Propositional Logic Syntactic Sugar

$\varphi \Leftrightarrow \psi := (\neg\varphi \vee \psi) \wedge (\neg\psi \vee \varphi)$ $\qquad \varphi \rightarrow \psi := \neg\varphi \vee \psi$

$\varphi \oplus \psi := (\varphi \wedge \neg\psi) \vee (\psi \wedge \neg\varphi)$ $\qquad \varphi \bar{\wedge} \psi := \neg(\varphi \wedge \psi)$

$(\alpha \Rightarrow \beta|\gamma) := (\neg\alpha \vee \beta) \wedge (\alpha \vee \gamma)$ $\qquad \varphi \bar{\vee} \psi := \neg(\varphi \vee \psi)$

Distributivity: $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$
$a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$
De Morgan: $\neg(a \vee b) \equiv (\neg a \wedge \neg b)$
$\neg(a \wedge b) \equiv (\neg a \vee \neg b)$
CNF: from truth table, take minterms that are 0. Each minterm is built as an OR of the negated variables. E.g., $(0,0,1) \rightarrow (x \vee y \vee \neg z)$.

###**SAT SOLVERS**
Satisfiability, Validity and Equivalence

$\mathrm{SAT}(\varphi) := \neg\mathrm{VALID}(\neg\varphi)$ $\quad \varphi \Leftrightarrow \psi := \mathrm{VALID}(\varphi \leftrightarrow \psi)$
$\mathrm{VALID}(\varphi) := (\varphi \Leftrightarrow 1)$ $\qquad \mathrm{SAT}(\varphi) := \neg(\varphi \Leftrightarrow 0)$.

Sequent Calculus:
-*Validity*: start with $\{\} \vdash \phi$; valid iff $\Gamma \cap \Delta \neq \{\}$ FOR ALL leaves.
-*Satisfiability*: start with $\{\phi\} \vdash \{\}$; satisfiable iff $\Gamma \cap \Delta = \{\}$ for AT LEAST ONE leaf.
-Counterexample/sat variable assignment: var is true, if $x \in \Gamma$; false otherwise; "don't care", if variable doesn't appear.

OPER.	LEFT	RIGHT
NOT	$\frac{\neg\phi,\Gamma\vdash\Delta}{\Gamma\vdash\phi,\Delta}$	$\frac{\Gamma\vdash\neg\phi,\Delta}{\phi,\Gamma\vdash\Delta}$
AND	$\frac{\phi\wedge\psi,\Gamma\vdash\Delta}{\phi,\psi,\Gamma\vdash\Delta}$	$\frac{\Gamma\vdash\phi\wedge\psi,\Delta}{\Gamma\vdash\phi,\Delta \qquad \Gamma\vdash\psi,\Delta}$
OR	$\frac{\phi\vee\psi,\Gamma\vdash\Delta}{\phi,\Gamma\vdash\Delta \qquad \psi,\Gamma\vdash\Delta}$	$\frac{\Gamma\vdash\phi\vee\psi,\Delta}{\Gamma\vdash\phi,\psi,\Delta}$

Resolution Calculus $\frac{\{\neg x\}\cup C_1 \qquad \{x\}\cup C_2}{C_1 \cup C_2}$

To prove unsatisfiability of given clauses in CNF: If we reach $\{\}$, the formula is unsatisfiable. E.g., $\{\{a\},\{\neg a,b\},\{\neg b\}\}$, we get:
$\{a\}+\{\neg a,b\} \rightarrow \{b\}; \{b\}+\{\neg b\} \rightarrow \{\}$ (unsatisfiable).
To prove validity, prove UNSAT of negated formula.

Davis Putnam Procedure - proves SAT; To prove validity: prove unsatisfiability of negated formula.
(1) Compute Linear Clause Form
(Don't forget to create the last clause $\{x_n\}$)
(2) Last variable has to be $\underline{1}$ (true) $\rightarrow$ find implied variables.
(3) For remaining variables: assume values and compute newly implied variables.
(4) If contradiction reached: backtrack.

Linear Clause Forms (Computes CNF) - Bottom up (inside out) in the syntax tree: convert "operators and variables" into new variable. E.g., $\neg a \vee b$ becomes $x_1 \leftrightarrow \neg a; x_2 \leftrightarrow x_1 \vee b$. Use rules below to find CNF. Create last clause {Xn}

$x \leftrightarrow \neg y \Leftrightarrow (\neg x \vee \neg y) \wedge (x \vee y)$
$x \leftrightarrow y_1 \wedge y_2 \Leftrightarrow$
$\quad (\neg x \vee y_1) \wedge (\neg x \vee y_2) \wedge (x \vee \neg y_1 \vee \neg y_2)$
$x \leftrightarrow y_1 \vee y_2 \Leftrightarrow$
$\quad (\neg x \vee y_1 \vee y_2) \wedge (x \vee \neg y_1) \wedge (x \vee \neg y_2)$
$x \leftrightarrow y_1 \rightarrow y_2 \Leftrightarrow$
$\quad (x \vee y_1) \wedge (x \vee \neg y_1 \vee \neg y_2) \wedge (\neg x \vee \neg y_1 \vee y_2)$
$x \leftrightarrow (y_1 \leftrightarrow y_2) \Leftrightarrow (x \vee y_1 \vee y_2) \wedge (x \vee \neg y_1 \vee \neg y_2) \wedge$
$\quad (\neg x \vee y_1 \vee \neg y_2) \wedge (\neg x \vee \neg y_1 \vee y_2)$
$x \leftrightarrow y_1 \oplus y_2 \Leftrightarrow (x \vee \neg y_1 \vee y_2) \wedge (x \vee y_1 \vee \neg y_2) \wedge$
$\quad (\neg x \vee y_1 \vee y_2) \wedge (\neg x \vee \neg y_1 \vee \neg y_2)$

```
Compose(int x, BddNode ψ, α)
 int m; BddNode h, l;
 if x>label(ψ) then
  return ψ;
 elseif x=label(ψ) then
  return ITE(α,high(ψ),
       low(ψ));
 else
  m=max{label(ψ),label(α)}
  (α0,α1):=Ops(α, m);
  (ψ0,ψ1):=Ops(ψ, m);
  h:=Compose(x,ψ1,α1);
  l:=Compose(x,ψ0,α0);
  return CreateNode(m,h,l)
 endif;end
```

```
ITE(BddNode i, j, k)
 int m; BddNode h, l;
 if i = 0 then return k
 elseif i=1 then
  return j
 elseif j=k then
  return k
 else
  m = max{label(i),
      label(j),label(k)}
  (i0,i1):=Ops(i,m);
  (j0,j1):=Ops(j,m);
  (k0,k1):=Ops(k,m);
  l:=ITE(i0,j0,k0);
  h:=ITE(i1, j1, k1);
  return CreateNode(m,h,l)
 endif;end
```

```
Constrain(Φ, β)
 if β=0 then
  ret 0
 elseif Φ ∈ {0,1}(β = 1)
     ret Φ
 else
  m=max{label(β),label(Φ)}
   (Φ0,Φ1):=Ops(Φ,m);
   (β0,β1):=Ops(β,m);
   if β0=0
  ret Constrain(Φ1,β1)
 elseif β1=0 then
  ret Constrain(Φ0,β0)
 else
  l:=Constrain(Φ0,β0);
  h:=Constrain(Φ1,β1);
  ret CreateNode(m,h,l)
endif;endif;end
```

```
Apply(⊙, Bddnode a, b)
 int m; BddNode h, l;
 if isLeaf(a)&isLeaf(b)
        then
  return Eval(⊙,label(a),
        label(b));
 else
  m=max{label(a),label(b)}
  (a0,a1):=Ops(a,m);
  (b0,b1):=Ops(b,m);
  h:=Apply(⊙,a1,b1);
  l:=Apply(⊙,a0,b0);
  return CreateNode(m,h,l)
 end;
end
```

```
Restrict(Φ,β)
 if β=0
  return 0
 elseif
       Φ ∈ {0,1} ∨ (β = 1)

  return Φ
 else
  m=max{label(β),label(Φ)}
  (Φ0,Φ1):=Ops(Φ,m);
  (β0,β1):=Ops(β,m)
 if β0=0
  return Restrict(Φ1,β1)
 elseif β1=0
  return Restrict(Φ0,β0)
 elseif m=label(Φ)
  return CreateNode(m,
   Restrict (Φ1,β1),
   Restrict (Φ0,β0))
 else
  return Restrict(Φ,
  Apply(∨,β0,β1))
 endif;endif;end
--------------------
Ops(v,m)
 x:=label(v);
 if m=degree(x)
  return (low(v),high(v))
 else return(v, v)
end;end
```

```
Exists(BddNode e, φ)
 if isLeaf(φ)∨isLeaf(e)
  return φ;
 elseif label(e)>label(φ)
  return Exist(high(e),φ)
 elseif label(e)=label(φ)
  h=Exist(high(e),high(φ)
  l=Exist(high(e),low(φ))
  return Apply(∨,l,h)
 else (label(e)<label(φ))
  h:=Exists(e,high(φ))
  l:=Exists(e,low(φ))
  return CreateNode(label(
     φ),h,l)
 endif; end function.
----------------------
ZDD: If positive cofactor
      = 0, redirect edge
      to negative
      cofactor.
If variable not in the
      formula, add with
      both edges pointing
      to same node.
----------------------
FDD: Positive Davio
      Decomposition. (
      Keep both edges to
      1 if happens!)
```

$\varphi = [\varphi]_x^0 \oplus x \wedge (\partial\varphi/\partial x)$
$(\partial\varphi/\partial x) := [\varphi]_x^0 \oplus [\varphi]_x^1$

Local Model Checking

$\frac{s\vdash_\Phi \varphi\wedge\psi}{\{s\vdash_\Phi\varphi\} \quad \{s\vdash_\Phi\psi\}}\wedge$		$\frac{s\vdash_\Phi \varphi\vee\psi}{\{s\vdash_\Phi\varphi\} \quad \{s\vdash_\Phi\psi\}}\vee$	
$\frac{s\vdash_\Phi \Box\varphi}{\{s_1\vdash_\Phi\varphi\}\ldots\{s_n\vdash_\Phi\varphi\}}\wedge$		$\frac{s\vdash_\Phi \Diamond\varphi}{\{s_1\vdash_\Phi\varphi\}\ldots\{s_n\vdash_\Phi\varphi\}}\vee$	
$\frac{s\vdash_\Phi \overleftarrow{\Box}\varphi}{\{s'_1\vdash_\Phi\varphi\}\ldots\{s'_n\vdash_\Phi\varphi\}}\wedge$		$\frac{s\vdash_\Phi \overleftarrow{\Diamond}\varphi}{\{s'_1\vdash_\Phi\varphi\}\ldots\{s'_n\vdash_\Phi\varphi\}}\vee$	
$\frac{s\vdash_\Phi \mu x.\varphi}{s\vdash_\Phi\varphi}$	$\frac{s\vdash_\Phi \nu x.\varphi}{s\vdash_\Phi\varphi}$	$\frac{s\vdash_\Phi x}{s\vdash_\Phi \mathfrak{D}_\Phi(x)}$	$\mathfrak{D}_\Phi$(replace w. initial form.)
$\{s_1 \ldots s_n\} = suc_\exists^{\mathcal{R}}(s)$ and $\{s'_1 \ldots s'_n\} = pre_\exists^{\mathcal{R}}(s)$			

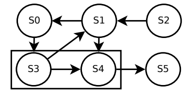
Approximations and Ranks

If (s,$\mu x.\varphi$) repeats$\rightarrow$return 1	$apx_0(\mu x.\varphi) := 0$
If (s,$\nu x.\varphi$) repeats$\rightarrow$return 0	$apx_0(\nu x.\varphi) := 1$

Tarski-Knaster Theorem: $\mu :=$ starts $\perp \rightharpoonup$ least fixpoint $\spadesuit$ $\nu :=$ starts $\top \rightharpoonup$ greatest fixpoint

Quantif. $\exists x.\varphi := [\varphi]_x^1 \vee [\varphi]_x^0$ $\clubsuit$ $\forall x.\varphi := [\varphi]_x^1 \wedge [\varphi]_x^0$

Predecessor and Successor
$\Diamond := pre_\exists^{\mathcal{R}}(Q) := \exists x'_1,...,x'_n.\varphi_{\mathcal{R}} \wedge [\varphi_Q]_{x_1,...,x_n}^{x'_1,...,x'_n}$
$\overleftarrow{\Diamond} := suc_\exists^{\mathcal{R}}(Q) := [\exists x_1,...,x_n.\varphi_{\mathcal{R}} \wedge \varphi_Q]_{x'_1,...,x'_n}^{x_1,...,x_n}$
$\Box = pre_\forall^{\mathcal{R}}(Q) := \forall x'_1,...,x'_n.\varphi_{\mathcal{R}} \rightarrow [\varphi_Q]_{x_1,...,x_n}^{x'_1,...,x'_n}$
$\overleftarrow{\Box} := suc_\forall^{\mathcal{R}}(Q) := [\forall x_1,...,x_n.\varphi_{\mathcal{R}} \rightarrow \varphi_Q]_{x'_1,...,x'_n}^{x_1,...,x_n}$



Example: $\Box/\overleftarrow{\Box}$
$pre_\forall^{\mathcal{R}}(\{S3,S4\}) = \{S0,S5\}$
$suc_\forall^{\mathcal{R}}(\{S3,S4\}) = \{S2,S5\}$

```
pre∀R(Q = {S1,...,Sn})
for each node n in K:
 if(n points to a node
      that is not in Q)
  n ∉ pre∀R(Q)
 else
  n ∈ pre∀R(Q)
```

```
suc∀R(Q = {S1,...,Sn})
for each node n in K:
 if(n is pointed by a node
      that is not in Q)
  n ∉ suc∀R(Q)
 else
  n ∈ suc∀R(Q)
```

###**AUTOMATA**
Automata types: G$\rightarrow$Safety; F$\rightarrow$Liveness; FG$\rightarrow$Persistence/Co-Buchi; GF$\rightarrow$Fairness/Buchi.
Automaton Determinization
$\mathbf{NDet_G \rightarrow Det_G}$**:** 1.Remove all states/edges that do not satisfy acceptance condition; 2.Use Subset construction (Rabin-Scott); 3.Acceptance condition will be the states where {} is never reached.
$\mathbf{\{NDet_F (partial)\ or\ NDet_{prefix}\} \rightarrow Det_{FG}}$**:** Breakpoint Construction.
$\mathbf{NDet_F\ (total) \rightarrow Det_F}$**:** Subset Construction.
$\mathbf{NDet_{FG} \rightarrow Det_{FG}}$**:** Breakpoint Construction.
$\mathbf{NDet_{GF} \rightarrow \{Det_{Rabin}\ or\ Det_{Streett}\}}$**:** Safra Algorithm.

Boolean Operations on ω-Automata
Complement
$$\neg\mathcal{A}_\forall(Q,\mathcal{I},\mathcal{R},\mathcal{F}) = \mathcal{A}_\exists(Q,\mathcal{I},\mathcal{R},\neg\mathcal{F})$$
$$\neg\mathcal{A}_\exists(Q,\mathcal{I},\mathcal{R},\mathcal{F}) = \mathcal{A}_\forall(Q,\mathcal{I},\mathcal{R},\neg\mathcal{F})$$
Conjunction
$$(\mathcal{A}_\exists(Q_1,\mathcal{I}_1,\mathcal{R}_1,\mathcal{F}_1) \wedge \mathcal{A}_\exists(Q_2,\mathcal{I}_2,\mathcal{R}_2,\mathcal{F}_2)) = \mathcal{A}_\exists(Q_1 \cup Q_2, \mathcal{I}_1 \wedge \mathcal{I}_2, \mathcal{R}_1 \wedge \mathcal{R}_2, \mathcal{F}_1 \wedge \mathcal{F}_2)$$
Disjunction
$$(\mathcal{A}_\exists(Q_1,\mathcal{I}_1,\mathcal{R}_1,\mathcal{F}_1) \vee \mathcal{A}_\exists(Q_2,\mathcal{I}_2,\mathcal{R}_2,\mathcal{F}_2)) = \mathcal{A}_\exists\begin{pmatrix} Q_1 \cup Q_2 \cup \{q\}, \\ (\neg q \wedge \mathcal{I}_1) \vee (q \wedge \mathcal{I}_2), \\ (\neg q \wedge \mathcal{R}_1 \wedge \neg q') \vee (q \wedge \mathcal{R}_2 \wedge q'), \\ (\neg q \wedge \mathcal{F}_1) \vee (q \wedge \mathcal{F}_2) \end{pmatrix}$$
If both automata are totally defined,
$$(\mathcal{A}_\exists(Q_1,\mathcal{I}_1,\mathcal{R}_1,\mathcal{F}_1) \vee \mathcal{A}_\exists(Q_2,\mathcal{I}_2,\mathcal{R}_2,\mathcal{F}_2)) = \mathcal{A}_\exists(Q_1 \cup Q_2, \mathcal{I}_1 \wedge \mathcal{I}_2, \mathcal{R}_1 \wedge \mathcal{R}_2, \mathcal{F}_1 \vee \mathcal{F}_2)$$
Eliminate Nesting - Acceptance condition **must** be an automata of the same type
$$\mathcal{A}_\exists(Q^1,\mathcal{I}_1^1,\mathcal{R}_1^1,\mathcal{A}_\exists(Q^2,\mathcal{I}_1^2,\mathcal{R}_1^2,\mathcal{F}_1))$$
$$= \mathcal{A}_\exists(Q^1 \cup Q^2, \mathcal{I}_1^1 \wedge \mathcal{I}_1^2, \mathcal{R}_1^1 \wedge \mathcal{R}_1^2, \mathcal{F}_1))$$
Boolean Operations of G
(1)$\neg G\varphi = F\neg\varphi$ $\qquad$ (2)$G\varphi \wedge G\psi = G[\varphi \wedge \psi]$
(3)$G\varphi \vee G\psi = \mathcal{A}_\exists(\{p,q\}, p \wedge q, [p' \leftrightarrow p \wedge \varphi] \wedge [q' \leftrightarrow q \wedge \psi], G[p \vee q])$
Boolean Operations of F
(1)$\neg F\varphi = G\neg\varphi$ $\qquad$ (2)$F\varphi \vee F\psi = F[\varphi \vee \psi]$
(3)$F\varphi \wedge F\psi = \mathcal{A}_\exists(\{p,q\}, \neg p \wedge \neg q, [p' \leftrightarrow p \vee \varphi] \wedge [q' \leftrightarrow q \vee \psi], F[p \wedge q])$
Boolean Operations of FG
(1)$\neg FG\varphi = GF\neg\varphi$ $\qquad$ (2)$FG\varphi \wedge FG\psi = FG[\varphi \wedge \psi]$
(3)$FG\varphi \vee FG\psi = \mathcal{A}_\exists(\{q\}, \neg q, q' \leftrightarrow (q \Rightarrow \psi|\neg\varphi), FG[\neg q \vee \psi])$
Boolean Operations of GF
(1)$\neg GF\varphi = FG\neg\varphi$ $\qquad$ (2)$GF\varphi \vee GF\psi = GF[\varphi \vee \psi]$
(3)$GF\varphi \wedge GF\psi = \mathcal{A}_\exists(\{q\}, \neg q, q' \leftrightarrow (q \Rightarrow \neg\psi|\varphi), GF[q \wedge \psi])$

Transformation of Acceptance Conditions
Reduction of G
$G\varphi = \mathcal{A}_\exists(\{q\}, q, \varphi \wedge q \wedge q', Fq))$
$G\varphi = \mathcal{A}_\exists(\{q\}, q, q' \leftrightarrow q \wedge \varphi, FGq)$
$G\varphi = \mathcal{A}_\exists(\{q\}, q, q' \leftrightarrow q \wedge \varphi, GFq)$
Reduction of F
$F\varphi$ can **not** be expressed by $NDet_G$
$F\varphi = \mathcal{A}_\exists(\{q\}, \neg q, q' \leftrightarrow q \vee \varphi, FGq)$
$F\varphi = \mathcal{A}_\exists(\{q\}, \neg q, q' \leftrightarrow q \vee \varphi, GFq)$
Reduction of FG
$FG\varphi$ can **not** be expressed by $NDet_G$
$FG\varphi = \mathcal{A}_\exists(\{q\}, \neg q, q \rightarrow \varphi \wedge q', Fq)$
$$FG\varphi = \mathcal{A}_\exists\left(\begin{matrix} \{p,q\}, \qquad \neg p \wedge \neg q, \\ \left[\begin{matrix}(p \rightarrow p') \wedge (p' \rightarrow p \vee \neg q) \wedge \\ (q' \leftrightarrow (p \wedge \neg q \vee \neg\varphi) \vee (p \wedge q))\end{matrix}\right], \\ G\neg q \wedge Fp \end{matrix}\right)$$
$$FG\varphi = \mathcal{A}_\exists\left(\begin{matrix} \{p,q\}, \qquad \neg p \wedge \neg q, \\ \left[\begin{matrix}(p \rightarrow p') \wedge (p' \rightarrow p \vee \neg q) \wedge \\ (q' \leftrightarrow (p \wedge \neg q \vee \neg\varphi) \vee (p \wedge q))\end{matrix}\right], \\ GF[p \wedge \neg q] \end{matrix}\right)$$

###**TEMPORAL LOGICS**
(S1)Pure LTL: AFGa
(S2)LTL + CTL: AFa
(S3)Pure CTL: AGEFa
(S4)CTL*: AFGa $\vee$ AGEFa
Remarks *Beware of Finite Paths*
E and A quantify over infinite paths.
Aφ holds on every state that has no infinite path;
Eφ is false on every state that has no infinite path;
A0 holds on states with only finite paths;
E1 is false on state with only finite paths;
$\Box$0 holds on states with no successor states;
$\Diamond$1 holds on states with successor states.
$F\varphi = \varphi \vee XF\varphi$ $\qquad G\varphi = \varphi \wedge XG\varphi$
$[\varphi\ U\ \psi] = \psi \vee (\varphi \wedge X[\varphi\ U\ \psi])$
$[\varphi\ B\ \psi] = \neg\psi \wedge (\varphi \vee X[\varphi\ B\ \psi])$
$[\varphi\ W\ \psi] = (\psi \wedge \varphi) \vee (\neg\psi \wedge X[\varphi\ W\ \psi])$

LTL Syntactic Sugar: analog for past operators
$G\varphi = \neg[1\ \underline{U}\ (\neg\varphi)]$ $\qquad F\varphi = [1\ \underline{U}\ \varphi]$
$[\varphi\ W\ \psi] = \neg[(\neg\varphi \vee \neg\psi)\ \underline{U}\ (\neg\varphi \wedge \psi)]$
$[\varphi\ \underline{W}\ \psi] = [(\neg\psi)\ \underline{U}\ (\varphi \wedge \psi)]$ *($\neg\psi$ holds until $\varphi \wedge \psi$)*
$[\varphi\ B\ \psi] = \neg[(\neg\varphi)\ \underline{U}\ \psi)]$
$[\varphi\ \underline{B}\ \psi] = [(\neg\psi)\ \underline{U}\ (\varphi \wedge \neg\psi)]$*($\psi$ can't hold when φ holds)*
$[\varphi\ U\ \psi] = \neg[(\neg\psi)\ \underline{U}\ (\neg\varphi \wedge \neg\psi)]$
$[\varphi\ U\ \psi] = [\varphi\ \underline{U}\ \psi] \vee G\varphi$
$[\varphi\ \underline{U}\ \psi] = \neg[(\neg\psi)\ U\ (\neg\varphi \wedge \neg\psi)]$
$[\varphi\ \underline{U}\ \psi] = \neg[(\neg\psi)\ W\ (\varphi \rightarrow \psi)]$
$[\varphi\ \underline{U}\ \psi] = [\psi\ \underline{W}\ (\varphi \rightarrow \psi)]$
$[\varphi\ \underline{U}\ \psi] = \neg[(\neg\varphi)\ B\ \psi]$*($\varphi$ doesn't matter when ψ holds)*
$[\varphi\ \underline{U}\ \psi] = [\psi\ \underline{B}\ (\neg\varphi \wedge \neg\psi]$
CTL Syntactic Sugar: analog for past operators
Existential Operators
$EF\varphi = E[1\ \underline{U}\ \varphi]$
$EG\varphi = E[\varphi\ U\ 0]$
$E[\varphi\ U\ \psi] = E[\varphi\ \underline{U}\ \psi] \vee EG\varphi$
$E[\varphi\ B\ \psi] = E[(\neg\psi)\ \underline{U}\ (\varphi \wedge \neg\psi)] \vee EG\neg\psi$
$E[\varphi\ B\ \psi] = E[(\neg\psi)\ U\ (\varphi \wedge \neg\psi)]$
$E[\varphi\ \underline{B}\ \psi] = E[(\neg\psi)\ \underline{U}\ (\varphi \wedge \neg\psi)]$
$E[\varphi\ \underline{B}\ \psi] = E[(\neg\psi\ \underline{U}\ (\varphi \wedge \neg\psi)]$
$E[\varphi\ W\ \psi] = E[(\neg\psi)\ \underline{U}\ (\varphi \wedge \psi)] \vee EG\neg\psi$
$E[\varphi\ W\ \psi] = E[(\neg\psi)\ U\ (\varphi \wedge \psi)]$
$E[\varphi\ \underline{W}\ \psi] = E[(\neg\psi)\ \underline{U}\ (\varphi \wedge \psi)]$
Universal Operators
$AX\varphi = \neg EX\neg\varphi$
$AG\varphi = \neg E[1\ \underline{U}\ \neg\varphi]$

$AF\varphi = \neg EG\neg\varphi$
$AF\varphi = \neg E[(\neg\varphi)\ U\ 0]$
$A[\varphi\ U\ \psi] = \neg E[(\neg\psi)\ \underline{U}\ (\neg\varphi \wedge \neg\psi)]$
$A[\varphi\ \underline{U}\ \psi] = \neg E[(\neg\psi)\ \underline{U}\ (\neg\varphi \wedge \neg\psi)] \wedge \neg EG\neg\psi$
$A[\varphi\ \underline{U}\ \psi] = \neg E[(\neg\psi)\ U\ (\neg\varphi \wedge \neg\psi)]$
$A[\varphi\ B\ \psi] = \neg E[(\neg\varphi)\ \underline{U}\ \psi]$
$A[\varphi\ \underline{B}\ \psi] = \neg E[(\neg\varphi)\ U\ \psi]$
$A[\varphi\ \underline{B}\ \psi] = \neg E[(\neg\varphi \vee \psi)\ \underline{U}\ \psi] \wedge \neg EG(\neg\varphi \vee \psi)$
$A[\varphi\ W\ \psi] = \neg E[(\neg\psi)\ \underline{U}\ (\neg\varphi \wedge \psi)]$
$A[\varphi\ \underline{W}\ \psi] = \neg E[(\neg\psi)\ \underline{U}\ (\neg\varphi \wedge \psi)] \wedge \neg EG\neg\psi$
$A[\varphi\ \underline{W}\ \psi] = \neg E[(\neg\psi)\ U\ (\neg\varphi \wedge \psi)]$

CTL* to CTL - Existential Operators

$EX\varphi = EXE\varphi$
$EF\varphi = EFE\varphi$ $\quad EFG\varphi \equiv EFEG\varphi$
$E[\varphi\ W\ \psi] = E[(E\varphi)\ W\ \psi]$
$E[\varphi\ \underline{W}\ \psi] = E[(E\varphi)\ \underline{W}\ \psi]$
$E[\psi\ U\ \varphi] = E[\psi\ U\ E(\varphi)]$
$E[\psi\ \underline{U}\ \varphi] = E[\psi\ \underline{U}\ E(\varphi)]$
$E[\varphi\ B\ \psi] = E[(E\varphi)\ B\ \psi]$
$E[\varphi\ \underline{B}\ \psi] = E[(E\varphi)\ \underline{B}\ \psi]$
obs. $EGF\varphi \neq EGEF\varphi \rightarrow$ can't be converted

CTL* to CTL - Universal Operators

$AX\varphi = AXA\varphi$
$AG\varphi = AGA\varphi$
$A[\varphi\ W\ \psi] = A[(A\varphi)\ W\ \psi]$
$A[\varphi\ \underline{W}\ \psi] = A[(A\varphi)\ \underline{W}\ \psi]$
$A[\varphi\ U\ \psi] = A[A(\varphi)\ U\ \psi]$
$A[\varphi\ \underline{U}\ \psi] = A[A(\varphi)\ \underline{U}\ \psi]$
$A[\psi\ B\ \varphi] = A[\psi\ B\ (E(\varphi)]$
$A[\psi\ \underline{B}\ \varphi] = A[\psi\ \underline{B}\ (E(\varphi)]$

Weak Equivalences

$*[\varphi U\psi] := [\varphi\underline{U}\psi] \vee G\varphi$ $\quad *[\varphi B\psi] := [\varphi\underline{B}\psi] \vee G\neg\psi$
*same to past version $\quad [\varphi W\psi] := \neg[(\neg\varphi)\underline{W}\psi]$
$\overleftarrow{X}\varphi := \neg\overleftarrow{\underline{X}}\neg\varphi$ *(at t0 : weak true. strong false)*

Negation Normal Form

$\neg(\varphi\wedge\psi) = \neg\varphi\vee\neg\psi$ $\quad \neg(\varphi\vee\psi) = \neg\varphi\wedge\neg\psi$
$\neg\neg\varphi = \varphi$ $\quad \neg X\varphi = X\neg\varphi$
$\neg G\varphi = F\neg\varphi$ $\quad \neg F\varphi = G\neg\varphi$
$\neg[\varphi\ U\ \psi] = [(\neg\varphi)\ \underline{B}\ \psi]$ $\quad \neg[\varphi\ \underline{U}\ \psi] = [(\neg\varphi)\ B\ \psi]$
$\neg[\varphi\ B\ \psi] = [(\neg\varphi)\ \underline{U}\ \psi]$ $\quad \neg[\varphi\ \underline{B}\ \psi] = [(\neg\varphi)\ U\ \psi]$
$\neg A\varphi = E\neg\varphi$ $\quad \neg E\varphi = A\neg\varphi$
$\neg\overleftarrow{X}\varphi = \overleftarrow{\underline{X}}\neg\varphi$ $\quad \neg\overleftarrow{\underline{X}}\varphi = \overleftarrow{X}\neg\varphi$
$\neg\overleftarrow{G}\varphi = \overleftarrow{F}\neg\varphi$ $\quad \neg\overleftarrow{F}\varphi = \overleftarrow{G}\neg\varphi$
$\neg[\varphi\ \overleftarrow{U}\ \psi] = [(\neg\varphi)\ \overleftarrow{\underline{B}}\ \psi]$ $\quad \neg[\varphi\ \overleftarrow{\underline{U}}\ \psi] = [(\neg\varphi)\ \overleftarrow{B}\ \psi]$
$\neg[\varphi\ \overleftarrow{B}\ \psi] = [(\neg\varphi)\ \overleftarrow{\underline{U}}\ \psi]$ $\quad \neg[\varphi\ \overleftarrow{\underline{B}}\ \psi] = [(\neg\varphi)\ \overleftarrow{U}\ \psi]$

Equivalences and Tips

$[\varphi\underline{U}\psi] \equiv \varphi$ *don't matter when* ψ *hold*
$[\varphi\underline{B}\psi] \equiv \psi$ *can't hold when* φ *hold*
$[\varphi\underline{W}\psi] \equiv \neg\psi$ *hold until* $\varphi\ \wedge\ \psi$
$[\varphi U\psi] \equiv [\varphi\underline{U}\psi] \vee G\varphi$
$[a\underline{U}Fb] \equiv Fb$
$F[a\underline{U}b] \equiv Fb \equiv [Fa\underline{U}Fb]$
$[\varphi B\psi] \equiv [\varphi\underline{B}\psi] \vee G\neg\psi$
$F[a\underline{B}b] \equiv F[a\wedge\neg b]$
$[\varphi W\psi] \equiv \neg[\neg\varphi\underline{W}\psi]$
$AEA \equiv A$ $\quad \bullet\ GFX \equiv GXF$
$FF\varphi \equiv F\varphi$ $\quad \bullet\ GG\varphi \equiv G\varphi$
$GF\varphi \equiv XGF\varphi \equiv FGF\varphi \equiv GGF\varphi \equiv GFGF\varphi \equiv FGGF\varphi$
$FG\varphi \equiv XFG\varphi \equiv FFG\varphi \equiv GFG\varphi \equiv GFFG\varphi \equiv FGFG\varphi$
$GF(x\vee y) \equiv GFx \vee GFy$
$E(\varphi\wedge\psi) \equiv E\varphi\wedge E\psi(Careful!\ Only\ sometimes!)$
$E(\varphi\vee\psi) \equiv E\varphi\vee E\psi(Careful!\ Only\ sometimes!)$
$E[(a\underline{U}b)\wedge(c\underline{U}d)] \equiv$
$\quad E[(a\wedge c)\underline{U}(b\wedge E(c\underline{U}d)\vee d\wedge E(a\underline{U}b))]$

$AG(\varphi\wedge\psi) \equiv AG\varphi\wedge AG\psi$
$A(G[aUb] \equiv G(a\vee b))$
$A(G[aBb] \equiv G(\neg b))$
$A(G[aWb] \equiv G(b\rightarrow a))$
$A(G[a\underline{U}b] \equiv G(a\vee b)\wedge GFb)$
$A(G[a\underline{B}b] \equiv G(\neg b)\wedge GFa)$
▶ *The following are initially but not generally valid*
$A(G\overleftarrow{X}a \equiv Ga)$ $\quad \bullet\ A(G\overleftarrow{\underline{X}}a \equiv false)$
$A(G\overleftarrow{G}a \equiv Ga)$ $\quad \bullet\ A(G\overleftarrow{F}a \equiv a)$
$A(G[a\overleftarrow{U}b] \equiv G(a\vee b)$ $\quad \bullet\ A(G[a\overleftarrow{\underline{U}}b] \equiv b\wedge G(a\vee b)$
$A(G[a\overleftarrow{B}b] \equiv G(\neg b)$ $\quad \bullet\ A(G[a\overleftarrow{\underline{B}}b] \equiv a\wedge G(\neg b)$
$A(G[a\overleftarrow{W}b] \equiv G(b\rightarrow a)$ $\quad \bullet\ A(G[a\overleftarrow{\underline{W}}b] \equiv b\wedge G(b\rightarrow a)$
$A(F\overleftarrow{X}a \equiv true)$ $\quad \bullet\ A(F\overleftarrow{\underline{X}}a \equiv Fa)$
$A(F\overleftarrow{G}a \equiv a)$ $\quad \bullet\ A(F\overleftarrow{F}a \equiv Fa)$
$A(F[a\overleftarrow{U}b] \equiv Fb\vee FPGa)$ $\quad \bullet\ A(F[a\overleftarrow{\underline{U}}b] \equiv Fb)$
$A(F[a\overleftarrow{B}b] \equiv F(a\wedge\neg b)\vee FPG(\neg a\wedge\neg b)$
$A(F[a\overleftarrow{\underline{B}}b] \equiv F(a\wedge\neg b)$
$A(F[a\overleftarrow{W}b] \equiv F(a\wedge b)\vee FPG\neg b)$
$A(F[a\overleftarrow{\underline{W}}b] \equiv F(a\wedge b)$

Eliminate boolean op. after path quantify

$[\varphi_1\ \underline{U}\ \psi_1]\wedge[\varphi_2\ \underline{U}\ \psi_2] =$
$$\left[(\varphi_1\wedge\varphi_2)\ \underline{U}\ \begin{pmatrix}\psi_1\wedge[\varphi_2\ \underline{U}\psi_2]\vee\\ \psi_2\wedge[\varphi_1\ \underline{U}\psi_1]\end{pmatrix}\right]$$
$[\varphi_1\ \underline{U}\ \psi_1]\wedge[\varphi_2\ U\ \psi_2] =$
$$\left[(\varphi_1\wedge\varphi_2)\ \underline{U}\ \begin{pmatrix}\psi_1\wedge[\varphi_2\ U\psi_2]\vee\\ \psi_2\wedge[\varphi_1\ \underline{U}\psi_1]\end{pmatrix}\right]$$
$[\varphi_1\ U\ \psi_1]\wedge[\varphi_2\ U\ \psi_2] =$
$$\left[(\varphi_1\wedge\varphi_2)\ \underline{U}\ \begin{pmatrix}\psi_1\wedge[\varphi_2\ U\psi_2]\vee\\ \psi_2\wedge[\varphi_1\ U\psi_1]\end{pmatrix}\right]$$

###MONADIC PREDICATE

LO2

first order terms are defined as:
$-t\in V_{\sum}|typ_{\sum}(t) = \mathbb{N}\subseteq Term^{LO2}_{\sum}$
formulas LO2 are defined as:
$-t1 < t2 \in L_{LO2}$
$-p^{(t)}\in L_{LO2}$
$-\neg\varphi,\varphi\wedge\psi\in L_{LO2}$
$-\exists t.\varphi\in L_{LO2}$
$-\exists p.\varphi\in L_{LO2}$
where:
$-t,t_1,t_2\tau\in V_{\sum}\ with\ typ_{\sum}(_t) = typ_{\sum}(_{t1}) = typ_{\sum}(_{t2}) = \mathbb{N}$
$-\varphi,\psi\in\zeta_{LO2}$
$-t\in V_{\sum}\ with\ typ_{\sum}(t) = \mathbb{N}$
$-p\in V_{\sum}\ with\ typ_{\sum}(p) = \mathbb{N}\rightarrow\mathbb{B}$

###TRANSLATIONS

CTL* Modelchecking to LTL model checking

Let's φ_i be a pure path formula (without path quantifiers), Ψ be a propositional formula, abbreviate subformulas $E\varphi$ and $A\psi$ working bottom-up the syntax tree to obtain the following normal form: $\phi = \text{let} \begin{bmatrix}x_1 = A\varphi_1\\ \vdots\\ x_n = A\varphi_n\end{bmatrix}$ in Ψ end
Use LTL model checking to compute $Q_i := [\![A\varphi_i]\!]_{\mathcal{K}_{i-1}}$, where $\mathcal{K}_0 := \mathcal{K}$ and $\mathcal{K}_{i+1}$ is obtained from $\mathcal{K}_i$ by labelling the states Q_i with x_i. Finally compute $[\![\Psi]\!]_{\mathcal{K}_n}$

function LO2_S1S(Φ)
case Φ **of**
$t1 < t2$: **return** $\exists p.[\forall t.p^{(t)}\rightarrow p^{(SUC(t))}]\wedge\neg p^{(t1)}\wedge p^{(t2)}$:
$p^{(t)}$: **return** $p^{(t)}$;
$\neg\varphi$: **return** $\neg LO2\_S1S(\varphi)$;
$\varphi\wedge\psi$: **return** $LO2\_S1S(\varphi)\wedge LO2\_S1S(\psi)$;
$\exists t.\varphi$: **return** $\exists t.LO2\_S1S(\varphi)$;
$\exists p.\varphi$: **return** $\exists p.LO2\_S1S(\varphi)$;
end
end
function S1S_LO2(Φ)
case Φ **of**
$p^{(n)}$:
return $\exists t0...tn.p^{(tn)}\wedge zero(t0)\wedge\bigwedge_{i=0}^{n-1} succ(ti,ti{+}1)$;
$p^{(t0+n)}$:
return $\exists t1...tn.p^{(tn)}\wedge\bigwedge_{i=0}^{n-1} succ(ti,ti+1)$;
$\neg\varphi$: **return** $\neg S1S\_LO2(\varphi)$;
$\varphi\wedge\psi$: **return** $S1S\_LO2(\varphi)\wedge S1S\_LO2(\psi)$;
$\exists t.\varphi$: **return** $\exists t.S1S\_LO2(\varphi)$;
$\exists p.\varphi$: **return** $\exists p.S1S\_LO2(\varphi)$;
end
end
function Tp2Od(t0, Φ) *temporal to LO1*
case Φ **of**
$is\_var(\Phi)$: $\Psi^{(t0)}$;
$\neg\varphi$: **return** $\neg Tp2Od(\varphi)$;
$\varphi\wedge\psi$: **return** $Tp2Od(\varphi)\wedge Tp2Od(\psi)$;
$\varphi\vee\psi$: **return** $Tp2Od(\varphi)\vee Tp2Od(\psi)$;
$X\varphi$: $\Psi := \exists t1.(t0 < t1)\wedge$
$\quad\forall t2.t0 < t2\rightarrow t1\leq t2)\wedge Tp2Od(t1,\varphi)$;
$[\varphi\underline{U}\psi]$: $\Psi := \exists t1.t0\leq t1\wedge Tp2Od(t1,\psi)\wedge$
$\quad interval((t0,1,t1,0),\varphi)$;
$[\varphi B\psi]$: $\Psi := \forall t1.t0\leq t1\wedge$
$\quad interval((t0,1,t1,0),\neg\varphi)\rightarrow Tp2Od(t1,\neg\psi)$;
$\overleftarrow{X}\varphi$: $\Psi := \forall t1.(t1 < t0)\wedge$
$\quad(\forall t2.t2 < t0\rightarrow t2\leq t1)\rightarrow Tp2Od(t1,\varphi)$;
$\overleftarrow{\underline{X}}\varphi$: $\Psi := \exists t1.(t1 < t0)\wedge$
$\quad(\forall t2.t2 < t0\rightarrow t2\leq t1)\wedge Tp2Od(t1,\varphi)$;
$[\varphi\overleftarrow{\underline{U}}\psi]$: $\Psi := \exists t1.t1\leq t0\wedge Tp2Od(t1,\psi)\wedge$
$\quad interval((t1,0,t0,1),\varphi)$;
$[\varphi\overleftarrow{B}\psi]$: $\Psi := \forall t1.t1\leq t0\wedge$
$\quad interval((t1,0,t0,1),\neg\varphi)\rightarrow Tp2Od(t1,\neg\psi)$;
end
return Ψ
end
function interval(l, φ)
case Φ **of**
$(t0,0,t1,0)$:
return $\forall t2.t0 < t2\wedge t2 < t1\rightarrow Tp2Od(t2,\varphi)$;
$(t0,0,t1,1)$:
return $\forall t2.t0 < t2\wedge t2\leq t1\rightarrow Tp2Od(t2,\varphi)$;
$(t0,1,t1,0)$:
return $\forall t2.t0\leq t2\wedge t2 < t1\rightarrow Tp2Od(t2,\varphi)$;
$(t0,1,t1,1)$:
return $\forall t2.t0\leq t2\wedge t2\leq 3t1\rightarrow Tp2Od(t2,\varphi)$;
end
end

ω-Automaton to LO2

$A_{\exists}(q1,...,qn,\psi I,\psi R,\psi F)$ *(input automaton)*
$\exists q1..qn,\Theta LO2(0,\psi I)\wedge(\forall t.\Theta LO2(t,\psi R))\wedge$
$(\forall .t1\exists t2.t1 < t2\wedge\Theta LO2(t2,\psi F))$

Where ΘLO2(t, Φ) is:

-$\Theta LO2(t,p) := p(t)$ *for variable p*
-$\Theta LO2(t,X\psi) := \Theta LO2(t+1,\psi)$
-$\Theta LO2(t,\neg\psi) := \neg\Theta LO2(t,\psi)$
-$\Theta LO2(t,\varphi\wedge\psi) := \Theta LO2(t,\varphi)\wedge\Theta LO2(t,\psi)$
-$\Theta LO2(t,\varphi\vee\psi) := \Theta LO2(t,\varphi)\vee\Theta LO2(t,\psi)$

LTL to ω-automata

$\phi\langle X\varphi\rangle_x\Leftrightarrow\mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow X\varphi,\phi\langle q\rangle_x)$
$\phi\langle X\varphi\rangle_x\Leftrightarrow$
$\quad\mathcal{A}_{\exists}(\{q_0,q_1\},1,(q_0\leftrightarrow\varphi)\wedge(q_1\leftrightarrow Xq_0),\phi\langle q_1\rangle_x)$
$\phi\langle G\varphi\rangle_x\Leftrightarrow$
$\quad\mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow\varphi\wedge Xq,\phi\langle q\rangle_x\wedge GF[\varphi\rightarrow q])$
$\phi\langle F\varphi\rangle_x\Leftrightarrow$
$\quad\mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow\varphi\wedge Xq,\phi\langle q\rangle_x\wedge GF[q\rightarrow\varphi])$
$\phi\langle[\varphi\ U\ \psi]\rangle_x\Leftrightarrow$
$\quad\mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow\psi\vee\varphi\wedge Xq,\phi\langle q\rangle_x\wedge GF[\varphi\rightarrow q])$
$\phi\langle[\varphi\ \underline{U}\ \psi]\rangle_x\Leftrightarrow$
$\quad\mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow\psi\vee\varphi\wedge Xq,\phi\langle q\rangle_x\wedge GF[q\rightarrow\psi])$
$\phi\langle[\varphi\ B\ \psi]\rangle_x\Leftrightarrow$
$\quad\mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow\neg\psi\wedge(\varphi\vee Xq),\phi\langle q\rangle_x\wedge GF[q\vee\psi])$
$\phi\langle[\varphi\ \underline{B}\ \psi]\rangle_x\Leftrightarrow$
$\quad\mathcal{A}_{\exists}(\{q\},1,q\leftrightarrow\neg\psi\wedge(\varphi\vee Xq),\phi\langle q\rangle_x\wedge GF[q\rightarrow\varphi])$
$\phi\langle\overleftarrow{X}\varphi\rangle_x\Leftrightarrow\mathcal{A}_{\exists}(\{q\},q,Xq\leftrightarrow\varphi,\phi\langle q\rangle_x)$
$\phi\langle\overleftarrow{\underline{X}}\varphi\rangle_x\Leftrightarrow\mathcal{A}_{\exists}(\{q\},\neg q,Xq\leftrightarrow\varphi,\phi\langle q\rangle_x)$
$\phi\langle\overleftarrow{G}\varphi\rangle_x\Leftrightarrow\mathcal{A}_{\exists}(\{q\},q,Xq\leftrightarrow\varphi\wedge q,\phi\langle\varphi\wedge q\rangle_x)$
$\phi\langle\overleftarrow{F}\varphi\rangle_x\Leftrightarrow\mathcal{A}_{\exists}(\{q\},\neg q,Xq\leftrightarrow\varphi\vee q,\phi\langle\varphi\vee q\rangle_x)$
$\phi\langle[\varphi\ \overleftarrow{U}\ \psi]\rangle_x\Leftrightarrow$
$\quad\mathcal{A}_{\exists}(\{q\},q,Xq\leftrightarrow\psi\vee\varphi\wedge q,\phi\langle\psi\vee\varphi\wedge q\rangle_x)$
$\phi\langle[\varphi\ \overleftarrow{\underline{U}}\ \psi]\rangle_x\Leftrightarrow$
$\quad\mathcal{A}_{\exists}(\{q\},\neg q,Xq\leftrightarrow\psi\vee\varphi\wedge q,\phi\langle\psi\vee\varphi\wedge q\rangle_x)$
$\phi\langle[\varphi\ \overleftarrow{B}\ \psi]\rangle_x\Leftrightarrow$
$\quad\mathcal{A}_{\exists}(\{q\},q,Xq\leftrightarrow\neg\psi\wedge(\varphi\vee q),\phi\langle\neg\psi\wedge(\varphi\vee q)\rangle_x)$
$\phi\langle[\varphi\ \overleftarrow{\underline{B}}\ \psi]\rangle_x\Leftrightarrow$
$\quad\mathcal{A}_{\exists}(\{q\},\neg q,Xq\leftrightarrow\neg\psi\wedge(\varphi\vee q),\phi\langle\neg\psi\wedge(\varphi\vee q)\rangle_x)$

CTL to $\mu-Calculus(\Phi_{inf} = \nu y.\Diamond y)$

$EX\varphi = \Diamond(\Phi_{inf}\wedge\varphi)$
$EG\varphi = \nu x.\varphi\wedge\Diamond x$
$EF\varphi = \mu x.\Phi_{inf}\wedge\varphi\vee\Diamond x$
$E[\varphi\underline{U}\psi] = \mu x.(\Phi_{inf}\wedge\psi)\vee\varphi\wedge\Diamond x$
$E[\varphi U\psi] = \nu x.(\Phi_{inf}\wedge\psi)\vee\varphi\wedge\Diamond x$
$E[\varphi\underline{B}\psi] = \mu x.\neg\psi\wedge(\Phi_{inf}\wedge\varphi\vee\Diamond x)$
$E[\varphi B\psi] = \nu x.\neg\psi\wedge(\Phi_{inf}\wedge\varphi\vee\Diamond x)$
$AX\varphi = \Box(\Phi_{inf}\rightarrow\varphi)$
$AG\varphi = \nu x.(\Phi_{inf}\rightarrow\varphi)\wedge\Box x$
$AF\varphi = \mu x.\varphi\vee\Box x$
$A[\varphi\underline{U}\psi] = \mu x.\psi\vee(\Phi_{inf}\rightarrow\varphi)\wedge\Box x$
$A[\varphi U\psi] = \nu x.\psi\vee(\Phi_{inf}\rightarrow\varphi)\wedge\Box x$
$A[\varphi\underline{B}\psi] = \mu x.(\Phi_{inf}\rightarrow\neg\psi)\wedge(\varphi\vee\Box x)$
$A[\varphi B\psi] = \nu x.(\Phi_{inf}\rightarrow\neg\psi)\wedge(\varphi\vee\Box x)$

G and μ-calculus (safety property)

-$[\nu x.\varphi\wedge\Diamond x]_K$
-Contains states s where an infinite path π starts with $\forall t.\pi^{(t)}\in[\varphi]_K$
-φ holds always on π

F and μ-calculus (liveness property)

-$[\mu x.\varphi\vee\Diamond x]_K$
-Contains states s where a (possibly finite) path π starts with $\exists t.\pi^{(t)}\in[\varphi]_K$
-φ holds at least once on π

FG and μ-calculus (persistence property)

-$[\mu y.[\nu x.\varphi\wedge\Diamond x]\vee\Diamond y]_K$
-Contains states s where an infinite path π starts with $\exists t1.\forall t2.\pi^{(t1+t2)}\in[\varphi]_K$
-φ holds after some point on π

GF and μ-calculus (fairness property)

-$[\nu y.[\mu x.(y\wedge\varphi)\vee\Diamond x]]_K$
-Contains states s where an infinite path π starts with
$\forall t1.\exists t2.\pi^{(t1+t2)}\in[\varphi]_K ?????t1+t2ort1+t0?????$
-φ holds infinitely often on π