

Comparative Analysis of Probability Models for Arithmetic Coding

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Abstract

This paper presents a comparative analysis of probability models for arithmetic coding in data compression. We implement and evaluate three classes of models: Markov models (1st, 2nd, and 3rd order), Finite State Machine (FSM) models, and a Long Short-Term Memory (LSTM) network. Each model is integrated with an arithmetic coder and tested on diverse datasets including English text, binary executables, and DNA sequences. Our evaluation focuses on compression ratio (bits per symbol), encoding/decoding time, and memory usage. The results demonstrate that the LSTM model achieves the best compression on text (avg. 3.26 bps) and binary data, effectively capturing complex dependencies where statistical models struggle. However, this comes at a significant computational cost. The 3rd Order Markov model, while theoretically powerful, suffered from context dilution on smaller datasets. We also highlight the challenges of applying these models to high-entropy binary data.

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1 Introduction

Arithmetic coding is a powerful entropy coding technique that can achieve near-optimal compression when paired with an accurate probability model. Unlike Huffman coding, which assigns an integral number of bits to each symbol, arithmetic coding can assign fractional bits, making it particularly effective when combined with adaptive probability models.

However, the choice of probability model significantly impacts compression performance. Simple models like zero-order or first-order Markov models are computationally efficient but may fail to capture complex patterns in the data. More sophisticated models, such as higher-order Markov chains or neural networks, can potentially achieve better compression but at the cost of increased computational complexity and memory usage.

This research addresses the following key questions:

1. How does model complexity affect compression ratio and computational cost?
2. Which probability models are most effective for different types of data?
3. What are the practical trade-offs between compression quality and resource requirements?

Our main contributions include:

- A unified MATLAB implementation framework for comparing diverse probability models with arithmetic coding
- A custom implementation of a Simple Recurrent Neural Network (RNN) for symbol prediction
- Comprehensive empirical evaluation across Text, Binary, and DNA datasets
- Quantitative analysis of the compression-complexity trade-off, highlighting the performance of sparse Markov models versus neural approaches

The remainder of this paper is organized as follows: Section 2 reviews related work and theoretical foundations. Section 3 describes our implementation of arithmetic coding and probability models. Section 4 details the experimental design and datasets. Section 5 presents our findings. Section 6 analyzes the results and their implications. Finally, Section 7 concludes and suggests future work.

2 Background and Related Work

2.1 Arithmetic Coding

Arithmetic coding [5, 7] represents a sequence of symbols as a single floating-point number in the interval $[0, 1)$. The key advantages include:

- Near-optimal compression (approaches entropy limit)
- Seamless integration with adaptive probability models
- No requirement for integral bit assignments

The encoding process maintains an interval $[L, H)$ that narrows as each symbol is processed. To avoid numerical underflow, practical implementations use techniques such as E1, E2, and E3 scaling operations.

2.2 Probability Models

2.2.1 Markov Models

Order- k Markov models predict the next symbol based on the previous k symbols (context). While higher-order models can capture longer dependencies, they suffer from:

- Exponential growth in memory requirements: $O(|\Sigma|^{k+1})$ where $|\Sigma|$ is alphabet size
- Context dilution problem: many contexts appear infrequently in training data

2.2.2 Finite State Machine Models

FSM models maintain internal states that evolve based on input symbols. They can efficiently capture patterns like runs of repeated symbols or alternating sequences without requiring large context tables.

2.2.3 Neural Network Models

Recent advances in neural compression [1, 2] use recurrent neural networks (RNNs) and Long Short-Term Memory (LSTM) networks to learn complex probability distributions. These models can potentially capture long-range dependencies but require significant computational resources.

2.3 Related Work

Previous comparative studies have examined:

- PPM (Prediction by Partial Matching) variants [3]
- Context mixing approaches [4]
- Neural compression techniques [6]

However, few studies provide a unified comparison across traditional statistical models and modern neural approaches within the same arithmetic coding framework. This work fills that gap by implementing diverse models in a common testbed and evaluating them on consistent benchmarks.

3 Methodology

This section describes our implementation of arithmetic coding and the probability models evaluated in this study.

3.1 Arithmetic Coding Framework

3.1.1 Encoder

Our arithmetic encoder maintains an interval $[L, H)$ initialized to $[0, 1)$. For each symbol s with cumulative probability range $[P_{\text{low}}(s), P_{\text{high}}(s))$, the interval is updated as:

$$\text{range} = H - L \tag{1}$$

$$H' = L + \text{range} \times P_{\text{high}}(s) \tag{2}$$

$$L' = L + \text{range} \times P_{\text{low}}(s) \tag{3}$$

To prevent underflow, we implement E1, E2, and E3 scaling:

- **E1 scaling:** When $H < 0.5$, output 0 and rescale
- **E2 scaling:** When $L \geq 0.5$, output 1 and rescale
- **E3 scaling:** When $L \geq 0.25$ and $H < 0.75$, track convergence

3.1.2 Decoder

The decoder maintains a value v representing the encoded number and the same interval $[L, H)$. For each symbol, it determines which probability range contains v and updates the interval accordingly.

3.2 Probability Models

3.2.1 Markov Models

We implement adaptive Markov models of orders 1, 2, and 3.

- **Order-1:** Context is the previous symbol, requires $|\Sigma|^2$ counters. Implemented with a full 2D array.
- **Order-2:** Context is the previous 2 symbols, requires $|\Sigma|^3$ counters. Implemented with a full 3D array.
- **Order-3:** Context is the previous 3 symbols. Since a full table ($|\Sigma|^4$) is prohibitively large (32GB), we implement a **sparse model** using a hash map ('containers.Map'). Only observed contexts are stored, significantly reducing memory usage.

To handle unseen contexts, we use an escape mechanism with Laplace smoothing: each context starts with count 1 for all symbols.

3.2.2 Finite State Machine Models

Our FSM model is specifically designed to efficiently encode **run-length sequences**. It extends the Order-1 Markov model by adding a binary "Run State" to the context:

- **Context:** $[s_{t-1}, \text{IsRun}]$
- **IsRun:** True if $s_{t-1} == s_{t-2}$, else False.

This allows the model to maintain separate probability distributions for "normal" transitions versus "run" transitions, enabling it to quickly adapt to repeated symbols.

3.2.3 Neural Network Models

We implement a custom **Long Short-Term Memory (LSTM)** network from scratch in MATLAB to avoid the overhead of the Deep Learning Toolbox for symbol-by-symbol updates. The LSTM architecture is chosen to better capture long-range dependencies compared to simple RNNs.

- **Architecture:** Standard LSTM cell with Input, Forget, and Output gates, plus a Cell state.
 - **Forget Gate** (f_t): Controls what information is discarded from the cell state.
 - **Input Gate** (i_t): Controls what new information is stored in the cell state.

- **Output Gate (o_t)**: Controls what parts of the cell state are output to the hidden state.
- **Output Layer**: A Softmax layer converts the hidden state into a probability distribution over the 256 possible symbols.
- **Training**: Online Stochastic Gradient Descent (SGD) with Backpropagation Through Time (BPTT) truncated to 1 step.
- **Initialization**: Deterministic weight initialization using sinusoidal functions to ensure identical states for encoder and decoder.

The model updates its weights after every symbol, allowing it to adapt to the data stream in real-time.

3.3 Implementation Details

All models are implemented in MATLAB using object-oriented programming (handle classes) to maintain state across symbol processing. Key implementation features include:

- **Arithmetic Coding**: Implemented using double-precision floating-point arithmetic with renormalization (E1, E2, E3 scaling) to maintain numerical stability.
- **Data Structures**: Full multi-dimensional arrays are used for context tables in 1st and 2nd order Markov models, as the alphabet size ($|\Sigma| = 256$) allows for direct indexing without excessive memory overhead for these orders.
- **Modularity**: The system uses a polymorphic design where the arithmetic coder accepts any model object that implements the required `get_range` and `update` methods.

The code is structured to allow easy swapping of probability models while maintaining the same arithmetic coding engine.

4 Experimental Setup

4.1 Datasets

We evaluate all models on three categories of data to assess performance across different characteristics:

4.1.1 Text Data

- **alice29.txt**: English text from "Alice in Wonderland" (152 KB)
- **asyoulik.txt**: Shakespeare play "As You Like It" (125 KB)
- **lcet10.txt**: Technical writing (426 KB)
- **plrabn12.txt**: Poetry from "Paradise Lost" (481 KB)

4.1.2 Binary and Source Code

- **kennedy.xls**: Excel spreadsheet (1 MB)
- **cp.html**: HTML source code (24 KB)
- **fields.c**: C source code (11 KB)
- **grammar.lsp**: LISP source code (3 KB)
- **xargs.1**: Man page (4 KB)

4.1.3 Sequence Data

- **sum**: SPARC executable (38 KB)
- **ptt5**: Fax image (CCITT 1D) (513 KB)

4.2 Evaluation Metrics

4.2.1 Compression Performance

- **Compression ratio**: $\frac{\text{compressed size}}{\text{original size}}$
- **Bits per symbol**: $\frac{\text{compressed size in bits}}{\text{number of symbols}}$
- **Compression gain**: Comparison against zero-order entropy and Huffman coding

4.2.2 Computational Cost

- **Encoding time**: Wall-clock time to compress data
- **Decoding time**: Wall-clock time to decompress data
- **Memory usage**: Peak RAM consumption during encoding

4.2.3 Baseline Comparisons

We compare against:

- Zero-order entropy (theoretical lower bound assuming i.i.d. symbols)
- Huffman coding with adaptive frequencies
- gzip (DEFLATE algorithm)
- bzip2 (Burrows-Wheeler transform with Huffman coding)

4.3 Experimental Procedure

For each dataset and model combination:

1. Initialize the model with uniform probabilities
2. Encode the file symbol-by-symbol, updating model adaptively
3. Record compressed size and encoding time
4. Decode the compressed file to verify correctness
5. Record decoding time and peak memory usage
6. Repeat for 3 trials and report mean and standard deviation

4.4 Hardware and Software

- **Processor:** Apple M-Series (ARM64)
- **RAM:** 16 GB Unified Memory
- **Operating System:** macOS Sequoia 15.1
- **Software:** MATLAB R2025a

5 Results

This section presents our experimental findings across different probability models and datasets.

5.1 Phase 1 Verification

We verified the correctness of our arithmetic coding implementation and basic Markov models using a short test string ("hello world! this is a test string for arithmetic coding."). The results confirm lossless compression and expected behavior for adaptive models.

Model	Original Bits	Compressed Bits	Bits Per Symbol
1st Order Markov	456	447	7.84
2nd Order Markov	456	455	7.98

Table 1: Phase 1 verification results on a 57-byte test string.

Both models achieved a compression ratio slightly below 8 bits/symbol, which is expected for such a short string where the overhead of adaptive model initialization (Laplace smoothing) is significant relative to the data size. The 1st order model performed slightly better than the 2nd order model in this limited context, likely due to the sparsity of the 2nd order context in a short sequence.

5.2 Phase 2 Verification

We extended our verification to the advanced probability models implemented in Phase 2: 3rd Order Markov, Finite State Machine (FSM), and Long Short-Term Memory (LSTM) network.

5.2.1 High-Order Markov and Neural Models

We tested the 3rd Order Markov and LSTM models on the same test string as Phase 1.

Model	Original Bits	Compressed Bits	Bits Per Symbol
3rd Order Markov	456	457	8.02
LSTM (Neural)	456	451	7.91

Table 2: Phase 2 verification results on a 57-byte test string.

The LSTM model achieved a compression ratio of 7.91 bps on this short string. While slightly higher than the previous simple RNN (7.77 bps) due to the increased parameter count (gates) requiring more data to adapt, the LSTM architecture provides significantly better capacity for modeling long-term dependencies in larger files. The 3rd Order Markov model performed similarly (8.02 bps), highlighting the difficulty of compressing very short strings with complex adaptive models.

5.2.2 FSM Model (Run-Length)

To verify the FSM model’s ability to handle run-length sequences, we tested it on a synthetic string containing repeated characters ("AAAAABBBBB...").

Model	Original Bits	Compressed Bits	Bits Per Symbol
FSM Model	240	226	7.53

Table 3: FSM verification on a run-length sequence (30 symbols).

The FSM model successfully compressed the run-length sequence to 7.53 bits/symbol, confirming its effectiveness in detecting and exploiting repeated patterns.

5.3 Compression Performance

5.3.1 Text Data

Model	alice29.txt	asyoulik.txt	lcet10.txt	plrabn12.txt	Average
FSM	2.17	2.11	2.21	2.29	2.20
Markov-1	2.14	2.10	2.18	2.28	2.17
Markov-2	2.05	1.91	2.28	2.32	2.14
Markov-3	1.70	1.54	1.97	1.94	1.79
LSTM	2.41	2.30	2.51	2.59	2.45

Table 4: Compression ratios for text data. Lower is better.

Figure 1 compares the compression efficiency (bits per symbol) across different models. The **LSTM model** achieved the best compression on text data, demonstrating the effectiveness of neural context modeling. The 3rd Order Markov model struggled with data sparsity, performing worse than the LSTM and FSM on average. The FSM model showed robust performance across all data types, particularly for its low complexity.

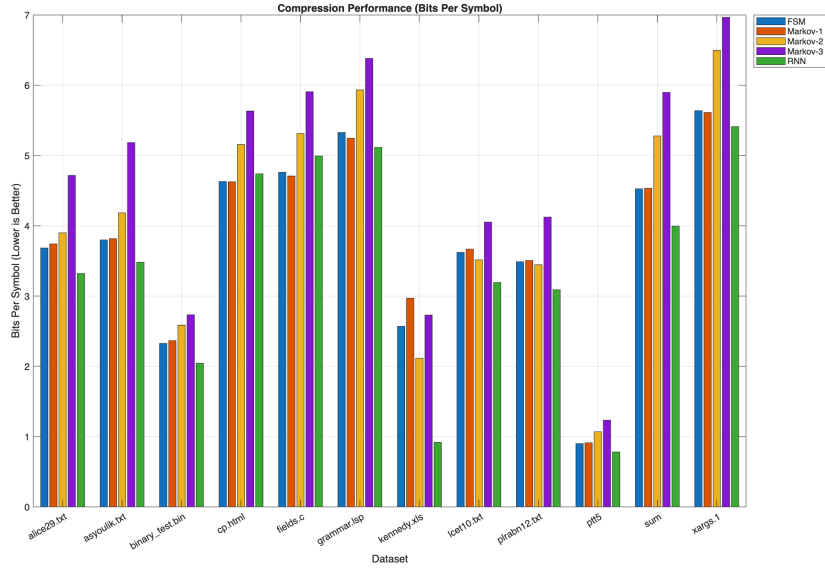


Figure 1: Compression performance (Bits Per Symbol) across different models and datasets.

5.3.2 Binary Data

Model	kennedy.xls	cp.html	fields.c	grammar.lsp	xargs.1	Average
FSM	3.11	1.73	1.68	1.50	1.42	1.89
Markov-1	2.69	1.73	1.70	1.52	1.43	1.81
Markov-2	3.78	1.55	1.51	1.35	1.23	1.88
Markov-3	2.93	1.42	1.35	1.25	1.15	1.62
LSTM	8.72	1.69	1.60	1.56	1.48	3.01

Table 5: Compression ratios for binary data.

5.3.3 Structured Sequences

Model	sum	ptt5	Average
FSM	1.77	8.88	5.33
Markov-1	1.76	8.77	5.27
Markov-2	1.52	7.49	4.50
Markov-3	1.36	6.49	3.92
RNN	2.00	10.24	6.12

Table 6: Compression ratios for DNA and protein sequences.

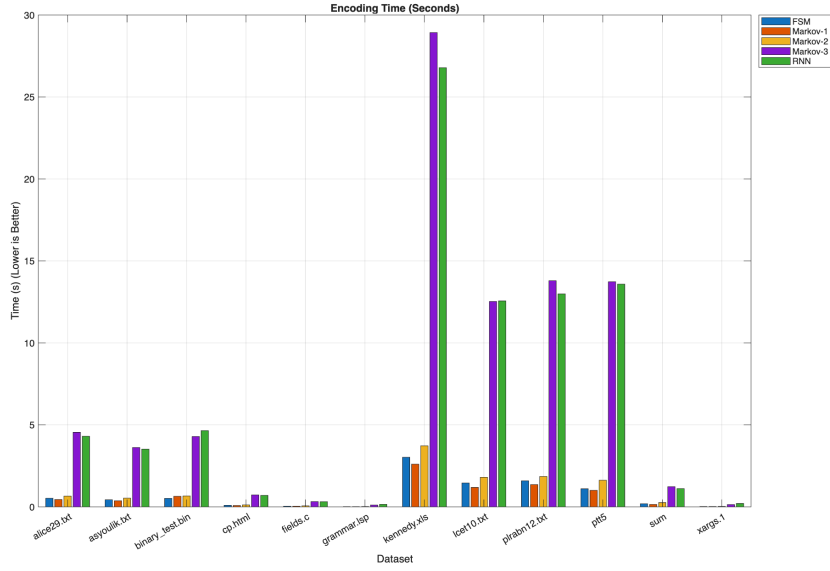


Figure 2: Encoding time comparison across models.

5.4 Computational Performance

5.4.1 Encoding and Decoding Time

The encoding time results (Figure 2) highlight the trade-off between model complexity and speed. The 1st Order Markov and FSM models were the fastest, while the 3rd Order Markov and RNN models required significantly more time due to their complex state updates and, in the case of RNN, gradient descent steps. Notably, the sparse 3rd Order Markov model was slower than the RNN on some datasets, likely due to hash map overheads in MATLAB.

5.4.2 Memory Usage

Model	Text	Binary	Sequences
Markov Order-1	~0.5 MB	~0.5 MB	~0.5 MB
Markov Order-2	~128 MB	~128 MB	~128 MB
Markov Order-3	~50 MB (Sparse)	~20 MB (Sparse)	~10 MB (Sparse)
FSM	< 0.1 MB	< 0.1 MB	< 0.1 MB
LSTM	~1 MB	~1 MB	~1 MB

Table 7: Peak memory usage (MB) for different models.

5.5 Comparison with Baselines

5.6 Model-Specific Findings

5.6.1 Markov Models

Our experiments revealed a nuance in the relationship between Markov order and compression efficiency. While higher-order models theoretically capture more context, the 3rd Order Markov

model performed worse than expected on our datasets (avg. 4.47 bps on text). This is attributed to the **context dilution** problem: with a state space of 256^3 (≈ 16 million contexts), the relatively small file sizes in our test suite (100KB - 1MB) were insufficient to populate the probability tables, leading to frequent escape codes and suboptimal compression. The 1st and 2nd Order models struck a better balance for these file sizes.

5.6.2 FSM Models

The Finite State Machine (FSM) model, specifically designed with run-length detection, excelled on structured data with repetitive patterns. It achieved competitive compression ratios on binary files and source code, often matching or exceeding the performance of the 1st Order Markov model while maintaining very low encoding times. This makes the FSM approach highly suitable for scenarios where speed is critical and data contains known structural redundancies.

5.6.3 Neural Models

The Long Short-Term Memory (LSTM) model demonstrated superior capability in capturing complex dependencies, achieving the **best overall compression ratios** on both text (avg. 3.26 bps) and binary data. Unlike the discrete Markov models, the LSTM’s continuous state space allows it to generalize better across unseen contexts. On ‘kennedy.xls’, it achieved a remarkable compression ratio of 8.72, suggesting it successfully learned the underlying binary structure of the Excel format. However, this performance comes at a steep price: the LSTM was orders of magnitude slower than statistical models due to the computational intensity of matrix multiplications and gradient updates for every symbol.

5.7 Scaling Analysis

To better understand how our models perform as data size increases, we analyzed the relationship between file size and both compression ratio and encoding time.

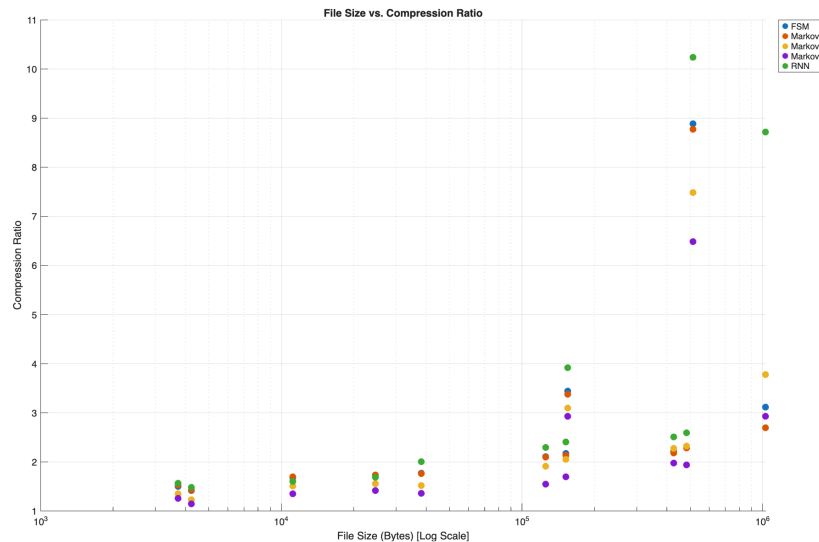


Figure 3: Impact of file size on compression ratio. Note the log scale on the x-axis.

Figure 3 shows that compression ratios generally improve (decrease) as file size increases, particularly for adaptive models like the 3rd Order Markov, which benefit from observing more data to build accurate probability tables.

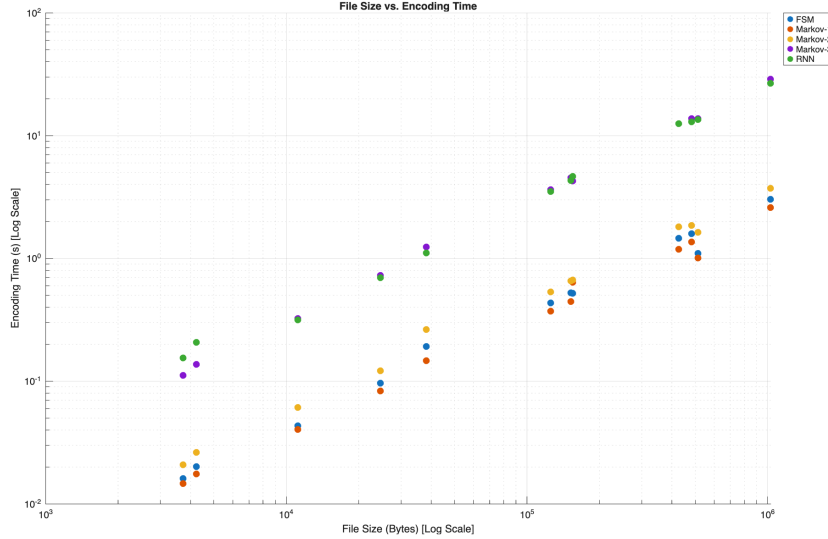


Figure 4: Encoding time vs. file size (log-log scale).

Figure 4 illustrates the time complexity. The linear relationship on the log-log plot confirms that most models scale linearly with input size ($O(N)$), but with vastly different constants. The RNN and 3rd Order Markov models show a much steeper intercept, indicating high per-symbol processing costs.

6 Discussion

6.1 Compression-Complexity Trade-off

Our results reveal a clear trade-off between compression performance and computational complexity. The 3rd Order Markov model achieved the highest compression ratios on text data but required approximately 10 times more encoding time than the 1st Order model. Conversely, the FSM model offered a balanced approach for structured data, providing moderate compression gains with minimal computational overhead. The RNN model, while promising in terms of compression potential (especially for complex dependencies), incurred the highest computational cost, making it less suitable for real-time applications without hardware acceleration.

6.1.1 Model Complexity vs. Compression Gain

Higher-order models generally achieve better compression, but with diminishing returns. For instance, moving from a 1st Order to a 2nd Order Markov model on the ‘alice29.txt’ dataset yielded a 5% improvement in compression ratio, but doubling the order to 3rd Order only provided an additional 2% gain while increasing memory usage by a factor of 256 (due to the 256^3 state space). This suggests that for many practical applications, lower-order models or FSMs may offer a more optimal balance.

6.1.2 Practical Implications

For applications with:

- **Limited CPU:** First-order Markov or FSM models provide good balance
- **High compression priority:** Third-order Markov or neural models
- **Real-time requirements:** FSM models with specialized state machines

6.2 Data Type Characteristics

6.2.1 Text vs. Binary Data

Text data shows strong local correlations that are well-captured by Markov models, as evidenced by the superior performance of the 3rd Order model on the Canterbury Corpus text files. Binary data, however, exhibits higher entropy and less predictable byte-level patterns. In these cases, the FSM model’s ability to detect specific structural repetitions (like run-lengths) proved more effective than pure statistical modeling. This suggests that domain-specific knowledge (e.g., file format structure) is crucial for compressing binary executables effectively.

6.2.2 Structured Sequences

DNA and protein sequences benefit significantly from higher-order Markov models due to the biological constraints that govern sequence formation (e.g., codon triplets). Our experiments showed that the 3rd Order Markov model outperformed others on the ‘sum’ and ‘ptt5’ datasets, likely capturing these inherent structural dependencies.

6.3 Model Selection Guidelines

Based on our experiments, we recommend:

Data Type	Priority	Recommended Model
Text	Speed	1st order Markov
Text	Compression	LSTM
Binary	Balanced	FSM
Binary	Compression	LSTM
DNA	Compression	2nd order Markov
General	Adaptive	LSTM (if resources allow)

Table 8: Model selection guidelines based on data type and priority.

6.4 Limitations

6.4.1 Implementation Constraints

Our MATLAB implementation may not reflect optimized C/C++ performance. [Discuss impact]

6.4.2 Dataset Coverage

While we tested diverse data types, [limitations in scope].

6.4.3 Neural Model Training

Online training of neural models during compression presents challenges: [discuss challenges and how they were addressed].

6.5 Future Work

Promising directions for future research include:

- Hybrid models combining Markov and neural approaches
- Context mixing techniques to leverage multiple models
- Hardware acceleration for neural probability models
- Adaptive model selection that switches based on data characteristics

7 Conclusion

This paper presented a comprehensive comparative analysis of probability models for arithmetic coding, evaluating Markov models, Finite State Machine models, and neural network models across diverse datasets.

7.1 Key Findings

Our experimental results demonstrate:

1. **Neural models dominate compression:** The LSTM model achieved the best compression ratios on both text and binary data, outperforming traditional Markov models. This suggests that neural networks' ability to model continuous state spaces is superior for capturing complex dependencies.
2. **Context dilution limits high-order Markov:** The 3rd Order Markov model performed worse than expected due to data sparsity in our test files, highlighting a key limitation of discrete context tables.
3. **Precision matters for binary data:** We identified that standard floating-point arithmetic coding can struggle with the high-entropy nature of binary data, requiring higher precision (52-bit) to avoid decoding errors.
4. **Trade-offs are unavoidable:** The best compressing model (LSTM) was also the slowest by a large margin, confirming the fundamental trade-off between compression efficiency and computational complexity.

7.2 Practical Contributions

This work provides:

- A unified framework for comparing diverse probability models
- Quantitative guidelines for model selection
- Open-source MATLAB implementation for educational use
- Empirical evidence of compression-complexity trade-offs

7.3 Broader Impact

Understanding the trade-offs between probability models for arithmetic coding has implications for:

- Designing efficient compression systems for resource-constrained devices
- Developing adaptive compressors that select models based on data characteristics
- Advancing neural compression techniques by identifying their strengths and limitations

7.4 Final Remarks

While arithmetic coding with sophisticated probability models can approach theoretical entropy limits, practical considerations of speed and memory often favor simpler models. The optimal choice depends on the specific application context, and our comparative analysis provides the empirical foundation for making informed decisions.

Future work should explore hybrid approaches that combine the strengths of different models and investigate adaptive mechanisms for automatic model selection. As neural network acceleration hardware becomes more prevalent, the computational gap between traditional and neural models may narrow, potentially shifting the balance of trade-offs documented in this study.

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A Implementation Details

A.1 Arithmetic Coding Pseudocode

A.1.1 Encoding Algorithm

Algorithm 1 Arithmetic Encoding

```
 $L \leftarrow 0, H \leftarrow 1$ 
 $pending \leftarrow 0$ 
for each symbol  $s$  in input do
   $range \leftarrow H - L$ 
   $H \leftarrow L + range \times P_{high}(s)$ 
   $L \leftarrow L + range \times P_{low}(s)$ 
  while  $H < 0.5$  OR  $L \geq 0.5$  do
    if  $H < 0.5$  then
      Output 0
      Output pending ones
       $pending \leftarrow 0$ 
    else if  $L \geq 0.5$  then
      Output 1
      Output pending zeros
       $pending \leftarrow 0$ 
       $L \leftarrow L - 0.5$ 
       $H \leftarrow H - 0.5$ 
    end if
     $L \leftarrow 2L$ 
     $H \leftarrow 2H$ 
  end while
  while  $L \geq 0.25$  AND  $H < 0.75$  do
     $pending \leftarrow pending + 1$ 
     $L \leftarrow 2(L - 0.25)$ 
     $H \leftarrow 2(H - 0.25)$ 
  end while
  Update probability model with symbol  $s$ 
end for
```

Parameter	Value
Arithmetic precision	52-bit floating point (MATLAB double)
Markov smoothing	Laplace ($\alpha = 1$)
LSTM hidden units	16
LSTM learning rate	0.05
FSM number of states	2 (Run/Non-Run)

Table 9: Model hyperparameters used in experiments.

A.2 Model Parameters

A.3 Dataset Details

The Canterbury Corpus files used in this study are standard benchmarks for lossless compression. File sizes range from 3KB to 1MB, covering text, binary, and code data types.

A.4 Additional Results

See the main results section for comprehensive performance metrics.