Propositional Logic Syntactic Sugar

$\varphi \Leftrightarrow \psi := (\neg\varphi \vee \psi) \wedge (\neg\psi \vee \varphi)$ $\varphi \to \psi := \neg\varphi \vee \psi$
$\varphi \oplus \psi := (\varphi \wedge \neg\psi) \vee (\psi \wedge \neg\varphi)$ $\varphi \bar{\wedge} \psi := \neg(\varphi \wedge \psi)$
$(\alpha \Rightarrow \beta | \gamma) := (\neg\alpha \vee \beta) \wedge (\alpha \vee \gamma)$ $\varphi \bar{\vee} \psi := \neg(\varphi \vee \psi)$

Distributivity: $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$
$a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$
De Morgan: $\neg(a \vee b) \equiv (\neg a \wedge \neg b)$
$\neg(a \wedge b) \equiv (\neg a \vee \neg b)$
CNF: from truth table, take minterms that are 0. Each minterm is built as an OR of the negated variables. E.g., $(0,0,1) \to (x \vee y \vee \neg z)$.
###**SAT SOLVERS**
Satisfiability, Validity and Equivalence

$\mathrm{SAT}(\varphi) := \neg\mathrm{VALID}(\neg\varphi)$ $\varphi \Leftrightarrow \psi := \mathrm{VALID}(\varphi \leftrightarrow \psi)$
$\mathrm{VALID}(\varphi) := (\varphi \Leftrightarrow 1)$ $\mathrm{SAT}(\varphi) := \neg(\varphi \Leftrightarrow 0)$.

Sequent Calculus:
-*Validity*: start with $\{\} \vdash \phi$; valid iff $\Gamma \cap \Delta \neq \{\}$ FOR ALL leaves.
-*Satisfiability*: start with $\{\phi\} \vdash \{\}$; satisfiable iff $\Gamma \cap \Delta = \{\}$ for AT LEAST ONE leaf.
-Counterexample/sat variable assignment: var is true, if $x \in \Gamma$; false otherwise; "don't care", if variable doesn't appear.

OPER.	LEFT	RIGHT
NOT	$\frac{\neg\phi,\Gamma\vdash\Delta}{\Gamma\vdash\phi,\Delta}$	$\frac{\Gamma\vdash\neg\phi,\Delta}{\phi,\Gamma\vdash\Delta}$
AND	$\frac{\phi\wedge\psi,\Gamma\vdash\Delta}{\phi,\psi,\Gamma\vdash\Delta}$	$\frac{\Gamma\vdash\phi\wedge\psi,\Delta}{\Gamma\vdash\phi,\Delta \qquad \Gamma\vdash\psi,\Delta}$
OR	$\frac{\phi\vee\psi,\Gamma\vdash\Delta}{\phi,\Gamma\vdash\Delta \qquad \psi,\Gamma\vdash\Delta}$	$\frac{\Gamma\vdash\phi\vee\psi,\Delta}{\Gamma\vdash\phi,\psi,\Delta}$

Resolution Calculus $\frac{\{\neg x\}\cup C_1 \qquad \{x\}\cup C_2}{C_1 \cup C_2}$

To prove unsatisfiability of given clauses in CNF: If we reach $\{\}$, the formula is unsatisfiable. E.g., $\{\{a\},\{\neg a, b\},\{\neg b\}\}$, we get: $\{a\}+\{\neg a, b\} \to \{b\}; \{b\}+\{\neg b\} \to \{\}$ (unsatisfiable). To prove validity, prove UNSAT of negated formula.

Davis Putnam Procedure - proves SAT; To prove validity: prove unsatisfiability of negated formula.
(1) Compute Linear Clause Form
(Don't forget to create the last clause $\{x_n\}$)
(2) Last variable has to be $\underline{1}$ (true) $\to$ find implied variables.
(3) For remaining variables: assume values and compute newly implied variables.
(4) If contradiction reached: backtrack.

Linear Clause Forms (Computes CNF) - Bottom up (inside out) in the syntax tree: convert "operators and variables" into new variable. E.g., $\neg a \vee b$ becomes $x_1 \leftrightarrow \neg a; x_2 \leftrightarrow x_1 \vee b$. Use rules below to find CNF. Create last clause {Xn}

$x \leftrightarrow \neg y \Leftrightarrow (\neg x \vee \neg y) \wedge (x \vee y)$
$x \leftrightarrow y_1 \wedge y_2 \Leftrightarrow$
$(\neg x \vee y_1) \wedge (\neg x \vee y_2) \wedge (x \vee \neg y_1 \vee \neg y_2)$
$x \leftrightarrow y_1 \vee y_2 \Leftrightarrow$
$(\neg x \vee y_1 \vee y_2) \wedge (x \vee \neg y_1) \wedge (x \vee \neg y_2)$
$x \leftrightarrow y_1 \to y_2 \Leftrightarrow$
$(x \vee y_1) \wedge (x \vee \neg y_1 \vee \neg y_2) \wedge (\neg x \vee \neg y_1 \vee y_2)$
$x \leftrightarrow (y_1 \leftrightarrow y_2) \Leftrightarrow (x \vee y_1 \vee y_2) \wedge (x \vee \neg y_1 \vee \neg y_2) \wedge$
$(\neg x \vee y_1 \vee \neg y_2) \wedge (\neg x \vee \neg y_1 \vee y_2)$
$x \leftrightarrow y_1 \oplus y_2 \Leftrightarrow (x \vee \neg y_1 \vee y_2) \wedge (x \vee y_1 \vee \neg y_2) \wedge$
$(\neg x \vee y_1 \vee y_2) \wedge (\neg x \vee \neg y_1 \vee \neg y_2)$

```
Compose(int x, BddNode ψ, α)
 int m; BddNode h, l;
 if x>label(ψ) then
  return ψ;
 elseif x=label(ψ) then
  return ITE(α,high(ψ),
        low(ψ));
 else
  m=max{label(ψ),label(α)}
  (α0,α1):=Ops(α, m);
  (ψ0,ψ1):=Ops(ψ, m);
  h:=Compose(x,ψ1,α1);
  l:=Compose(x,ψ0,α0);
  return CreateNode(m,h,l)
 endif;end
```

```
ITE(BddNode i, j, k)
 int m; BddNode h, l;
 if i = 0 then return k
 elseif i=1 then
  return j
 elseif j=k then
  return k
 else
  m = max{label(i),
      label(j),label(k)}
  (i0,i1):=Ops(i,m);
  (j0,j1):=Ops(j,m);
  (k0,k1):=Ops(k,m);
  l:=ITE(i0,j0,k0);
  h:=ITE(i1, j1, k1);
  return CreateNode(m,h,l)
 end;end
```

```
Constrain(Φ, β)
 if β=0 then
  ret 0
 elseif Φ ∈ {0,1}(β = 1)
     ret Φ
 else
  m=max{label(β),label(Φ)}
    (Φ0,Φ1):=Ops(Φ,m);
    (β0,β1):=Ops(β,m);
    if β0=0
  ret Constrain(Φ1,β1)
 elseif β1=0 then
  ret Constrain(Φ0,β0)
 else
  l:=Constrain(Φ0,β0);
  h:=Constrain(Φ1,β1);
  ret CreateNode(m,h,l)
endif;endif;end
```

```
Apply(⊙, Bddnode a, b)
 int m; BddNode h, l;
 if isLeaf(a)&isLeaf(b)
      then
  return Eval(⊙,label(a),
       label(b));
 else
  m=max{label(a),label(b)}
  (a0,a1):=Ops(a,m);
  (b0,b1):=Ops(b,m);
  h:=Apply(⊙,a1,b1);
  l:=Apply(⊙,a0,b0);
  return CreateNode(m,h,l)
 end;
end
```

```
Restrict(Φ,β)
 if β=0
  return 0
 elseif
      Φ ∈ {0,1} ∨ (β = 1)

  return Φ
 else
  m=max{label(β),label(Φ)}
  (Φ0,Φ1):=Ops(Φ,m);
  (β0,β1):=Ops(β,m)
 if β0=0
  return Restrict(Φ1,β1)
 elseif β1=0
  return Restrict(Φ0,β0)
 elseif m=label(Φ)
  return CreateNode(m,
   Restrict (Φ1,β1),
   Restrict (Φ0,β0))
 else
  return Restrict(Φ,
  Apply(∨,β0,β1))
 endif;endif;end
--------------------------
Ops(v,m)
 x:=label(v);
 if m=degree(x)
  return (low(v),high(v))
 else return(v, v)
end;end
```

```
Exists(BddNode e, φ)
 if isLeaf(φ)∨isLeaf(e)
  return φ;
 elseif label(e)>label(φ)
  return Exist(high(e),φ)
 elseif label(e)=label(φ)
  h=Exist(high(e),high(φ)
  l=Exist(high(e),low(φ))
  return Apply(∨,l,h)
 else (label(e)<label(φ))
  h:=Exists(e,high(φ))
  l:=Exists(e,low(φ))
  return CreateNode(label(
      φ),h,l)
 endif; end function.
--------------------------
ZDD: If positive cofactor
     = 0, redirect edge
     to negative
     cofactor.
If variable not in the
     formula, add with
     both edges pointing
     to same node.
--------------------------
FDD: Positive Davio
     Decomposition. (
     Keep both edges to
     1 if happens!)
```
$\varphi = [\varphi]_x^0 \oplus x \wedge (\partial\varphi/\partial x)$
$(\partial\varphi/\partial x) := [\varphi]_x^0 \oplus [\varphi]_x^1$

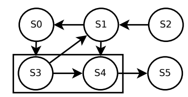
Local Model Checking

$\frac{s\vdash_\Phi \varphi\wedge\psi}{\{s\vdash_\Phi\varphi\} \quad \{s\vdash_\Phi\psi\}}\wedge$	$\frac{s\vdash_\Phi \varphi\vee\psi}{\{s\vdash_\Phi\varphi\} \quad \{s\vdash_\Phi\psi\}}\vee$
$\frac{s\vdash_\Phi \Box\varphi}{\{s_1\vdash_\Phi\varphi\}\ldots\{s_n\vdash_\Phi\varphi\}}\wedge$	$\frac{s\vdash_\Phi \Diamond\varphi}{\{s_1\vdash_\Phi\varphi\}\ldots\{s_n\vdash_\Phi\varphi\}}\vee$
$\frac{s\vdash_\Phi \overleftarrow{\Box}\varphi}{\{s_1'\vdash_\Phi\varphi\}\ldots\{s_n'\vdash_\Phi\varphi\}}\wedge$	$\frac{s\vdash_\Phi \overleftarrow{\Diamond}\varphi}{\{s_1'\vdash_\Phi\varphi\}\ldots\{s_n'\vdash_\Phi\varphi\}}\vee$
$\frac{s\vdash_\Phi \mu x.\varphi}{s\vdash_\Phi\varphi}$ $\quad$ $\frac{s\vdash_\Phi \nu x.\varphi}{s\vdash_\Phi\varphi}$	$\frac{s\vdash_\Phi x}{s\vdash_\Phi \mathfrak{D}_\Phi(x)}$ $\quad$ $\mathfrak{D}_\Phi$(replace w. initial form.)
$\{s_1 \ldots s_n\} = suc_\exists^\mathcal{R}(s)$ and $\{s_1' \ldots s_n'\} = pre_\exists^\mathcal{R}(s)$	

Approximations and Ranks

If (s,$\mu x.\varphi$) repeats$\to$return 1	$apx_0(\mu x.\varphi) := 0$
If (s,$\nu x.\varphi$) repeats$\to$return 0	$apx_0(\nu x.\varphi) := 1$

Tarski-Knaster Theorem: $\mu :=$ starts $\perp \to$ least fixpoint ♠ $\nu :=$ starts $\top \to$ greatest fixpoint
Quantif. $\exists x.\varphi := [\varphi]_x^1 \vee [\varphi]_x^0$ ♣ $\forall x.\varphi := [\varphi]_x^1 \wedge [\varphi]_x^0$
Predecessor and Successor
$\Diamond := pre_\exists^\mathcal{R}(Q) := \exists x_1', \ldots, x_n'.\varphi_\mathcal{R} \wedge [\varphi_Q]_{x_1,\ldots,x_n}^{x_1',\ldots,x_n'}$
$\overleftarrow{\Diamond} := suc_\exists^\mathcal{R}(Q) := [\exists x_1, \ldots, x_n.\varphi_\mathcal{R} \wedge \varphi_Q]_{x_1',\ldots,x_n'}^{x_1,\ldots,x_n}$
$\Box = pre_\forall^\mathcal{R}(Q) := \forall x_1', \ldots, x_n'.\varphi_\mathcal{R} \to [\varphi_Q]_{x_1,\ldots,x_n}^{x_1',\ldots,x_n'}$
$\overleftarrow{\Box} := suc_\forall^\mathcal{R}(Q) := [\forall x_1, \ldots, x_n.\varphi_\mathcal{R} \to \varphi_Q]_{x_1',\ldots,x_n'}^{x_1,\ldots,x_n}$



Example: $\Box/\overleftarrow{\Box}$
$pre_\forall^\mathcal{R}(\{S3, S4\}) = \{S0, S5\}$
$suc_\forall^\mathcal{R}(\{S3, S4\}) = \{S2, S5\}$

```
pre∀R(Q = {S1,...,Sn})
for each node n in K:
 if(n points to a node
       that is not in Q)
  n ∉ pre∀R(Q)
 else
  n ∈ pre∀R(Q)
```

```
suc∀R(Q = {S1,...,Sn})
for each node n in K:
 if(n is pointed by a node
       that is not in Q)
  n ∉ suc∀R(Q)
 else
  n ∈ suc∀R(Q)
```

###**AUTOMATA**
Automata types: G$\to$Safety; F$\to$Liveness; FG$\to$Persistence/Co-Buchi; GF$\to$Fairness/Buchi.
Automaton Determinization
$\mathbf{NDet_G \to Det_G}$: 1.Remove all states/edges that do not satisfy acceptance condition; 2.Use Subset construction (Rabin-Scott); 3.Acceptance condition will be the states where {} is never reached.
$\mathbf{\{NDet_F (partial)\ or\ NDet_{prefix}\} \to Det_{FG}}$: Breakpoint Construction.
$\mathbf{NDet_F}$ **(total)**$\mathbf{\to Det_F}$: Subset Construction.
$\mathbf{NDet_{FG} \to Det_{FG}}$: Breakpoint Construction.
$\mathbf{NDet_{GF} \to \{Det_{Rabin}\ or\ Det_{Streett}\}}$: Safra Algorithm.
Boolean Operations on ω-Automata
Complement
$$\neg\mathcal{A}_\forall(Q,\mathcal{I},\mathcal{R},\mathcal{F}) = \mathcal{A}_\exists(Q,\mathcal{I},\mathcal{R},\neg\mathcal{F})$$
$$\neg\mathcal{A}_\exists(Q,\mathcal{I},\mathcal{R},\mathcal{F}) = \mathcal{A}_\forall(Q,\mathcal{I},\mathcal{R},\neg\mathcal{F})$$
Conjunction
$$(\mathcal{A}_\exists(Q_1,\mathcal{I}_1,\mathcal{R}_1,\mathcal{F}_1) \wedge \mathcal{A}_\exists(Q_2,\mathcal{I}_2,\mathcal{R}_2,\mathcal{F}_2)) = \mathcal{A}_\exists(Q_1 \cup Q_2, \mathcal{I}_1 \wedge \mathcal{I}_2, \mathcal{R}_1 \wedge \mathcal{R}_2, \mathcal{F}_1 \wedge \mathcal{F}_2)$$
Disjunction
$$(\mathcal{A}_\exists(Q_1,\mathcal{I}_1,\mathcal{R}_1,\mathcal{F}_1) \vee \mathcal{A}_\exists(Q_2,\mathcal{I}_2,\mathcal{R}_2,\mathcal{F}_2)) =$$
$$\mathcal{A}_\exists\left(\begin{array}{c} Q_1 \cup Q_2 \cup \{q\}, \\ (\neg q \wedge \mathcal{I}_1) \vee (q \wedge \mathcal{I}_2), \\ (\neg q \wedge \mathcal{R}_1 \wedge \neg q') \vee (q \wedge \mathcal{R}_2 \wedge q'), \\ (\neg q \wedge \mathcal{F}_1) \vee (q \wedge \mathcal{F}_2) \end{array}\right)$$
If both automata are totally defined,
$$(\mathcal{A}_\exists(Q_1,\mathcal{I}_1,\mathcal{R}_1,\mathcal{F}_1) \vee \mathcal{A}_\exists(Q_2,\mathcal{I}_2,\mathcal{R}_2,\mathcal{F}_2)) = \mathcal{A}_\exists(Q_1 \cup Q_2, \mathcal{I}_1 \wedge \mathcal{I}_2, \mathcal{R}_1 \wedge \mathcal{R}_2, \mathcal{F}_1 \vee \mathcal{F}_2)$$
Eliminate Nesting - Acceptance condition **must** be an automata of the same type
$$\mathcal{A}_\exists(Q^1,\mathcal{I}_1^1,\mathcal{R}_1^1,\mathcal{A}_\exists(Q^2,\mathcal{I}_1^2,\mathcal{R}_1^2,\mathcal{F}_1))$$
$$= \mathcal{A}_\exists(Q^1 \cup Q^2, \mathcal{I}_1^1 \wedge \mathcal{I}_1^2, \mathcal{R}_1^1 \wedge \mathcal{R}_1^2, \mathcal{F}_1))$$
Boolean Operations of G
(1)$\neg G\varphi = F\neg\varphi$ (2)$G\varphi \wedge G\psi = G[\varphi \wedge \psi]$
(3)$G\varphi \vee G\psi = \mathcal{A}_\exists(\{p,q\}, p \wedge q, [p' \leftrightarrow p \wedge \varphi] \wedge [q' \leftrightarrow q \wedge \psi], G[p \vee q])$
Boolean Operations of F
(1)$\neg F\varphi = G\neg\varphi$ (2)$F\varphi \vee F\psi = F[\varphi \vee \psi]$
(3)$F\varphi \wedge F\psi = \mathcal{A}_\exists(\{p,q\}, \neg p \wedge \neg q, [p' \leftrightarrow p \vee \varphi] \wedge [q' \leftrightarrow q \vee \psi], F[p \wedge q])$
Boolean Operations of FG
(1)$\neg FG\varphi = GF\neg\varphi$ (2)$FG\varphi \wedge FG\psi = FG[\varphi \wedge \psi]$
(3)$FG\varphi \vee FG\psi = \mathcal{A}_\exists(\{q\}, \neg q, q' \leftrightarrow (q \Rightarrow \psi | \neg\varphi), FG[\neg q \vee \psi])$
Boolean Operations of GF
(1)$\neg GF\varphi = FG\neg\varphi$ (2)$GF\varphi \vee GF\psi = GF[\varphi \vee \psi]$
(3)$GF\varphi \wedge GF\psi = \mathcal{A}_\exists(\{q\}, \neg q, q' \leftrightarrow (q \Rightarrow \neg\psi | \varphi), GF[q \wedge \psi])$

Transformation of Acceptance Conditions
Reduction of G
$G\varphi = \mathcal{A}_\exists(\{q\}, q, \varphi \wedge q \wedge q', Fq))$
$G\varphi = \mathcal{A}_\exists(\{q\}, q, q' \leftrightarrow q \wedge \varphi, FGq)$
$G\varphi = \mathcal{A}_\exists(\{q\}, q, q' \leftrightarrow q \wedge \varphi, GFq)$
Reduction of F
$F\varphi$ can **not** be expressed by $NDet_G$
$F\varphi = \mathcal{A}_\exists(\{q\}, \neg q, q' \leftrightarrow q \vee \varphi, FGq)$
$F\varphi = \mathcal{A}_\exists(\{q\}, \neg q, q' \leftrightarrow q \vee \varphi, GFq)$
Reduction of FG
$FG\varphi$ can **not** be expressed by $NDet_G$
$FG\varphi = \mathcal{A}_\exists(\{q\}, \neg q, q \to \varphi \wedge q', Fq)$
$$FG\varphi = \mathcal{A}_\exists\left(\begin{array}{c} \{p,q\}, \qquad \neg p \wedge \neg q, \\ \left[\begin{array}{c} (p \to p') \wedge (p' \to p \vee \neg q) \wedge \\ (q' \leftrightarrow (p \wedge \neg q \vee \neg \varphi) \vee (p \wedge q)) \end{array}\right], \\ G\neg q \wedge Fp \end{array}\right)$$
$$FG\varphi = \mathcal{A}_\exists\left(\begin{array}{c} \{p,q\}, \qquad \neg p \wedge \neg q, \\ \left[\begin{array}{c} (p \to p') \wedge (p' \to p \vee \neg q) \wedge \\ (q' \leftrightarrow (p \wedge \neg q \vee \neg \varphi) \vee (p \wedge q)) \end{array}\right], \\ GF[p \wedge \neg q] \end{array}\right)$$

###**TEMPORAL LOGICS**
(S1)Pure LTL: AFGa
(S2)LTL + CTL: AFa
(S3)Pure CTL: AGEFa
(S4)CTL*: AFGa $\vee$ AGEFa
Remarks *Beware of Finite Paths*
E and A quantify over infinite paths.
Aφ holds on every state that has no infinite path;
Eφ is false on every state that has no infinite path;
A0 holds on states with only finite paths;
E1 is false on state with only finite paths;
$\Box$0 holds on states with no successor states;
$\Diamond$1 holds on states with successor states.
$F\varphi = \varphi \vee XF\varphi$ $G\varphi = \varphi \wedge XG\varphi$
$[\varphi\ U\ \psi] = \psi \vee (\varphi \wedge X[\varphi\ U\ \psi])$
$[\varphi\ B\ \psi] = \neg\psi \wedge (\varphi \vee X[\varphi\ B\ \psi])$
$[\varphi\ W\ \psi] = (\psi \wedge \varphi) \vee (\neg\psi \wedge X[\varphi\ W\ \psi])$

LTL Syntactic Sugar: analog for past operators
$G\varphi = \neg[1\ \underline{U}\ (\neg\varphi)]$ $F\varphi = [1\ \underline{U}\ \varphi]$
$[\varphi\ W\ \psi] = \neg[(\neg\varphi \vee \neg\psi)\ \underline{U}\ (\neg\varphi \wedge \psi)]$
$[\varphi\ \underline{W}\ \psi] = [(\neg\psi)\ \underline{U}\ (\varphi \wedge \psi)]$ *($\neg\psi$ holds until $\varphi \wedge \psi$)*
$[\varphi\ B\ \psi] = \neg[(\neg\varphi)\ \underline{U}\ \psi)]$
$[\varphi\ \underline{B}\ \psi] = [(\neg\psi)\ \underline{U}\ (\varphi \wedge \neg\psi)]$*($\psi$ can't hold when φ holds)*
$[\varphi\ U\ \psi] = \neg[(\neg\psi)\ \underline{U}\ (\neg\varphi \wedge \neg\psi)]$
$[\varphi\ U\ \psi] = [\varphi\ \underline{U}\ \psi] \vee G\varphi$
$[\varphi\ \underline{U}\ \psi] = \neg[(\neg\psi)\ U\ (\neg\varphi \wedge \neg\psi)]$
$[\varphi\ \underline{U}\ \psi] = \neg[(\neg\psi)\ W\ (\varphi \to \psi)]$
$[\varphi\ \underline{U}\ \psi] = [\psi\ \underline{W}\ (\varphi \to \psi)]$
$[\varphi\ \underline{U}\ \psi] = \neg[(\neg\varphi)\ B\ \psi]$*($\varphi$ doesn't matter when ψ holds)*
$[\varphi\ \underline{U}\ \psi] = [\psi\ \underline{B}\ (\neg\varphi \wedge \neg\psi]$
CTL Syntactic Sugar: analog for past operators
Existential Operators
$EF\varphi = E[1\ \underline{U}\ \varphi]$
$EG\varphi = E[\varphi\ U\ 0]$
$E[\varphi\ U\ \psi] = E[\varphi\ \underline{U}\ \psi] \vee EG\varphi$
$E[\varphi\ B\ \psi] = E[(\neg\psi)\ \underline{U}\ (\varphi \wedge \neg\psi)] \vee EG\neg\psi$
$E[\varphi\ B\ \psi] = E[(\neg\psi)\ U\ (\varphi \wedge \neg\psi)]$
$E[\varphi\ \underline{B}\ \psi] = E[(\neg\psi)\ \underline{U}\ (\varphi \wedge \neg\psi)]$
$E[\varphi\ \underline{B}\ \psi] = E[(\neg\psi\ \underline{U}\ (\varphi \wedge \neg\psi)]$
$E[\varphi\ W\ \psi] = E[(\neg\psi)\ \underline{U}\ (\varphi \wedge \psi)] \vee EG\neg\psi$
$E[\varphi\ W\ \psi] = E[(\neg\psi)\ U\ (\varphi \wedge \psi)]$
$E[\varphi\ \underline{W}\ \psi] = E[(\neg\psi)\ \underline{U}\ (\varphi \wedge \psi)]$
Universal Operators
$AX\varphi = \neg EX\neg\varphi$
$AG\varphi = \neg E[1\ \underline{U}\ \neg\varphi]$

$AF\varphi = \neg EG\neg\varphi$
$AF\varphi = \neg E[(\neg\varphi)\ U\ 0]$
$A[\varphi\ U\ \psi] = \neg E[(\neg\psi)\ \underline{U}\ (\neg\varphi\wedge\neg\psi)]$
$A[\varphi\ \underline{U}\ \psi] = \neg E[(\neg\psi)\ \underline{U}\ (\neg\varphi\wedge\neg\psi)]\wedge\neg EG\neg\psi$
$A[\varphi\ \underline{U}\ \psi] = \neg E[(\neg\psi)\ U\ (\neg\varphi\wedge\neg\psi)]$
$A[\varphi\ B\ \psi] = \neg E[(\neg\varphi)\ \underline{U}\ \psi]$
$A[\varphi\ \underline{B}\ \psi] = \neg E[(\neg\varphi)\ U\ \psi]$
$A[\varphi\ \underline{B}\ \psi] = \neg E[(\neg\varphi\vee\psi)\ \underline{U}\ \psi]\wedge\neg EG(\neg\varphi\vee\psi)$
$A[\varphi\ W\ \psi] = \neg E[(\neg\psi)\ \underline{U}\ (\neg\varphi\wedge\psi)]$
$A[\varphi\ \underline{W}\ \psi] = \neg E[(\neg\psi)\ \underline{U}\ (\neg\varphi\wedge\psi)]\wedge\neg EG\neg\psi$
$A[\varphi\ \underline{W}\ \psi] = \neg E[(\neg\psi)\ U\ (\neg\varphi\wedge\psi)]$

CTL* to CTL - Existential Operators
$EX\varphi = EXE\varphi$
$EF\varphi = EFE\varphi$ $\qquad$ $EFG\varphi \equiv EFEG\varphi$
$E[\varphi\ W\ \psi] = E[(E\varphi)\ W\ \psi]$
$E[\varphi\ \underline{W}\ \psi] = E[(E\varphi)\ \underline{W}\ \psi]$
$E[\psi\ U\ \varphi] = E[\psi\ U\ E(\varphi)]$
$E[\psi\ \underline{U}\ \varphi] = E[\psi\ \underline{U}\ E(\varphi)]$
$E[\varphi\ B\ \psi] = E[(E\varphi)\ B\ \psi]$
$E[\varphi\ \underline{B}\ \psi] = E[(E\varphi)\ \underline{B}\ \psi]$
obs. $EGF\varphi \neq EGEF\varphi \rightarrow$ can't be converted

CTL* to CTL - Universal Operators
$AX\varphi = AXA\varphi$
$AG\varphi = AGA\varphi$
$A[\varphi\ W\ \psi] = A[(A\varphi)\ W\ \psi]$
$A[\varphi\ \underline{W}\ \psi] = A[(A\varphi)\ \underline{W}\ \psi]$
$A[\varphi\ U\ \psi] = A[A(\varphi)\ U\ \psi]$
$A[\varphi\ \underline{U}\ \psi] = A[A(\varphi)\ \underline{U}\ \psi]$
$A[\psi\ B\ \varphi] = A[\psi\ B\ (E(\varphi)]$
$A[\psi\ \underline{B}\ \varphi] = A[\psi\ \underline{B}\ (E(\varphi)]$

Weak Equivalences
$*[\varphi U\psi] := [\varphi\underline{U}\psi]\vee G\varphi$ $\qquad$ $*\ [\varphi B\psi] := [\varphi\underline{B}\psi]\vee G\neg\psi$
*same to past version $\qquad$ $[\varphi W\psi] := \neg[(\neg\varphi)\underline{W}\psi]$
$\overleftarrow{X}\varphi := \neg\underline{\overleftarrow{X}}\neg\varphi$ *(at t0 : weak true. strong false)*

Negation Normal Form
$\neg(\varphi\wedge\psi) = \neg\varphi\vee\neg\psi$ $\qquad$ $\neg(\varphi\vee\psi) = \neg\varphi\wedge\neg\psi$
$\neg\neg\varphi = \varphi$ $\qquad$ $\neg X\varphi = X\neg\varphi$
$\neg G\varphi = F\neg\varphi$ $\qquad$ $\neg F\varphi = G\neg\varphi$
$\neg[\varphi\ U\ \psi] = [(\neg\varphi)\ \underline{B}\ \psi]$ $\qquad$ $\neg[\varphi\ \underline{U}\ \psi] = [(\neg\varphi)\ B\ \psi]$
$\neg[\varphi\ B\ \psi] = [(\neg\varphi)\ \underline{U}\ \psi]$ $\qquad$ $\neg[\varphi\ \underline{B}\ \psi] = [(\neg\varphi)\ U\ \psi]$
$\neg A\varphi = E\neg\varphi$ $\qquad$ $\neg E\varphi = A\neg\varphi$
$\neg\overleftarrow{X}\varphi = \underline{\overleftarrow{X}}\neg\varphi$ $\qquad$ $\neg\underline{\overleftarrow{X}}\varphi = \overleftarrow{X}\neg\varphi$
$\neg\overleftarrow{G}\varphi = \overleftarrow{F}\neg\varphi$ $\qquad$ $\neg\overleftarrow{F}\varphi = \overleftarrow{G}\neg\varphi$
$\neg[\varphi\ \overleftarrow{U}\ \psi] = [(\neg\varphi)\ \underline{\overleftarrow{B}}\ \psi]$ $\qquad$ $\neg[\varphi\ \underline{\overleftarrow{U}}\ \psi] = [(\neg\varphi)\ \overleftarrow{B}\ \psi]$
$\neg[\varphi\ \overleftarrow{B}\ \psi] = [(\neg\varphi)\ \underline{\overleftarrow{U}}\ \psi]$ $\qquad$ $\neg[\varphi\ \underline{\overleftarrow{B}}\ \psi] = [(\neg\varphi)\ \overleftarrow{U}\ \psi]$

Equivalences and Tips
$[\varphi\underline{U}\psi] \equiv \varphi\ don't\ matter\ when\ \psi\ hold$
$[\varphi\underline{B}\psi] \equiv \psi\ can't\ hold\ when\ \varphi\ hold$
$[\varphi\underline{W}\psi] \equiv \neg\psi\ hold\ until\ \varphi\ \wedge\ \psi$
$[\varphi U\psi] \equiv [\varphi\underline{U}\psi]\vee G\varphi$
$[a\underline{U}Fb] \equiv Fb$
$F[a\underline{U}b] \equiv Fb \equiv [Fa\underline{U}Fb]$
$[\varphi B\psi] \equiv [\varphi\underline{B}\psi]\vee G\neg\psi$
$F[a\underline{B}b] \equiv F[a\wedge\neg b]$
$[\varphi W\psi] \equiv \neg[\neg\varphi\underline{W}\psi]$
$AEA \equiv A$ $\qquad$ $\bullet\ GFX \equiv GXF$
$FF\varphi \equiv F\varphi$ $\qquad$ $\bullet\ GG\varphi \equiv G\varphi$
$GF\varphi \equiv XGF\varphi \equiv FGF\varphi \equiv GGF\varphi \equiv GFGF\varphi \equiv FGGF\varphi$
$FG\varphi \equiv XFG\varphi \equiv FFG\varphi \equiv GFG\varphi \equiv GFFG\varphi \equiv FGFG\varphi$
$GF(x\vee y) \equiv GFx\vee GFy$
$E(\varphi\wedge\psi) \equiv E\varphi\wedge E\psi(in\ general)$
$E(\varphi\vee\psi) \equiv E\varphi\vee E\psi$
$E[(a\underline{U}b)\wedge(c\underline{U}d)] \equiv$
$\quad E[(a\wedge c)\underline{U}(b\wedge E(c\underline{U}d)\vee d\wedge E(a\underline{U}b))]$
$AG(\varphi\wedge\psi) \equiv AG\varphi\wedge AG\psi$

Eliminate boolean op. after path quantify
$$[\varphi_1\ \underline{U}\ \psi_1]\wedge[\varphi_2\ \underline{U}\ \psi_2] = \left[(\varphi_1\wedge\varphi_2)\ \underline{U}\ \begin{pmatrix}\psi_1\wedge[\varphi_2\ \underline{U}\psi_2]\vee\\ \psi_2\wedge[\varphi_1\ \underline{U}\psi_1]\end{pmatrix}\right]$$
$$[\varphi_1\ \underline{U}\ \psi_1]\wedge[\varphi_2\ U\ \psi_2] = \left[(\varphi_1\wedge\varphi_2)\ \underline{U}\ \begin{pmatrix}\psi_1\wedge[\varphi_2\ U\psi_2]\vee\\ \psi_2\wedge[\varphi_1\ \underline{U}\psi_1]\end{pmatrix}\right]$$
$$[\varphi_1\ U\ \psi_1]\wedge[\varphi_2\ U\ \psi_2] = \left[(\varphi_1\wedge\varphi_2)\ \underline{U}\ \begin{pmatrix}\psi_1\wedge[\varphi_2\ U\psi_2]\vee\\ \psi_2\wedge[\varphi_1\ U\psi_1]\end{pmatrix}\right]$$

###**MONADIC PREDICATE**
S1S
First order terms are defined as follows:
$-0\in Term^{S1S}_{\sum}$
$-t\in V_{\sum}|typ_{\sum}(t)=\mathbb{N}\subseteq Term^{S1S}_{\sum}$
$-SUC(\tau)\in Term^{S1S}_{\sum} if \tau\in Term^{S1S}_{\sum}$
Formulas ζ_{S1S} are defined as:
$-p^{(t)}\in L_{S1S}$ (predicate p at time t)
$-\neg\varphi,\varphi\wedge\psi\in L_{S1S}$
$-\exists t.\varphi\in L_{S1S}$
$-\exists p.\varphi\in L_{S1S}$
where:
$-\tau\in Term^{S1S}_{\sum}$
$-\varphi,\psi\in\zeta_{S1S}$
$-t\in V_{\sum}\ with\ typ_{\sum}(t)=\mathbb{N}$
$-p\in V_{\sum}\ with\ typ_{\sum}(p)=\mathbb{N}\rightarrow\mathbb{B}$
LO2
first order terms are defined as:
$-t\in V_{\sum}|typ_{\sum}(t)=\mathbb{N}\subseteq Term^{LO2}_{\sum}$
formulas LO2 are defined as:
$-t1<t2\in L_{LO2}$
$-p^{(t)}\in L_{LO2}$
$-\neg\varphi,\varphi\wedge\psi\in L_{LO2}$
$-\exists t.\varphi\in L_{LO2}$
$-\exists p.\varphi\in L_{LO2}$
where:
$-t,t_1,t_2\tau\in V_{\sum}\ with\ typ_{\sum}(_t)=typ_{\sum}(_{t1})=typ_{\sum}(t_{t2})=\mathbb{N}$
$-\varphi,\psi\in\zeta_{LO2}$
$-t\in V_{\sum}\ with\ typ_{\sum}(t)=\mathbb{N}$
$-p\in V_{\sum}\ with\ typ_{\sum}(p)=\mathbb{N}\rightarrow\mathbb{B}$
LO2' Consider the following set $\zeta_{LO2'}$ of formulas:
$-Subset(p,q),Sing(p),andPSUC(p,q)belongto\zeta_{LO2'}$
$-\neg\varphi,\varphi\wedge\psi$
$-\exists p.\varphi$
where $-\varphi,\psi\in\zeta_{LO2'}$
$-p\in V_{\sum}\ with\ typ_{\sum}(p)=\mathbb{N}\rightarrow\mathbb{B}$
$\zeta_{LO2'}\ hasnonumericvariables$
numeric variable t is replaced by a singleton set p_t
$\zeta_{LO2'}$ is as expressive as LO2 and S1S

###**TRANSLATIONS**
CTL* Modelchecking to LTL model checking
Let's φ_i be a pure path formula (without path quantifiers), Ψ be a propositional formula, abbreviate subformulas $E\varphi$ and $A\psi$ working bottom-up the syntax tree to obtain the following
normal form: $\phi = \text{let}\begin{bmatrix}x_1=A\varphi_1\\ \vdots\\ x_n=A\varphi_n\end{bmatrix}\text{in }\Psi\text{ end}$
Use LTL model checking to compute $Q_i := [\![A\varphi_i]\!]_{\mathcal{K}_{i-1}}$, where $\mathcal{K}_0 := \mathcal{K}$ and $\mathcal{K}_{i+1}$ is obtained from $\mathcal{K}_i$ by labelling the states Q_i with x_i.
Finally compute $[\![\Psi]\!]_{\mathcal{K}_n}$

function LO2_S1S(Φ)
  case Φ **of**
    $t1<t2$: **return** $\exists p.[\forall t.p^{(t)}\rightarrow p^{(SUC(t))}]\wedge\neg p^{(t1)}\wedge p^{(t2)}$;
    $p^{(t)}$: **return** $p^{(t)}$;
    $\neg\varphi$: **return** $\neg LO2\_S1S(\varphi)$;
    $\varphi\wedge\psi$: **return** $LO2\_S1S(\varphi)\wedge LO2\_S1S(\psi)$;
    $\exists t.\varphi$: **return** $\exists t.LO2\_S1S(\varphi)$;
    $\exists p.\varphi$: **return** $\exists p.LO2\_S1S(\varphi)$;
  end
end
function S1S_LO2(Φ)
  case Φ **of**
    $p^{(n)}$:
return $\exists t0...tn.p^{(tn)}\wedge zero(t0)\wedge\bigwedge_{i=0}^{n-1} succ(ti,ti+1)$;
    $p^{(t0+n)}$:
return $\exists t1...tn.p^{(tn)}\wedge\bigwedge_{i=0}^{n-1} succ(ti,ti+1)$;
    $\neg\varphi$: **return** $\neg S1S\_LO2(\varphi)$;
    $\varphi\wedge\psi$: **return** $S1S\_LO2(\varphi)\wedge S1S\_LO2(\psi)$;
    $\exists t.\varphi$: **return** $\exists t.S1S\_LO2(\varphi)$;
    $\exists p.\varphi$: **return** $\exists p.S1S\_LO2(\varphi)$;
  end
end
function Tp2Od(t0, Φ) *temporal to LO1*
  case Φ **of**
    $is\_var(\Phi)$: $\Psi^{(t0)}$;
    $\neg\varphi$: **return** $\neg Tp2Od(\varphi)$;
    $\varphi\wedge\psi$: **return** $Tp2Od(\varphi)\wedge Tp2Od(\psi)$;
    $\varphi\vee\psi$: **return** $Tp2Od(\varphi)\vee Tp2Od(\psi)$;
    $X\varphi : \Psi := \exists t1.(t0<t1)\wedge$
      $\forall t2.t0<t2\rightarrow t1\leq t2)\wedge Tp2Od(t1,\varphi)$;
    $[\varphi\underline{U}\psi] : \Psi := \exists t1.t0\leq t1\wedge Tp2Od(t1,\psi)\wedge$
      $interval((t0,1,t1,0),\varphi)$;
    $[\varphi B\psi] : \Psi := \forall t1.t0\leq t1\wedge$
      $interval((t0,1,t1,0),\neg\varphi)\rightarrow Tp2Od(t1,\neg\psi)$;
    $\overleftarrow{X}\varphi : \Psi := \forall t1.(t1<t0)\wedge$
      $(\forall t2.t2<t0\rightarrow t2\leq t1)\rightarrow Tp2Od(t1,\varphi)$;
    $\underline{\overleftarrow{X}}\varphi : \Psi := \exists t1.(t1<t0)\wedge$
      $(\forall t2.t2<t0\rightarrow t2\leq t1)\wedge Tp2Od(t1,\varphi)$;
    $[\varphi\underline{\overleftarrow{U}}\psi] : \Psi := \exists t1.t1\leq t0\wedge Tp2Od(t1,\psi)\wedge$
      $interval((t1,0,t0,1),\varphi)$;
    $[\varphi\overleftarrow{B}\psi] : \Psi := \forall t1.t1\leq t0\wedge$
      $interval((t1,0,t0,1),\neg\varphi)\rightarrow Tp2Od(t1,\neg\psi)$;
  end
  return Ψ
end
function interval(l, φ)
  case Φ **of**
    $(t0,0,t1,0)$:
      return $\forall t2.t0<t2\wedge t2<t1\rightarrow Tp2Od(t2,\varphi)$;
    $(t0,0,t1,1)$:
      return $\forall t2.t0<t2\wedge t2\leq t1\rightarrow Tp2Od(t2,\varphi)$;
    $(t0,1,t1,0)$:
      return $\forall t2.t0\leq t2\wedge t2<t1\rightarrow Tp2Od(t2,\varphi)$;
    $(t0,1,t1,1)$:
      return $\forall t2.t0\leq t2\wedge t2\leq 3t1\rightarrow Tp2Od(t2,\varphi)$;
  end
end

ω-Automaton to LO2
$A_{\exists}(q1,...,qn,\psi I,\psi R,\psi F)(input\ automaton)$
$\exists q1..qn,\Theta LO2(0,\psi I)\wedge(\forall t.\Theta LO2(t,\psi R))\wedge(\forall.t1\exists t2.t1<t2\wedge\Theta LO2(t2,\psi F))$
Where ΘLO2(t, Φ) is:
-$\Theta LO2(t,p) := p(t)\ for\ variable\ p$
-$\Theta LO2(t,X\psi) := \Theta LO2(t+1,\psi)$
-$\Theta LO2(t,\neg\psi) := \neg\Theta LO2(t,\psi)$
-$\Theta LO2(t,\varphi\wedge\psi) := \Theta LO2(t,\varphi)\wedge\Theta LO2(t,\psi)$
-$\Theta LO2(t,\varphi\vee\psi) := \Theta LO2(t,\varphi)\vee\Theta LO2(t,\psi)$

LTL to ω-automata
$\phi\langle X\varphi\rangle_x \Leftrightarrow \mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow X\varphi,\phi\langle q\rangle_x)$
$\phi\langle X\varphi\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q_0,q_1\},1,(q_0\leftrightarrow\varphi)\wedge(q_1\leftrightarrow Xq_0),\phi\langle q_1\rangle_x)$
$\phi\langle G\varphi\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow\varphi\wedge Xq,\phi\langle q\rangle_x\wedge GF[\varphi\rightarrow q])$
$\phi\langle F\varphi\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow\varphi\wedge Xq,\phi\langle q\rangle_x\wedge GF[q\rightarrow\varphi])$
$\phi\langle[\varphi\ U\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow\psi\vee\varphi\wedge Xq,\phi\langle q\rangle_x\wedge GF[\varphi\rightarrow q])$
$\phi\langle[\varphi\ \underline{U}\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow\psi\vee\varphi\wedge Xq,\phi\langle q\rangle_x\wedge GF[q\rightarrow\psi])$
$\phi\langle[\varphi\ B\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow\neg\psi\wedge(\varphi\vee Xq),\phi\langle q\rangle_x\wedge GF[q\vee\psi])$
$\phi\langle[\varphi\ \underline{B}\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow\neg\psi\wedge(\varphi\vee Xq),\phi\langle q\rangle_x\wedge GF[q\rightarrow\varphi])$
$\phi\langle\overleftarrow{X}\varphi\rangle_x \Leftrightarrow \mathcal{A}_{\exists}(\{q\},q,Xq\leftrightarrow\varphi,\phi\langle q\rangle_x)$
$\phi\langle\overleftarrow{X}\varphi\rangle_x \Leftrightarrow \mathcal{A}_{\exists}(\{q\},\neg q,Xq\leftrightarrow\varphi,\phi\langle q\rangle_x)$
$\phi\langle\overleftarrow{G}\varphi\rangle_x \Leftrightarrow \mathcal{A}_{\exists}(\{q\},q,Xq\leftrightarrow\varphi\wedge q,\phi\langle\varphi\wedge q\rangle_x)$
$\phi\langle\overleftarrow{F}\varphi\rangle_x \Leftrightarrow \mathcal{A}_{\exists}(\{q\},\neg q,Xq\leftrightarrow\varphi\vee q,\phi\langle\varphi\vee q\rangle_x)$
$\phi\langle[\varphi\ \overleftarrow{U}\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\},q,Xq\leftrightarrow\psi\vee\varphi\wedge q,\phi\langle\psi\vee\varphi\wedge q\rangle_x)$
$\phi\langle[\varphi\ \underline{\overleftarrow{U}}\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\},\neg q,Xq\leftrightarrow\psi\vee\varphi\wedge q,\phi\langle\psi\vee\varphi\wedge q\rangle_x)$
$\phi\langle[\varphi\ \overleftarrow{B}\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\},q,Xq\leftrightarrow\neg\psi\wedge(\varphi\vee q),\phi\langle\neg\psi\wedge(\varphi\vee q)\rangle_x)$
$\phi\langle[\varphi\ \underline{\overleftarrow{B}}\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\},\neg q,Xq\leftrightarrow\neg\psi\wedge(\varphi\vee q),\phi\langle\neg\psi\wedge(\varphi\vee q)\rangle_x)$

CTL to $\mu - Calculus(\Phi_{inf}=\nu y.\Diamond y)$
$EX\varphi = \Diamond(\Phi_{inf}\wedge\varphi)$
$EG\varphi = \nu x.\varphi\wedge\Diamond x$
$EF\varphi = \mu x.\Phi_{inf}\wedge\varphi\vee\Diamond x$
$E[\varphi\underline{U}\psi] = \mu x.(\Phi_{inf}\wedge\psi)\vee\varphi\wedge\Diamond x$
$E[\varphi U\psi] = \nu x.(\Phi_{inf}\wedge\psi)\vee\varphi\wedge\Diamond x$
$E[\varphi\underline{B}\psi] = \mu x.\neg\psi\wedge(\Phi_{inf}\wedge\varphi\vee\Diamond x)$
$E[\varphi B\psi] = \nu x.\neg\psi\wedge(\Phi_{inf}\wedge\varphi\vee\Diamond x)$
$AX\varphi = \Box(\Phi_{inf}\rightarrow\varphi)$
$AG\varphi = \nu x.(\Phi_{inf}\rightarrow\varphi)\wedge\Box x$
$AF\varphi = \mu x.\varphi\vee\Box x$
$A[\varphi\underline{U}\psi] = \mu x.\psi\vee(\Phi_{inf}\rightarrow\varphi)\wedge\Box x$
$A[\varphi U\psi] = \nu x.\psi\vee(\Phi_{inf}\rightarrow\varphi)\wedge\Box x$
$A[\varphi\underline{B}\psi] = \mu x.(\Phi_{inf}\rightarrow\neg\psi)\wedge(\varphi\vee\Box x)$
$A[\varphi B\psi] = \nu x.(\Phi_{inf}\rightarrow\neg\psi)\wedge(\varphi\vee\Box x)$

G and μ-calculus (safety property)
-$[\nu x.\varphi\wedge\Diamond x]_K$
-Contains states s where an infinite path π starts with $\forall t.\pi^{(t)}\in[\varphi]_K$
-φ holds always on π
F and μ-calculus (liveness property)
-$[\mu x.\varphi\vee\Diamond x]_K$
-Contains states s where a (possibly finite) path π starts with $\exists t.\pi^{(t)}\in[\varphi]_K$
-φ holds at least once on π
FG and μ-calculus (persistence property)
-$[\mu y.[\nu x.\varphi\wedge\Diamond x]\vee\Diamond y]_K$
-Contains states s where an infinite path π starts with $\exists t1.\forall t2.\pi^{(t1+t2)}\in[\varphi]_K$
-φ holds after some point on π
GF and μ-calculus (fairness property)
-$[\nu y.[\mu x.(y\wedge\varphi)\vee\Diamond x]]_K$
-Contains states s where an infinite path π starts with
$\forall t1.\exists t2.\pi^{(t1+t2)}\in[\varphi]_K ?????t1+t2ort1+t0?????$
-φ holds infinitely often on π