

CS 260 – Homework #2

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Problem 1 (1.13 from text)

Show that the following statements are true.

a. 17 is $O(1)$

b. $\frac{n(n-1)}{2}$ is $O(n^2)$

c. $\max(n^3, 10n^2)$ is $O(n^3)$

d. $\sum_{i=0}^n i^k$ is $O(n^{k+1})$ and $\Omega(n^{k+1})$ for integer k

e. If $p(x)$ is any k^{th} degree polynomial with a positive leading coefficient, then $p(n)$ is $O(n^k)$ and $\Omega(n^k)$

Solution

$\begin{aligned} &a. \text{ 17 is } O(1) \\ &c(1) \geq 17, \forall n > n_0 \\ &c = 34; n_0 = 1 \\ &34 \geq 17 \\ &\therefore 17 \text{ is } O(1) \forall n \geq 1 \end{aligned}$	$\begin{aligned} &b. \frac{n(n-1)}{2} \text{ is } O(n^2) \\ &cn^2 \geq \frac{n(n-1)}{2}, \forall n > n_0 \\ &c = 2; n_0 = 1 \\ &4n^2 \geq n^2 - n \\ &3n^2 + n \geq 0 \\ &n(3n+1) \geq 0 \\ &\therefore \frac{n(n-1)}{2} \text{ is } O(n^2) \forall n \geq 1 \end{aligned}$
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c. $\max(n^3, 10n^2)$ is $O(n^3)$

$$\underline{n \leq 10}$$

$$cn^3 \geq 10n^2$$

$$c = 1; n_0 = 1$$

$$n^3 \geq 10n^2$$

$$n^3 - 10n^2 \geq 0$$

$$n^2(n - 10) \geq 0$$

$$\underline{n > 10}$$

$$cn^3 \geq n^3$$

$$c = 1; n_0 = 1$$

$$2n^3 \geq n^3$$

$$n^3 \geq 0$$

$\therefore \max(n^3, 10n^2)$ is $O(n^3)$

d. $\sum_{i=0}^n i^k$ is $O(n^{k+1})$ and $\Omega(n^{k+1})$ for integer k

$$\sum_{i=0}^n i^k \text{ is } O(n^{k+1})$$

$$k = 1$$

$$\text{Summation} = \frac{n(n+1)}{2}$$

$$cn^2 \geq \frac{n(n+1)}{2}$$

$$c = 2, n_0 = 1$$

$$2n^2 \geq \frac{n^2 + n}{2}$$

$$4n^2 \geq n^2 + n$$

$$3n^2 - n \geq 0$$

$$n(3n - 1) \geq 0$$

$$\sum_{i=0}^n i^k \text{ is } \Omega(n^{k+1})$$

$$k = 1$$

$$c * \frac{n(n+1)}{2} \geq n^2$$

$$c = 4; n_0 = 1$$

$$4n^2 + 4n \geq 2n^2$$

$$2n^2 + 4n \geq 0$$

$$2n(2n + 2) \geq 0$$

$\therefore \sum_{i=0}^n i^k$ is $O(n^{k+1})$ and $\Omega(n^{k+1})$ for integer k

e. If $p(x)$ is any k^{th} degree polynomial with a positive leading coefficient,
then $p(n)$ is $O(n^k)$ and $\Omega(n^k)$

$$O(n^k)$$

$$cn^k \geq an^k + bn^{k-1} + dn^{k-2}$$

$$(c - a)n^k - bn^{k-1} - dn^{k-2} \geq 0$$

$$\text{if } k = 2; c = 2; a = 1; b = 1; d = 1$$

$$(2 - 1)n^2 - 1n - 1n^0 \geq 0$$

$$n^2 - n \geq 1$$

$$n(n - 1) \geq 1$$

$$\therefore p(x) \text{ is } O(n^k)$$

$$\Omega(n^k)$$

$$\text{if } k = 2; c = 1; a = 3; b = 1; d = 1$$

$$c(3n^k + n^{k-1} + 1n^{k-2}) \geq n^k$$

$$3n^2 + n + 1 \geq n^2$$

$$2n^2 + n + 1 \geq 0$$

$$\therefore p(x) \text{ is } \Omega(n^k)$$

Problem 2 (1.16 from text)

Order the following functions by growth rate: (a) n , (b) \sqrt{n} , (c) $\log n$, (d) $\log \log n$, (e) $\log_2 n$, (f) $n/\log n$, (g) $\sqrt{n} \log_2 n$, (h) $(1/3)^n$, (i) $(3/2)^n$, (j) 17.

Solution

In decreasing order:

j. 17

i. $(3/2)^n$

a. n

f. $n/\log(n)$

g. $\sqrt{n} \log_2 n$

b. \sqrt{n}

e. $\log_2 n$

c. $\log n$

d. $\log(\log(n))$

h. $(1/3)^n$

Problem 3 (1.18 from text)

Here is a function $\text{max}(i, n)$ that returns the largest element in positions i through $i+n-1$ of an integer array A . You may assume for convenience that n is a power of 2.

```
function max ( i, n: integer ): integer;
var
m1, m2: integer;
begin
    if n = 1 then
        return (A[i])
    else begin
        m1 := max(i, n div 2);
        m2 := max(i+n div 2, n div 2);
        if m1 < m2 then
            return (m2)
        else
            return (m1)
        end
    end
end
```

- a. Let $T(n)$ be the worst-case time taken by max with second argument n . That is, n is the number of elements of which the largest is found. Write an equation expressing $T(n)$ in terms of $T(j)$ for one or more values of j less than n and a constant or constants that represent the times taken by individual statements of the max program.
- b. Give a tight big oh upper bound on $T(n)$. Your answer should be equal to the big omega lower bound, and be as simple as possible.

Solution

- a. $n=2$... 2 more $\text{max}()$ calls in addition to the first
 $n=4$... 4 more $\text{max}()$ calls
 $n=8$... 8 more $\text{max}()$ calls
 Mapped to a function, $T(j)=2*T(j/2)+1$
- b. $T(n)$ is $O(n)$

Problem 4 (2.9 from text)

The following procedure was intended to remove all occurrences of element x from list L . Explain why it doesn't always work and suggest a way to repair the procedure so it performs its intended task.

```
procedure delete ( x: elementtype; var L: LIST );
var
    p: position; begin
    p := FIRST(L);
    while p <> END(L) do begin
        if RETRIEVE(p, L) = x then
            DELETE(p, L);
        p := NEXT(p, L)
    end
end; { delete }
```

Solution

An issue arises with this particular delete procedure when there are consecutive occurrences of the element x . After deleting an x , the position moves to the next one. However, the procedure fails to check if the new element in the position that was just deleted, if it contains the element x . This can be fixed by including an else statement to the if of the retrieve. This will make sure consecutive occurrences of x won't get missed.

Problem 5 (2.11 from text)

Suppose L is a LIST and p , q , and r are positions. As a function of n , the length of list L , determine how many times the functions FIRST, END, and NEXT are executed by the following program.

```
 $p := \text{FIRST}(L);$   $O(1)$   
while  $p \neq \text{END}(L)$  do begin  $O(n)$   
     $q := p;$   
    while  $q \neq \text{END}(L)$  do begin  $O(n)$   
         $q := \text{NEXT}(q, L);$   
         $r := \text{FIRST}(L);$   
        while  $r \neq q$  do  $O(n)$   
             $r := \text{NEXT}(r, L)$   
    end;  
     $p := \text{NEXT}(p, L)$   
end;
```

Solution

FIRST $\Rightarrow 1 + n \cdot n = 1 + n^2$

END $\Rightarrow n \cdot n = n^2$

NEXT $\Rightarrow n \cdot n \cdot n = n^3$