

## Question 1

```
maxSum(A, j):
    if(j == 0):
        return A[0]
    prev_max = maxSum(A, j - 1) + A[j]
    return max(prev_max, A[j])
```

## Question 2

(a)

```
maxSumMemo(A, j T):
    if(j == 0):
        return A[0]
    if(T[j] != -infinity):
        return T[j]
    prev_max = maxSumMemo(A, j - 1, T) + A[j]
    T[j] = max(prev_max, A[j])
    return T[j]

maxSum(A):
    n = A.length
    T is a table/array of length n, whose values are all initialised to -infinity
    max = maxSumMemo(A, n - 1, T)
    for(k = n - 2; k >= 0; k--):
        curr = maxSumMemo(A, k, T)
        if(curr > max):
            max = curr
    return max
```

(b)

In `maxSum()`, when `maxSumMemo()` is called for the first time outside the for-loop, it recursively iterates through all  $n$  entries of  $A$ , filling up the table  $T$  with their maximum subarray sums. This takes  $n$  steps.

Then in `maxSum()`, the for-loop iterates approximately  $n$  times, calling `maxSumMemo()` each time. However, this time `maxSumMemo()` will only take a fixed number of steps, since the table  $T$  has already been filled out. So, the for-loop takes  $n$  steps.

So, the number of steps is approximately  $n + n$ . Which is a  $\Theta(n)$  cost.

## Question 3

(a)

```
maxSumTabulate(A):
```

```

n = A.length
T is a table/array of length n
max = T[0]
for(k = 1; k < n; k++):
    T[k] = max(T[k - 1] + A[k], A[k])
    if(T[k] > max):
        max = T[k]
return max

```

(b)

There is one loop that iterates over all  $n$  entries, and the cost within the loop is fixed. So, the cost is  $\Theta(n)$ .

#### Question 4

(a)

The probability that  $\text{bad}(a, b)$  returns the correct answer can be split into two scenarios. Particularly, notice that  $\text{weightedCoin}(a)$  has a  $\frac{a-1}{a}$  probability of returning 0 and a  $\frac{1}{a}$  probability of returning 1.

Case 1: When  $\text{weightedCoin}(a)$  returns 0, the correct answer  $a + b$  is returned.

Case 2: When  $\text{weightedCoin}(a)$  returns 1,  $\text{random}(b)$  is called, with a  $\frac{1}{b}$  probability of returning  $b$ , in which case the function would return  $a + b$ , the correct answer.

So, the overall probability that  $\text{bad}(a, b)$  returns the correct answer is the probability of case 1 occurring plus the probability that case 2 occurs AND returns the correct answer. Thus, the probability that  $\text{bad}(a, b)$  returns the correct answer is

$$\frac{a-1}{a} + \left(\frac{1}{a}\right)\left(\frac{1}{b}\right)$$

(b)

In case 1, the function can only return  $a + b$ . In case 2, the function can return  $a + 1, a + 2, \dots, a + b$ , each with a probability of  $\frac{1}{b}$ . So from this information, and the previously found probability of returning  $a + b$ , we derive the following probability distribution:

X	a+1	a+2	...	a + b - 1	a + b
P(X)	$\left(\frac{1}{a}\right)\left(\frac{1}{b}\right)$	$\left(\frac{1}{a}\right)\left(\frac{1}{b}\right)$	...	$\left(\frac{1}{a}\right)\left(\frac{1}{b}\right)$	$\frac{a-1}{a} + \left(\frac{1}{a}\right)\left(\frac{1}{b}\right)$

Now, we can calculate the expected value of  $X$ ,  $E(X)$ .

$$\begin{aligned}
E(X) &= \sum_{k=1}^{b-1} ((a+k)(\frac{1}{a})(\frac{1}{b})) + (a+b)(\frac{a-1}{a} + (\frac{1}{a})(\frac{1}{b})) \\
&= (\frac{1}{a})(\frac{1}{b}) \sum_{k=1}^b (a+k) + (a+b)(\frac{a-1}{a}) \\
&= (\frac{1}{a})(\frac{1}{b})(ab + \frac{b(b+1)}{2}) + (a-1) + (\frac{b(a-1)}{a}) \\
&= 1 + \frac{b+1}{2a} + a - 1 + \frac{2b(a-1)}{2a} \\
&= \frac{b+1+2ba-2b}{2a} + a \\
&= \frac{2ba-b+1}{2a} + a \\
&= a + b - \frac{b}{2a} + \frac{1}{2a}
\end{aligned}$$