Network Programming

- Network Programming as Programming across
 Machine Boundaries
- · The Sockets API
- Reliable Communication Channels: TCP
- Dangerous at any Speed; connectionless communication and UDP
- · Server Design
- Reading: R&R, Ch 18

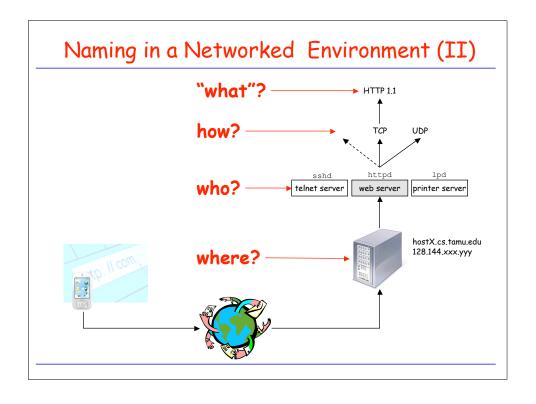
Naming in a Networked Environment

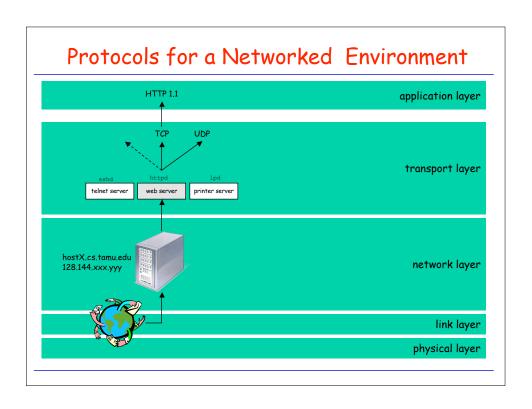
Reminder:

- · Naming within a single address space?
 - addresses, duh!
- · Naming across address spaces?
 - file descriptors
 - filenames (use file system as name space)
 - keys

New:

· Naming across machine boundaries?





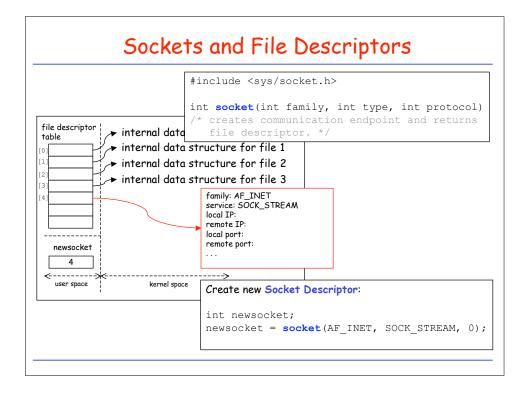
The Socket API

- What does an API need to support? [Comer&Stevens]
 - allocate local resources (buffers)
 - specify local and remote communication endpoints
 - initiate a connection
 - wait for an incoming connection
 - send or receive data
 - determine when data arrives
 - generate urgent data
 - handle incoming urgent data
 - terminate a connection gracefully
 - handle connection termination from the remote site
 - abort communication
 - handle error conditions or a connection abort
 - release local resources when communication finishes

- Existing TCP/IP APIs:
 - Berkeley Socket API
 - aka socket API, socket interface, sockets
 - adapted by Linux and others
 - Windows Sockets
 - System V Unix TLI (Transport Layer Interface)

Specifying a Protocol Interface

- Reality Check: "All" networks today are based on TCP/IP.
- · How to define a network API, then?!
 - Approach 1: Define functions that specifically support TCP/IP communication.
 - e.g. makeTCPconnection(int32 host, int16 portno);
 - Approach 2: Define functions that support network communication in general, and use parameters to handle TCP/IP as a special case.
- The socket API provides generalized functions that support network communication using many possible protocols.
- The programmer specifies the type of service required rather than the name of a specific network protocol.

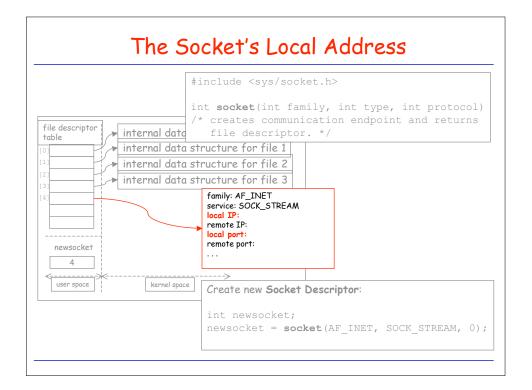


TCP? UDP?!!

- · Connection-oriented Style (SOCK_STREAM) of communication
- Implemented in TCP/IP as the Transport Control Protocol (TCP).
- TCP provides full reliability:
 - verifies that data arrives, with automatic retransmission
 - computes checksums to detect corruption
 - uses sequence numbers to guarantee ordering of received packets
 - automatically eliminates duplicated packets
 - provides flow control (ensures that sender does not send more packets than the receiver can handle)
 - informs sender if network becomes inoperable for any reason
- TCP protects network resources:
 - provides congestion control (throttles transmission when it detects that network is congested)
- · What is the cost of all this?! Connection establishment overhead.

TCP? UDP?!!

- Connectionless Style (SOCK_DGRAM) of communication
- Implemented in TCP/IP as the User Datagram Protocol (UDP).
- UDP provides no guarantee about reliable delivery:
 - packets can be lost, duplicated, delayed, or delivered out of order
- UDP works well if the underlying network works well, e.g. local network
- In practice, programmers use UDP only when:
 - The application requires that UDP must be used. (The application has been designed to handle reliability and delivery errors.)
 - 2. The application relies on hardware broadcast or multicast for delivery.
 - 3. Application runs in reliable, local environment, and overhead for reliability is unnecessary.



Defining the Socket's Local Address (server)

Host vs. Network Byte Order

- · Big-endian vs. Little-endian.
- Network representation requires big-endian.
- Portability?!

```
#include <arpa/inet.h>

uint32_t htonl(uint32_t hostlong);
uint32_t ntohl(uint32_t netlong);
uint16_t htons(uint16_t hostshort);
uint16_t ntohs(uint16_t netshort);
struct sockaddr_in server;

int sock = socket(AF_INET, SOCK_STREAM, 0);
server.sin_family = AF_INET;
server.sin_addr.s_addr = htonl(INADDR_ANY);
server.sin_port = htons((short)8652);

bind(sock, &server, sizeof(server));
```

Prepare to Accept Incoming Connections (server)

```
#include <sys/socket.h>
int listen(int socket, int backlog);
/* specify willingness to accept incoming
   connections at given socket, with given
   queue limit.
   Connections are then accepted with
   accept(). */
```

When request for connection from client comes in, both client and server execute hand-shake procedure to set up the connection.

When server is busy, two things can happen

Connection request is queued (as long as backlog not exceeded)

or

 Connection request is refused (client receives ECONNREFUSED error)

Handle Incoming Connections (server)

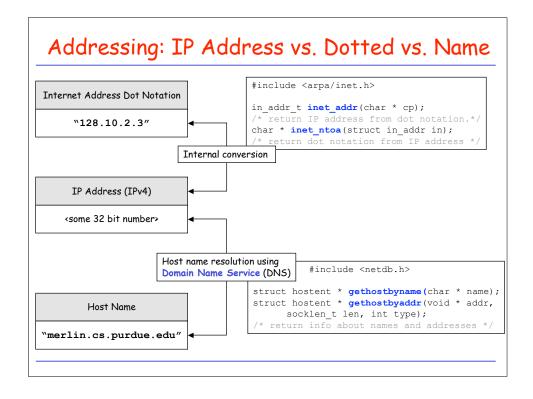
- Extract first connection request on queue of pending connections.
- · Create new socket with same properties of given socket.
- Returns new file descriptor for the socket.
- Blocks caller if no pending connections are present in queue.
- New socket may not be used to accept more connections.
- · Original socket socket remains open.
- Argument address (result parameter) is filled in with the address of connecting entity.
- Argument address_len (value-result parameter) initially contains length of space pointed to by address. On return it contains length of actual length of space.

In the meantime, at the Client's End (client)

Note difference between SOCK_DGRAM and SOCK_STREAM sockets:

- For SOCK_STREAM sockets, connect attempts to establish connection with other socket.
 - · Generally, stream sockets may connect only once.
- For SOCK_DGRAM sockets, connect specifies the peer with which to associate socket.
 - Datagram sockets may dissolve association by connecting to invalid address, such as null address.
- · For address, fill in family, address, and port number.

client socket() bind() connect() connect() read() / write() close() client socket() bind() connect() read() / write() close()



Example TCP Client: DAYTIME client [Comer]

```
#define LINELEN 128

/* forward */
int connectTCP(const char * host, const char * service);

/* main program */
int main(argc, char * argv) {

   char * host = "localhost";/* use local host if none supplied */
   char * service = "daytime"; /* default service port */

   if (argc > 1) host = argv[1];
   if (argc > 2) service = argv[2];

   int s = connectTCP(host, service);

while ( (int n = read(s, buf, LINELEN)) > 0) {
    buf[n] = "\0"; /* ensure null terminated */
        (void) fputs(buf, stdout);
   }
}
```

Example TCP Client: (cont...)

```
#define LINELEN 128
             int connectTCP(const char * host , const char * service) {
/* forward
               struct sockaddr_in sin; /* Internet endpoint address */
               memset(&sin, 0, sizeof(sin));
               sin.sin_family = AF_INET;
/* main pr
               /* Map service name to port number */
int main(a
               if (struct servent * pse = getservbyname(service, "tcp") )
                   sin.sin_port = pse->s_port;
  char * h
               else if ((sin.sin_port = htons((unsigned short)atoi(service))) == 0)
                   \verb|errexit("can't get < %s> service entry \n", service);|\\
               /* Map host name to IP address, allowing for dotted decimal */
  if (argc
               if (struct hostent * phe = gethostbyname(host) )
  if (argc
                   {\tt memcpy\,(\&sin.sin\_addr,\ phe->h\_addr,\ phe->h\_length)\;;}
               else if ( (sin.sin addr.s addr = inet addr(host)) == INADDR NONE )
  int s =
                   errexit("can't get <%s> host entry\n", host);
               /* Allocate socket */
  while (
               int s = socket(AF_INET, SOCK_STREAM, 0);
    buf[n]
               if (s < 0) errexit("can't create socket: %s\n", strerror(errno));</pre>
    (void)
  }
               if (connect(s, (struct sockaddr *)&sin, sizeof(sin)) < 0)</pre>
                   errexit("can't connect to %s.%s: %s\n", host, service, strerror(errno));
               return s;
```

Server Software Design: Issues [Comer]

· Concurrent vs. Iterative Servers:

The term concurrent server refers to whether the server permits multiple requests to proceed concurrently, not to whether the underlying implementation uses multiple, concurrent threads of execution.

Iterative server implementations are easier to build and understand, but may result in poor performance because they make clients wait for service.

Connection-Oriented vs. Connectionless Access:

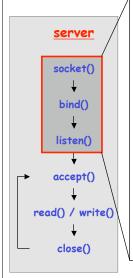
Connection-oriented (TCP, typically) servers are easier to implement, but have resources bound to connections.

Reliable communication over UDP is not easy!

· Stateful vs. Stateless Servers:

How much information should the server maintain about clients? (What if clients crash, and server does not know?)

Example: Iterative, Connection-Oriented Server



```
int passiveTCPsock(const char * service, int backlog) {
  struct sockaddr_in sin;
                                      /* Internet endpoint address */
                                     /* Zero out address */
 memset(&sin, 0, sizeof(sin));
sin.sin_family = AF_INET;
sin.sin_addr.s_addr = INADDR_ANY;
    Map service name to port number */
  if (struct servent * pse = getservbyname(service, "tcp") )
      sin.sin_port = pse->s_port;
  else if ((sin.sin_port = htons((unsigned short)atoi(service))) == 0)
      errexit("can't get <%s> service entry\n", service);
  /* Allocate socket */
 int s = socket(AF INET, SOCK STREAM, 0);
 if (s < 0) errexit("can't create socket: %s\n", strerror(errno));</pre>
  /* Bind the socket */
 if (bind(s, (struct sockaddr *)&sin, sizeof(sin)) < 0)</pre>
       errexit("can't bind to ...\n");
  /* Listen on socket */
 if (listen(s, backlog) < 0)
       errexit("can't listen on ...\n")
 return s:
```

