Propositional Logic Syntactic Sugar

$\varphi \Leftrightarrow \psi := (\neg\varphi \vee \psi) \wedge (\neg\psi \vee \varphi)$ $\quad \varphi \rightarrow \psi := \neg\varphi \vee \psi$

$\varphi \oplus \psi := (\varphi \wedge \neg\psi) \vee (\psi \wedge \neg\varphi)$ $\quad \varphi \bar{\wedge} \psi := \neg(\varphi \wedge \psi)$

$(\alpha \Rightarrow \beta|\gamma) := (\neg\alpha \vee \beta) \wedge (\alpha \vee \gamma)$ $\quad \varphi \bar{\vee} \psi := \neg(\varphi \vee \psi)$

Distributivity: $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$
$a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$
De Morgan: $\neg(a \vee b) \equiv (\neg a \wedge \neg b)$
$\neg(a \wedge b) \equiv (\neg a \vee \neg b)$
CNF: from truth table, take minterms that are 0. Each minterm is built as an OR of the negated variables. E.g., $(0,0,1) \rightarrow (x \vee y \vee \neg z)$.

###**SAT SOLVERS**

Satisfiability, Validity and Equivalence

$\text{SAT}(\varphi) := \neg\text{VALID}(\neg\varphi)$ $\quad \varphi \Leftrightarrow \psi := \text{VALID}(\varphi \leftrightarrow \psi)$

$\text{VALID}(\varphi) := (\varphi \Leftrightarrow 1)$ $\quad \text{SAT}(\varphi) := \neg(\varphi \Leftrightarrow 0)$.

Sequent Calculus:
-*Validity*: start with $\{\} \vdash \phi$; valid iff $\Gamma \cap \Delta \neq \{\}$ FOR ALL leaves.
-*Satisfiability*: start with $\{\phi\} \vdash \{\}$; satisfiable iff $\Gamma \cap \Delta = \{\}$ for AT LEAST ONE leaf.
-Counterexample/sat variable assignment: var is true, if $x \in \Gamma$; false otherwise; "don't care", if variable doesn't appear.

OPER.	LEFT	RIGHT
NOT	$\frac{\neg\phi, \Gamma \vdash \Delta}{\Gamma \vdash \phi, \Delta}$	$\frac{\Gamma \vdash \neg\phi, \Delta}{\phi, \Gamma \vdash \Delta}$
AND	$\frac{\phi \wedge \psi, \Gamma \vdash \Delta}{\phi, \psi, \Gamma \vdash \Delta}$	$\frac{\Gamma \vdash \phi \wedge \psi, \Delta}{\Gamma \vdash \phi, \Delta \quad \Gamma \vdash \psi, \Delta}$
OR	$\frac{\phi \vee \psi, \Gamma \vdash \Delta}{\phi, \Gamma \vdash \Delta \quad \psi, \Gamma \vdash \Delta}$	$\frac{\Gamma \vdash \phi \vee \psi, \Delta}{\Gamma \vdash \phi, \psi, \Delta}$

Resolution Calculus $\frac{\{\neg x\} \cup C_1 \quad \{x\} \cup C_2}{C_1 \cup C_2}$

To prove unsatisfiability of given clauses in CNF: If we reach $\{\}$, the formula is unsatisfiable. E.g., $\{\{a\}, \{\neg a, b\}, \{\neg b\}\}$, we get: $\{a\} + \{\neg a, b\} \rightarrow \{b\}; \{b\} + \{\neg b\} \rightarrow \{\}$ (unsatisfiable). To prove validity, prove UNSAT of negated formula.

Davis Putnam Procedure - proves SAT; To prove validity: prove unsatisfiability of negated formula.
(1) Compute Linear Clause Form
(Don't forget to create the last clause $\{x_n\}$)
(2) Last variable has to be 1 (true) → find implied variables.
(3) For remaining variables: assume values and compute newly implied variables.
(4) If contradiction reached: backtrack.

Linear Clause Forms (Computes CNF) - Bottom up in the syntax tree: convert "operators and variables" into new variable. E.g., $\neg a \vee b$ becomes $x_1 \leftrightarrow \neg a; x_2 \leftrightarrow x_1 \vee b$. Use rules below to find CNF.

$x \leftrightarrow \neg y \Leftrightarrow (\neg x \vee \neg y) \wedge (x \vee y)$
$x \leftrightarrow y_1 \wedge y_2 \Leftrightarrow$
$\quad (\neg x \vee y_1) \wedge (\neg x \vee y_2) \wedge (x \vee \neg y_1 \vee \neg y_2)$
$x \leftrightarrow y_1 \vee y_2 \Leftrightarrow$
$\quad (\neg x \vee y_1 \vee y_2) \wedge (x \vee \neg y_1) \wedge (x \vee \neg y_2)$
$x \leftrightarrow y_1 \rightarrow y_2 \Leftrightarrow$
$\quad (x \vee y_1) \wedge (x \vee \neg y_1 \vee \neg y_2) \wedge (\neg x \vee \neg y_1 \vee y_2)$
$x \leftrightarrow (y_1 \leftrightarrow y_2) \Leftrightarrow (x \vee y_1 \vee y_2) \wedge (x \vee \neg y_1 \vee \neg y_2) \wedge$
$\quad (\neg x \vee y_1 \vee \neg y_2) \wedge (\neg x \vee \neg y_1 \vee y_2)$
$x \leftrightarrow y_1 \oplus y_2 \Leftrightarrow (x \vee \neg y_1 \vee y_2) \wedge (x \vee y_1 \vee \neg y_2) \wedge$
$\quad (\neg x \vee y_1 \vee y_2) \wedge (\neg x \vee \neg y_1 \vee \neg y_2)$

```
Compose(int x, BddNode ψ, α)
 int m; BddNode h, l;
 if x>label(ψ) then
  return ψ;
 elseif x=label(ψ) then
  return ITE(α,high(ψ),
        low(ψ));
 else
  m=max{label(ψ),label(α)}
  (α0,α1):=Ops(α, m);
  (ψ0,ψ1):=Ops(ψ, m);
  h:=Compose(x,ψ1,α1);
  l:=Compose(x,ψ0,α0);
  return CreateNode(m,h,l)
 endif;end
```

```
ITE(BddNode i, j, k)
 int m; BddNode h, l;
 if i = 0 then return k
 elseif i=1 then
  return j
 elseif j=k then
  return k
 else
  m = max{label(i),
      label(j),label(k)}
  (i0,i1):=Ops(i,m);
  (j0,j1):=Ops(j,m);
  (k0,k1):=Ops(k,m);
  l:=ITE(i0,j0,k0);
  h:=ITE(i1, j1, k1);
  return CreateNode(m,h,l)
 end;end
```

```
Constrain(Φ, β)
 if β=0 then
  ret 0
 elseif Φ ∈ {0,1}(β = 1)
     ret Φ
 else
  m=max{label(β),label(Φ)}
    (Φ0,Φ1):=Ops(Φ,m);
    (β0,β1):=Ops(β,m);
    if β0=0
  ret Constrain(Φ1,β1)
 elseif β1=0 then
  ret Constrain(Φ0,β0)
 else
  l:=Constrain(Φ0,β0);
  h:=Constrain(Φ1,β1);
  ret CreateNode(m,h,l)
endif;endif;end
```

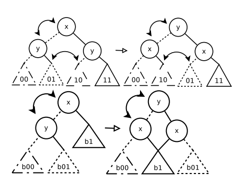
```
Apply(⊙, Bddnode a, b)
 int m; BddNode h, l;
 if isLeaf(a)&isLeaf(b)
      then
  return Eval(⊙,label(a),
       label(b));
 else
  m=max{label(a),label(b)}
  (a0,a1):=Ops(a,m);
  (b0,b1):=Ops(b,m);
  h:=Apply(⊙,a1,b1);
  l:=Apply(⊙,a0,b0);
  return CreateNode(m,h,l)
 end;
end
```

```
Restrict(Φ, β)
 if β=0
  return 0
 elseif
       Φ ∈ {0,1} ∨ (β = 1)

  return Φ
 else
  m=max{label(β),label(Φ)}
  (Φ0,Φ1):=Ops(Φ,m);
  (β0,β1):=Ops(β,m)
 if β0=0
  return Restrict(Φ1,β1)
 elseif β1=0
  return Restrict(Φ0,β0)
 elseif m=label(Φ)
  return CreateNode(m,
   Restrict (Φ1,β1),
   Restrict (Φ0,β0))
 else
  return Restrict(Φ,
  Apply(∨,β0,β1))
 endif;endif;end
------------------------
Ops(v,m)
 x:=label(v);
 if m=degree(x)
  return (low(v),high(v))
 else return(v, v)
end;end
```

FDD: Positive Davio Decomposition
$\varphi = [\varphi]_x^0 \oplus x \wedge (\partial\varphi/\partial x)$
$(\partial\varphi/\partial x) := [\varphi]_x^0 \oplus [\varphi]_x^1$

ZDD: If positive cofactor = 0, redirect edge to negative cofactor.
If variable not in the formula, add with both edges pointing to same node.



Local Model Checking

$\frac{s \vdash_\Phi \varphi \wedge \psi}{\{s \vdash_\Phi \varphi\} \quad \{s \vdash_\Phi \psi\}} \wedge$	$\frac{s \vdash_\Phi \varphi \vee \psi}{\{s \vdash_\Phi \varphi\} \quad \{s \vdash_\Phi \psi\}} \vee$
$\frac{s \vdash_\Phi \Box\varphi}{\{s_1 \vdash_\Phi \varphi\} \ldots \{s_n \vdash_\Phi \varphi\}} \wedge$	$\frac{s \vdash_\Phi \Diamond\varphi}{\{s_1 \vdash_\Phi \varphi\} \ldots \{s_n \vdash_\Phi \varphi\}} \vee$
$\frac{s \vdash_\Phi \overleftarrow{\Box}\varphi}{\{s'_1 \vdash_\Phi \varphi\} \ldots \{s'_n \vdash_\Phi \varphi\}} \wedge$	$\frac{s \vdash_\Phi \overleftarrow{\Diamond}\varphi}{\{s'_1 \vdash_\Phi \varphi\} \ldots \{s'_n \vdash_\Phi \varphi\}} \vee$

$\frac{s \vdash_\Phi \mu x.\varphi}{s \vdash_\Phi \varphi}$	$\frac{s \vdash_\Phi \nu x.\varphi}{s \vdash_\Phi \varphi}$	$\frac{s \vdash_\Phi x}{s \vdash_\Phi \mathfrak{D}_\Phi(x)}$	$\mathfrak{D}_\Phi$(replace w. initial form.)

$\{s_1 \ldots s_n\} = suc_\exists^\mathcal{R}(s)$ and $\{s'_1 \ldots s'_n\} = pre_\exists^\mathcal{R}(s)$

Approximations and Ranks

If (s,$\mu x.\varphi$) repeats→return 1	$apx_0(\mu x.\varphi) := 0$
If (s,$\nu x.\varphi$) repeats→return 0	$apx_0(\nu x.\varphi) := 1$
$apx_{n+1}(\mu x.\varphi) := [\varphi]_x^{apxn(\mu x.\varphi)}$	
$apx_{n+1}(\nu x.\varphi) := [\varphi]_x^{apxn(\nu x.\varphi)}$	

Tarski-Knaster Theorem: μ := starts $\perp \rightarrow$ least fixpoint ♠ ν := starts $\top \rightarrow$ greatest fixpoint

Quantif. $\exists x.\varphi := [\varphi]_x^1 \vee [\varphi]_x^0$ ♣ $\forall x.\varphi := [\varphi]_x^1 \wedge [\varphi]_x^0$

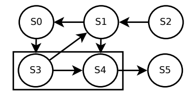
Predecessor and Successor

$\Diamond := pre_\exists^\mathcal{R}(Q) := \exists x'_1, \ldots, x'_n.\varphi_\mathcal{R} \wedge [\varphi_Q]_{x_1,\ldots,x_n}^{x'_1,\ldots,x'_n}$

$\overleftarrow{\Diamond} := suc_\exists^\mathcal{R}(Q) := [\exists x_1, \ldots, x_n.\varphi_\mathcal{R} \wedge \varphi_Q]_{x'_1,\ldots,x'_n}^{x_1,\ldots,x_n}$

$\Box = pre_\forall^\mathcal{R}(Q) := \forall x'_1, \ldots, x'_n.\varphi_\mathcal{R} \rightarrow [\varphi_Q]_{x_1,\ldots,x_n}^{x'_1,\ldots,x'_n}$

$\overleftarrow{\Box} := suc_\forall^\mathcal{R}(Q) := [\forall x_1, \ldots, x_n.\varphi_\mathcal{R} \rightarrow \varphi_Q]_{x'_1,\ldots,x'_n}^{x_1,\ldots,x_n}$



Example: $\Box/\overleftarrow{\Box}$

$pre_\forall^\mathcal{R}(\{S3, S4\}) = \{S0, S5\}$

$suc_\forall^\mathcal{R}(\{S3, S4\}) = \{S2, S5\}$

```
pre∀R(Q = {S1,...,Sn})
for each node n in K:
 if(n points to a node
       that is not in Q)
   n ∉ pre∀R(Q)
 else
   n ∈ pre∀R(Q)
```

```
suc∀R(Q = {S1,...,Sn})
for each node n in K:
 if(n is pointed by a node
       that is not in Q)
   n ∉ suc∀R(Q)
 else
   n ∈ suc∀R(Q)
```

###**AUTOMATA**

Automata types: G→Safety; F→Liveness; FG→Persistence/Co-Buchi; GF→Fairness/Buchi.

Automaton Determinization

$NDet_G \rightarrow Det_G$: 1.Remove all states/edges that do not satisfy acceptance condition; 2.Use Subset construction (Rabin-Scott); 3.Acceptance condition will be the states where {} is never reached.
$\{NDet_F$(partial) or $NDet_{prefix}\} \rightarrow Det_{FG}$: Breakpoint Construction.
$NDet_F$ (total)$\rightarrow Det_F$: Subset Construction.
$NDet_{FG} \rightarrow Det_{FG}$: Breakpoint Construction.
$NDet_{GF} \rightarrow \{Det_{Rabin}$ or $Det_{Streett}\}$: Safra Algorithm.

Boolean Operations on ω-Automata

Complement

$$\neg\mathcal{A}_\forall(Q, \mathcal{I}, \mathcal{R}, \mathcal{F}) = \mathcal{A}_\exists(Q, \mathcal{I}, \mathcal{R}, \neg\mathcal{F})$$
$$\neg\mathcal{A}_\exists(Q, \mathcal{I}, \mathcal{R}, \mathcal{F}) = \mathcal{A}_\forall(Q, \mathcal{I}, \mathcal{R}, \neg\mathcal{F})$$

Conjunction

$$(\mathcal{A}_\exists(Q_1, \mathcal{I}_1, \mathcal{R}_1, \mathcal{F}_1) \wedge \mathcal{A}_\exists(Q_2, \mathcal{I}_2, \mathcal{R}_2, \mathcal{F}_2)) =$$
$$\mathcal{A}_\exists(Q_1 \cup Q_2, \mathcal{I}_1 \wedge \mathcal{I}_2, \mathcal{R}_1 \wedge \mathcal{R}_2, \mathcal{F}_1 \wedge \mathcal{F}_2)$$

Disjunction

$$(\mathcal{A}_\exists(Q_1, \mathcal{I}_1, \mathcal{R}_1, \mathcal{F}_1) \vee \mathcal{A}_\exists(Q_2, \mathcal{I}_2, \mathcal{R}_2, \mathcal{F}_2)) =$$
$$\mathcal{A}_\exists \left(\begin{array}{c} Q_1 \cup Q_2 \cup \{q\}, \\ (\neg q \wedge \mathcal{I}_1) \vee (q \wedge \mathcal{I}_2), \\ (\neg q \wedge \mathcal{R}_1 \wedge \neg q') \vee (q \wedge \mathcal{R}_2 \wedge q'), \\ (\neg q \wedge \mathcal{F}_1) \vee (q \wedge \mathcal{F}_2) \end{array} \right)$$

If both automata are totally defined,

$$(\mathcal{A}_\exists(Q_1, \mathcal{I}_1, \mathcal{R}_1, \mathcal{F}_1) \vee \mathcal{A}_\exists(Q_2, \mathcal{I}_2, \mathcal{R}_2, \mathcal{F}_2)) =$$
$$\mathcal{A}_\exists(Q_1 \cup Q_2, \mathcal{I}_1 \wedge \mathcal{I}_2, \mathcal{R}_1 \wedge \mathcal{R}_2, \mathcal{F}_1 \vee \mathcal{F}_2)$$

Eliminate Nesting - Acceptance condition **must** be an automata of the same type

$$\mathcal{A}_\exists(Q^1, \mathcal{I}_1^1, \mathcal{R}_1^1, \mathcal{A}_\exists(Q^2, \mathcal{I}_1^2, \mathcal{R}_1^2, \mathcal{F}_1))$$
$$= \mathcal{A}_\exists(Q^1 \cup Q^2, \mathcal{I}_1^1 \wedge \mathcal{I}_1^2, \mathcal{R}_1^1 \wedge \mathcal{R}_1^2, \mathcal{F}_1))$$

Boolean Operations of G

(1)$\neg G\varphi = F\neg\varphi$ $\quad$ (2)$G\varphi \wedge G\psi = G[\varphi \wedge \psi]$
(3)$G\varphi \vee G\psi = \mathcal{A}_\exists(\{p, q\}, p \wedge q,$
$\quad [p' \leftrightarrow p \wedge \varphi] \wedge [q' \leftrightarrow q \wedge \psi], G[p \vee q])$

Boolean Operations of F

(1)$\neg F\varphi = G\neg\varphi$ $\quad$ (2)$F\varphi \vee F\psi = F[\varphi \vee \psi]$
(3)$F\varphi \wedge F\psi = \mathcal{A}_\exists(\{p, q\}, \neg p \wedge \neg q,$
$\quad [p' \leftrightarrow p \vee \varphi] \wedge [q' \leftrightarrow q \vee \psi], F[p \wedge q])$

Boolean Operations of FG

(1)$\neg FG\varphi = GF\neg\varphi$ $\quad$ (2)$FG\varphi \wedge FG\psi = FG[\varphi \wedge \psi]$
(3)$FG\varphi \vee FG\psi = \mathcal{A}_\exists(\{q\}, \neg q, q' \leftrightarrow (q \Rightarrow \psi|\neg\varphi),$
$\quad FG[\neg q \vee \psi])$

Boolean Operations of GF

(1)$\neg GF\varphi = FG\neg\varphi$ $\quad$ (2)$GF\varphi \vee GF\psi = GF[\varphi \vee \psi]$
(3)$GF\varphi \wedge GF\psi = \mathcal{A}_\exists(\{q\}, \neg q, q' \leftrightarrow (q \Rightarrow \neg\psi|\varphi),$
$\quad GF[q \wedge \psi])$

Transformation of Acceptance Conditions

Reduction of G

$G\varphi = \mathcal{A}_\exists(\{q\}, q, \varphi \wedge q \wedge q', Fq))$
$G\varphi = \mathcal{A}_\exists(\{q\}, q, q' \leftrightarrow q \wedge \varphi, FGq)$
$G\varphi = \mathcal{A}_\exists(\{q\}, q, q' \leftrightarrow q \wedge \varphi, GFq)$

Reduction of F

$F\varphi$ can **not** be expressed by $NDet_G$
$F\varphi = \mathcal{A}_\exists(\{q\}, \neg q, q' \leftrightarrow q \vee \varphi, FGq)$
$F\varphi = \mathcal{A}_\exists(\{q\}, \neg q, q' \leftrightarrow q \vee \varphi, GFq)$

Reduction of FG

$FG\varphi$ can **not** be expressed by $NDet_G$
$FG\varphi = \mathcal{A}_\exists(\{q\}, \neg q, q \rightarrow \varphi \wedge q', Fq)$

$$FG\varphi = \mathcal{A}_\exists \left(\begin{array}{c} \{p, q\}, \qquad \neg p \wedge \neg q, \\ \left[\begin{array}{c} (p \rightarrow p') \wedge (p' \rightarrow p \vee \neg q) \wedge \\ (q' \leftrightarrow (p \wedge \neg q \vee \neg\varphi) \vee (p \wedge q)) \end{array} \right], \\ G\neg q \wedge Fp \end{array} \right)$$

$$FG\varphi = \mathcal{A}_\exists \left(\begin{array}{c} \{p, q\}, \qquad \neg p \wedge \neg q, \\ \left[\begin{array}{c} (p \rightarrow p') \wedge (p' \rightarrow p \vee \neg q) \wedge \\ (q' \leftrightarrow (p \wedge \neg q \vee \neg\varphi) \vee (p \wedge q)) \end{array} \right], \\ GF[p \wedge \neg q] \end{array} \right)$$

###**TEMPORAL LOGICS**

(S1)Pure LTL: AFGa
(S2)LTL + CTL: AFa
(S3)Pure CTL: AGEFa
(S4)CTL*: AFGa ∨ AGEFa

Remarks *Beware of Finite Paths*
E and A quantify over infinite paths.
Aφ holds on every state that has no infinite path;
Eφ is false on every state that has no infinite path;
A0 holds on states with only finite paths;
E1 is false on state with only finite paths;
$\Box$0 holds on states with no successor states;
$\Diamond$1 holds on states with successor states.

$F\varphi = \varphi \vee XF\varphi$ $\quad G\varphi = \varphi \wedge XG\varphi$
$[\varphi\ U\ \psi] = \psi \vee (\varphi \wedge X[\varphi\ U\ \psi])$
$[\varphi\ B\ \psi] = \neg\psi \wedge (\varphi \vee X[\varphi\ B\ \psi])$
$[\varphi\ W\ \psi] = (\psi \wedge \varphi) \vee (\neg\psi \wedge X[\varphi\ W\ \psi])$

LTL Syntactic Sugar: analog for past operators

$G\varphi = \neg[1\ \underline{U}\ (\neg\varphi)]$ $\quad F\varphi = [1\ \underline{U}\ \varphi]$
$[\varphi\ W\ \psi] = \neg[(\neg\varphi \vee \neg\psi)\ \underline{U}\ (\neg\varphi \wedge \psi)]$
$[\varphi\ \underline{W}\ \psi] = [(\neg\psi)\ \underline{U}\ (\varphi \wedge \psi)]$ *($\neg\psi$ holds until $\varphi \wedge \psi$)*
$[\varphi\ B\ \psi] = \neg[(\neg\varphi)\ \underline{U}\ \psi)]$
$[\varphi\ \underline{B}\ \psi] = [(\neg\psi)\ \underline{U}\ (\varphi \wedge \neg\psi)]$*($\psi$ can't hold when φ holds)*
$[\varphi\ U\ \psi] = \neg[(\neg\psi)\ \underline{U}\ (\neg\varphi \wedge \neg\psi)]$
$[\varphi\ U\ \psi] = [\varphi\ \underline{U}\ \psi] \vee G\varphi$
$[\varphi\ \underline{U}\ \psi] = \neg[(\neg\psi)\ U\ (\neg\varphi \wedge \neg\psi)]$
$[\varphi\ \underline{U}\ \psi] = \neg[(\neg\psi)\ W\ (\varphi \rightarrow \psi)]$
$[\varphi\ \underline{U}\ \psi] = [\psi\ \underline{W}\ (\varphi \rightarrow \psi)]$
$[\varphi\ \underline{U}\ \psi] = \neg[(\neg\varphi)\ B\ \psi]$*($\varphi$ doesn't matter when ψ holds)*
$[\varphi\ \underline{U}\ \psi] = [\psi\ \underline{B}\ (\neg\varphi \wedge \neg\psi]$

CTL Syntactic Sugar: analog for past operators

Existential Operators

$EF\varphi = E[1\ \underline{U}\ \varphi]$
$EG\varphi = E[\varphi\ U\ 0]$
$E[\varphi\ U\ \psi] = E[\varphi\ \underline{U}\ \psi] \vee EG\varphi$
$E[\varphi\ B\ \psi] = E[(\neg\psi)\ \underline{U}\ (\varphi \wedge \neg\psi)] \vee EG\neg\psi$
$E[\varphi\ B\ \psi] = E[(\neg\psi)\ U\ (\varphi \wedge \neg\psi)]$
$E[\varphi\ \underline{B}\ \psi] = E[(\neg\psi)\ \underline{U}\ (\varphi \wedge \neg\psi)]$
$E[\varphi\ \underline{B}\ \psi] = E[(\neg\psi\ \underline{U}\ (\varphi \wedge \neg\psi)]$
$E[\varphi\ W\ \psi] = E[(\neg\psi)\ \underline{U}\ (\varphi \wedge \psi)] \vee EG\neg\psi$
$E[\varphi\ W\ \psi] = E[(\neg\psi)\ U\ (\varphi \wedge \psi)]$
$E[\varphi\ \underline{W}\ \psi] = E[(\neg\psi)\ \underline{U}\ (\varphi \wedge \psi)]$

Universal Operators

$AX\varphi = \neg EX\neg\varphi$
$AG\varphi = \neg E[1\ \underline{U}\ \neg\varphi]$
$AF\varphi = \neg EG\neg\varphi$
$AF\varphi = \neg E[(\neg\varphi)\ U\ 0]$
$A[\varphi\ U\ \psi] = \neg E[(\neg\psi)\ \underline{U}\ (\neg\varphi \wedge \neg\psi)]$
$A[\varphi\ \underline{U}\ \psi] = \neg E[(\neg\psi)\ \underline{U}\ (\neg\varphi \wedge \neg\psi)] \wedge \neg EG\neg\psi$
$A[\varphi\ \underline{U}\ \psi] = \neg E[(\neg\psi)\ U\ (\neg\varphi \wedge \neg\psi)]$
$A[\varphi\ B\ \psi] = \neg E[(\neg\varphi)\ \underline{U}\ \psi]$
$A[\varphi\ \underline{B}\ \psi] = \neg E[(\neg\varphi)\ U\ \psi]$
$A[\varphi\ \underline{B}\ \psi] = \neg E[(\neg\varphi \vee \psi)\ \underline{U}\ \psi] \wedge \neg EG(\neg\varphi \vee \psi)$
$A[\varphi\ W\ \psi] = \neg E[(\neg\psi)\ \underline{U}\ (\neg\varphi \wedge \psi)]$
$A[\varphi\ \underline{W}\ \psi] = \neg E[(\neg\psi)\ \underline{U}\ (\neg\varphi \wedge \psi)] \wedge \neg EG\neg\psi$
$A[\varphi\ \underline{W}\ \psi] = \neg E[(\neg\psi)\ U\ (\neg\varphi \wedge \psi)]$

CTL* to CTL - Existential Operators

$EX\varphi = EXE\varphi$
$EF\varphi = EFE\varphi$ $\qquad EFG\varphi \equiv EFEG\varphi$
$E[\varphi\ W\ \psi] = E[(E\varphi)\ W\ \psi]$
$E[\varphi\ \underline{W}\ \psi] = E[(E\varphi)\ \underline{W}\ \psi]$
$E[\psi\ U\ \varphi] = E[\psi\ U\ E(\varphi)]$
$E[\psi\ \underline{U}\ \varphi] = E[\psi\ \underline{U}\ E(\varphi)]$
$E[\varphi\ B\ \psi] = E[(E\varphi)\ B\ \psi]$
$E[\varphi\ \underline{B}\ \psi] = E[(E\varphi)\ \underline{B}\ \psi]$

obs. $EGF\varphi \neq EGEF\varphi \rightarrow$ can't be converted

CTL* to CTL - Universal Operators

$AX\varphi = AXA\varphi$
$AG\varphi = AGA\varphi$
$A[\varphi\ W\ \psi] = A[(A\varphi)\ W\ \psi]$
$A[\varphi\ \underline{W}\ \psi] = A[(A\varphi)\ \underline{W}\ \psi]$
$A[\varphi\ U\ \psi] = A[A(\varphi)\ U\ \psi]$
$A[\varphi\ \underline{U}\ \psi] = A[A(\varphi)\ \underline{U}\ \psi]$
$A[\psi\ B\ \varphi] = A[\psi\ B\ (E(\varphi)]$
$A[\psi\ \underline{B}\ \varphi] = A[\psi\ \underline{B}\ (E(\varphi)]$

Weak Equivalences

$*[\varphi U \psi] := [\varphi \underline{U} \psi] \vee G\varphi$ $\qquad * [\varphi B \psi] := [\varphi \underline{B} \psi] \vee G\neg\psi$
*same to past version $\qquad [\varphi W \psi] := \neg[(\neg\varphi)\underline{W}\psi]$
$\overleftarrow{X}\varphi := \neg\underline{\overleftarrow{X}}\neg\varphi$ *(at t0 : weak true. strong false)*

Negation Normal Form

$\neg(\varphi \wedge \psi) = \neg\varphi \vee \neg\psi$	$\neg(\varphi \vee \psi) = \neg\varphi \wedge \neg\psi$
$\neg\neg\varphi = \varphi$	$\neg X\varphi = X\neg\varphi$
$\neg G\varphi = F\neg\varphi$	$\neg F\varphi = G\neg\varphi$
$\neg[\varphi\ U\ \psi] = [(\neg\varphi)\ \underline{B}\ \psi]$	$\neg[\varphi\ \underline{U}\ \psi] = [(\neg\varphi)\ B\ \psi]$
$\neg[\varphi\ B\ \psi] = [(\neg\varphi)\ \underline{U}\ \psi]$	$\neg[\varphi\ \underline{B}\ \psi] = [(\neg\varphi)\ U\ \psi]$
$\neg A\varphi = E\neg\varphi$	$\neg E\varphi = A\neg\varphi$
$\neg\overleftarrow{X}\varphi = \underline{\overleftarrow{X}}\neg\varphi$	$\neg\underline{\overleftarrow{X}}\varphi = \overleftarrow{X}\neg\varphi$
$\neg\overleftarrow{G}\varphi = \overleftarrow{F}\neg\varphi$	$\neg\overleftarrow{F}\varphi = \overleftarrow{G}\neg\varphi$
$\neg[\varphi\ \overleftarrow{U}\ \psi] = [(\neg\varphi)\ \underline{\overleftarrow{B}}\ \psi]$	$\neg[\varphi\ \underline{\overleftarrow{U}}\ \psi] = [(\neg\varphi)\ \overleftarrow{B}\ \psi]$
$\neg[\varphi\ \overleftarrow{B}\ \psi] = [(\neg\varphi)\ \underline{\overleftarrow{U}}\ \psi]$	$\neg[\varphi\ \underline{\overleftarrow{B}}\ \psi] = [(\neg\varphi)\ \overleftarrow{U}\ \psi]$

Equivalences and Tips

$[\varphi\underline{B}\psi] \equiv \psi$ *can't hold when* φ *hold*
$[\varphi U\psi] \equiv [\varphi\underline{U}\psi] \vee G\varphi$
$[a\underline{U}Fb] \equiv Fb$
$F[a\underline{U}b] \equiv Fb \equiv [Fa\underline{U}Fb]$
$[\varphi B\psi] \equiv [\varphi\underline{B}\psi] \vee G\neg\psi$
$F[a\underline{B}b] \equiv F[a \wedge \neg b]$
$[\varphi W\psi] \equiv \neg[\neg\varphi\underline{W}\psi]$
$AEA \equiv A$ $\qquad \bullet\ GFX \equiv GXF$
$FF\varphi \equiv F\varphi$ $\qquad \bullet\ GG\varphi \equiv G\varphi$
$GF\varphi \equiv XGF\varphi \equiv FGF\varphi \equiv GGF\varphi \equiv GFGF\varphi \equiv FGGF\varphi$
$FG\varphi \equiv XFG\varphi \equiv FFG\varphi \equiv GFG\varphi \equiv GFFG\varphi \equiv FGFG\varphi$
$GF(x \vee y) \equiv GFx \vee GFy$
$E(\varphi \wedge \psi) \equiv E\varphi \wedge E\psi$ *(in general)*
$E(\varphi \vee \psi) \equiv E\varphi \vee E\psi$
$AG(\varphi \wedge \psi) \equiv AG\varphi \wedge AG\psi$

Eliminate boolean op. after path quantify

$$[\varphi_1\ \underline{U}\ \psi_1] \wedge [\varphi_2\ \underline{U}\ \psi_2] = \left[(\varphi_1 \wedge \varphi_2)\ \underline{U}\ \begin{pmatrix}\psi_1 \wedge [\varphi_2\ \underline{U}\psi_2] \vee \\ \psi_2 \wedge [\varphi_1\ \underline{U}\psi_1]\end{pmatrix}\right]$$

$$[\varphi_1\ \underline{U}\ \psi_1] \wedge [\varphi_2\ U\ \psi_2] = \left[(\varphi_1 \wedge \varphi_2)\ \underline{U}\ \begin{pmatrix}\psi_1 \wedge [\varphi_2\ U\psi_2] \vee \\ \psi_2 \wedge [\varphi_1\ \underline{U}\psi_1]\end{pmatrix}\right]$$

$$[\varphi_1\ U\ \psi_1] \wedge [\varphi_2\ U\ \psi_2] = \left[(\varphi_1 \wedge \varphi_2)\ \underline{U}\ \begin{pmatrix}\psi_1 \wedge [\varphi_2\ U\psi_2] \vee \\ \psi_2 \wedge [\varphi_1\ U\psi_1]\end{pmatrix}\right]$$

MONADIC PREDICATE

S1S

First order terms are defined as follows:
$-0 \in Term^{S1S}_{\sum}$
$-t \in V_{\sum} | typ_{\sum}(t) = \mathbb{N} \subseteq Term^{S1S}_{\sum}$
$-SUC(\tau) \in Term^{S1S}_{\sum} if \tau \in Term^{S1S}_{\sum}$

Formulas ζ_{S1S} are defined as:
$-p^{(t)} \in L_{S1S}$ (predicate p at time t)
$-\neg\varphi, \varphi \wedge \psi \in L_{S1S}$
$-\exists t.\varphi \in L_{S1S}$
$-\exists p.\varphi \in L_{S1S}$
where:
$-\tau \in Term^{S1S}_{\sum}$
$-\varphi, \psi \in \zeta_{S1S}$
$-t \in V_{\sum}$ *with* $typ_{\sum}(t) = \mathbb{N}$
$-p \in V_{\sum}$ *with* $typ_{\sum}(p) = \mathbb{N} \rightarrow \mathbb{B}$

LO2

first order terms are defined as:
$-t \in V_{\sum} | typ_{\sum}(t) = \mathbb{N} \subseteq Term^{LO2}_{\sum}$

formulas LO2 are defined as:
$-t1 < t2 \in L_{LO2}$
$-p^{(t)} \in L_{LO2}$
$-\neg\varphi, \varphi \wedge \psi \in L_{LO2}$
$-\exists t.\varphi \in L_{LO2}$
$-\exists p.\varphi \in L_{LO2}$
where:
$-t, t_1, t_2\tau \in V_{\sum}$ *with* $typ_{\sum}(_t) = typ_{\sum}(_{t1}) = typ_{\sum}(_{t2}) = \mathbb{N}$
$-\varphi, \psi \in \zeta_{LO2}$
$-t \in V_{\sum}$ *with* $typ_{\sum}(t) = \mathbb{N}$
$-p \in V_{\sum}$ *with* $typ_{\sum}(p) = \mathbb{N} \rightarrow \mathbb{B}$

LO2' Consider the following set $\zeta_{LO2'}$ of formulas:
$-Subset(p,q), Sing(p), and PSUC(p,q) belong to \zeta_{LO2'}$
$-\neg\varphi, \varphi \wedge \psi$
$-\exists p.\varphi$
where $-\varphi, \psi \in \zeta_{LO2'}$
$-p \in V_{\sum}$ *with* $typ_{\sum}(p) = \mathbb{N} \rightarrow \mathbb{B}$
$\zeta_{LO2'}$ *hasnonumericvariables*
numeric variable t is replaced by a singleton set p_t
$\zeta_{LO2'}$ is as expressive as LO2 and S1S

TRANSLATIONS

CTL* Modelchecking to LTL model checking

Let's φ_i be a pure path formula (without path quantifiers), Ψ be a propositional formula, abbreviate subformulas $E\varphi$ and $A\psi$ working bottom-up the syntax tree to obtain the following normal form: $\phi = $ let $\begin{bmatrix} x_1 = A\varphi_1 \\ \vdots \\ x_n = A\varphi_n \end{bmatrix}$ in Ψ end

Use LTL model checking to compute $Q_i := [\![A\varphi_i]\!]_{\mathcal{K}_{i-1}}$, where $\mathcal{K}_0 := \mathcal{K}$ and $\mathcal{K}_{i+1}$ is obtained from $\mathcal{K}_i$ by labelling the states Q_i with x_i. Finally compute $[\![\Psi]\!]_{\mathcal{K}_n}$

function LO2_S1S(Φ)
  case Φ **of**
    $t1 < t2$: **return** $\exists p.[\forall t.p^{(t)} \rightarrow p^{(SUC(t))}] \wedge \neg p^{(t1)} \wedge p^{(t2)}$:
    $p^{(t)}$: **return** $p^{(t)}$;
    $\neg\varphi$: **return** $\neg LO2\_S1S(\varphi)$;
    $\varphi \wedge \psi$: **return** $LO2\_S1S(\varphi) \wedge LO2\_S1S(\psi)$;
    $\exists t.\varphi$: **return** $\exists t.LO2\_S1S(\varphi)$;
    $\exists p.\varphi$: **return** $\exists p.LO2\_S1S(\varphi)$;
  end
end

function S1S_LO2(Φ)
  case Φ **of**
    $p^{(n)}$:
return $\exists t0...tn.p^{(tn)} \wedge zero(t0) \wedge \bigwedge_{i=0}^{n-1} succ(ti, ti{+}1)$;
    $p^{(t0+n)}$:
return $\exists t1...tn.p^{(tn)} \wedge \bigwedge_{i=0}^{n-1} succ(ti, ti+1)$;
    $\neg\varphi$: **return** $\neg S1S\_LO2(\varphi)$;
    $\varphi \wedge \psi$: **return** $S1S\_LO2(\varphi) \wedge S1S\_LO2(\psi)$;
    $\exists t.\varphi$: **return** $\exists t.S1S\_LO2(\varphi)$;
    $\exists p.\varphi$: **return** $\exists p.S1S\_LO2(\varphi)$;
  end
end

function Tp2Od(t0, Φ) *temporal to LO1*
  case Φ **of**
    $is\_var(\Phi) : \Psi^{(t0)}$;
    $\neg\varphi$: **return** $\neg Tp2Od(\varphi)$;
    $\varphi \wedge \psi$: **return** $Tp2Od(\varphi) \wedge Tp2Od(\psi)$;
    $\varphi \vee \psi$: **return** $Tp2Od(\varphi) \vee Tp2Od(\psi)$;
    $X\varphi : \Psi := \exists t1.(t0 < t1) \wedge$
        $\forall t2.t0 < t2 \rightarrow t1 \leq t2) \wedge Tp2Od(t1, \varphi)$;
    $[\varphi\underline{U}\psi] : \Psi := \exists t1.t0 \leq t1 \wedge Tp2Od(t1, \psi) \wedge$
        $interval((t0, 1, t1, 0), \varphi)$;
    $[\varphi B\psi] : \Psi := \forall t1.t0 \leq t1 \wedge$
        $interval((t0, 1, t1, 0), \neg\varphi) \rightarrow Tp2Od(t1, \neg\psi)$;
    $\overleftarrow{X}\varphi : \Psi := \forall t1.(t1 < t0) \wedge$
        $(\forall t2.t2 < t0 \rightarrow t2 \leq t1) \rightarrow Tp2Od(t1, \varphi)$;
    $\underline{\overleftarrow{X}}\varphi : \Psi := \exists t1.(t1 < t0) \wedge$
        $(\forall t2.t2 < t0 \rightarrow t2 \leq t1) \wedge Tp2Od(t1, \varphi)$;
    $[\varphi\underline{\overleftarrow{U}}\psi] : \Psi := \exists t1.t1 \leq t0 \wedge Tp2Od(t1, \psi) \wedge$
        $interval((t1, 0, t0, 1), \varphi)$;
    $[\varphi\overleftarrow{B}\psi] : \Psi := \forall t1.t1 \leq t0 \wedge$
        $interval((t1, 0, t0, 1), \neg\varphi) \rightarrow Tp2Od(t1, \neg\psi)$;
  end
  return Ψ
end

function interval(l, φ)
  case Φ **of**
    $(t0, 0, t1, 0)$:
      return $\forall t2.t0 < t2 \wedge t2 < t1 \rightarrow Tp2Od(t2, \varphi)$;
    $(t0, 0, t1, 1)$:
      return $\forall t2.t0 < t2 \wedge t2 \leq t1 \rightarrow Tp2Od(t2, \varphi)$;
    $(t0, 1, t1, 0)$:
      return $\forall t2.t0 \leq t2 \wedge t2 < t1 \rightarrow Tp2Od(t2, \varphi)$;
    $(t0, 1, t1, 1)$:
      return $\forall t2.t0 \leq t2 \wedge t2 \leq 3t1 \rightarrow Tp2Od(t2, \varphi)$;
  end
end

ω-Automaton to LO2

$A_{\exists}(q1, ..., qn, \psi I, \psi R, \psi F)$ *(input automaton)*
$\exists q1..qn, \Theta LO2(0, \psi I) \wedge (\forall t.\Theta LO2(t, \psi R)) \wedge (\forall .t1\exists t2.t1 < t2 \wedge \Theta LO2(t2, \psi F))$

Where ΘLO2(t, Φ) is:
-$\Theta LO2(t, p) := p(t)$ *for variable p*
-$\Theta LO2(t, X\psi) := \Theta LO2(t + 1, \psi)$
-$\Theta LO2(t, \neg\psi) := \neg\Theta LO2(t, \psi)$
-$\Theta LO2(t, \varphi \wedge \psi) := \Theta LO2(t, \varphi) \wedge \Theta LO2(t, \psi)$
-$\Theta LO2(t, \varphi \vee \psi) := \Theta LO2(t, \varphi) \vee \Theta LO2(t, \psi)$

LTL to ω-automata

$\phi\langle X\varphi\rangle_x \Leftrightarrow \mathcal{A}_{\exists}(\{q\}, 1, q \leftrightarrow X\varphi, \phi\langle q\rangle_x)$
$\phi\langle X\varphi\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q_0, q_1\}, 1, (q_0 \leftrightarrow \varphi) \wedge (q_1 \leftrightarrow Xq_0), \phi\langle q_1\rangle_x)$
$\phi\langle G\varphi\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\}, 1, q \leftrightarrow \varphi \wedge Xq, \phi\langle q\rangle_x \wedge GF[\varphi \rightarrow q])$
$\phi\langle F\varphi\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\}, 1, q \leftrightarrow \varphi \wedge Xq, \phi\langle q\rangle_x \wedge GF[q \rightarrow \varphi])$
$\phi\langle[\varphi\ U\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\}, 1, q \leftrightarrow \psi \vee \varphi \wedge Xq, \phi\langle q\rangle_x \wedge GF[\varphi \rightarrow q])$
$\phi\langle[\varphi\ \underline{U}\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\}, 1, q \leftrightarrow \psi \vee \varphi \wedge Xq, \phi\langle q\rangle_x \wedge GF[q \rightarrow \psi])$
$\phi\langle[\varphi\ B\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\}, 1, q \leftrightarrow \neg\psi \wedge (\varphi \vee Xq), \phi\langle q\rangle_x \wedge GF[q \vee \psi])$
$\phi\langle[\varphi\ \underline{B}\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\}, 1, q \leftrightarrow \neg\psi \wedge (\varphi \vee Xq), \phi\langle q\rangle_x \wedge GF[q \rightarrow \varphi])$
$\phi\langle\overleftarrow{X}\varphi\rangle_x \Leftrightarrow \mathcal{A}_{\exists}(\{q\}, q, Xq \leftrightarrow \varphi, \phi\langle q\rangle_x)$
$\phi\langle\overleftarrow{X}\varphi\rangle_x \Leftrightarrow \mathcal{A}_{\exists}(\{q\}, \neg q, Xq \leftrightarrow \varphi, \phi\langle q\rangle_x)$
$\phi\langle\overleftarrow{G}\varphi\rangle_x \Leftrightarrow \mathcal{A}_{\exists}(\{q\}, q, Xq \leftrightarrow \varphi \wedge q, \phi\langle\varphi \wedge q\rangle_x)$
$\phi\langle\overleftarrow{F}\varphi\rangle_x \Leftrightarrow \mathcal{A}_{\exists}(\{q\}, \neg q, Xq \leftrightarrow \varphi \vee q, \phi\langle\varphi \vee q\rangle_x)$
$\phi\langle[\varphi\ \overleftarrow{U}\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\}, q, Xq \leftrightarrow \psi \vee \varphi \wedge q, \phi\langle\psi \vee \varphi \wedge q\rangle_x)$
$\phi\langle[\varphi\ \underline{\overleftarrow{U}}\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\}, \neg q, Xq \leftrightarrow \psi \vee \varphi \wedge q, \phi\langle\psi \vee \varphi \wedge q\rangle_x)$
$\phi\langle[\varphi\ \overleftarrow{B}\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\}, q, Xq \leftrightarrow \neg\psi \wedge (\varphi \vee q), \phi\langle\neg\psi \wedge (\varphi \vee q)\rangle_x)$
$\phi\langle[\varphi\ \underline{\overleftarrow{B}}\ \psi]\rangle_x \Leftrightarrow$
  $\mathcal{A}_{\exists}(\{q\}, \neg q, Xq \leftrightarrow \neg\psi \wedge (\varphi \vee q), \phi\langle\neg\psi \wedge (\varphi \vee q)\rangle_x)$

CTL to $\mu - Calculus(\Phi_{inf} = \nu y.\Diamond y)$

$EX\varphi = \Diamond(\Phi_{inf} \wedge \varphi)$
$EG\varphi = \nu x.\varphi \wedge \Diamond x$
$EF\varphi = \mu x.\Phi_{inf} \wedge \varphi \vee \Diamond x$
$E[\varphi\underline{U}\psi] = \mu x.(\Phi_{inf} \wedge \psi) \vee \varphi \wedge \Diamond x$
$E[\varphi U\psi] = \nu x.(\Phi_{inf} \wedge \psi) \vee \varphi \wedge \Diamond x$
$E[\varphi\underline{B}\psi] = \mu x.\neg\psi \wedge (\Phi_{inf} \wedge \varphi \vee \Diamond x)$
$E[\varphi B\psi] = \nu x.\neg\psi \wedge (\Phi_{inf} \wedge \varphi \vee \Diamond x)$
$AX\varphi = \Box(\Phi_{inf} \rightarrow \varphi)$
$AG\varphi = \nu x.(\Phi_{inf} \rightarrow \varphi) \wedge \Box x$
$AF\varphi = \mu x.\varphi \vee \Box x$
$A[\varphi\underline{U}\psi] = \mu x.\psi \vee (\Phi_{inf} \rightarrow \varphi) \wedge \Box x$
$A[\varphi U\psi] = \nu x.\psi \vee (\Phi_{inf} \rightarrow \varphi) \wedge \Box x$
$A[\varphi\underline{B}\psi] = \mu x.(\Phi_{inf} \rightarrow \neg\psi) \wedge (\varphi \vee \Box x)$
$A[\varphi B\psi] = \nu x.(\Phi_{inf} \rightarrow \neg\psi) \wedge (\varphi \vee \Box x)$

G and μ-calculus (safety property)
-$[\nu x.\varphi \wedge \Diamond x]_K$
-Contains states s where an infinite path π starts with $\forall t.\pi^{(t)} \in [\varphi]_K$
-φ holds always on π

F and μ-calculus (liveness property)
-$[\mu x.\varphi \vee \Diamond x]_K$
-Contains states s where a (possibly finite) path π starts with $\exists t.\pi^{(t)} \in [\varphi]_K$
-φ holds at least once on π

FG and μ-calculus (persistence property)
-$[\mu y.[\nu x.\varphi \wedge \Diamond x] \vee \Diamond y]_K$
-Contains states s where an infinite path π starts with $\exists t1.\forall t2.\pi^{(t1+t2)} \in [\varphi]_K$
-φ holds after some point on π

GF and μ-calculus (fairness property)
-$[\nu y.[\mu x.(y \wedge \varphi) \vee \Diamond x]]_K$
-Contains states s where an infinite path π starts with
$\forall t1.\exists t2.\pi^{(t1+t2)} \in [\varphi]_K ?????t1 + t2 or t1 + t0?????$
-φ holds infinitely often on π