One-Time Passwords

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1

One-Time Password

- One-Time Password (OTP)
  - A password that is valid for only one login session or transaction
  - A.k.a. dynamic password, dynamic pin
- Pros
  - Not vulnerable to replay attack
  - Not vulnerable to password-reuse attack
- Cons
  - Hard to remember, so you need additional technology

2

Methods

- Based on time-synchronization
- Based on the previous password
- Based on a challenge

3

Time synchronization (→)

- Prover
  - Token, $clock_p$
- Verifier:
  - Authentication server, $clock_v$
- Problems
  - Clocks of prover and verifier are roughly synchronised
  - Network latency, user delay, clock skews

4

Time synchronization (→)

- Time Parameters
  - T0 = initial time
  - T = current time
  - X = time steps in a second
  - C = # of time-steps between T0 and T
    - C = (T – T0)/X
  - W = acceptance window
- Key
  - Key k shared between prover and verifier

Time synchronization

- The protocol
  - **Prover** **Authenticator**
    - $T_p \leftarrow clock_p()$
    - $C_p = (T_p - T_0)/X$
    - $HOTP = HMAC_k(C_p)$

----------HOTP---------->

$T_v \leftarrow clock_v()$
for all t in $[T_v – W/2, T_v + W/2]$ {
  $C_v = (t – T_0)/X$;
  if (HOTP == $H_k(C_v)$
    return TRUE;
}
return FALSE

< ----------TRUE|FALSE----------

Time synchronization

- For more details
  - D. M'Raihi, S. Machani, M. Pei, J. Rydell. TOTP: Time-Based One-Time Password Algorithm, RFC 6238, IETF, May 2011

7

Lamport's scheme

- Hash List
  - Setup
    - Seed $p_0 \leftarrow$ random()
    - $p_i = H(p_{i-1})$, i = 1, …, n
    - p_n is stored at the verifier by *offline means*
  - Password verification
    - Prover sends p_{n-1} to Verifier
    - Verifier returns ($p_n == H(p_{n-1})$)
    - *More in general*
      - Verifier returns ($p_i == H(p_{i-1})$) or ($p_i == H^i(p_0)$)
      - 2nd form in case p_i are not verified sequentially

8

Challenge-response

- Prover and Verifier share a key K
  - **Verifier** **Prover**

chl ← random()
send(Prover, chl)
----------->
res = H_k(chl)
send(Verifier, res)
< ----------
return (res == H_k(ch))