

Natural Deduction Rules for Propositional Logic

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Conjunction

The rules:

$$\frac{\varphi \quad \psi}{\varphi \wedge \psi} \wedge \mathbf{I} \qquad \frac{\varphi \wedge \psi}{\varphi} \wedge \mathbf{E}_l \qquad \frac{\varphi \wedge \psi}{\psi} \wedge \mathbf{E}_r$$

Representation in Fitch's style:

$$\begin{array}{c|c} n & \vdots \\ & \varphi \\ & \vdots \\ m & \psi \\ & \vdots \\ & \varphi \wedge \psi \quad \wedge \mathbf{I}(n, m) \end{array} \qquad \begin{array}{c|c} n & \vdots \\ & \varphi \wedge \psi \\ & \vdots \\ m & \varphi \quad \wedge \mathbf{E}_l(n) \end{array} \qquad \begin{array}{c|c} n & \vdots \\ & \varphi \wedge \psi \\ & \vdots \\ m & \psi \quad \wedge \mathbf{E}_r(n) \end{array}$$

Disjunction

The rules:

$$\frac{\varphi}{\varphi \vee \psi} \vee \mathbf{I}_l \qquad \frac{\psi}{\varphi \vee \psi} \vee \mathbf{I}_r \qquad \frac{\begin{array}{c} [\varphi] \\ \vdots \\ \theta \end{array} \quad \begin{array}{c} [\psi] \\ \vdots \\ \theta \end{array}}{\varphi \vee \psi \quad \theta} \vee \mathbf{E}$$

Representation in Fitch's style:

$$\begin{array}{c}
 \begin{array}{c}
 n \quad \left| \begin{array}{c} \vdots \\ \varphi \\ \vdots \\ \varphi \vee \psi \end{array} \right. \quad \vee \mathbf{I}_l(n)
 \end{array}
 \qquad
 \begin{array}{c}
 n \quad \left| \begin{array}{c} \vdots \\ \psi \\ \vdots \\ \varphi \vee \psi \end{array} \right. \quad \vee \mathbf{I}_r(n)
 \end{array}
 \qquad
 \begin{array}{c}
 \begin{array}{c}
 \vdots \\
 n \quad \varphi \vee \psi \\
 \vdots \\
 m_1 \quad \left| \begin{array}{c} \varphi \\ \vdots \end{array} \right. \\
 k_1 \quad \left| \begin{array}{c} \theta \end{array} \right. \\
 \vdots \\
 m_2 \quad \left| \begin{array}{c} \psi \\ \vdots \end{array} \right. \\
 k_2 \quad \left| \begin{array}{c} \theta \end{array} \right. \\
 \vdots \\
 \theta
 \end{array}
 \end{array}
 \qquad
 \vee \mathbf{E}(n, m_1 - k_1, m_2 - k_2)
 \end{array}$$

Negation

The rules:

$$\frac{\begin{array}{c} [\varphi] \\ \vdots \\ \bot \end{array}}{\neg \varphi} \neg \mathbf{I}
 \qquad
 \frac{\neg \neg \varphi}{\varphi} \neg \mathbf{E}$$

Representation in Fitch's style:

$$\begin{array}{c|c}
n & \vdots \\
& \hline
& \varphi \\
& \vdots \\
m & \perp \\
& \hline
& \neg\varphi
\end{array}
\quad \neg\mathbf{I}(n-m)$$

$$\begin{array}{c|c}
n & \vdots \\
& \neg\neg\varphi \\
& \vdots \\
& \varphi
\end{array}
\quad \neg\mathbf{E}(n)$$

False

The rules:

$$\frac{\varphi \quad \vdots \quad \neg\varphi}{\perp} \perp\mathbf{I}$$

$$\frac{\perp}{\varphi} \perp\mathbf{E}$$

Representation in Fitch's style:

$$\begin{array}{c|c}
n & \vdots \\
& \varphi \\
& \vdots \\
m & \neg\varphi \\
& \vdots \\
& \perp
\end{array}
\quad \perp\mathbf{I}(n, m)$$

$$\begin{array}{c|c}
n & \vdots \\
& \perp \\
& \vdots \\
& \varphi
\end{array}
\quad \perp\mathbf{E}(n)$$

Implication

The rules:

$$\frac{[\varphi] \quad \vdots \quad \psi}{\varphi \rightarrow \psi} \rightarrow\mathbf{I}$$

$$\frac{\varphi \rightarrow \psi \quad \varphi}{\psi} \rightarrow\mathbf{E}$$

Representation in Fitch's style:

$$\begin{array}{c} \vdots \\ n \quad | \quad \varphi \\ \hline \vdots \\ m \quad | \quad \psi \\ \varphi \rightarrow \psi \end{array} \quad \rightarrow \mathbf{I}(n-m)$$

$$\begin{array}{c|c} & \vdots \\ n & \varphi \rightarrow \psi \\ & \vdots \\ m & \varphi \\ & \vdots \\ & \psi \end{array} \rightarrow \mathbf{E}(n, m)$$