a_k - P_2 a_k - \dots - P_{k-1} a_k$$

This clearly follows the logic behind Gram-Schmidt: you expand to the next dimension an orthogonal basis by subtracting the projection of  $a$  on to the current vectors in the basis.

### 0.1 PART C

First, we note that  $(I - P)$  yields a projection matrix onto a space orthogonal to the space  $P$  projects onto. From this, it is easy to see why it represents the Modified Gram-Schmidt process. In the MGS process, when you build the orthogonal basis of dimension  $n$  with  $n$  linearly

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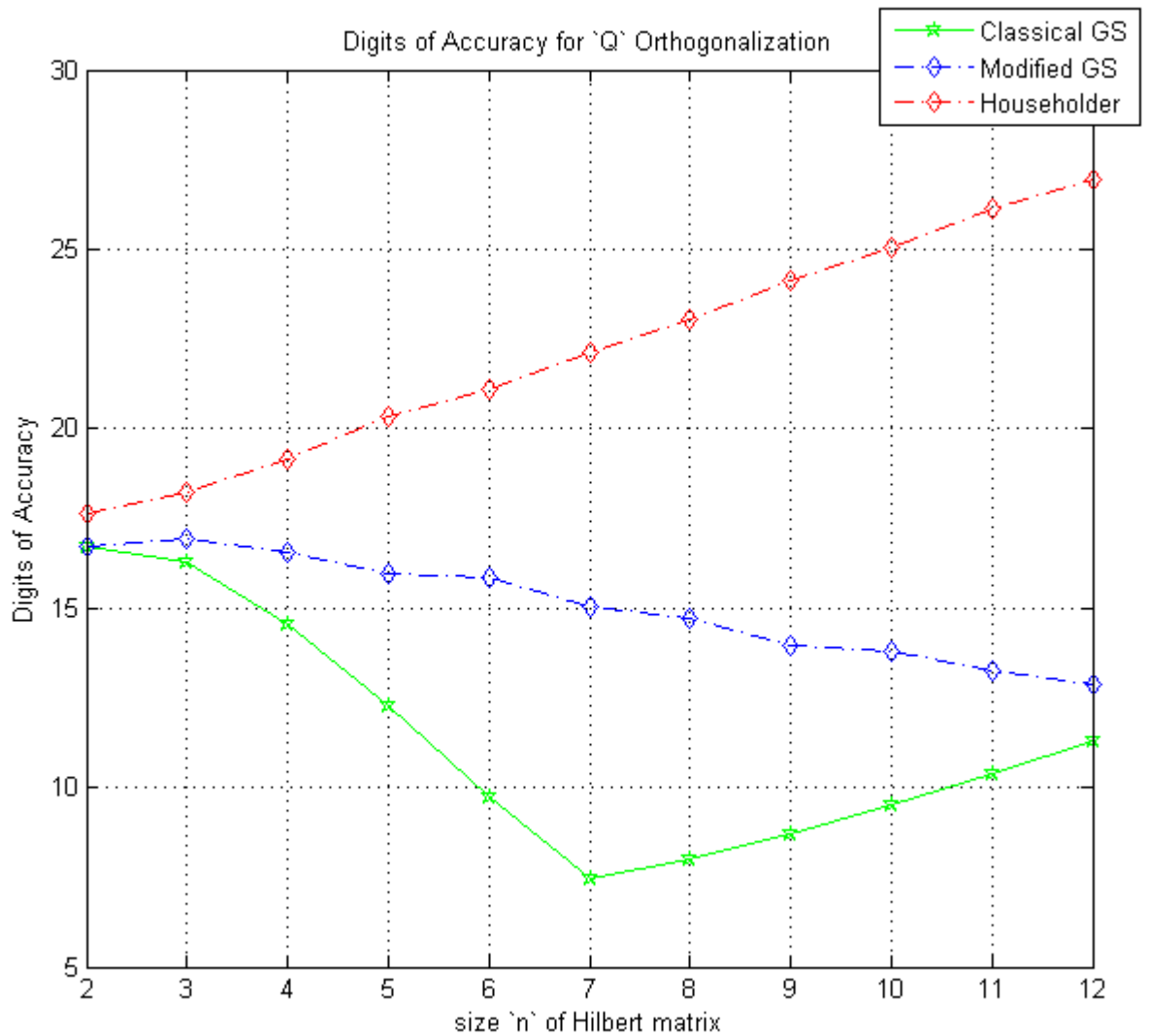
independent vectors, when you add a vector to the basis, you first subtract the component of the vector being added from the remaining vectors to be added. So if we start at  $a_1$ , no vectors have been added yet, so  $a_1$  remains unperturbed. But by the time we get to  $a_2$ , it has already had the component of  $a_1$  subtracted off. Interestingly, subtracting off the component of  $a_1$  from  $a_n$  for any  $n$ , is equivalent to projecting the vector  $a_n$  onto a space orthogonal to the space spanned by  $a_1$ . This can be computed with  $(I - P_1)$ . So by the time you start to add  $a_k$ , it has already been projected onto spaces orthogonal to  $a_1$  thru  $a_{k-1}$  respectively, thus yielding:  $(I - P_1) \dots (I - P_{k-1}) a_k$ .

## 0.2 PART D

The transformation in Part C is equivalent to the transformation in Part B, as shown in Part A. The formulation shown here is Equivalent to the ones in Part B and C because after one transformation by  $(I - P_1 - P_2 \dots P_n)$  produces a vector that is orthogonal to ALL of the spaces projected onto by  $P_1$  thru  $P_n$ , and as such, transforming an already transformed vector  $\vec{v}$  will result in  $P_k \vec{v} = 0$  for all  $k$ , leaving the only nonzero term in the original formulation to be  $I$ . As a result, any powers of  $m > 1$  are equivalent.

### PROBLEM 3.12

#### PART A AND B



In observing the above Graph, it is clear that the accuracy of Householder is the best, followed by the Modified Gram-Schmidt process, followed by the Classical Gram-Schmidt Process. Computationally, they are all bound by the number of columns in the matrix. The Classical and Modified Gram Schmidt processes are equal in run-time since they simply do things in a different order. However, the Modified process allows you to modify the A matrix in-place as it goes along, eliminating the need for separate storage for Q.

Finally, we have householder – householder is  $O(n^3)$  like G-S, and it also works on sets recursively smaller in size, so the runtime is more or less the same. However, it should be noted

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that it is more similar to the Modified G-S process in that it can modify the matrix in-place, eliminating the need for extra space for  $Q$ .