EXECUTION MONITOR AS PROGRAMMING LANGUAGE PRIMITIVE



Access Control Lists

Interpreter: OCAML code (eval)

$$\frac{env(x) = v}{env > Den x \Rightarrow v}$$

$$\frac{env \triangleright e_1 \Rightarrow v_1, env \triangleright e_2 \Rightarrow v_2}{env \triangleright Sum(e_1, e_2) \Rightarrow v_1 + v_2}$$

eval Den x, env ->lookup x env

eval Sum(e1, e2) env ->
eval e1 env + eval e2 env

Wrapped Interpreter: eval env acl

$$\frac{env(x) = v}{env \ acl \ \triangleright Den \ x \Rightarrow v}$$

```
\frac{add \in al \ env \rhd e_1 \Rightarrow v_1 \ , env \rhd e_2 \Rightarrow v_2}{env \ acl \ \rhd \ Sum(e_1, e_2) \Rightarrow v_1 + v_2}
```

eval Den x, env, al -> lookup x env

```
eval Sum(e1, e2) env al ->
if (emCheck alist "add")
then (eval e1 env al) + (eval e2 env al)
else failwith("Sum not allowed")
```

```
let rec eval exp env (alist:acl) = match exp with
  Eint(n) \rightarrow n
  Den x -> lookup env x
  Let(i,e1,e2) -> let v1 = eval e1 env acl in eval e2 (bind env i v1) acl
  Sum(e1,e2) -> if (emCheck alist "add")
    then (eval e1 env alist) + (eval e2 env alist)
    else failwith("Sum not allowed")
 | Times(e1,e2) -> if (emCheck alist "prod")
   then (eval e1 env alist) * (eval e2 env alist)
   else failwith("times not allowed")
 | Minus(e1,e2) -> if (emCheck alist "sub")
   then (eval e1 env alist) - (eval e2 env alist)
   else failwith("Minus not allowed");;
```

ACL as local Policies within Execution Monitor. #take 1

letEM x = e1 with al in e2

Add a local access policy for the evaluation of e1

FUN+EM

OCAML type for arithmetic expressions with variables and execution monitor

```
type ide = string
type iexp =
    | Eint of int
    | Let of string *iexp * iexp
    | Den of ide
    | Sum of iexp * iexp
    | Times of iexp * iexp
    | Minus of iexp * iexp
    | LetEM of string * iexp * acl * iexp;;
```

Extending access control lists

```
let rec extend al1 al2 =
  match al2 with
  | Empty -> al1
  | AC(aop, als) -> AC(aop, (extend al1 als));;
```

From Operational semantics to OCAML Code

```
\frac{env\ acl[al] \rhd e1 \Rightarrow v1 \ env[x=v1], acl \rhd e2 \Rightarrow v}{env, acl \rhd LetEM\ x=e1\ with\ al\ in\ e2 \Rightarrow v}
```

```
LetEM(i,e1,al,e2) ->
let newal = (extend alist al) in
let v = (eval e1 env newal) in
(eval e2 (bind env i v) alist)
```

let rec eval iexp env (alist:acl) = match iexp with

Execution Monitor take 2

letEM x = e1 with al in e2

Add a local access policy for the evaluation of e2

| Eint(n) -> n

let rec eval iexp env (alist:acl) = match iexp with

```
| Den x -> lookup env x
| Sum(e1,e2) -> if (emCheck alist "add")
then (eval e1 env alist) + (eval e2 env alist)
else failwith("Sum not allowed")
| Times(e1,e2) -> if (emCheck alist "prod")
then (eval e1 env alist) * (eval e2 env alist)
else failwith("times not allowed")
| Minus(e1,e2) -> if (emCheck alist "sub")
then (eval e1 env alist) - (eval e2 env alist)
else failwith("Minus not allowed")
| LetEM(i,e1,al,e2) -> let v = (eval e1 env acl) in
let newal = (extend alist al) in
in (eval e2 (bind env i v) newal)
```

Execution Monitor take 3

letEM x = e1 with al in e2

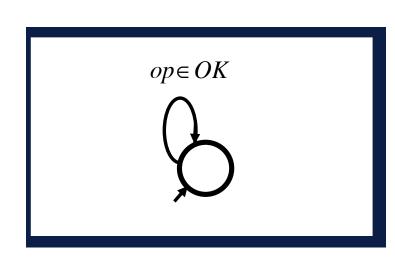
Add a local access policy for the evaluation of e 1 and e2

| Eint(n) -> n

let rec eval iexp env (alist:acl) = match iexp with

Discussion

Simple Acces Control Policies are just Automata



... but EM

- Performs arbitrary computation to decide whether to allow event or halt
 - Can have side effects
 - Can change program flow

Exercize

- 1.Extend the simple language with execution monitor to include functions and function calls
- 2.Extende the language to include a variety of security policies rather than ACL.