HW5 EE 4940 code

November 23, 2024

```
[1]: #Import libraries
import sympy as sp
import numpy as np
from decimal import *
np.set_printoptions(legacy='1.25')
```

0.1 Problem 1

You can ignore the work in here, this is just for my own personal understanding.

```
[]: a = 5
p = 19
n = 11
for x in range(0,p) :
    a_x = a**x % p
    a_n_x = a**(n*x) % p
    print(f'x = {x} | {a_n_x}')
```

```
x = 0 | 1 | 1
x = 1 | 5 | 6
x = 2 | 6 | 17
x = 3 | 11 | 7
x = 4 | 17 | 4
x = 5 | 9 | 5
x = 6 | 7 | 11
x = 7 | 16 | 9
x = 8 | 4 | 16
x = 9 | 1 | 1
x = 10 | 5 | 6
x = 11 | 6 | 17
x = 12 | 11 | 7
x = 13 | 17 | 4
x = 14 | 9 | 5
x = 15 | 7 | 11
x = 16 | 16 | 9
x = 17 | 4 | 16
x = 18 | 1 | 1
```

```
[ ]: p = 5
     X = []
     for a in range(1,p):
         i=1
         for x in range(1,p):
              a_x = a**x \% p
              \#print(f'x = \{x\} \mid \{a_x\} \mid \{a_n_x\}')
              if a_x == 1 :
                  X.append(i)
                  break
              i+=1
     print(X)
     res = []
     for val in X :
         if val not in res:
              res.append(val)
     print(res)
     print(f'{len(res)}')
    [1, 4, 4, 2]
    [1, 4, 2]
```

0.2 Problem 3

3

```
[]: def find modulo inverse(input, modulo): #Finds inverse in modulo.
         #e.g. 2 inverse in mod 3 is 2 b/c 2*2 = 4 \pmod{3} = 1
         \#==Inputs==
         #input: value to find the inverse of
         #modulo: value of the modulus
         i=modulo
         while i>0 :
             if np.mod(input*i,modulo) == 1 :
                 return i
             i-=1
         return 0
     def Elliptic_Add_GF(point_P, point_Q,a,b,p) : #
         if point_P == [0,0]: return point_Q #0 + Q = Q
         if point_Q == [0,0]: return point_P #0 + P = P
         if point_P == point_Q: # P + P = 2P
             inv = find_modulo_inverse((2*point_P[1]),p) #denominator of S
```

```
PB = [16,5]
for i in range(1,nB) :
    PB = Elliptic_Add_GF(PB,G,a,b,p)
    #print(PB)

print('')
print(f'PB = {PB}')
```

```
PB = [12, 17]
```

```
[]: xA = 4
M = [4,5]

C1 = [16,5]
for i in range(1,xA) :
        C1 = Elliptic_Add_GF(C1,G,a,b,p)

public_product = PB #xA*PB
for i in range(1,xA) :
        public_product = Elliptic_Add_GF(public_product,PB,a,b,p)
        print(public_product)

print(f'pub_product = {public_product}')

C2 = Elliptic_Add_GF(M,public_product,a,b,p)

C = [C1,C2]
    print(C)
```

[17, 0]

```
[12, 6]
    [0, 0]
    pub_product = [0, 0]
    [[19, 20], [4, 5]]
[]: #Alice public key
     C1 = C[0]
     C2 = C[1]
     print(C1)
     nB_x_C1 = C1
     for i in range(1,nB) :
         nB_x_C1 = Elliptic_Add_GF(nB_x_C1,C1,a,b,p)
         print(nB_x_C1)
     nB_x_C1[1] = -nB_x_C1[1]
     print(nB_x_C1)
     M = Elliptic_Add_GF(C2, nB_x_C1,a,b,p)
     print(M)
    [19, 20]
```

[12, 17]

[18, 10]

[17, 0]

[18, 13]

[12, 6]

[19, 3]

[0, 0]

[0, 0]

[4, 5]

0.3 Problem 5

(a) Write a computer program to compute a power of a number in \mathbb{Z}_p . Your program should take a, n, and p and compute $a^n \pmod{p}$. Note that a and n may be very large so that an is infeasible to be evaluated. Recall the efficient approach based on the binary expansion of n, we discussed in class.

```
[2]: def exp(a,n,p) : # computes a^n (mod p)
    output = 1
    while n > 0:
        if (n & 1) == 1 : output = (output * a) % p # if leftmost bit = 2
        a = (a * a) % p
        n >>= 1 # find next leftmost bit.
    return output
```

(b) Implement an encryptor for an RSA system with public key (n,e).

```
[]: def RSA_encrypt(M,e,n): #encrypts RSA message M by doing M^e (mod n)
    output = exp(M,e,n)
    return output
```

(c) Let M be your UMN student ID. Use your program to encrypt M with n=31189420800514467447616631563 and e=2887920783636036798964123603.

```
[]: id = 5631519
n = 31189420800514467447616631563
e = 2887920783636036798964123603

C = RSA_encrypt(id,e,n)
print(f'C = {C}')
```

C = 31027703950070711403380330146