

$\text{reverse} :: [a] \rightarrow [a]$   
 $\text{reverse } [] = []$   
 $\text{reverse } (x:xs) = \text{reverse } xs \# [x]$

This spec is guaranteed to be well-typed.

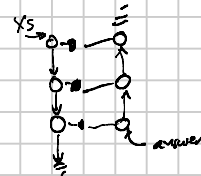
$\text{revcat} :: [a] \rightarrow [a] \rightarrow [a]$   
 $\text{revcat } xs\ ys = \text{reverse } xs \# ys$   
 if we implement  $\text{revcat}$ , we can then implement  $\text{reverse}$  in terms of  $\text{revcat}$ :  
 $\text{reverse } xs = \text{revcat } xs []$   
 we need a definition of  $\text{revcat}$  without reference to  $\text{reverse}$ ,  
 e.g. a direct recursive definition:

$\text{revcat } []\ ys$   
 $\quad \{ \text{spec. of revcat} \}$   
 $= \text{reverse } [] \# ys$   
 $\quad \{ \text{spec. of reverse} \}$   
 $= [] \# ys$   
 $\quad \{ \text{def. of } \# \}$   
 $= ys$

$\text{revcat } (x:xs)\ ys$   
 $= \{ \text{spec. of revcat} \}$   
 $\text{reverse } (x:xs) \# ys$   
 $= \{ \text{spec. of reverse} \}$   
 $(\text{reverse } xs \# [x]) \# ys$   
 $= \{ \text{associativity of } \# \}$   
 $\text{reverse } xs \# ([x] \# ys)$   
 $> \{ \text{by def. of } \# \}$   
 $\text{reverse } xs \# (x:ys)$   
 $\quad \{ \text{by spec. of revcat} \}$   
 $= \text{revcat } xs\ (x:ys)$

The new definition for  $\text{reverse}$ :

$\text{reverse} :: [a] \rightarrow [a]$   
 $\text{reverse } xs = \text{revcat } xs\ []$   
 where  
 $\text{revcat} :: [a] \rightarrow [a] \rightarrow [a]$   
 $\text{revcat } []\ ys = ys$   
 $\text{revcat } (x:xs)\ ys = \text{revcat } xs\ (x:ys)$



# Argued trees with size information

data ATree a = Leaf a  
| Fork Int (ATree a) (ATree a)

invariant will be that in a Fork n xt yt,  $n = \text{size } xt$

The invariant must be enforced when we build the tree,  
so we will only build Fork nodes using a *fork* fork

fork :: ATree a → ATree a → ATree a  
fork xt yt = Fork (size xt) xt yt  
size :: ATree a → Int  
size (Leaf x) = 1  
size (Fork n xt yt) = n + size yt

mkATree :: [a] → Maybe (ATree a)

mkATree xs =  
case xs of  
[] → Nothing  
[x] → Just (Leaf x)  
\_ → case mkATree ys of  
Nothing → Nothing  
Just yt → case mkATree zs of  
Nothing → Nothing  
Just zt → Just (fork yt zt)

where

$m = \text{length } xs \div 2$   
 $(ys, zs) = \text{splitAt } m \text{ } xs$

Exercise:  
→ Try with error  
instead of the Maybe  
→ Does it actually work?

spec. { retrieve :: ATree a → Int → a  
retrieve xt k = (flatten xt) !! k

retrieve (Leaf x) 0 = x  
retrieve (Fork m xt yt) k  
|  $k < m$  = retrieve xt k  
|  $k \geq m$  = retrieve yt (k-m)