

MATH 381 HW 2

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1. Show that $(p \rightarrow q) \vee (p \rightarrow r)$ is logically equivalent to $p \rightarrow (q \vee r)$.

$$\begin{aligned} & (p \rightarrow q) \vee (p \rightarrow r) \\ \equiv & \neg p \vee q \vee \neg p \vee r && [\text{Conditional-disjunction equivalence}] \\ \equiv & \neg p \vee \neg p \vee q \vee r && [\text{Commutative law}] \\ \equiv & \neg p \vee (q \vee r) && [\text{Idempotent law}] \\ \equiv & p \rightarrow (q \vee r) && [\text{Conditional-disjunction equivalence}] \end{aligned}$$

2. Show that $(q \wedge (p \rightarrow \neg q)) \rightarrow \neg p$ is a tautology using propositional equivalence and the laws of logic. Do not use a truth table.

$$\begin{aligned} & (q \wedge (p \rightarrow \neg q)) \rightarrow \neg p \\ \equiv & \neg(q \wedge (p \rightarrow \neg q)) \vee \neg p && [\text{Conditional-disjunction equivalence}] \\ \equiv & (\neg q \vee \neg(p \rightarrow \neg q)) \vee \neg p && [\text{De Morgan's law}] \\ \equiv & (\neg q \vee \neg(\neg p \vee \neg q)) \vee \neg p && [\text{Conditional-disjunction equivalence}] \\ \equiv & (\neg q \vee (p \wedge q)) \vee \neg p && [\text{De Morgan's law}] \\ \equiv & (\neg p \vee \neg q) \vee (\neg p \vee (p \wedge q)) && [\text{Distributive law}] \\ \equiv & \neg p \vee \neg q \vee ((\neg p \vee p) \wedge (\neg p \vee q)) && [\text{Distributive law}] \\ \equiv & \neg p \vee \neg q \vee (T \wedge (\neg p \vee q)) && [\text{Negation law}] \\ \equiv & \neg p \vee \neg q \vee \neg p \vee q && [\text{Identity law}] \\ \equiv & \neg p \vee \neg p \vee \neg q \vee q && [\text{Commutative law}] \\ \equiv & \neg p \vee (\neg q \vee q) && [\text{Idempotent law}] \\ \equiv & \neg p \vee T && [\text{Negation law}] \\ \equiv & T && [\text{Domination law}] \end{aligned}$$

3. Create a compound logical expression composed of at least two propositions (p, q , e.g.) that is a contradiction. Show that your expression really is a contradiction.

$$\begin{aligned}
& (\neg q \wedge \neg p) \wedge (p \vee q) \\
\equiv & (\neg q \wedge (p \vee q)) \wedge (\neg p \wedge (p \vee q)) && [\text{Distributive law}] \\
\equiv & ((\neg q \wedge p) \vee (\neg q \wedge q)) \wedge ((\neg p \wedge p) \vee (\neg p \wedge q)) && [\text{Distributive law}] \\
\equiv & ((\neg q \wedge p) \vee F) \wedge (F \vee (\neg p \wedge q)) && [\text{Negation law}] \\
\equiv & (\neg q \wedge p) \wedge (\neg p \wedge q) && [\text{Identity law}] \\
\equiv & \neg q \wedge (p \wedge \neg p) \wedge q && [\text{Associative law}] \\
\equiv & \neg q \wedge F \wedge q && [\text{Negation law}] \\
\equiv & F && [\text{Domination law}]
\end{aligned}$$

4. Determine whether each of the following propositions is true or false, with explanation.

- (a) $\forall x P(x) \rightarrow \exists x P(x)$, where $P(x)$ is an arbitrary predicate function.

This is necessarily **true** because if $P(x)$ holds for every x in the domain, then there certainly exists an x for which $P(x)$ holds. Only a single solution satisfies the conclusion, and the hypothesis states that every single input satisfies the conclusion.

- (b) $\exists x P(x) \rightarrow \forall x P(x)$, where $P(x)$ is an arbitrary predicate function.

This is not necessarily true and therefore **false** in the general case because the hypothesis only supposes that at least one input exists such that $P(x)$ holds, so it is not certain that every element satisfies the conclusion. While it is true that it is *possible* that every input satisfies the conclusion, the hypothesis only guarantees one.

- (c) $\forall x (2x \geq x)$, where the domain consists of all real numbers.

$$2x \geq x \implies \begin{cases} 2 \geq 1 & x \neq 0 \\ 2(0) \geq 0 & x = 0 \end{cases}$$

The predicate $2x \geq x$ is equivalent to $2 \geq 1$ (which is obviously true) by the division property of equality for $x \in \mathbb{R}$, $x \neq 0$. For

the case that $x = 0$, $2(0) \geq 0 \implies 0 \geq 0$, which is true.

Therefore, this proposition is **true**.

- (d) $\exists n(n^2 < n)$, where the domain consists of all natural numbers (positive integers).

$$n^2 < n \implies n^2 - n < 0$$

$$\text{Let } f(n) = n^2 - n = n(n - 1)$$

$$\implies f(n) = 0 \text{ for } n \in \{0, 1\}$$

$$f(2) = 2^2 - 2 = 4 - 2 = 2 > 0$$

$$\therefore f(n) \geq 0 \text{ for } n \in \mathbb{N} = \mathbb{Z}^+$$

We can rearrange $n^2 < n$ to get the equivalent proposition $n^2 - n < 0$ by the subtraction property of equality. If we declare a function $f(n) = n^2 - n$, we can find the roots to get any possible interval for which $n^2 - n < 0$ and therefore $n^2 < n$, where n would fulfill the proposition. There are only two real roots ($x \in \{0, 1\}$), which is the maximum number for a quadratic function by the fundamental theorem of algebra. Since we are only considering $n \geq 1$ since $n \in \mathbb{N} = \mathbb{Z}^+$, we only need to find whether $f(n)$ is positive or negative after the root $n = 1$, as we know the function won't cross the x -axis again for $x > 1$. Accordingly, $f(2) > 0$, so a natural number does not exist such that the proposition will hold, meaning it is **false**.

- (e) $\exists! x(x^3 = -1)$, where the domain consists of all real numbers.

$$x^3 = -1 \implies x = \sqrt[3]{-1} = -1$$

The quantified proposition can be rearranged algebraically to yield a single solution $x = -1$. We also know that this because the function $f(x) = x^3$ is an injective function i.e. a one-to-one mapping of elements $f : \mathbb{R} \rightarrow \mathbb{R}$. If we look at the graph of f , it passes a “horizontal line test,” where every single number y is only the output that corresponds to a single input x . Therefore, there exists a unique x that fulfills the predicate, so the proposition is **true**.

5. Translate the following English statements to logical expressions, introducing notation as needed for predicates. Then, give a negation of the logical expression as well as a negation of the English statement.

- (a) Some of the students in the class are not here today.

$$\exists s(\neg P(s))$$

The domain of s is the set of students in the class.

$P(s)$: s is present today.

$$\neg[\exists s(\neg P(s))] = \forall s P(s)$$

“Every student in the class is here today.”

- (b) The number \sqrt{x} is rational if x is an integer.

$$\forall x(\sqrt{x} \in \mathbb{Q}) \qquad x \in \mathbb{Z}$$

$$\neg[\forall x(\sqrt{x} \in \mathbb{Q})] = \exists x(\sqrt{x} \notin \mathbb{Q}) \qquad x \in \mathbb{Z}$$

“There is an integer x such that \sqrt{x} is irrational”