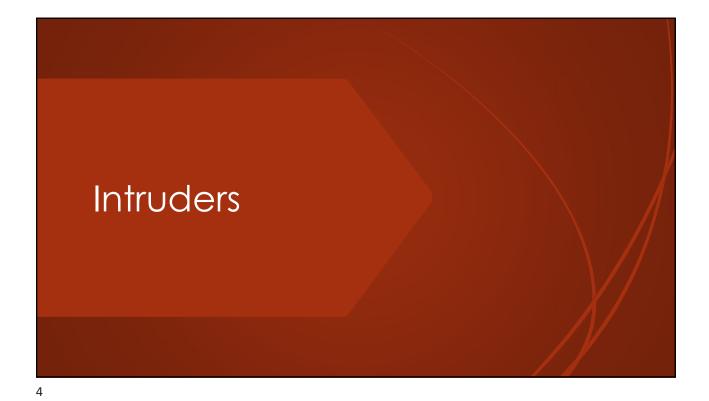




Learning objectives

- types of intruder behavior patterns;
- principles of and requirements for intrusion detection;
- key features of host-based intrusion detection;
- distributed host-based intrusion detection;
- key features of network-based intrusion detection;
- intrusion detection exchange format;
- honeypots;
- Snort



One of the key threats to security is the use of some form of hacking by an intruder (AKA hacker or cracker).
 Several reports indicate also:

 A general increase in malicious hacking activity
 An increase in attacks specifically targeted at individuals in organizations and the IT systems they use.

 Breaches (source: Verizon)
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Are either individuals, usually working as insiders, or members of a larger group of outsider attackers, who are motivated by social or political causes

Also know as hacktivists

Alm of their attacks is often to promote and publicize their cause typically through:

Website defacement Denical of service attacks in negative publicity or compromise of their targets.

Classes of Intruders – State-Sponsored Organizations

Groups of hackers sponsored by governments to conduct espionage or sabotage activities

Also known as Advanced Persistent Threats (APTs) due to the covert nature and persistence over extended periods involved with any attacks in this class

Widespread nature and scope of these activities by a wide range of countries from China, Russia, USA, UK, and their intelligence allies

Hackers with motivations other than those previously listed

Include classic hackers or crackers who are motivated by technical challenge or by peer-group esteem and reputation

Many of those responsible for discovering new categories of buffer overflow vulnerabilities could be regarded as members of this class

Given the wide availability of attack toolkits, there is a pool of "hobby hackers" using them to explore system and network security

Intruder Skill Levels – Apprentice

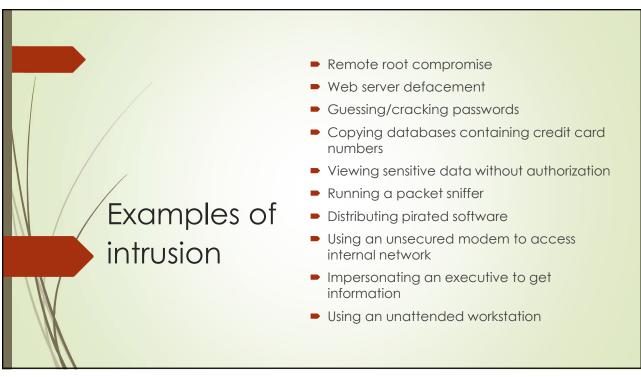
- Hackers with minimal technical skill who primarily use existing attack toolkits
- They likely comprise the largest number of attackers, including many criminal and activist attackers
- Given their use of existing known tools, these attackers are the easiest to defend against
- Also known as "script-kiddies" due to their use of existing scripts (tools)

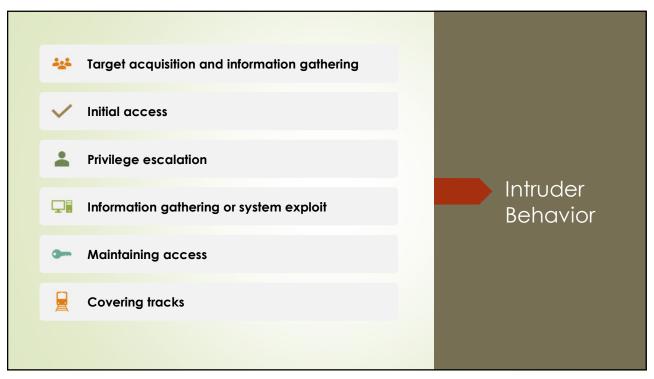
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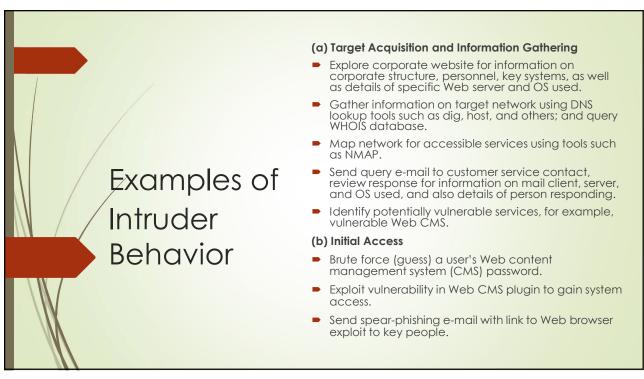
Intruder Skill Levels – Journeyman

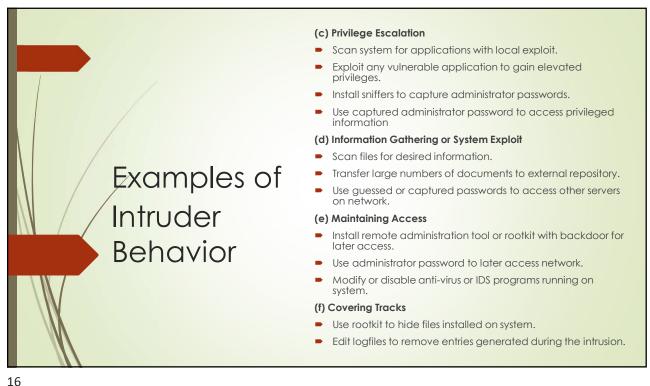
- Hackers with sufficient technical skills to modify and extend attack toolkits to use newly discovered, or purchased, vulnerabilities
- They may be able to locate new vulnerabilities to exploit that are similar to some already known
- Hackers with such skills are likely found in all intruder classes
- Adapt tools for use by others

Intruder Skill Levels – Master Hackers with high-level technical skills capable of discovering brand new categories of vulnerabilities Write new powerful attack toolkits Some of the better known classical hackers are of this level Some are employed by state-sponsored organizations Defending against these attacks is of the highest difficulty



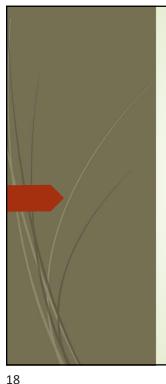






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Why intrusion detection?

- Authentication facilities, access control facilities, and firewalls all play a role in countering intrusions.
- Intrusion detection is another line of defense:
 - 1. If an intrusion is detected quickly enough, the intruder can be identified and ejected from the system before any damage is done or any data are
 - 2. An effective intrusion detection system can serve as a deterrent, thus acting to prevent intrusions.
 - 3. Intrusion detection enables the collection of information about intrusion techniques that can be used to strengthen intrusion prevention measures.

Assumptions:

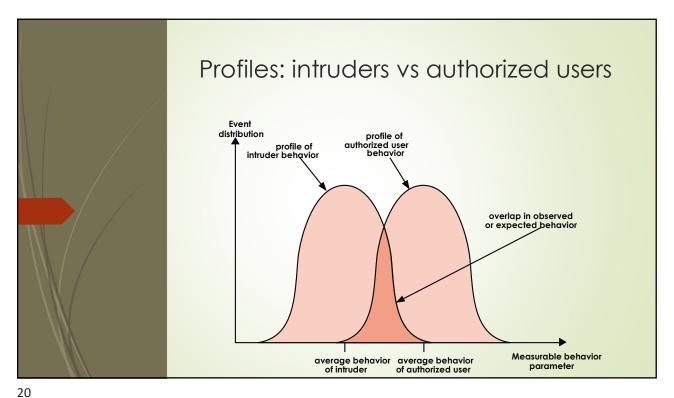
- The behavior of the intruder differs from that of a legitimate user in ways that can be quantified.
- Not a crisp, exact distinction between an attack by an intruder and the normal use of resources by an authorized user.

Intrusion Detection System (IDS)

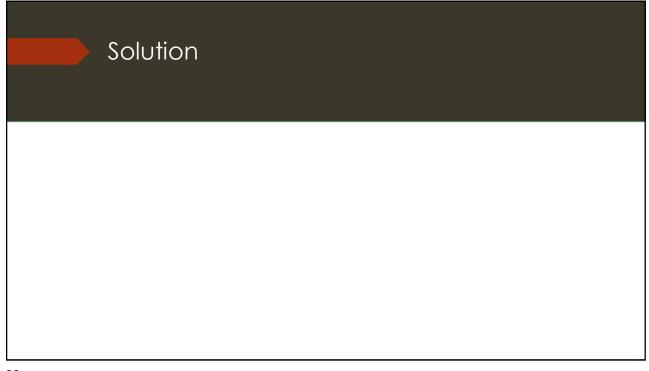
- Host-based IDS (HIDS)
 - Monitors the characteristics of a single host for suspicious activity
- Network-based IDS (NIDS)
 - Monitors network traffic and analyzes network, transport, and application protocols to identify suspicious activity
- Distributed or hybrid IDS
 - Combines information from a number of sensors, often both host and network based, in a central analyzer that is able to better identify and respond to intrusion activity

Comprises three logical components:

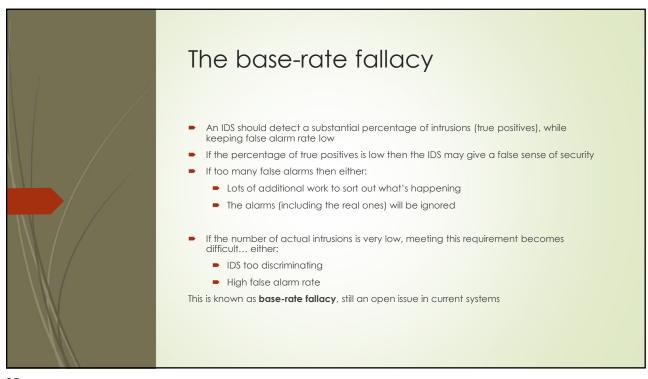
- Sensors collect data
- Analyzers determine if intrusion has occurred
- User interface view output or control system behavior

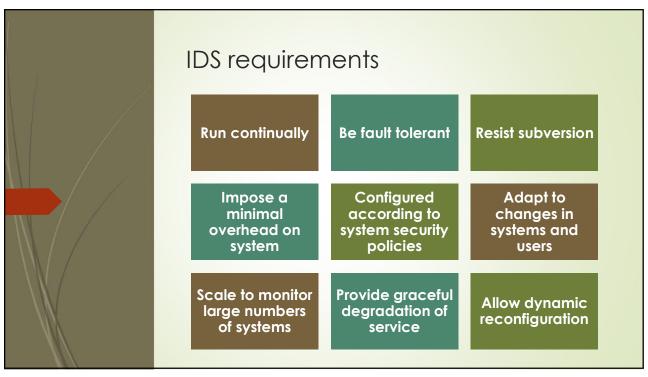


Suppose there are 2 actual intrusions for every 1000 authorized users, and the overlapping area covers 1% of the authorized users and 50% of the intruders. a) Sketch the event distribution and argue that this is not an unreasonable depiction. b) What is the probability that an event that occurs in this region (the overlapping area) is that of an authorized user? Keep in mind that 50% of all intrusions fall in this region.



Profiles of Behavior of Intruders and Authorized Users An early study of intrusion (Anderson 1980, still valid) postulated that: it is possible to distinguish between an outside attacker and a legitimate user with reasonable confidence. Patterns of legitimate user behavior can be established by observing past history, and significant deviation from such patterns can be detected. Detecting an inside attacker (a legitimate user acting in an unauthorized fashion) is more difficult: the distinction between abnormal and normal behavior may be small: such violations would be undetectable solely through the search for anomalous behavior; however, insider behavior might be detectable by an intelligent definition of the class of conditions that suggest unauthorized use.

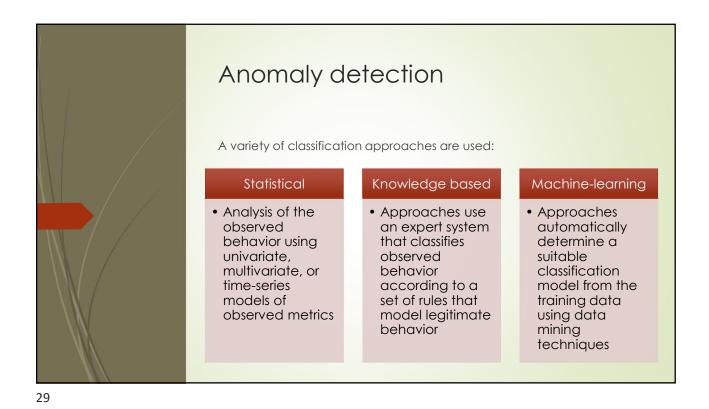




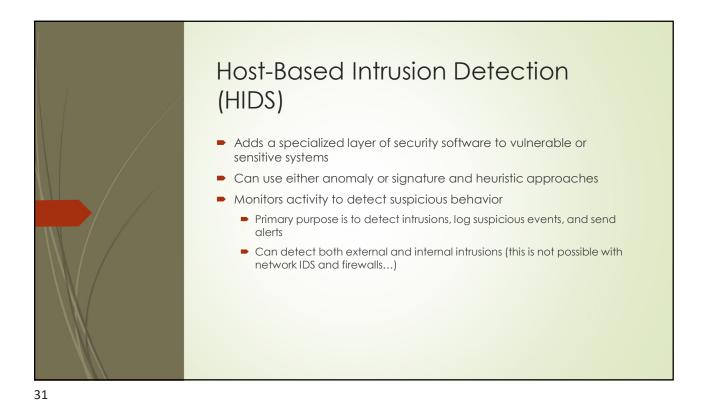
Review question What is an IDS in the end? Is it a hardware appliance? A piece of software? A process? A component of an O.S. or of a firewall?

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Analysis approaches in intrusion detection Signature/heuristic Anomaly detection detection Uses a set of known malicious Involves the collection of data relating to the behavior of data patterns or attack rules that legitimate users over a period of are compared with current time behavior Also known as misuse detection Current observed behavior is analyzed to determine whether Can only identify known attacks this behavior is that of a legitimate for which it has patterns or rules user or that of an intruder Can detect even unknown, zeroday attacks



Rule-based heuristic Signature identification approaches Match a large collection of known Signature or Involves the use of rules for identifying patterns of malicious data against data known penetrations or penetrations that stored on a system or in transit over a would exploit known weaknesses Heuristic network Detection The signatures need to be large enough Rules can also be defined that identify to minimize the false alarm rate, while suspicious behavior, even when the still detecting a sufficiently large fraction behavior is within the bounds of of malicious data established patterns of usage Widely used in anti-virus products, Typically rules used are specific to the network traffic scanning proxies, and in machine and to the operating system NIDS SNORT is an example of a rule-based



Data Sources and Sensors

Common data sources include:

System call traces

A fundamental component of intrusion detection is the sensor that collects data

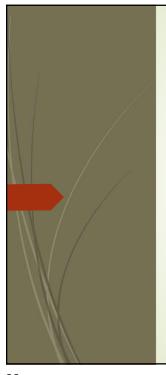
Common data sources include:

System call traces

Audit (log file) records

File integrity checksums

Registry access (Windows)



Anomaly HIDS

- In UNIX/Linux most effective results based on the analysis of system call traces
 - analyses sequences of system calls invoked by a process over time
 - System call traces can be produced by inserting hooks in the OS itself (e.g.BSM audit module)
 - Anomaly detection often on machine learning:
- In Windows The analysis of system calls activations does not work well:
 - Extensive use of Dynamic Link Libraries (DLL) that often hide the system calls
 - Even the analysis of the registry or of audit logs does not work very well
 - Best approaches analyze traces of DLL functions invocations (similar to Linux for the system calls)
- In all OSs: cryptographic checksums to verify integrity of files
 - Problematic that files can be legally modified (and they are many, even in the OS)

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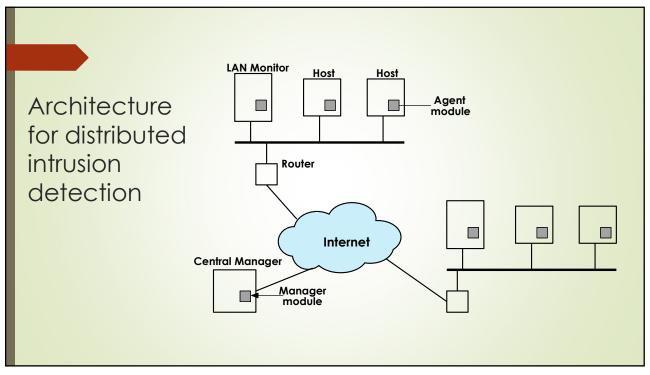
Linux System Calls and Windows DLLs Monitored

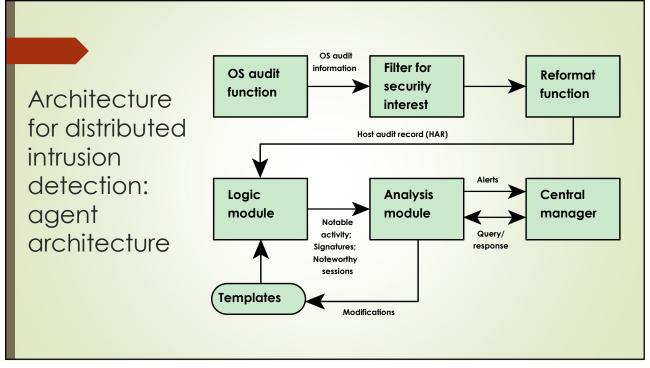
Ubuntu Linux system calls

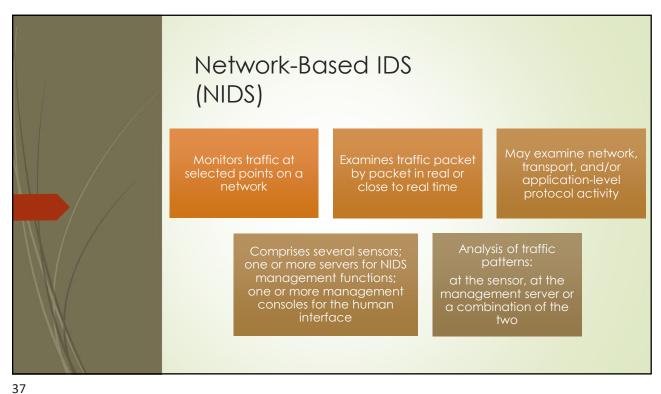
accept, access, acct, adjtime, aiocancel, aioread, aiowait, aiowrite, alarm, async_daemon, auditsys, bind, chdir, chmod, chown, chroot, close, connect, creat, dup, dup2, execv, execve, exit, exportfs, fchdir, fchmod, fchown, fchroot, fcntl, flock, fork, fpathconf, fstat, fstat, fstatfs, fsync, ftime, ftruncate, getdents, getdirentries, getdomainname, getdopt, getdatblesize, getfh, getgid, getgroups, gethostid, gethostname, getilimer, getmsg, getpagesize, getpeername, getpgrp, getpid, getpriority, getrlimit, getrusage, getsockname, getsockopt, gettimeofday, getuid, gtty, ioctl, kill, killpg, link, listen, Iseek, Istat, madvise, mctl, mincore, mkdir, mknod, mmap, mount, mount, mprotect, mpxchan, msgsys, msync, munmap, nfs_mount, nfssvc, nice, open, pathconf, pause, pcfs_mount, phys, pipe, poll, profil, ptrace, putmsg, quota, quotactl, read, readlink, readv, reboot, recv, recvfrom, recvmsg, rename, resuba, rfssys, rmdir, sbreak, sbrk, select, semsys, send, sendmsg, sendto, setdomainname, setdopt, setgid, setgroups, sethostid, sethostname, setitimer, setpgid, setggrp, setpgrp, setpriority, setquota, setregid, setreuid, setrlimit, setsid, setsockopt, settimeofday, setuid, shmsys, shutdown, sigblock, sigpause, sigpending, sigsetmask, sigstack, sigsys, sigvec, socket, socketaddr, socketpair, sstk, stat, stat, statfs, stime, stty, swapon, symlink, sync, sysconf, time, times, truncate, umask, umount, uname, unlink, unmount, ustat, utime, utimes, vadvise, vfork, vhangup, vlimit, ypixsys, vread, vtimes, vtrace, vwrite, wait, wait3, wait4, write, writev

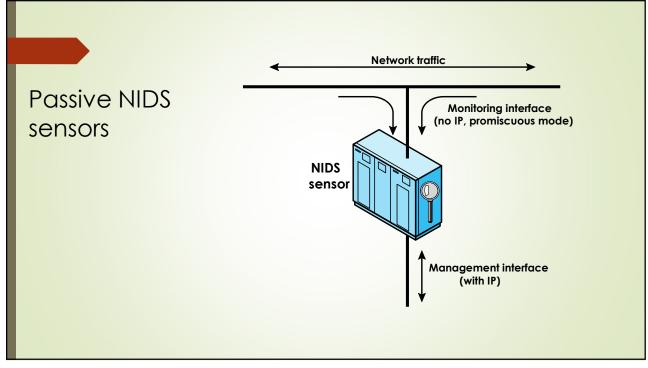
Key Windows DLL and executables

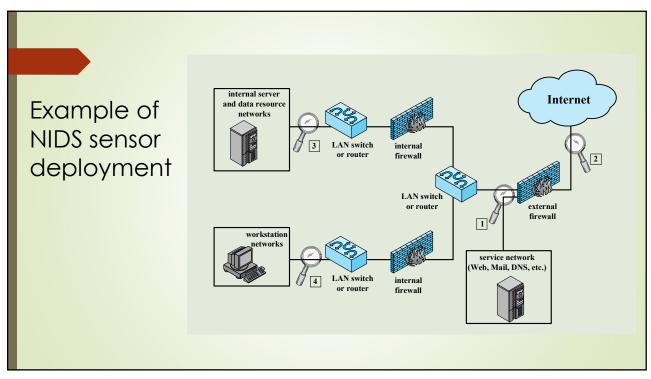
kernel32 msvcpp msvcrt mswsock ntdll ntoskrnl user32 ws2_32

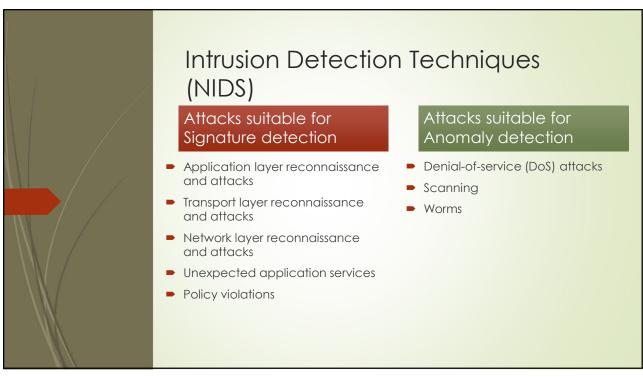


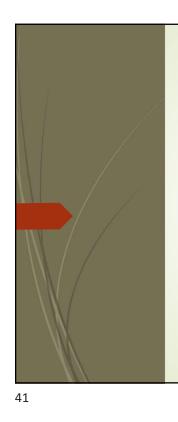












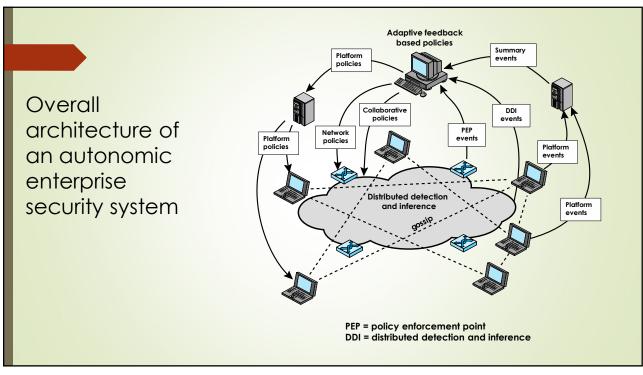
Stateful Protocol Analysis (SPA)

Subset of anomaly detection:

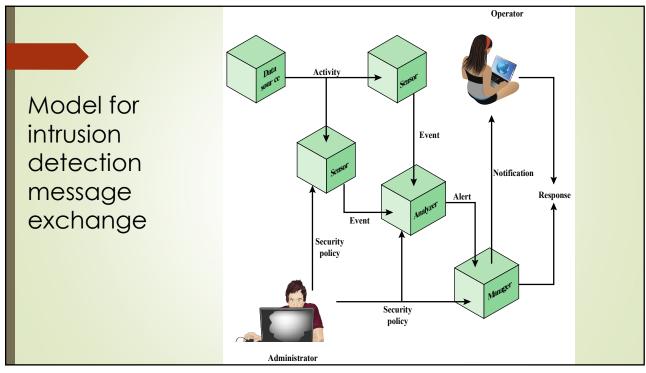
- Compares observed network traffic against predetermined universal vendor supplied profiles of benign protocol traffic
 - Different than anomaly techniques trained with organizationspecific traffic protocols
- Understands and tracks network, transport, and application protocol states to ensure they progress as expected
- A key disadvantage is the high resource use it requires

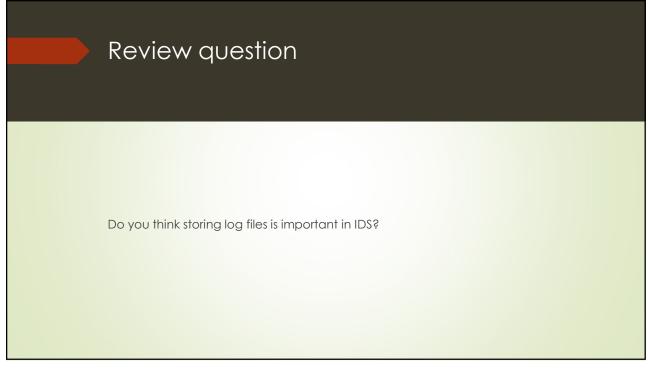
Logging of Alerts

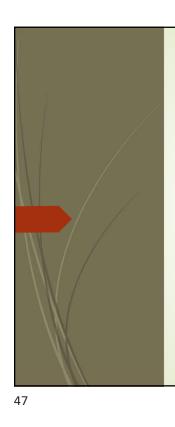
- Typical information logged by a NIDS sensor includes:
 - Timestamp
 - Connection or session ID
 - Event or alert type
 - Rating (e.g., priority, severity, impact, confidence)
 - Network, transport, and application layer protocols
 - Source and destination IP addresses
 - Source and destination TCP or UDP ports, or ICMP types and codes
 - Number of bytes transmitted over the connection
 - Decoded payload data, such as application requests and responses
 - State-related information (e.g., authenticated username)



Intrusion Detection Message Exchange Requirements (RFC 4766) • defines requirements for the Intrusion Detection Message **IETF Intrusion** Exchange Format (IDMEF) specifies requirements for a communication protocol for communicating IDMEF Detection The Intrusion Detection Message Exchange Working Format (RFC 4765) describes a data model to represent information exported Group by intrusion detection systems and explains the rationale for presents an implementation of the data model in the Extensible Markup Language (XML) and provides an XML Document Type Definition Defines data formats and exchange procedures for sharing The Intrusion Detection Exchange Protocol information of interest to intrusion detection and response systems describes the Intrusion Detection Exchange Protocol (IDXP): and to management systems that an application-level protocol for exchanging data between may need to interact with them intrusion detection entities • IDXP supports mutual authentication, integrity, and The working group issued several confidentiality over a connection-oriented protocol RFCs in 2007:

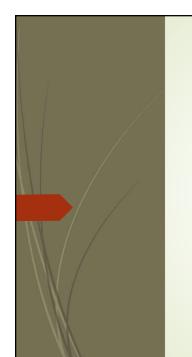






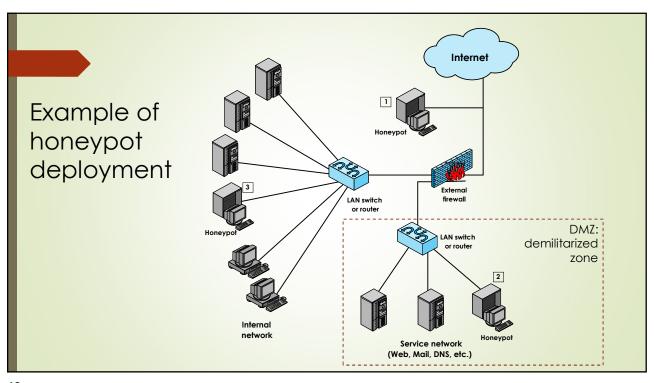
Honeypots

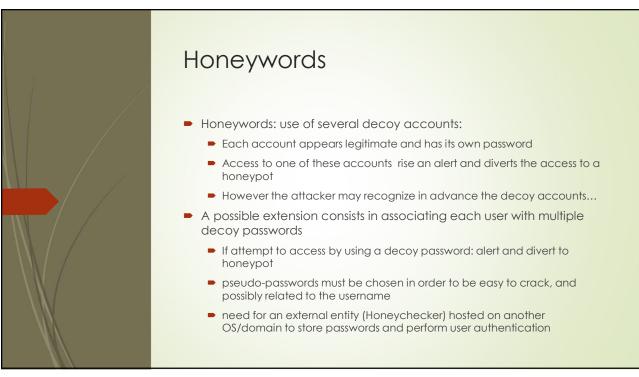
- Decoy systems designed to:
 - Lure a potential attacker away from critical systems
 - Collect information about the attacker's activity
 - Encourage the attacker to stay on the system long enough for administrators to respond
- Systems are filled with fabricated information that a legitimate user of the system wouldn't access
- Resources that have no production value
 - Hence any incoming communication is most likely a probe, scan, or an attack
 - Initiated outbound communication suggests that the system has probably been compromised

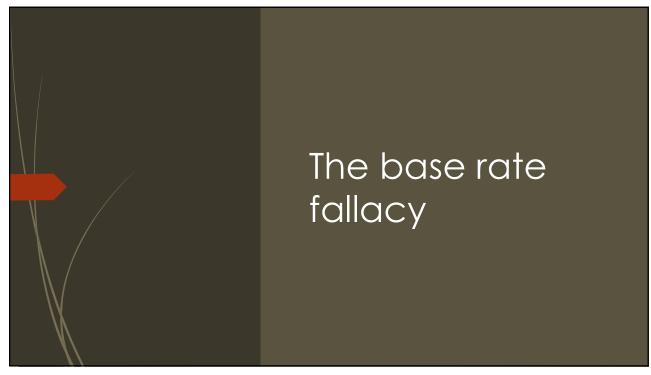


Honeypot Classifications

- Low interaction honeypot
 - Consists of a software package that emulates particular IT services or systems well
 enough to provide a realistic initial interaction, but does not execute a full version of
 those services or systems
 - Provides a less realistic target
 - Often sufficient for use as a component of a distributed IDS to warn of imminent attack
- High interaction honeypot
 - A real system, with a full operating system, services and applications, which are instrumented and deployed where they can be accessed by attackers
 - Is a more realistic target that may occupy an attacker for an extended period
 - However, it requires significantly more resources
 - If compromised could be used to initiate attacks on other systems







Conditional probability

- We often want to know a probability that is conditional on some event.
- The effect of the condition is to remove some of the outcomes from the sample space.
- Example: what is the probability of getting a sum of 8 on the roll of two dice if we know that the face of at least one die is an even number?
- 1. Because one die is even and the sum is even, the second die must show an even number.
- 2. Thus, there are three equally likely successful outcomes: (2, 6), (4, 4), and (6, 2)
- 3. While the total set of possibilities is [36 (number of events with both faces odd)] = 36 (3 * 3) = 27.
- 4. The resulting probability is 3/27 = 1/9.

Conditional probability

the conditional probability of an event A assuming the event B has occurred, denoted by Pr[A | B], is defined as the ratio:

$$\Pr[A \mid B] = \frac{\Pr[AB]}{\Pr[B]}$$

- \blacksquare where we assume Pr[B] is not zero.
- In our example:
 - \rightarrow A = {sum of 8}
 - \blacksquare B = {at least one die even}.

- The quantity Pr[A | B] encompasses all outcomes in which the sum is 8 **and** at least one die is even
 - As we have seen, there are three such outcomes.
- \blacksquare Thus, Pr[AB] = 3/36 = 1/12.
- A moment's thought should convince you that Pr[B] = 3/4.
- Follows that

$$\Pr[A \mid B] = \frac{1/12}{3/4} = \frac{1}{9}$$

Which confirms the previous reasoning.

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Conditional probability

Two events A and B are called independent if

$$Pr[AB] = Pr[A] \cdot Pr[B]$$

Recall that if A and B are independent,

$$Pr[A \mid B] = Pr[A]$$

and

$$Pr[B \mid A] = Pr[B]$$

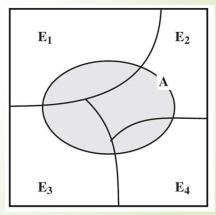
Conditional probability – Total probability

Total probability

- **given** a set of mutually exclusive events $E_1, E_2, ..., E_n$,
- such that the union of these events covers all possible outcomes,
- and given an arbitrary event A, then it can be shown that:

$$\Pr[A] = \sum_{i=1}^{n} (\Pr[A \mid E_i] \cdot \Pr[E_i])$$

Total probability theorem illustrated



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Conditional probability – Bayes theorem

Bayes' theorem:

- it is used to calculate "posterior odds"...
- ... the probability that something really is the case, given evidence in favor of it.

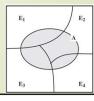
Example:

- if A happens,
- what is the probability that a given E_i is true?

The theorem may also be stated as follows:

$$\Pr[E_i \mid A] = \frac{\Pr[A \mid E_i] \cdot \Pr[E_i]}{\Pr[A]} =$$

$$= \frac{\Pr[A \mid E_i] \cdot \Pr[E_i]}{\sum_{i=1}^{n} (\Pr[A \mid E_i] \cdot \Pr[E_i])}$$



Conditional probability – Bayes theorem

Example:

- Suppose we are transmitting a sequence of zeroes and ones over a noisy transmission line.
 - Let SO and S1 be, at a given time, the events a 0 is sent and a 1 is sent, respectively,
 - ▶ Let RO and R1 be the events that a 0 is received and a 1 is received.
- Suppose we know the probabilities of the source:

$$Pr[S1] = p \text{ and } Pr[S0] = 1 - p$$

 ... and we observe the line to determine how frequently an error occurs when a one is sent and when a zero is sent, so that the following probabilities are calculated:

$$Pr[R0 \mid S1] = p_A \text{ and } Pr[R1 \mid S0] = p_B$$

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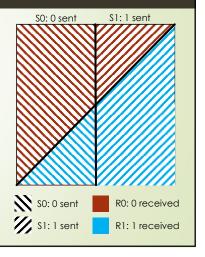
Conditional probability – Bayes theorem

If a zero is received (R0), we can then calculate the conditional probability of an error, namely the conditional probability that a one was sent (\$1) given that a zero was received, using Bayes' theorem:

$$\Pr[S1 \mid R0] = \frac{\Pr[R0 \mid S1] \cdot \Pr[S1]}{\Pr[R0 \mid S1] \cdot \Pr[S1] + \Pr[R0 \mid S0] \cdot \Pr[S0]} =$$

$$=\frac{p_A \cdot p}{p_A \cdot p + (1-p_B) \cdot (1-p)}$$

Bayes: $\Pr[E_i \mid A] = \frac{\Pr[A \mid E_i] \cdot \Pr[E_i]}{\sum_{i=1}^{n} (\Pr[A \mid E_i] \cdot \Pr[E_i])}$





Consider a patient that has a test for some disease that comes back positive (indicating he has the disease). You know that:

- The accuracy of the test is 87%:
 - if a patient has the disease, 87% of the time, the test yields the correct result,
 - if the patient does not have the disease, 87% of the time, the test yields the correct result.
- The incidence of the disease in the population is 1%.

Given that the test is positive, how probable is it that the patient does not have the disease?

That is, what is the probability that this is a false alarm?

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The base rate fallacy

- We need Bayes' theorem to get the correct answer:
- The accuracy of the test is 87%
- The incidence of the disease in the population is 1%.

$$\Pr[well \mid positive] = \frac{\Pr[positive \mid well] \cdot \Pr[well]}{\Pr[positive \mid disease] \cdot \Pr[disease] + \Pr[positive \mid well] \cdot \Pr[well]} = \frac{\Pr[positive \mid disease]}{\Pr[positive \mid disease]} = \frac{\Pr[positive \mid well] \cdot \Pr[well]}{\Pr[positive \mid disease]} = \frac{\Pr[positive \mid disease]}{\Pr[positive \mid diseas$$

$$= \frac{0.13 \cdot 0.99}{0.87 \cdot 0.01 + 0.13 \cdot 0.99} = 0.937$$

■ Which means that in most cases it's a false alarm

The base rate fallacy

The problem is that, when proposed to people, the answer is:

- Many subjects gave the answer 13%.
- The vast majority, including many physicians, gave a number below 50%.
- Many physicians who guessed wrong lamented,
 - "If you are right, there is no point in making clinical tests!"
- The reason most people get it wrong is that they do not consider the basic rate of incidence (the base rate) when intuitively solving the problem.
- This error is known as the base rate fallacy.

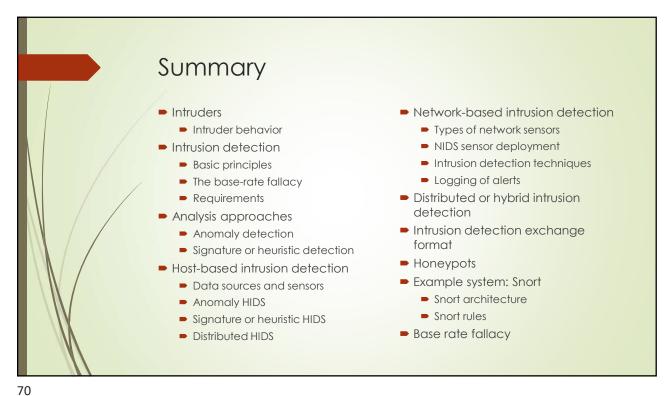
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The base rate fallacy

What happens when probabilities change?

accuracy	0,87	0,99	0,999	0,99
incidence	0,01	0,01	0,01	0,001
Pr[well positive] (false alarm rate)	0,94	0,5	0,09	0,91

In actual situations it was found that the probabilities associated with IDSs were such that the false alarm rate was unsatisfactory.



Threat Level	Signature		
Low	1 P1 + 1 P2		
Medium	1 P3 + 1 P4		
High	2 P4		
Exercise			
A decentralized NIDS is operating with two nodes in the network monitoring anomalous inflows of traffic. In addition, a central node is present, to generate an alarm signal upon receiving input signals from the two distributed nodes.			
The signatures of traffic inflow into the two IDS nodes follow one of four patterns: P1, P2, P3, and P4 (all equiprobable).			
The threat levels are classified by the central node based upon the observed traffic by the two NIDS at a given time and are given by the above table			
If, at a given time instance, at least one distributed node generates an alarm signal P4, what is the probability that the observed traffic in the network will be classified at threat level "Medium" or "High"?			

Solution

- The signatures of traffic inflow into the two IDS
- nodes follow one of four patterns: P1, ..., P4.

 If,at least one node generates an alarm P4, what is the probability that the observed traffic will be classified at threat level "Medium" or "High"?

Threat Level	Signature
Low	1 P1 + 1 P2
Medium	1 P3 + 1 P4
High	2 P4

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Exercise

The network of an organization has two intrusion detection sensors aimed at detecting cyberattacks in real-time by means of anomaly detection. The two sensors are based on a different technology, and they have the following accuracy in the detection of DoS, worms or scan attacks:

Accuracy	DoS	Scan	Worm
Sensor1	-	0.75	0.82
Sensor2	0.79	0.91	_

- Assume that, from historical records, 10% of the attacks are DoS, 50% are Scan and 40% are worms.
- If Sensor 2 raises an alarm for a DoS attack. What is the probability this is a false positive?

Solution

Accuracy	DoS	Scan	Worm
Sensor1	-	0.75	0.82
Sensor2	0.79	0.91	-

10% of the attacks are DoS, 50% are Scan and 40% are worms.

If Sensor 2 raises an alarm for a DoS attack, what is the probability that it is a false positive?

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Exercise

The network of an organization has two intrusion detection sensors aimed at detecting cyberattacks in real-time by means of anomaly detection. The two sensors are based on a different technology, and they have the following accuracy in the detection of DoS, worms or scan attacks:

Accuracy	DoS	Scan	Worm
Sensor1	-	0.75	0.82
Sensor2	0.79	0.91	_

- Assume that, from historical records, 10% of the attacks are DoS, 50% are Scan and 40% are worms.
- If Sensor 1 raises an alarm for a worm attack, what is the probability that it is a false positive?

Solution

Accuracy	DoS	Scan	Worm
Sensor1	-	0.75	0.82
Sensor2	0.79	0.91	-

10% of the attacks are DoS, 50% are Scan and 40% are worms.

If Sensor 1 raises an alarm for a worm attack, what is the probability that it is a false positive?

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Exercise

A taxicab was involved in a fatal hit-and-run accident at night. Two cab companies, the Green and the Blue, operate in the city. You are told that:

- 85% of the cabs in the city are Green and 15% are Blue.
- A witness identified the cab as Blue.

The court tested the reliability of the witness under the same circumstances that existed on the night of the accident and concluded that the witness was correct in identifying the color of the cab 80% of the time. What is the probability that the cab involved in the incident was Blue rather than Green?

85% of the cabs are Green; 15% are Blue. A witness identified the cab as Blue. the witness was correct 80% of the times. What is the probability that the cab involved in the incident was Blue rather than Green?

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