## 4. Average Time and Probabilistic Programs [WiP]

José Proença

Algorithms (CC4010) 2023/2024

CISTER - U.Porto, Porto, Portugal

https://cister-labs.github.io/alg2324





### **Overview**

- Measuring precisely performance of algorithms
- Measuring asymptotically performance of algorithms
- Analysing recursive functions
- Measuring precisely the average time of algorithms
- Possibly: sorting algorithms bubbleSort, swapSort, insertionSort, mergeSort, quickSort
- Next: analysis of sequences of operations (amortised analysis)

José Proença 2 / 11

### Recall goal

```
int count = 0;
for (int i=0; i<n; i++)
  if (v[i] == 0) count++</pre>
```

#### **RAM**

- worst-case: T(n) = 5 + 5n
- best-case: T(n) = 5 + 4n

### #array-accesses + #count-increments

- worst-case: T(n) = 2n
- best-case: T(n) = n
- average-case:

$$\overline{T}(n) = n + \sum_{0 \le r < n} P(v[r] = 0)$$

José Proença 3 / 11

## Preliminaries: series

### Recall arithmetic series

...

José Proença Preliminaries: series  $4 \ / \ 11$ 

## Recall geometric series

. . .

José Proença Preliminaries: series 5 / 11

# Calculating average cases

## Binary search

. .

José Proença Calculating average cases  $6 \ / \ 11$ 

## Two's complement

. .

José Proença Calculating average cases  $7 \ / \ 11$ 

### **Exercises**

. . .

### **Quicksort** analysis

. .

(See animation at https://visualgo.net/en/sorting.)

José Proença Calculating average cases  $9 \ / \ 11$ 

# slides by Pedro Ribeiro, slides 4 pages 9-13

### **Randomized Algorithms**

### Randomized algorithms

We call an algorithm **randomized** if its behavior is determined not only by its input but also by values produced by a **random-number generator** 

- Most programming environments offer a (deterministic)
   pseudorandom-number generator: it returns numbers that "look" statistically random
- We typically refer to the analysis of randomized algorithms by talking about the expected cost (ex: the expected running time)
- We can use probabilistic analysis to analyse randomized algorithms

- Consider rolling two dice and observing the results.
- We call this an experiment.
- It has **36 possible outcomes**: 1-1, 1-2, 1-3, 1-4, 1-5, 1-6, 2-1, 2-2, 2-3, ..., 6-4, 6-5, 6-6
- Each of these outcomes has probability 1/36 (assuming fair dice)
- What is the probability of the sum of dice being 7?

**Add** the probabilities of all the outcomes satisfying this condition: 1-6, 2-5, 3-4, 4-3, 5-2, 6-1 (probability is 1/6)



In the language of probability theory, this setting is characterized by a sample space S and a probability measure p.

- Sample Space is constituted by all possible outcomes, which are called elementary events
- In a **discrete probability distribution** (d.p.d.), the probability measure is a function p(e) (or Pr(e)) over elementary events e such that:
  - p(e) > 0 for all  $e \in S$
  - $\sum_{e \in S} p(e) = 1$
- An event is a subset of the sample space.
- For a d.p.d. the probability of an event is just the **sum** of the probabilities of its elementary events.

 A random variable is a function from elementary events to integers or reals:

Ex: let  $X_1$  be a random variable representing result of first die and  $X_2$  representing the second die.

 $X = X_1 + X_2$  would represent the sum of the two We could now ask: what is the probability that X = 7?

• One property of a random variable we care is **expectation**:

### **Expectation**

For a discrete random variable X over sample space S, the expected value of X is:

$$\mathbf{E}[X] = \sum_{e \in S} Pr(e)X(e)$$

• In words: the expectation of a random variable X is just its average value over S, where each elementary event e is weighted according to its probability.

Ex: If we roll a single die, the expected value is 3.5 (all six elementary events have equal probability).

 One possible rewrite of the previous equation, grouping elementary events:

### **Expectation (possible rewrite)**

$$\mathbf{E}[X] = \sum_{a} Pr(X = a)a$$

### Las Vegas vs. Monte Carlo

- QuickSort always returns a correct result (a sorted array) but its runtime is a random variable (with  $\mathcal{O}(n \log n)$  in expectation)
- Some randomized algorithms are not guaranteed to be correct, but their runtime is fixed.

### Las Vegas Algorithms

Randomized algorithms that always output the correct answer, and whose runtimes are random variables.

### **Monte Carlo Algorithms**

Randomized algorithms that always terminate in a given time bound, but are correct with at least some (high) probability.

José Proença Calculating average cases  $11 \ / \ 11$