## 1

## $\ell_1$ -norm minimization

## Charles J. Garfinkle

**Theorem 1** (Noiseless recovery). Suppose the  $n \times m$  matrix  $\mathbf{A}$  has  $\|\mathbf{A}_j\|_2 = 1$  for  $j = 1, \ldots, m$  and  $L_2(\mathbf{A}) > 0$ . Let  $\mathbf{x}_i \in \mathbb{R}^m$  be such that  $\mathbf{x}_i = c_i \mathbf{e}_i$  ( $c_i \neq 0$ ) for  $i = 1, \ldots, m$ . Suppose  $n \times m'$  matrix  $\mathbf{B}$  with  $\|\mathbf{B}_j\|_2 = 1$  for  $j = 1, \ldots, m'$  and vectors  $\overline{\mathbf{x}}_i, \ldots, \overline{\mathbf{x}}_m$  together solve:

$$\min \sum_{i=1}^{m} \|\overline{\mathbf{x}}_i\|_1 \quad subject \ to \quad \mathbf{B}\overline{\mathbf{x}}_i = \mathbf{A}\mathbf{x}_i \quad for \ i = 1, \dots, m. \tag{1}$$

Then  $A = B_S P$  for some  $S \subseteq [m']$  of size m and  $m \times m$  permutation matrix P.

*Proof.* Fixing i, we have

$$\|\mathbf{x}_i\|_1 = |c_i| \|\mathbf{A}\mathbf{e}_i\|_2 = \|\mathbf{A}\mathbf{x}_i\|_2 = \|\mathbf{B}\overline{\mathbf{x}}_i\|_2 \le \sum_j |\bar{c}_j| \|\mathbf{B}\mathbf{e}_j\|_2 = \|\overline{\mathbf{x}}_i\|_1$$
 (2)

so that  $\|\overline{\mathbf{x}}_i\|_1 = \|\mathbf{x}_i\|_1 + \varepsilon_i$  for some  $\varepsilon_i \geq 0$ . But also,

$$\sum_{i} \|\mathbf{x}_i\|_1 \ge \sum_{i} \|\overline{\mathbf{x}}_i\|_1 = \sum_{i} (\|\mathbf{x}_i\|_1 + \varepsilon_i)$$
(3)

Hence  $\sum_i \varepsilon_i \le 0$ . But since every  $\varepsilon_i$  is non-negative, it must but that  $\varepsilon_i = 0$  for i = 1, ..., m. Thus,

$$\|\overline{\mathbf{x}}_i\|_1 = \|\mathbf{x}_i\|_1 = \|\mathbf{B}\overline{\mathbf{x}}_i\|_2.$$
 (4)

**Lemma 1.** If  $\|\mathbf{T}\mathbf{v}\|_2 = \|\mathbf{v}\|_1$  for  $\mathbf{T}$  with  $\|\mathbf{T}_j\|_2 = 1$  for all j and  $|\mathbf{T}_i \cdot \mathbf{T}_j| < 1$  for all i, j then  $\mathbf{v}$  has at most one non-zero entry.

*Proof.* Let  $\mathbf{v} = \sum_{i} c_{i} e_{j}$ . Then,

$$\mathbf{T}\mathbf{v} = \sum_{i} c_{j} \mathbf{T}_{j} = \sum_{i} c_{j} \sum_{i} \mathbf{T}_{ij} \mathbf{e}_{i} = \sum_{i} \sum_{j} c_{j} \mathbf{T}_{ij} \mathbf{e}_{i} = \sum_{i} d_{i} \mathbf{e}_{i} \quad \text{where } d_{i} = \sum_{j} c_{j} \mathbf{T}_{ij}$$
 (5)

So that

$$\|\mathbf{T}\mathbf{v}\|_{2}^{2} = \sum_{i} d_{i}^{2} = \sum_{i} \left(\sum_{j} c_{j} \mathbf{T}_{ij}\right)^{2} \tag{6}$$

$$= \sum_{i} \left( \sum_{j} \left( c_{j} \mathbf{T}_{ij} \right)^{2} + \sum_{j} \sum_{k \neq j} \left( c_{j} \mathbf{T}_{ij} \right) \left( c_{k} \mathbf{T}_{ik} \right) \right)$$
(7)

$$= \sum_{i} c_j^2 \sum_{i} \mathbf{T}_{ij}^2 + \sum_{i} \sum_{k \neq i} c_j c_k \sum_{i} \mathbf{T}_{ij} \mathbf{T}_{ik}$$

$$\tag{8}$$

$$= \sum_{j} c_j^2 + \sum_{j} \sum_{k \neq j} c_j c_k \left( \mathbf{T}_j \cdot \mathbf{T}_k \right) \tag{9}$$

where we have applied the formula  $(\sum_i a_i)^2 = \sum_i a_i^2 + \sum_i \sum_{j \neq i} a_i a_j$ . Next, observe

$$\|\mathbf{v}\|_{1}^{2} = \left(\sum_{j} |c_{j}|\right)^{2} = \sum_{j} c_{j}^{2} + \sum_{j} \sum_{k \neq j} |c_{j}| |c_{k}|$$
(10)

so that  $\|\mathbf{T}\mathbf{v}\|_2^2 = \|\mathbf{v}\|_1$  implies  $\sum_i \sum_{j \neq i} \alpha_{ij} = 0$  for  $\alpha_{ij} = |c_i||c_j| - c_i c_j (\mathbf{T}_i \cdot \mathbf{T}_j)$ . Now,

$$\alpha_{ij} \ge |c_i||c_j| - |c_ic_j\left(\mathbf{T}_i \cdot \mathbf{T}_j\right)| = |c_i||c_j|\left(1 - |\mathbf{T}_i \cdot \mathbf{T}_j|\right) \ge 0 \tag{11}$$

since  $\|\mathbf{T}_j\|_2 = 1$  for all j. Thus  $\alpha_{ij} = 0$  for all i, j and we have  $|c_i||c_j| = |c_i||c_j||\mathbf{T}_i \cdot \mathbf{T}_j|$ , or  $|c_i||c_j|(1 - |\mathbf{T}_i \cdot \mathbf{T}_j|) = 0$ . Since  $|\mathbf{T}_i \cdot \mathbf{T}_j| < 1$ , this implies  $|c_i||c_j| = 0$  for all i, j. Thus,  $\mathbf{v}$  has at most one nonzero entry, since if  $c_i \neq 0$  then  $c_j = 0$  for all  $j \neq i$ .

By Lemma 1, either B has colinear columns or the  $\bar{\mathbf{x}}_i$  are all 1-sparse vectors.

It is easy to see that if  $\mathbf{B}$  has colinear columns, there exists some  $S \subset [m']$  such that the submatrix  $\mathbf{B}_S$  has no colinear columns and satisfies  $\mathbf{B}_S \overline{\mathbf{x}}_i' = \mathbf{B} \overline{\mathbf{x}}_i$  for some  $\overline{\mathbf{x}}_i' \in \mathbb{R}^{|S|}$  with  $\|\overline{\mathbf{x}}_i'\|_1 = \|\overline{\mathbf{x}}_i\|_1$  for all i. Simply set S = [m'] and  $\overline{\mathbf{x}}_i' = \overline{\mathbf{x}}_i$  and let i = 1. Iterating through  $j \neq i$ , if  $|\mathbf{B}_i \cdot \mathbf{B}_j| = 1$ , let  $S \to S \setminus \{j\}$  and  $\overline{c}_i' \to \overline{c}_i' + \overline{c}_j \cdot \operatorname{sgn}(\mathbf{B}_i \cdot \mathbf{B}_j)$ . Induct on i.

By the above claim, there is some  $S \subseteq [m']$  such that  $\mathbf{B}_S$  has no colinear columns (i.e.  $L_2(\mathbf{B}) > 0$ ) and satisfies the assumptions of the theorem for some  $\overline{\mathbf{x}}_i' \in \mathbb{R}^{|S|}$  for which  $\|\overline{\mathbf{x}}_i'\|_1 = \|\overline{\mathbf{x}}_i\|_1$  for all i. By the lemma, then, the  $\overline{\mathbf{x}}_i'$  must all have at most one non-zero entry (i.e. they are 1-sparse) and the result follows by application of the  $\ell_0$ -norm theorem.

**Theorem 2** (Noisy recovery). Suppose the  $n \times m$  matrix  $\mathbf{A}$  has  $\|\mathbf{A}_j\|_2 = 1$  for j = 1, ..., m and  $L_2(\mathbf{A}) > 0$ . Let  $\mathbf{x}_i \in \mathbb{R}^m$  be such that  $\mathbf{x}_i = c_i \mathbf{e}_i$  ( $c_i \neq 0$ ) for i = 1, ..., m. Suppose  $n \times m'$  matrix  $\mathbf{B}$  with  $\|\mathbf{B}_j\|_2 = 1$  for j = 1, ..., m' and vectors  $\overline{\mathbf{x}}_i, ..., \overline{\mathbf{x}}_m$  together solve:

$$\min \sum_{i=1}^{m} \|\overline{\mathbf{x}}_i\|_1 \quad subject \ to \quad \|\mathbf{B}\overline{\mathbf{x}}_i - \mathbf{A}\mathbf{x}_i\|_2 \le \varepsilon \quad for \ i = 1, \dots, m. \tag{12}$$

Then  $A = B_S P$  for some  $S \subseteq [m']$  of size m and  $m \times m$  permutation matrix P.

*Proof.* By the reverse triangle inequality, we have for all i:

$$\varepsilon \ge \|\mathbf{A}\mathbf{x}_i - \mathbf{B}\overline{\mathbf{x}}_i\|_2 \ge \|\mathbf{A}\mathbf{x}_i\| - \|\mathbf{B}\overline{\mathbf{x}}_i\|_2 |. \tag{13}$$

So,

$$\|\overline{\mathbf{x}}_i\|_1 \ge \|\mathbf{B}\overline{\mathbf{x}}_i\|_2 \ge \|\mathbf{A}\mathbf{x}_i\|_2 - \varepsilon = |c_i| - \varepsilon. \tag{14}$$

It is trivial to show that letting  $\mathbf{B} = \mathbf{A}$  and  $\overline{\mathbf{x}}_i = \left(1 - \frac{\varepsilon}{|c_i|}\right)\mathbf{x}_i$  for all i is a particular solution satisfying the constraints, since then  $\|\mathbf{B}\overline{\mathbf{x}}_i - \mathbf{A}\mathbf{x}_i\|_2 = \varepsilon$ . Thus, a solution to the minimization problem must satisfy:

$$\sum_{i=1}^{m} \|\overline{\mathbf{x}}_i\|_1 \le \sum_{i} \left(|c_i| - \varepsilon\right). \tag{15}$$

Taken together with Eq. 14, this implies that  $\|\overline{\mathbf{x}}_i\|_1 = |c_i| - \varepsilon$  for all i, thus  $\|\overline{\mathbf{x}}_i\|_1 = \|\mathbf{B}\overline{\mathbf{x}}_i\|_2$  for all i. By Lemma 1, either  $\mathbf{B}$  has colinear columns or the  $\overline{\mathbf{x}}_i$  are all 1-sparse vectors.

It is easy to see that if  $\mathbf{B}$  has colinear columns, there exists some  $S \subset [m']$  such that the submatrix  $\mathbf{B}_S$  has no colinear columns and satisfies  $\mathbf{B}_S \overline{\mathbf{x}}_i' = \mathbf{B} \overline{\mathbf{x}}_i$  for some  $\overline{\mathbf{x}}_i' \in \mathbb{R}^{|S|}$  with  $\|\overline{\mathbf{x}}_i'\|_1 = \|\overline{\mathbf{x}}_i\|_1$  for all i. Simply set S = [m'] and  $\overline{\mathbf{x}}_i' = \overline{\mathbf{x}}_i$  and let i = 1. Iterating through  $j \neq i$ , if  $|\mathbf{B}_i \cdot \mathbf{B}_j| = 1$ , let  $S \to S \setminus \{j\}$  and  $\overline{c}_i' \to \overline{c}_i' + \overline{c}_j \cdot \operatorname{sgn}(\mathbf{B}_i \cdot \mathbf{B}_j)$ . Induct on i.

By the above claim, there is some  $S \subseteq [m']$  such that  $\mathbf{B}_S$  has no colinear columns (i.e.  $L_2(\mathbf{B}) > 0$ ) and satisfies the assumptions of the theorem for some  $\overline{\mathbf{x}}_i' \in \mathbb{R}^{|S|}$  for which  $\|\overline{\mathbf{x}}_i'\|_1 = \|\overline{\mathbf{x}}_i\|_1$  for all i. By the lemma, then, these  $\overline{\mathbf{x}}_i'$  must all have at most one non-zero entry (i.e. they are 1-sparse). Therefore, there exist  $\overline{c}_1, \ldots, \overline{c}_{\overline{m}}$  and a map  $\pi : [m] \to [\overline{m}]$  such that

$$\|c_i \mathbf{A}_i - \bar{c}_i \mathbf{B}_{\pi(i)}\|_2 \le \varepsilon \quad \text{for all } i$$
 (16)

We could end here by applying the noisy  $\ell_0$ -norm theorem; this would imply dictionary recovery up to an error commensurate with  $\varepsilon$ . Instead, we wil go further by using the fact that in this case we know  $|\overline{c}_i| = |c_i| - \varepsilon$  for all i.

**Lemma 2.**  $\|\mathbf{u} - \mathbf{v}\|_2 \le \varepsilon$  for  $\|\mathbf{v}\|_2 = \|\mathbf{u}\|_2 - \varepsilon \implies \|\mathbf{v}\|_2 \mathbf{u} = \|\mathbf{u}\|_2 \mathbf{v}$ .

Proof. Observe that

$$\varepsilon^2 \ge \|\mathbf{u}\|_2^2 + \|\mathbf{v}\|_2^2 - 2\langle \mathbf{u}, \mathbf{v} \rangle \tag{17}$$

$$= \|\mathbf{u}\|_{2}^{2} + (\|\mathbf{u}\|_{2} - \varepsilon)^{2} - 2\langle \mathbf{u}, \mathbf{v}\rangle \tag{18}$$

$$=2\|\mathbf{u}\|_{2}^{2}-2\|\mathbf{u}\|_{2}\varepsilon+\varepsilon^{2}-2\langle\mathbf{u},\mathbf{v}\rangle\tag{19}$$

$$\implies \langle \mathbf{u}, \mathbf{v} \rangle \ge \|\mathbf{u}\|_2 (\|\mathbf{u}\|_2 - \varepsilon) = \|\mathbf{u}\|_2 \|\mathbf{v}\|_2 \tag{20}$$

But  $\langle \mathbf{u}, \mathbf{v} \rangle \leq \|\mathbf{u}\|_2 \|\mathbf{v}\|_2$  always, thus  $\langle \mathbf{u}, \mathbf{v} \rangle = \|\mathbf{u}\|_2 \|\mathbf{v}\|_2$ . Therefore,

$$\|\|\mathbf{u}\|_{2}\mathbf{v} - \|\mathbf{v}\|_{2}\mathbf{u}\|_{2}^{2} = (\|\mathbf{u}\|_{2}\|\mathbf{v}\|_{2})^{2} + (\|\mathbf{v}\|_{2}\|\mathbf{u}\|_{2})^{2} - 2\|\mathbf{u}\|_{2}\|\mathbf{v}\|_{2}\langle\mathbf{u},\mathbf{v}\rangle = 0$$
(21)

So 
$$\|\mathbf{u}\|_2 \mathbf{v} = \|\mathbf{v}\|_2 \mathbf{u}$$
.

By Lemma 2 (i.e. with  $c_i \mathbf{A}_i = \mathbf{u}$  and  $\bar{c}_i \mathbf{B}_{\pi(i)} = \mathbf{v}$ ), we have for all i that  $|c_i|\bar{c}_i \mathbf{B}_{\pi(i)} = |\bar{c}_i|c_i \mathbf{A}_i$ , or

$$\overline{c}_i' \mathbf{B}_{\pi(i)} = c_i \mathbf{A}_i \quad \text{for all } i$$
 (22)

for  $\overline{c}_i' = \operatorname{sign}(\overline{c}_i) \cdot |c_i|$ , and the result follows by application of the noiseless  $\ell_0$ -norm theorem. (Alternatively, we may directly infer  $\mathbf{B}_{\pi(i)} = \operatorname{sign}(c_i\overline{c}_j)\mathbf{A}_i$ ) and apply only the argument of the the  $\ell_0$ -norm proof establishing that  $\pi$  is a permutation.)