

# Incorrect Systems: It's not the Problem, It's the Solution

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Universität Salzburg



DREAMS Seminar, Berkeley, July 2012

Software

Software/  
Hardware

Hardware

Software

Software/  
Hardware

Hardware

Krishna Palem  
Rice

Software

Software/  
Hardware

Probabilistic or  
Approximate  
Computing

Krishna Palem  
Rice

Software

Software/  
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Probabilistic or  
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Computing

Rakesh Kumar  
UIUC

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Software

Stochastic  
Processors

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Software

Martin Rinard  
MIT

Stochastic  
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Probabilistic or  
Approximate  
Computing

Krishna Palem  
Rice

Program  
Transformation

Martin Rinard  
MIT

Stochastic  
Processors

Rakesh Kumar  
UIUC

Probabilistic or  
Approximate  
Computing

Krishna Palem  
Rice

## Program Transformation

1. memory leaks
2. addressing errors
3. infinite loops

## Stochastic Processors

Rakesh Kumar  
UIUC

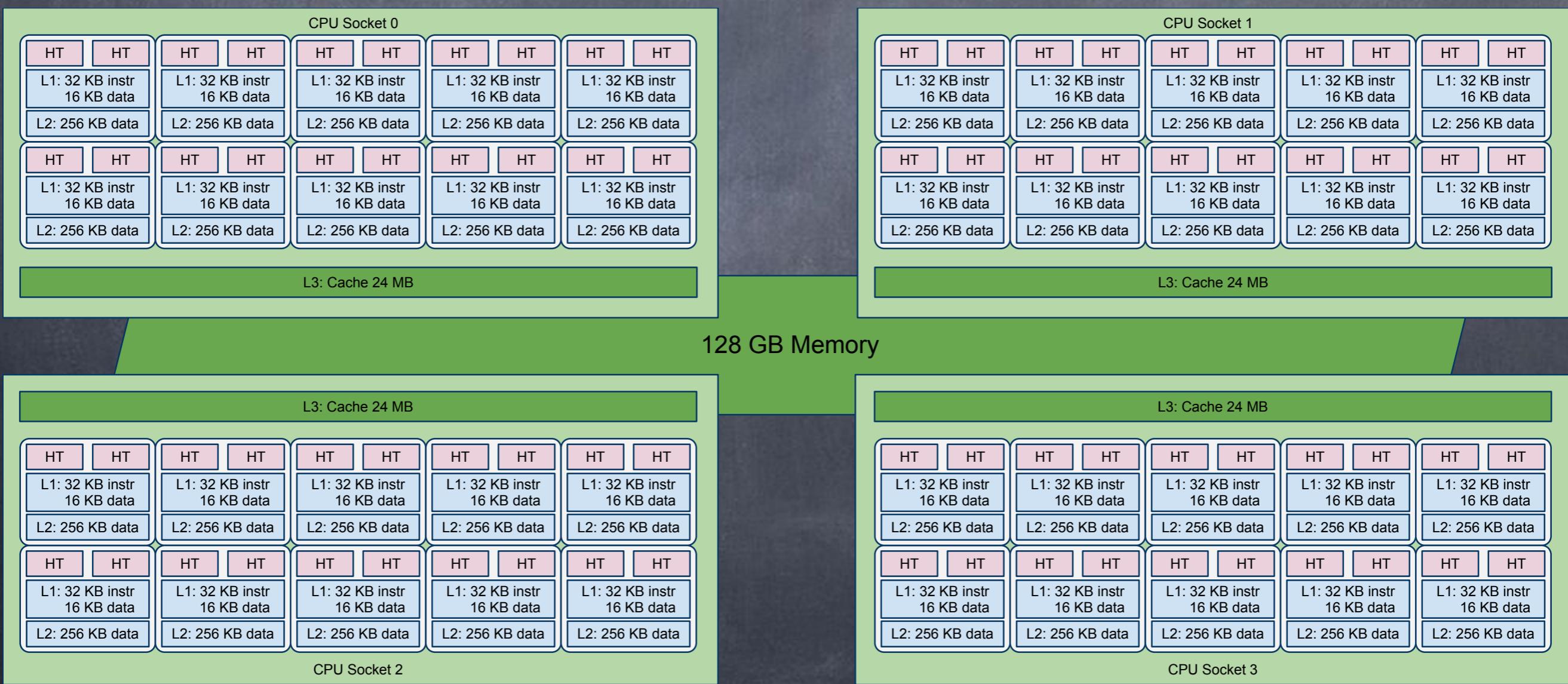
## Probabilistic or Approximate Computing

Krishna Palem  
Rice

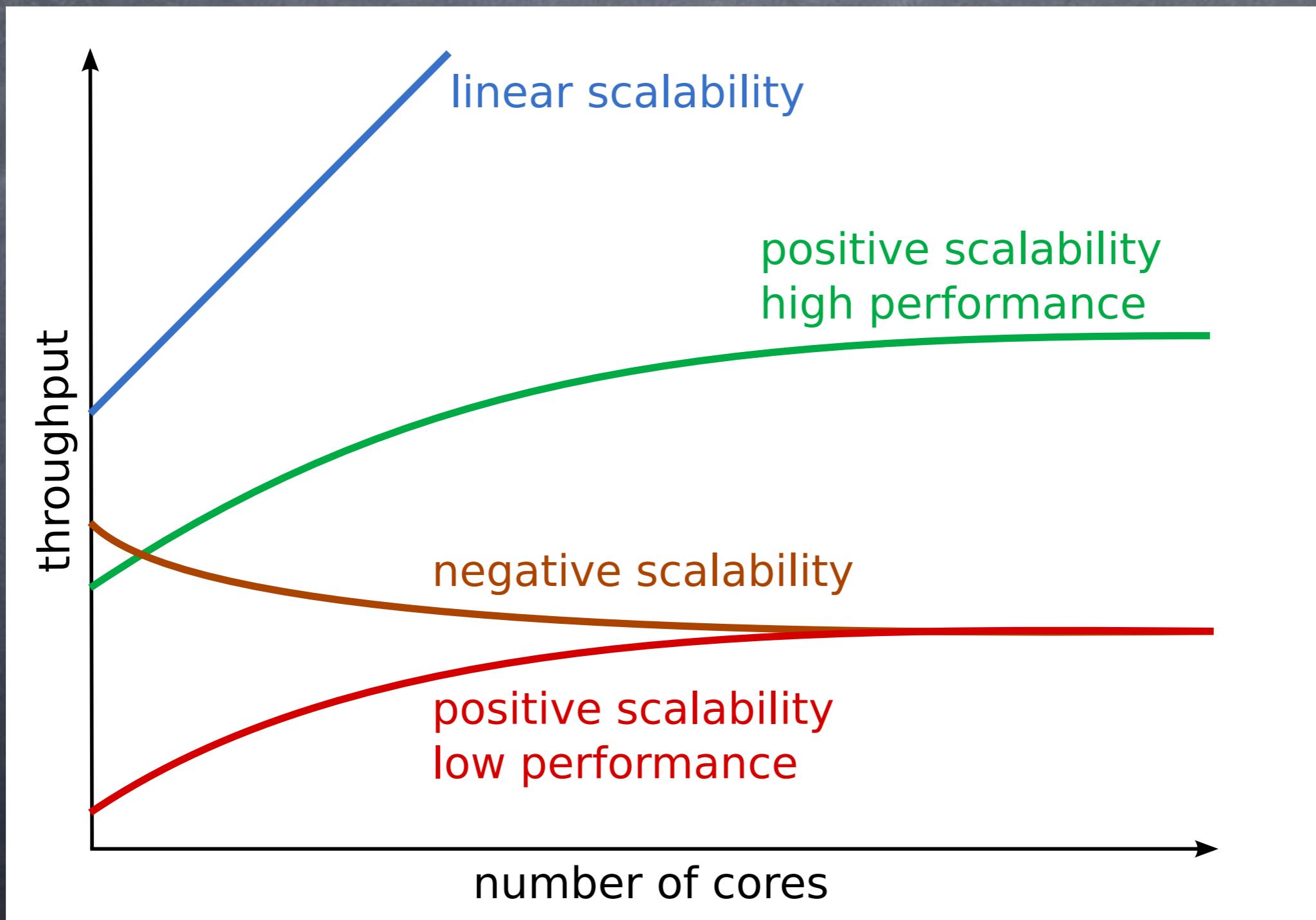
# Metrics of Correctness in Systems Engineering

Joint work w/ A. Haas,  
M. Lippautz, H. Payer,  
H. Röck, A. Sokolova and our  
collaborators at IST Austria  
T. Henzinger, A. Sezgin

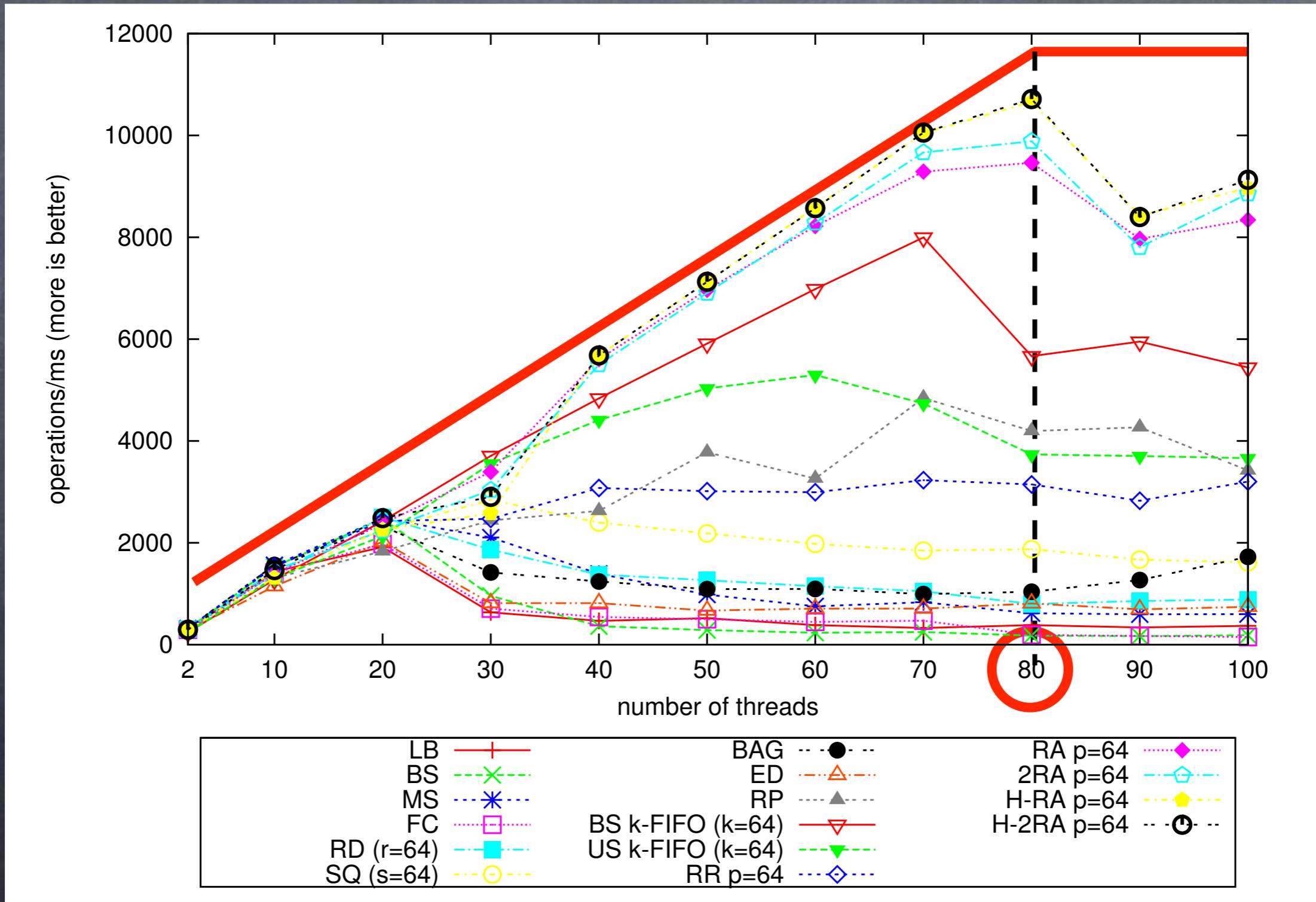
4 processors × 10 cores ×  
 2 hardware threads =  
 80 hardware threads



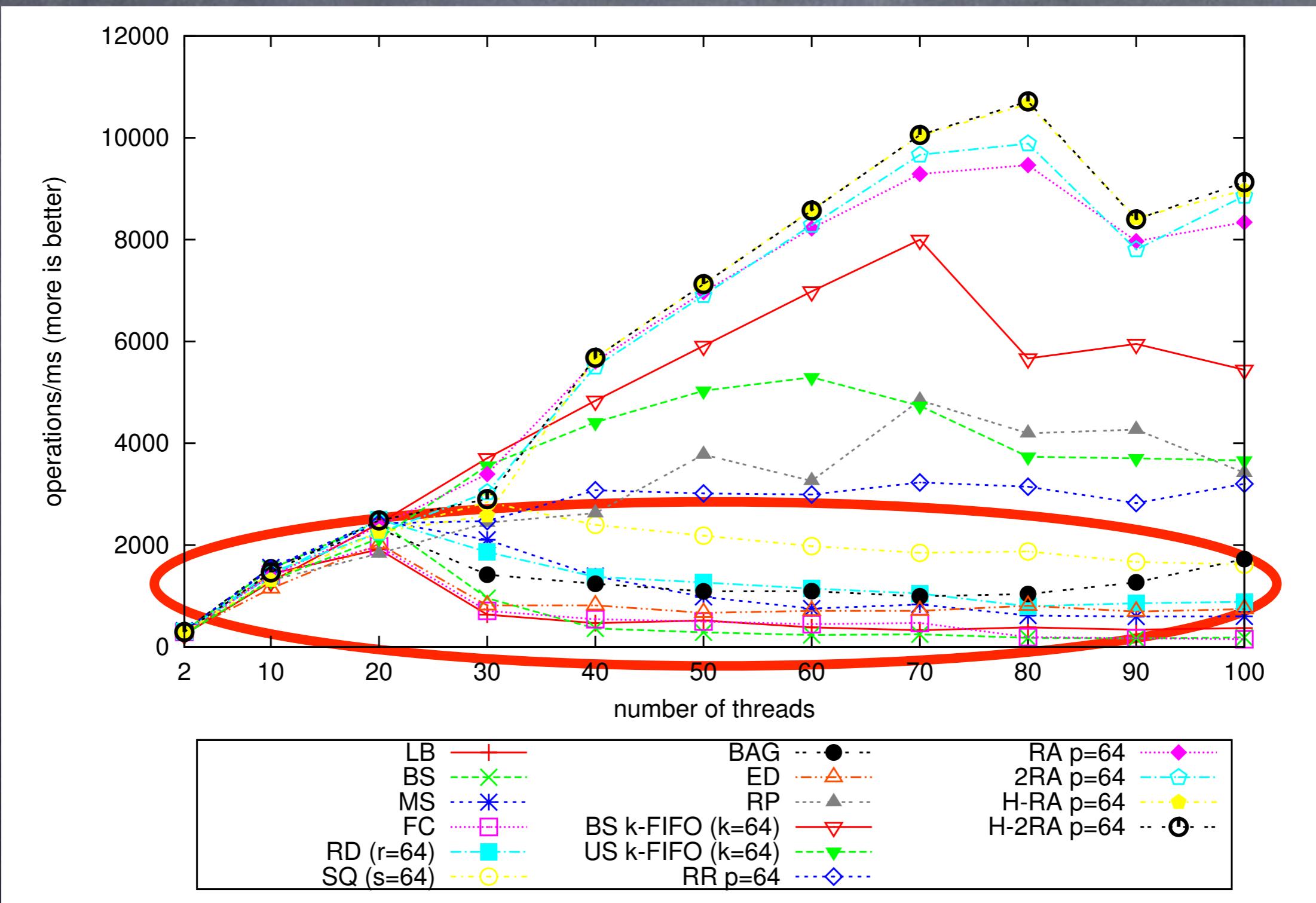
# Performance & Scalability



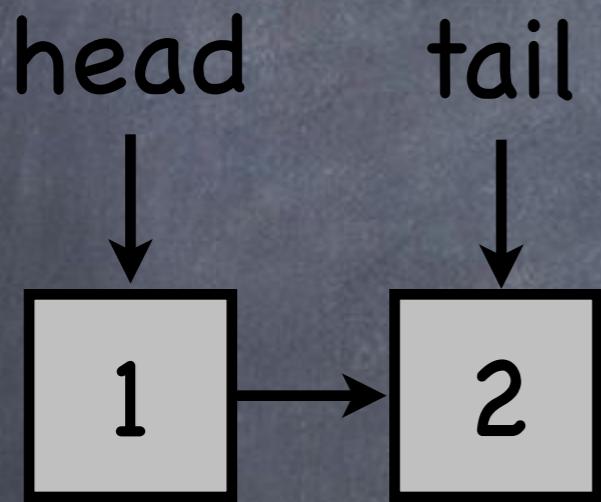
# Ideal 80-Thread Performance



# Regular FIFO Queues



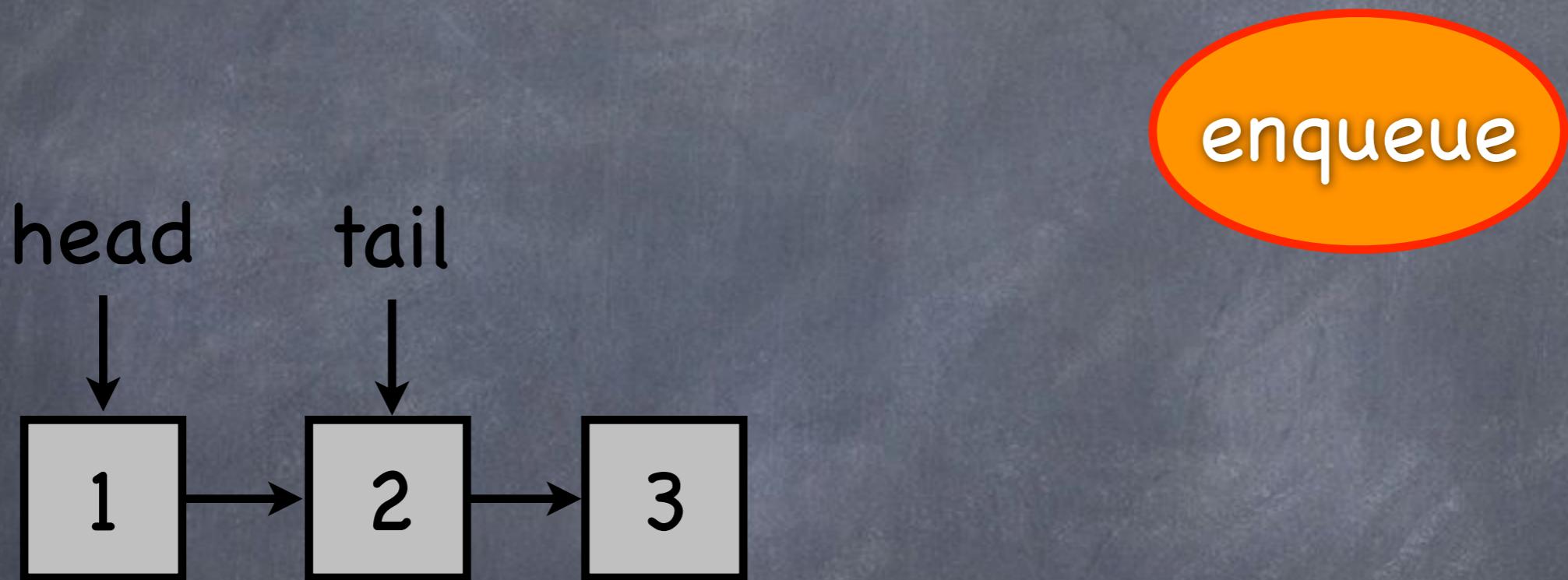
# Concurrent First-in- First-out (FIFO) Queue



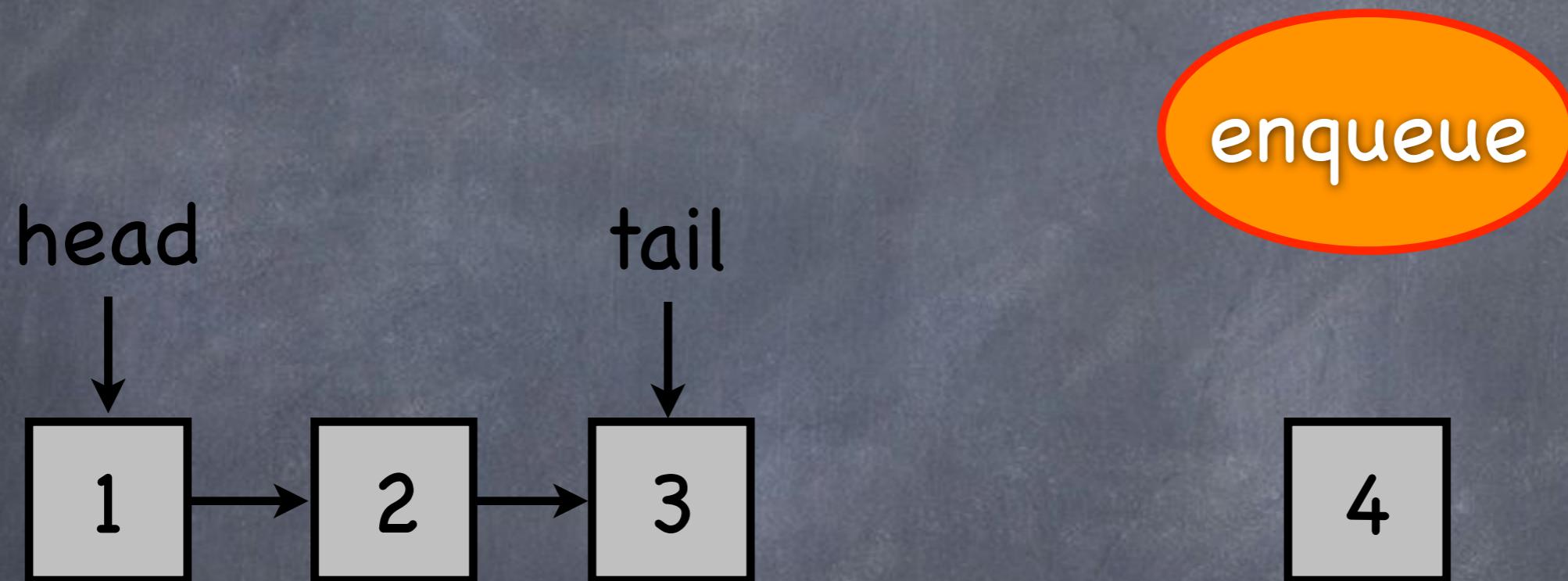
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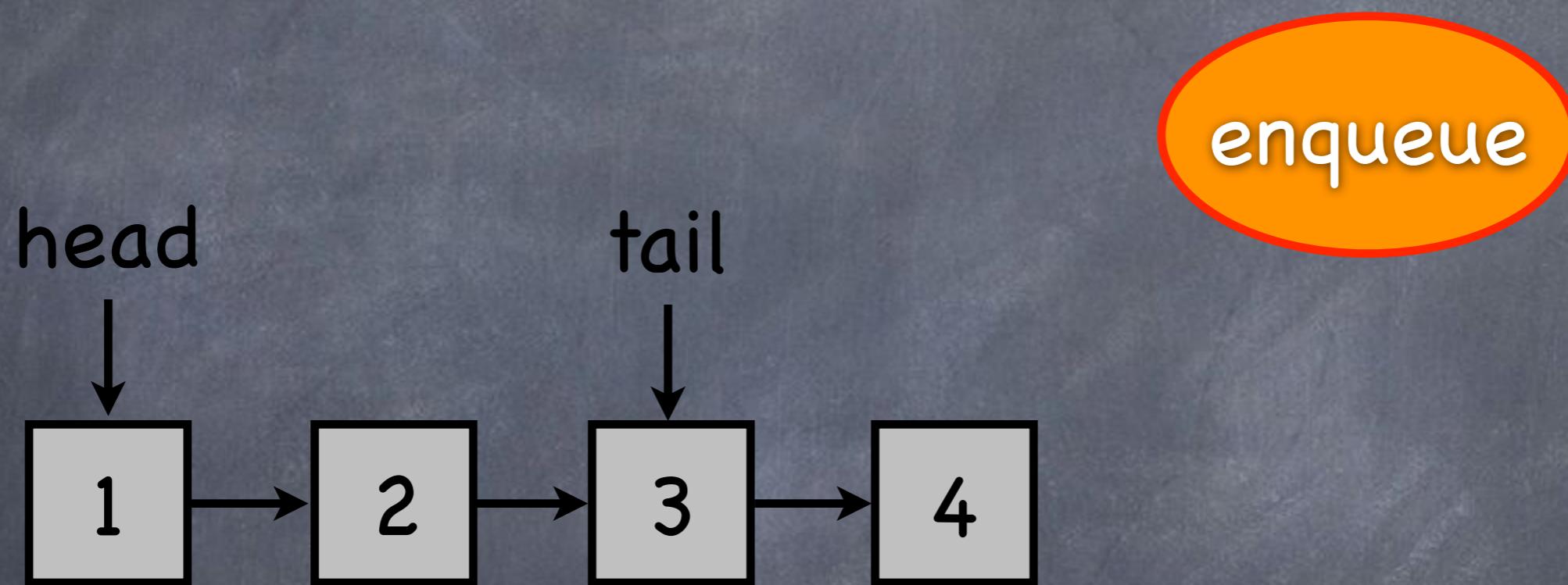
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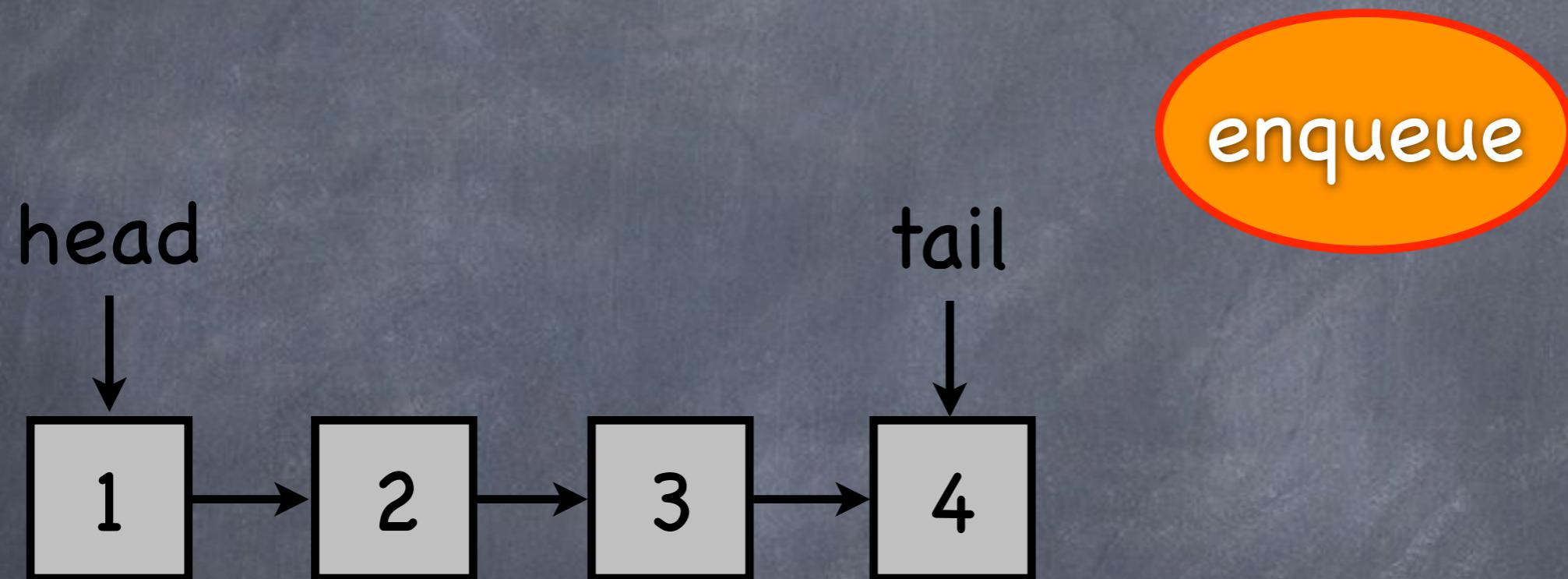
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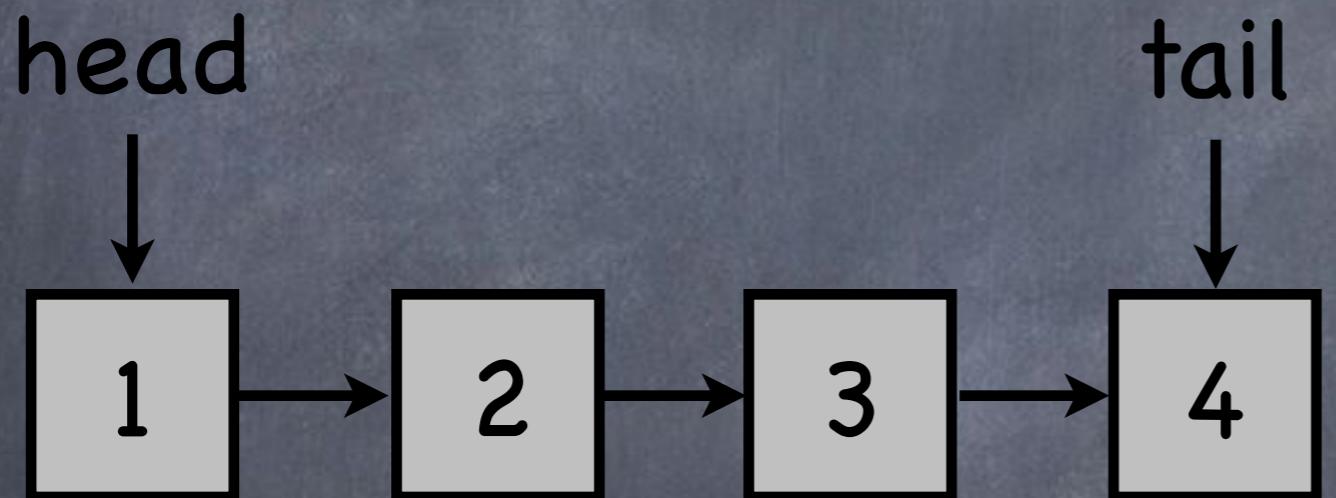
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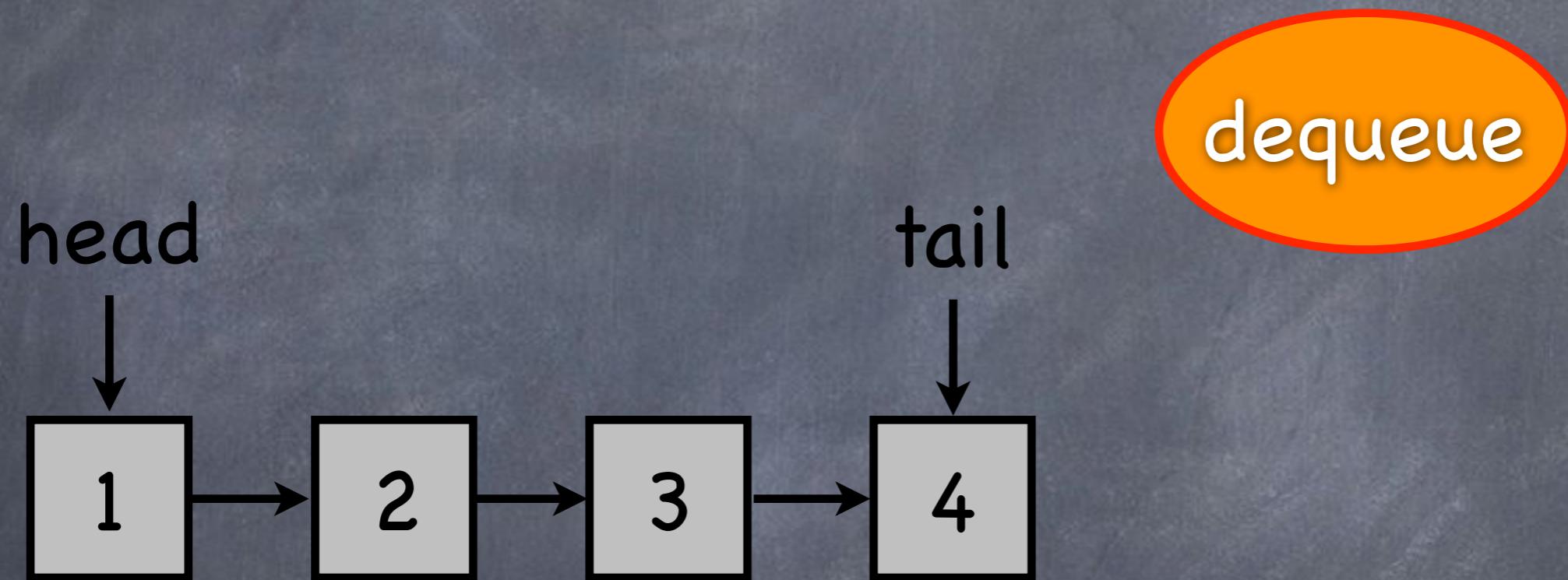
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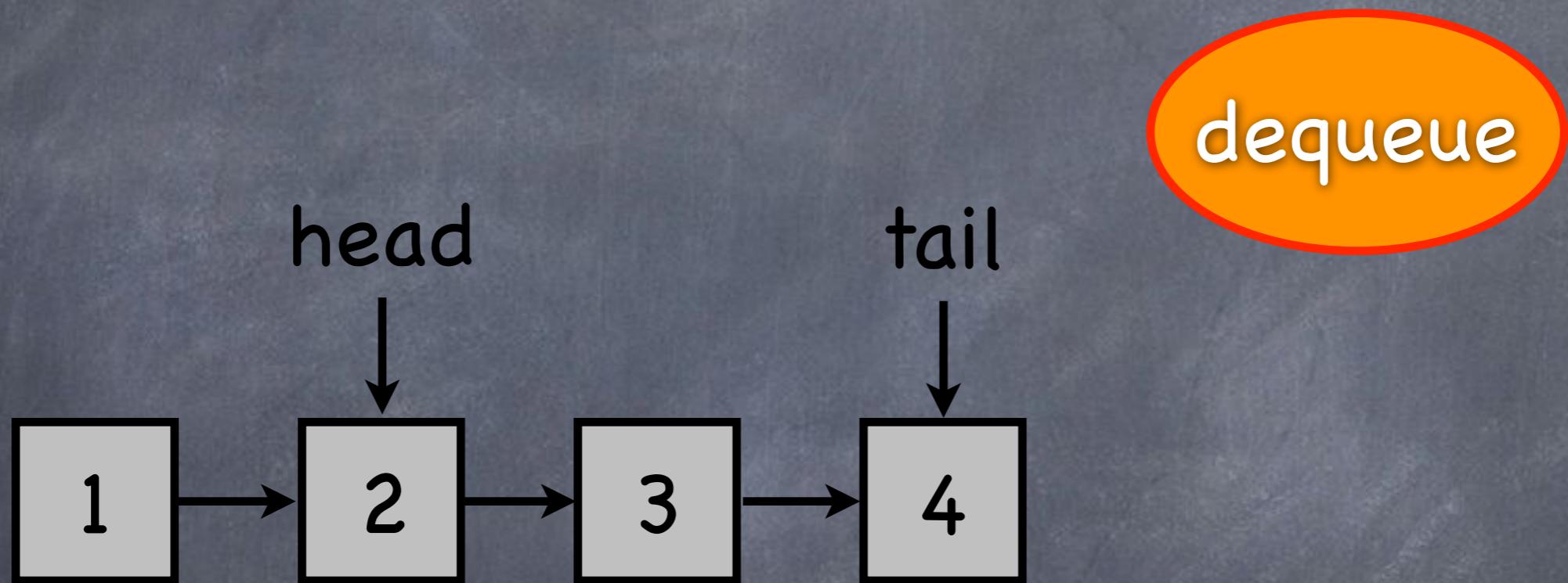
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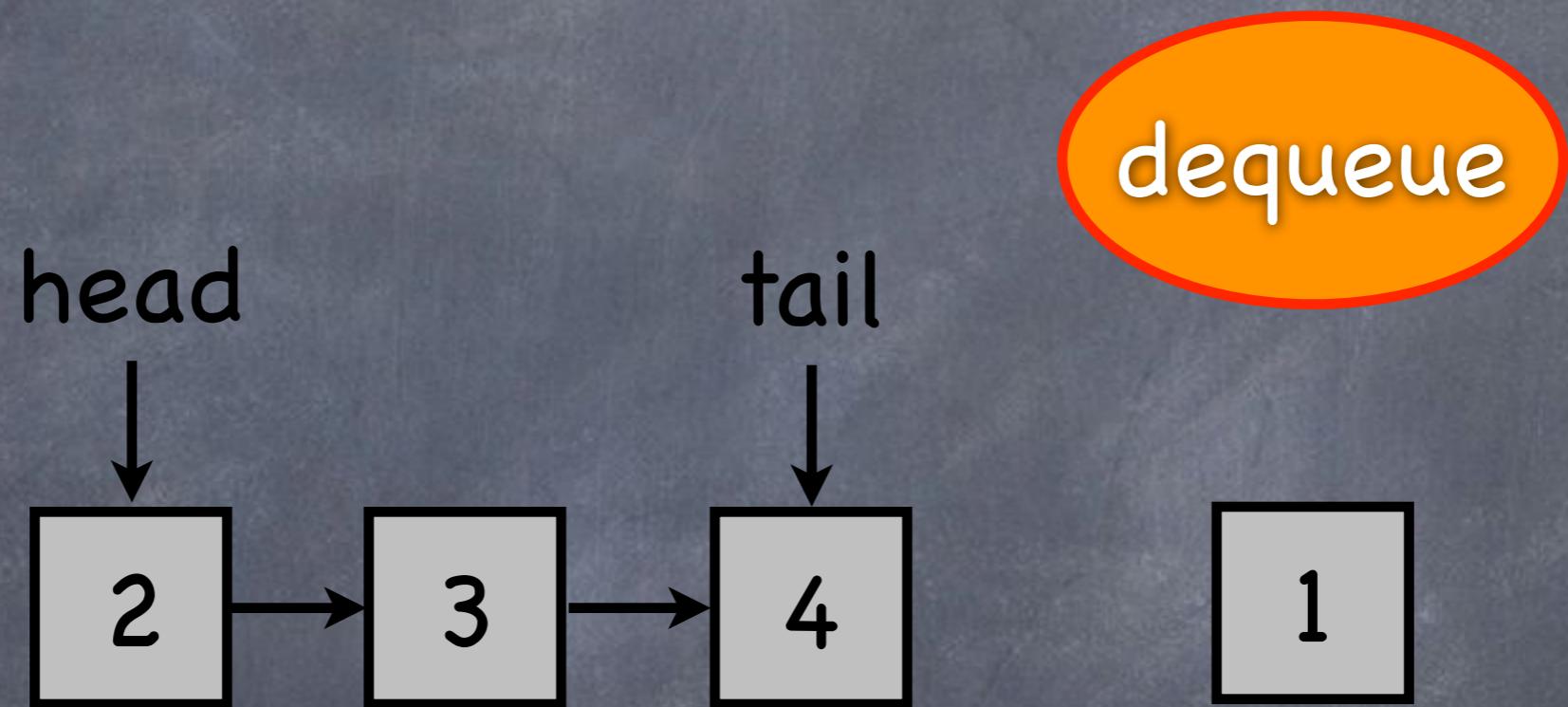
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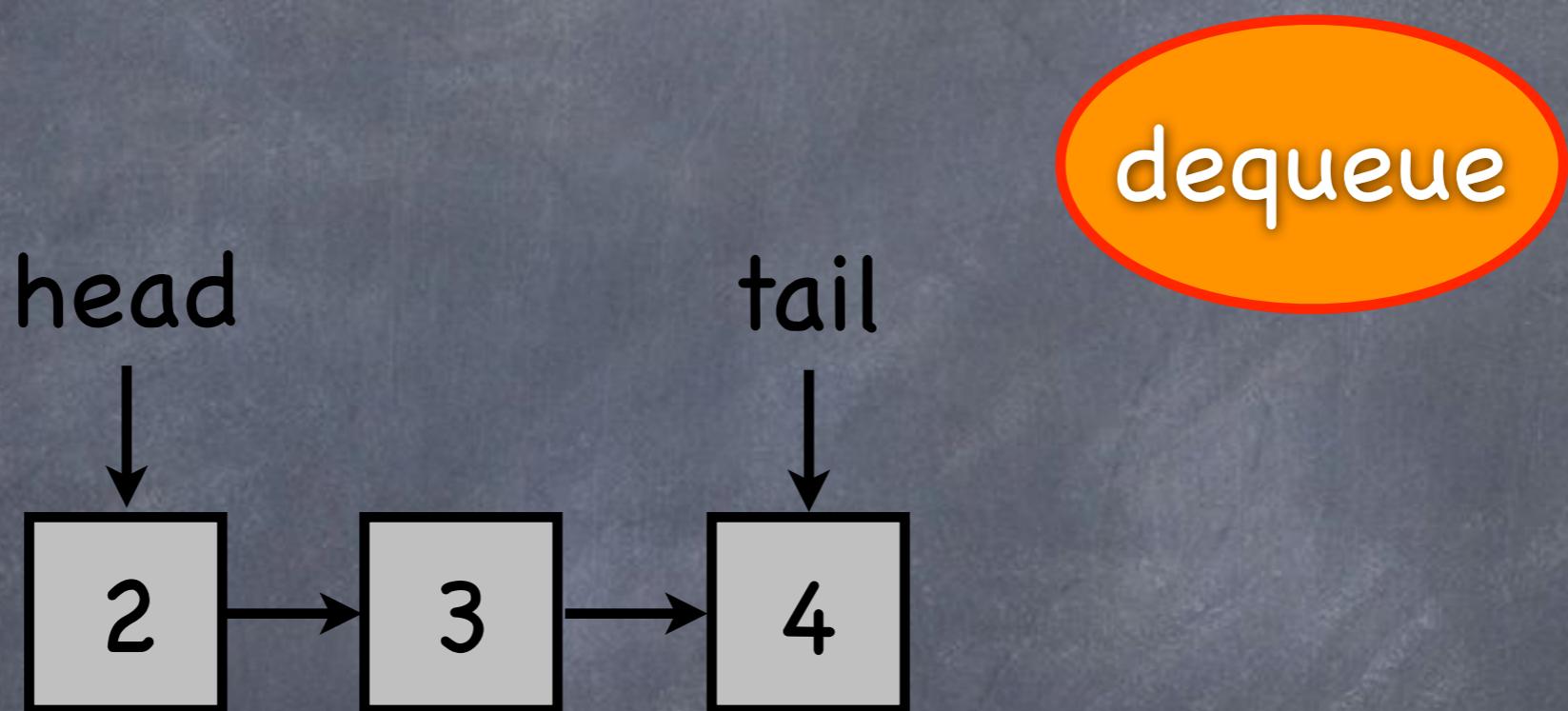
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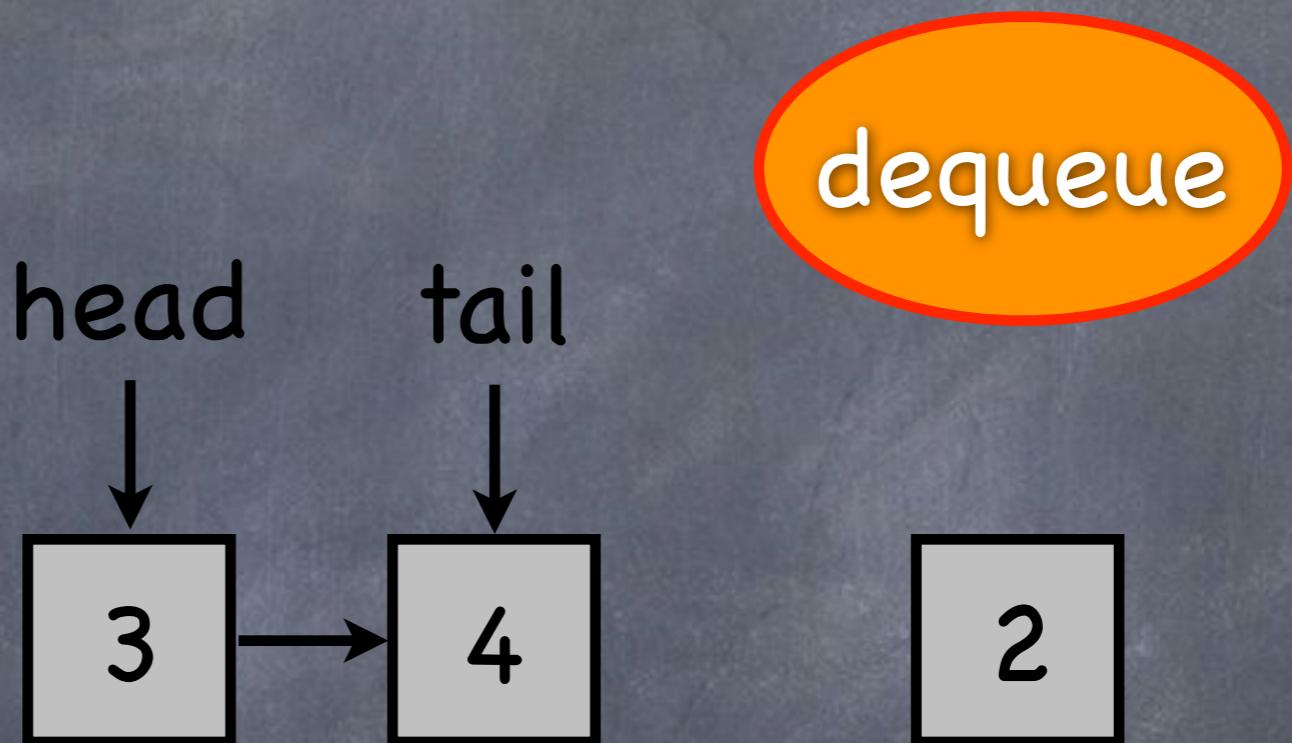
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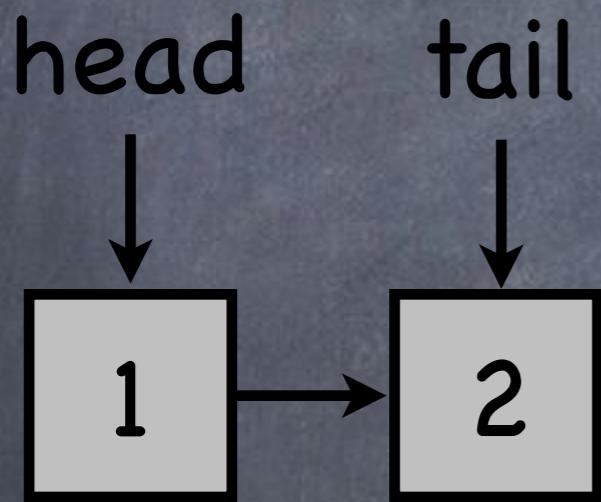
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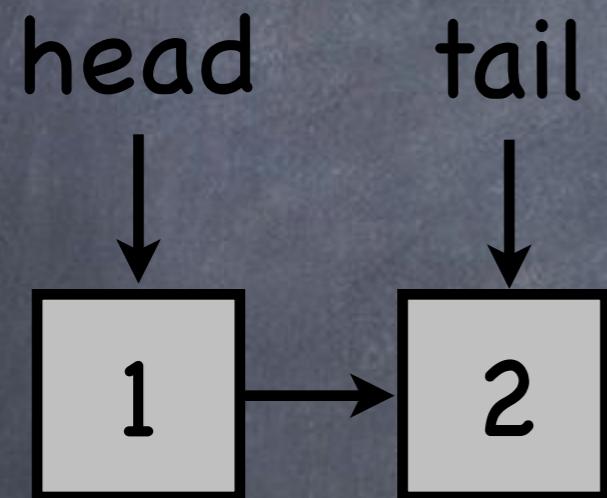
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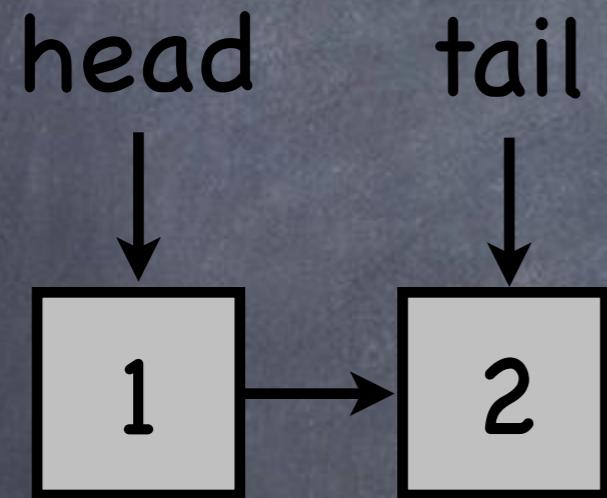


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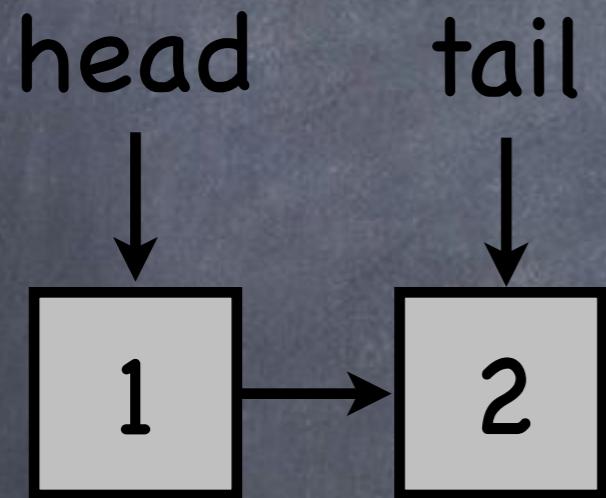
-> 1 lock

# Concurrent First-in- First-out (FIFO) Queue



-> 1 lock -> 2 locks

# Concurrent First-in-First-out (FIFO) Queue



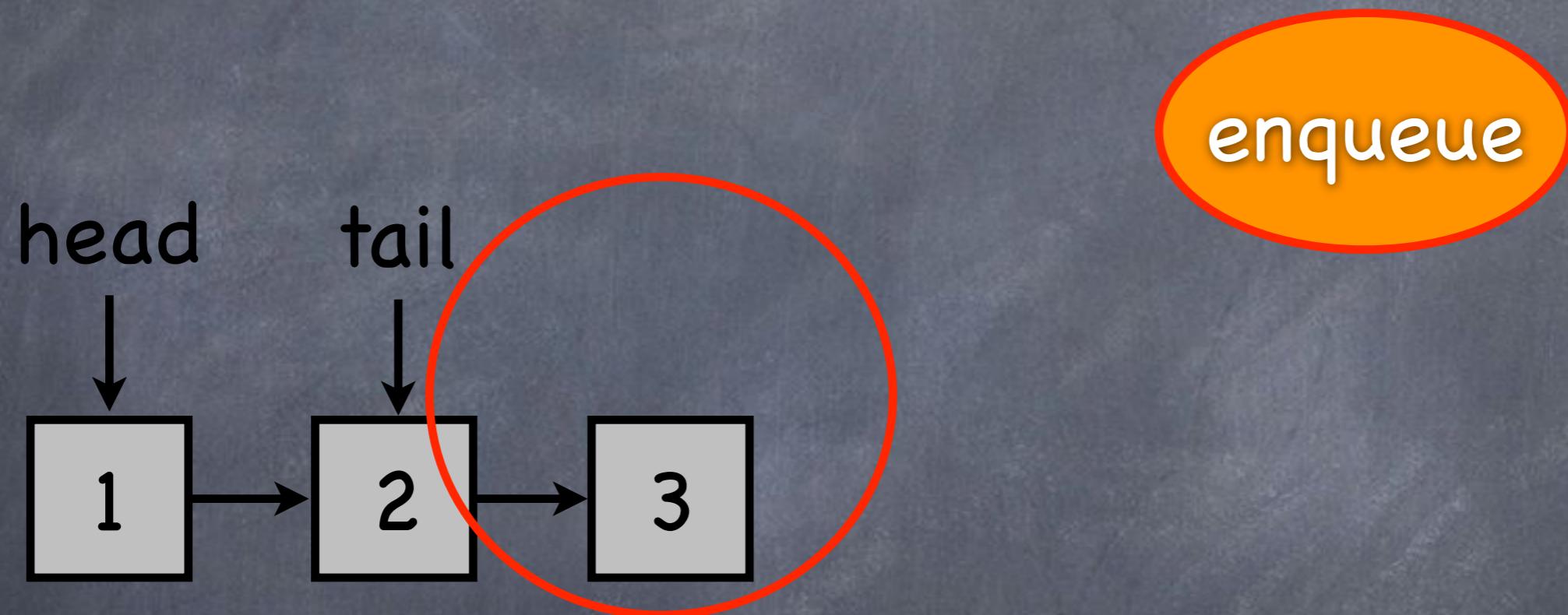
-> 1 lock -> 2 locks -> 0 locks

# Concurrent First-in-First-out (FIFO) Queue



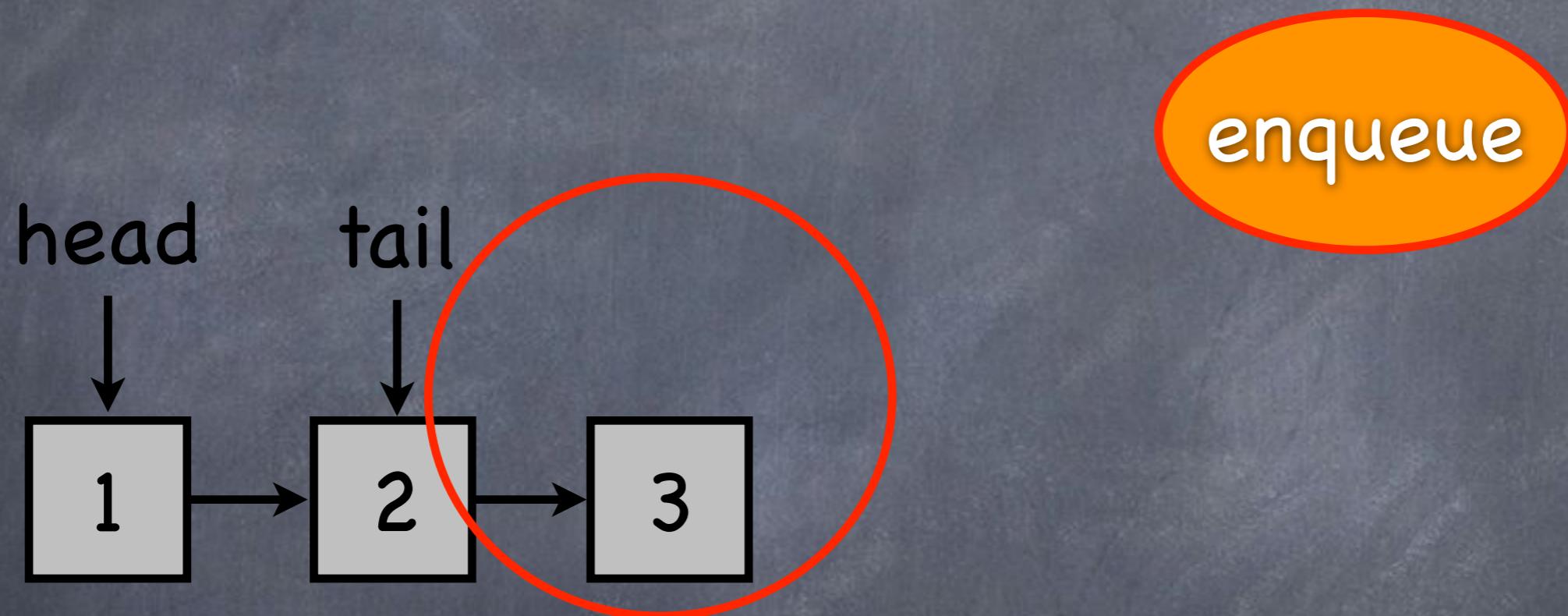
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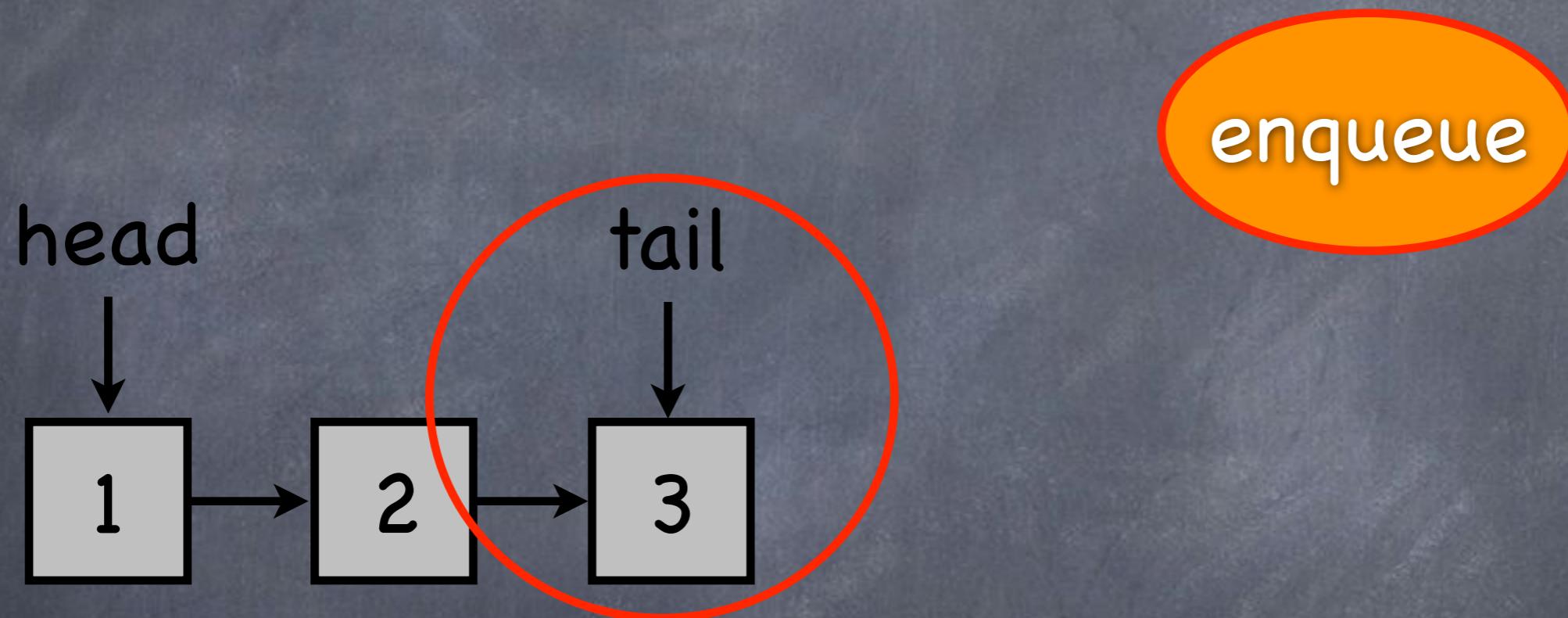
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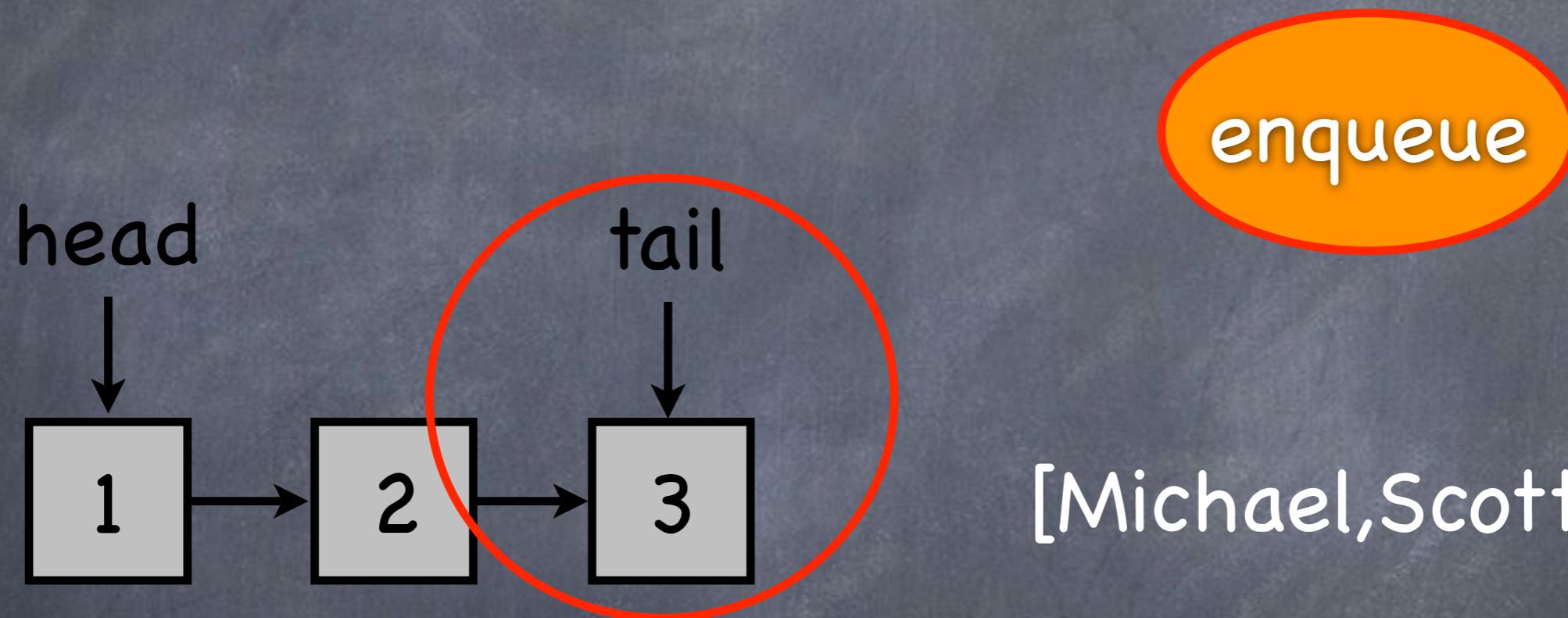
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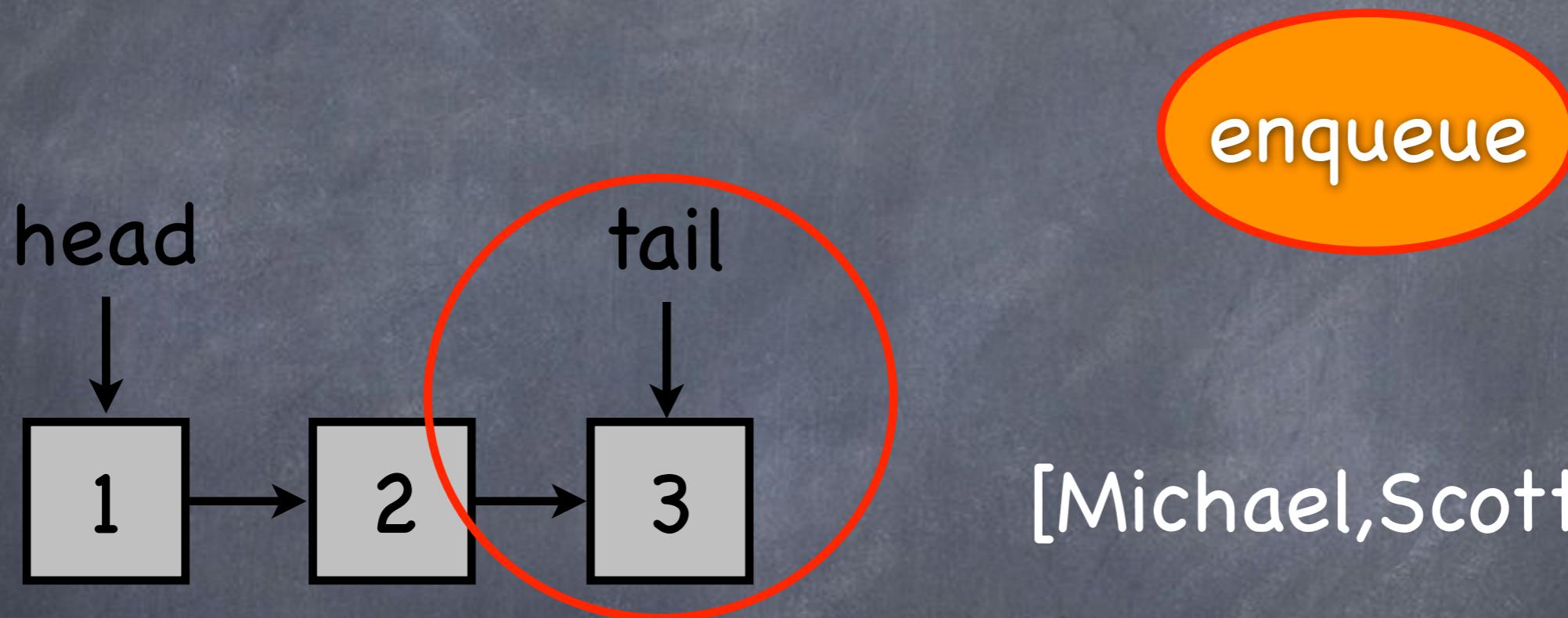
# Concurrent First-in-First-out (FIFO) Queue



[Michael, Scott '96]

-> 1 lock -> 2 locks -> 0 locks -> compare & swap

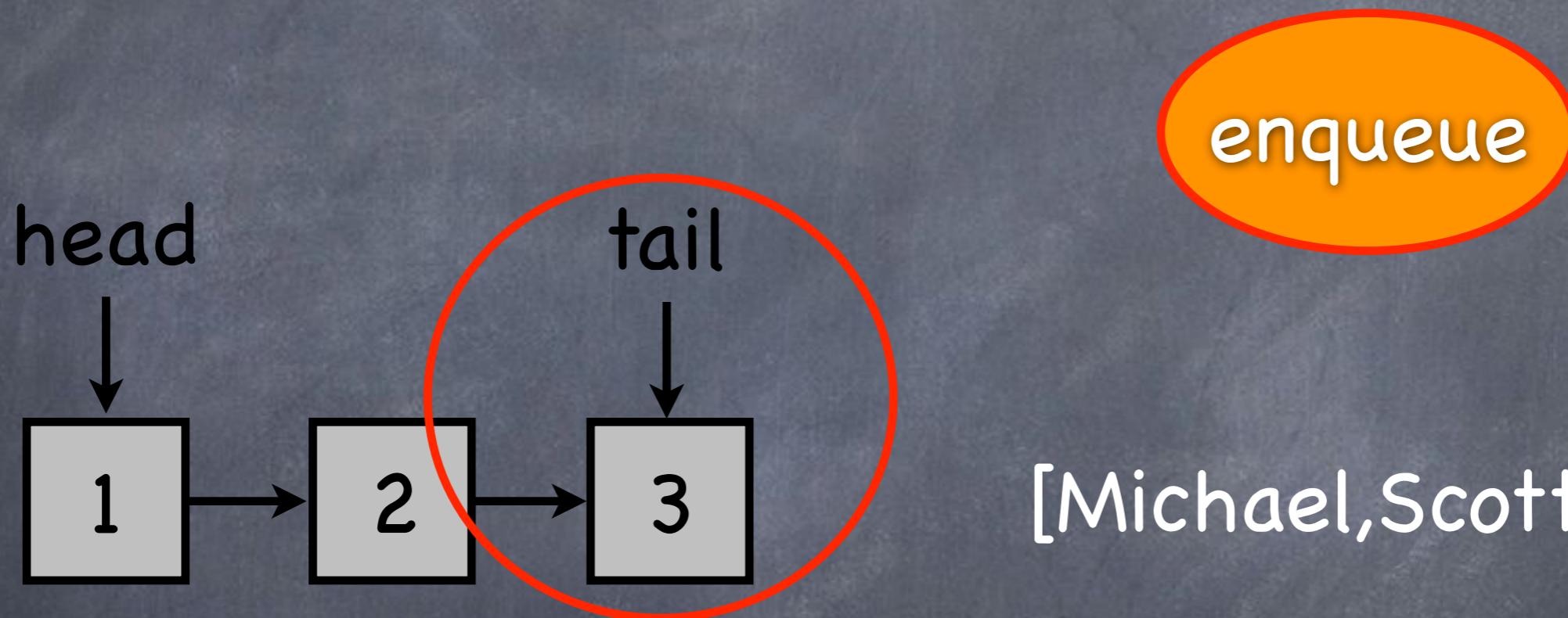
# Concurrent First-in-First-out (FIFO) Queue



[Michael, Scott '96]

- > 1 lock
- > 2 locks
- > 0 locks
- > compare & swap
- > lock-based vs. **lock-free** vs. wait-free?

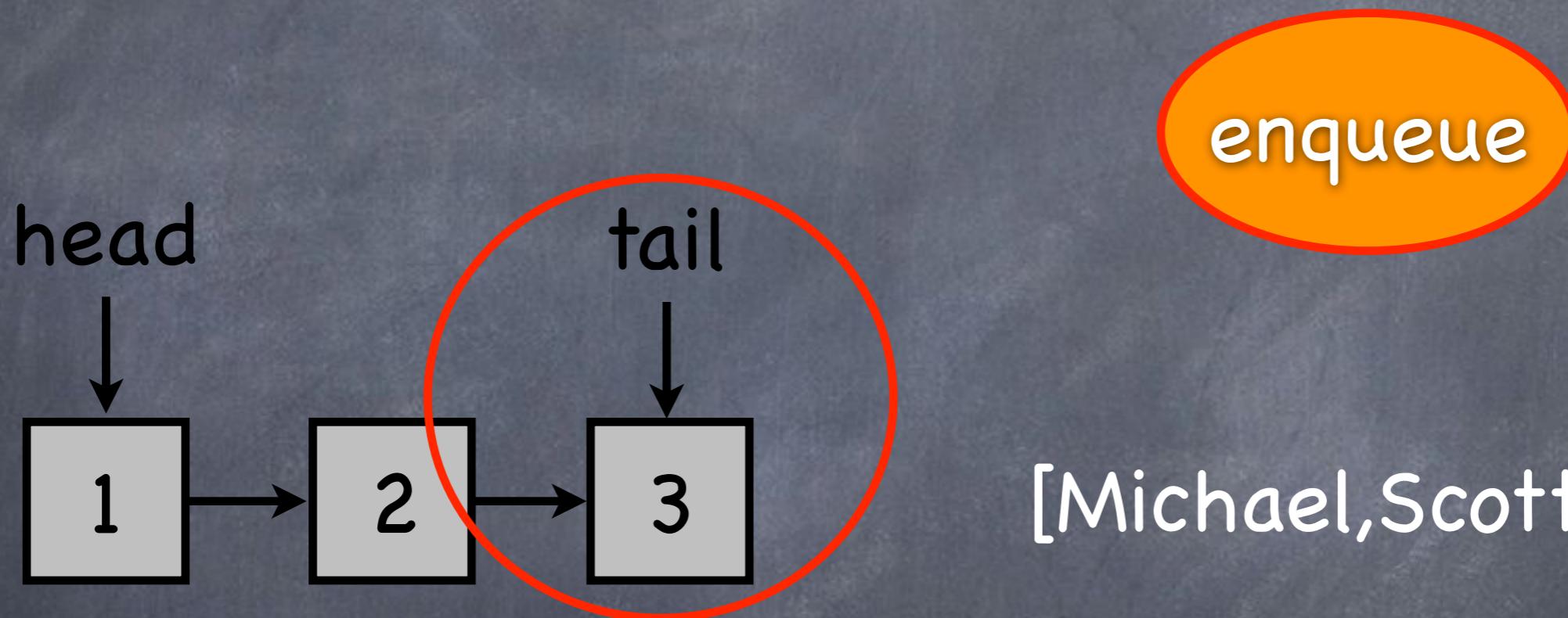
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[Michael, Scott '96]

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- > memory contention on **head** and **tail** pointers!

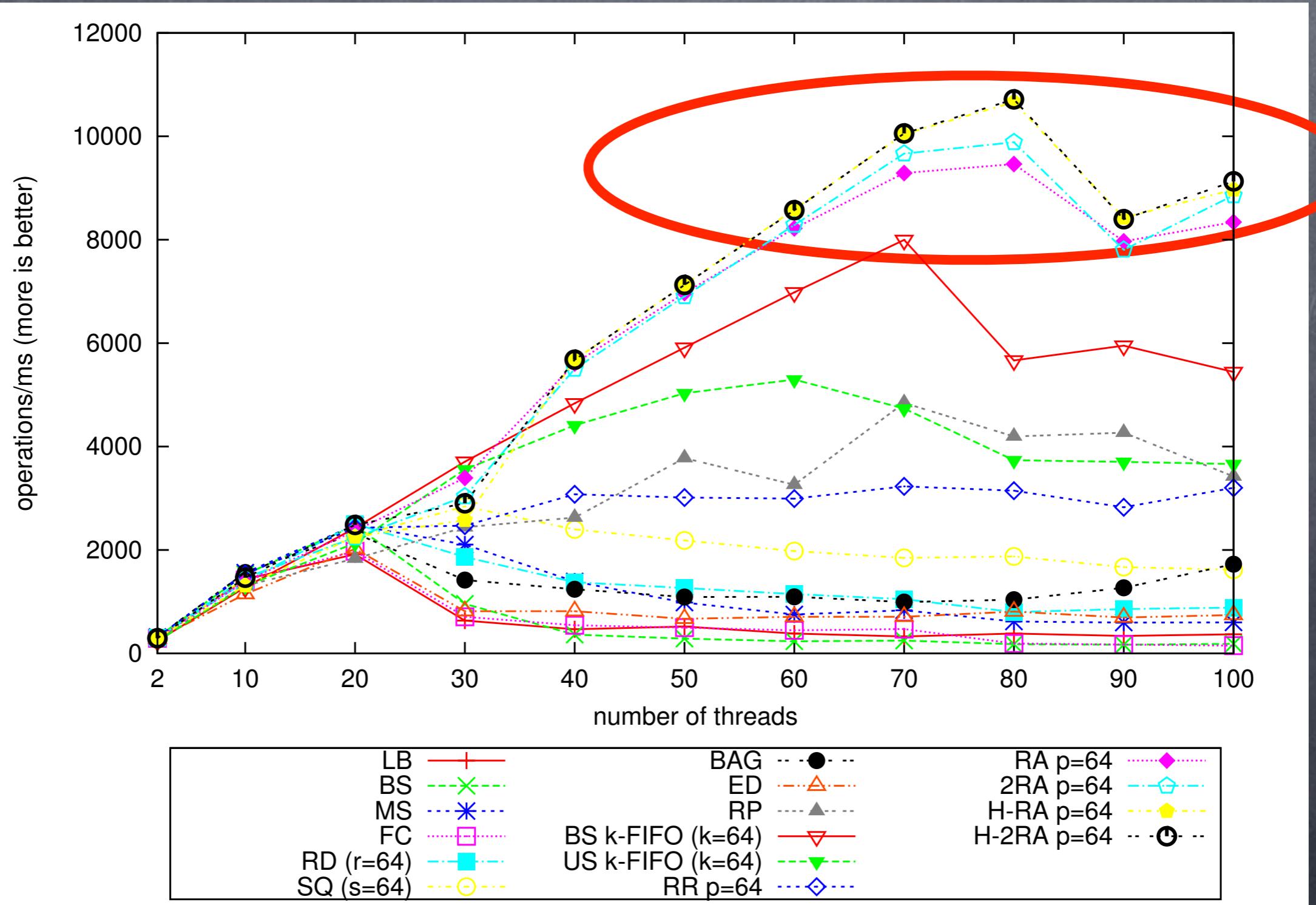
# Concurrent First-in-First-out (FIFO) Queue



[Michael, Scott '96]

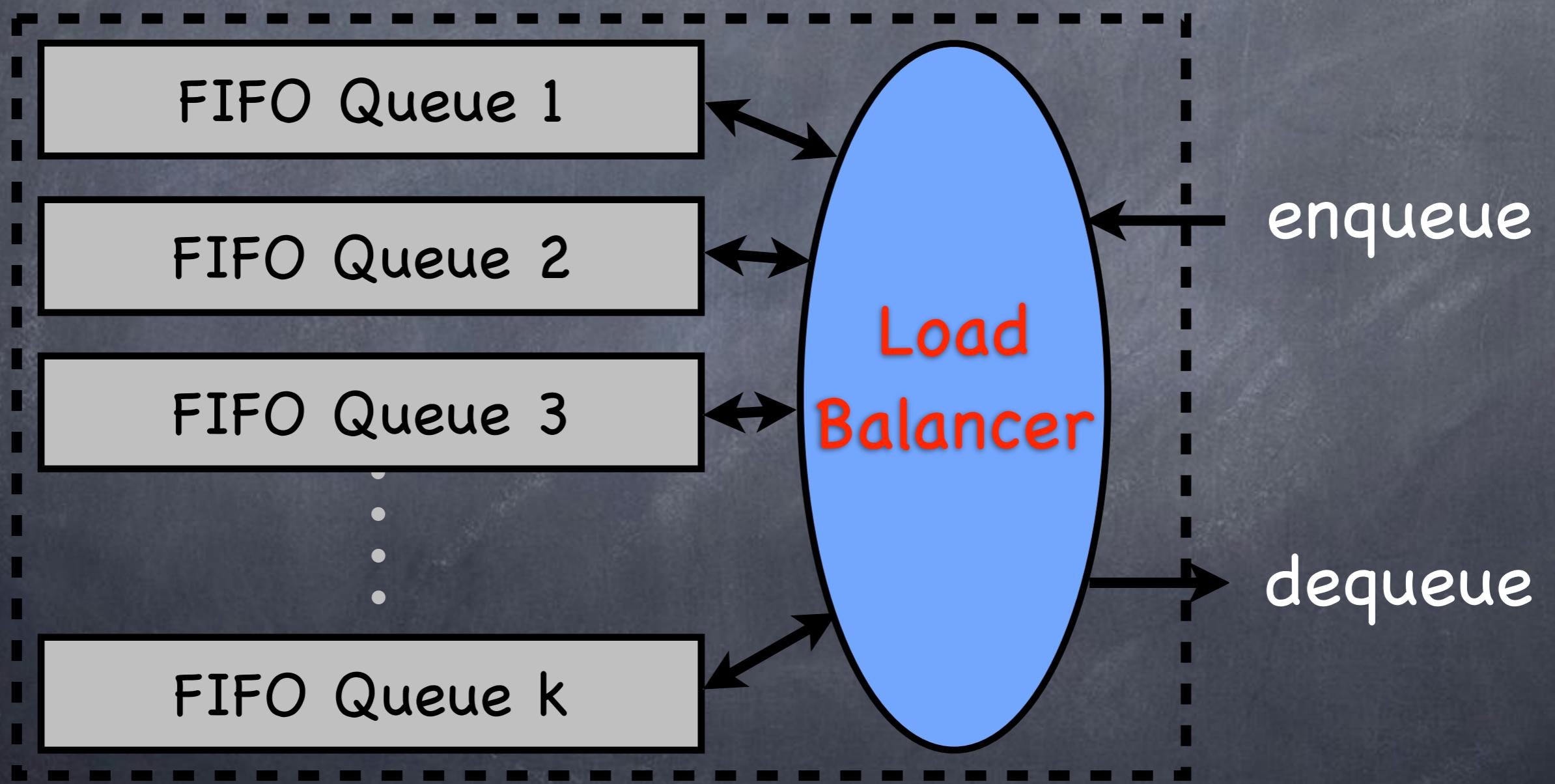
- > 1 lock -> 2 locks -> 0 locks -> compare & swap
- > lock-based vs. **lock-free** vs. wait-free?
- > memory contention on **head** and **tail** pointers!
- > and on **next** pointers!

# Distributed Queues

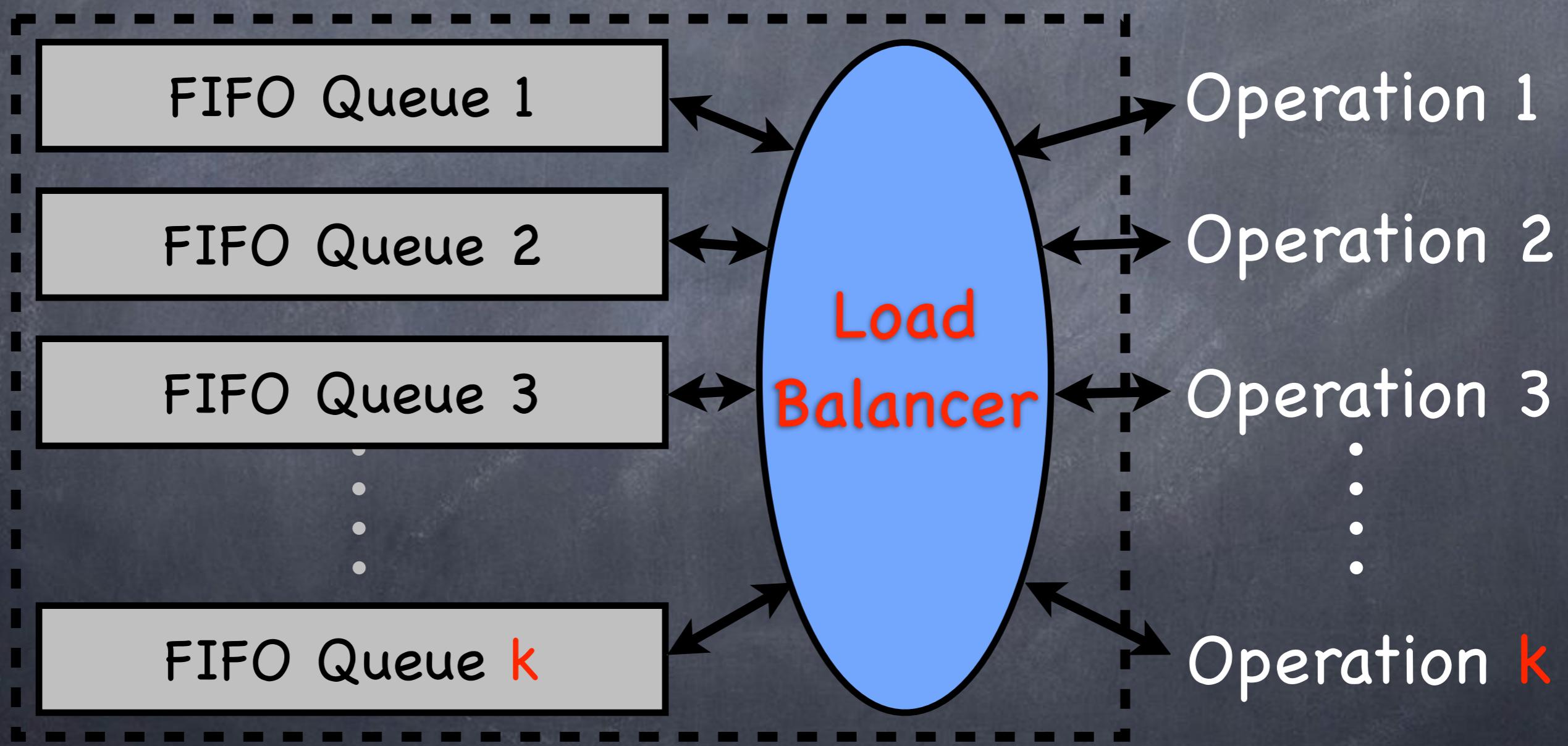


# Distributed Queues

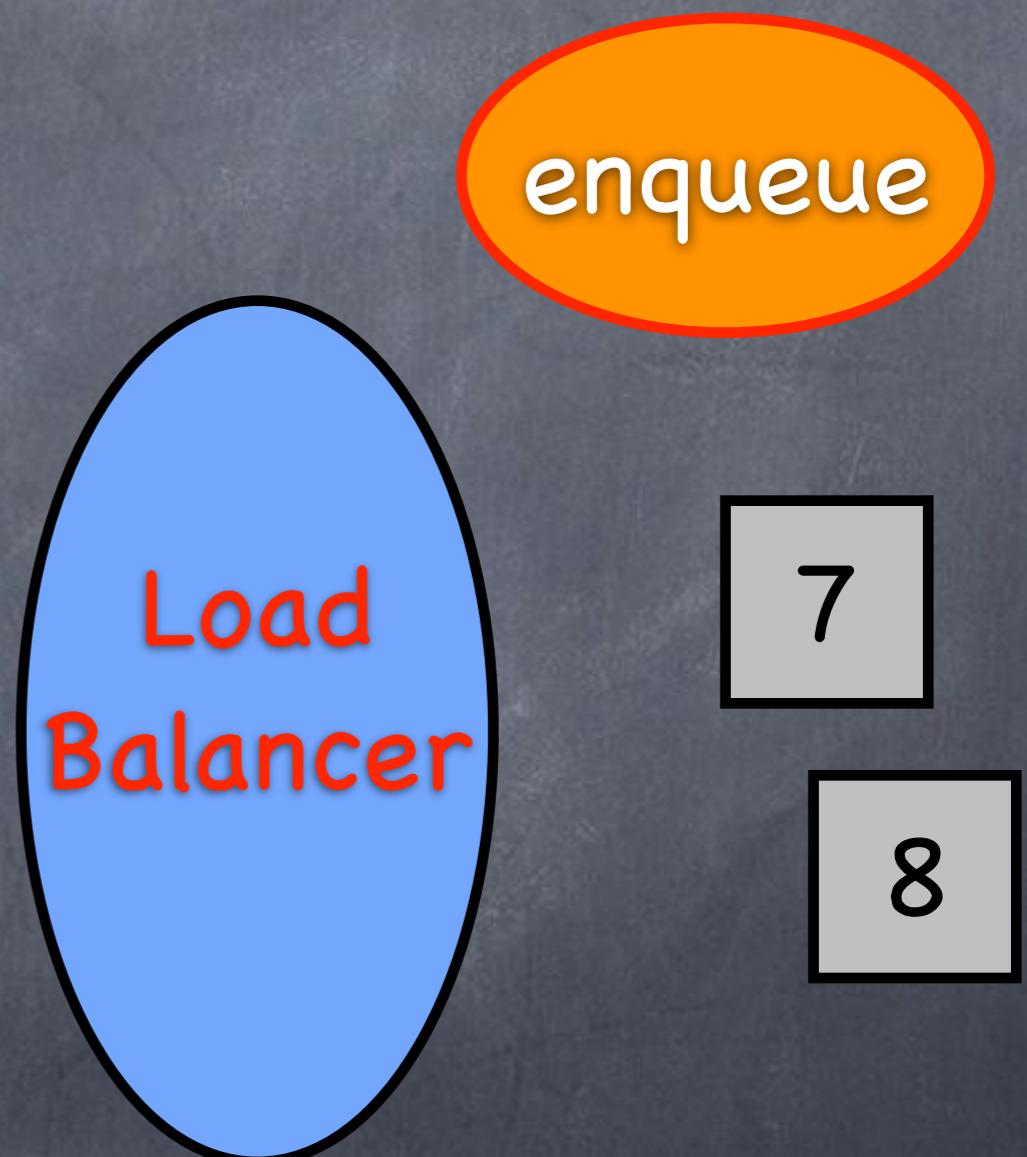
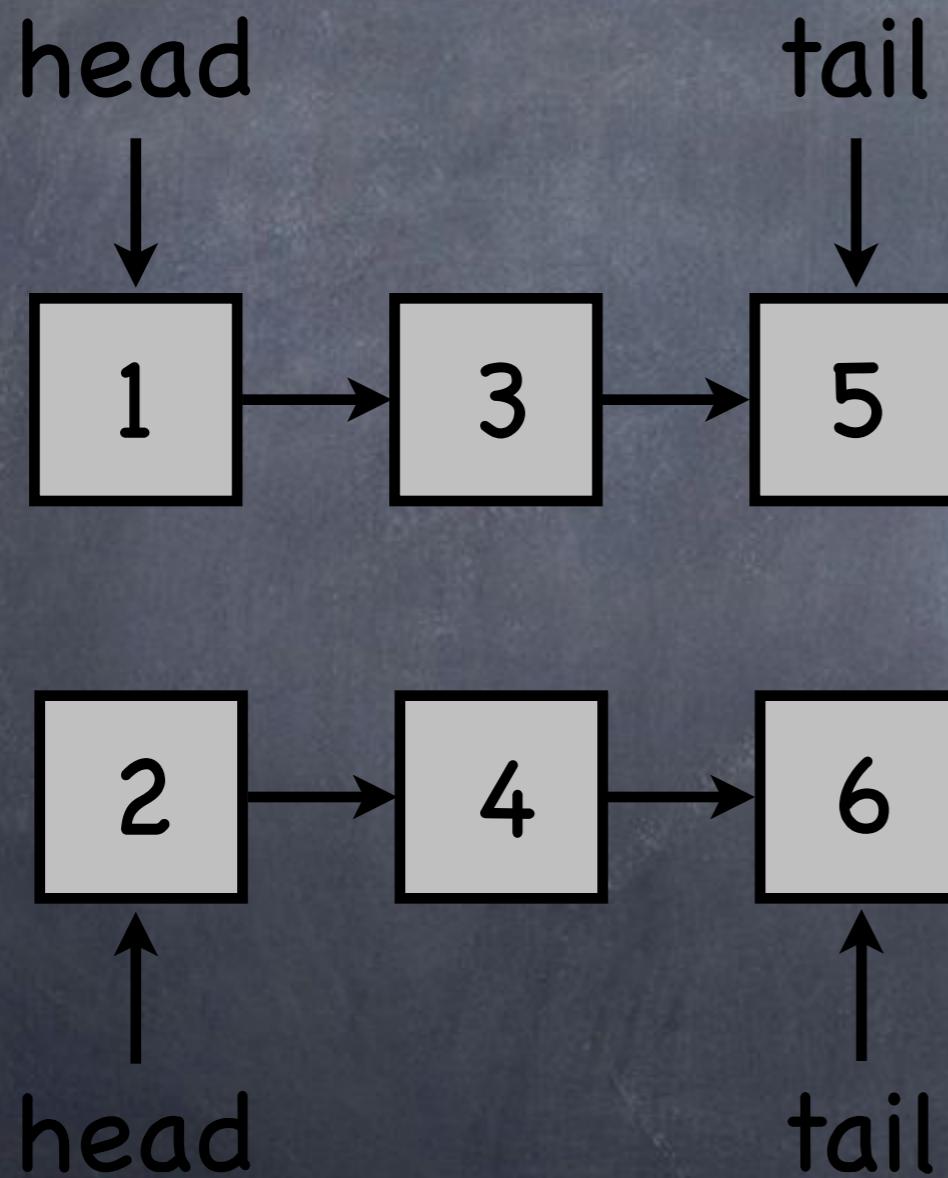
[\_,Payer,Röck,Sokolova'12]



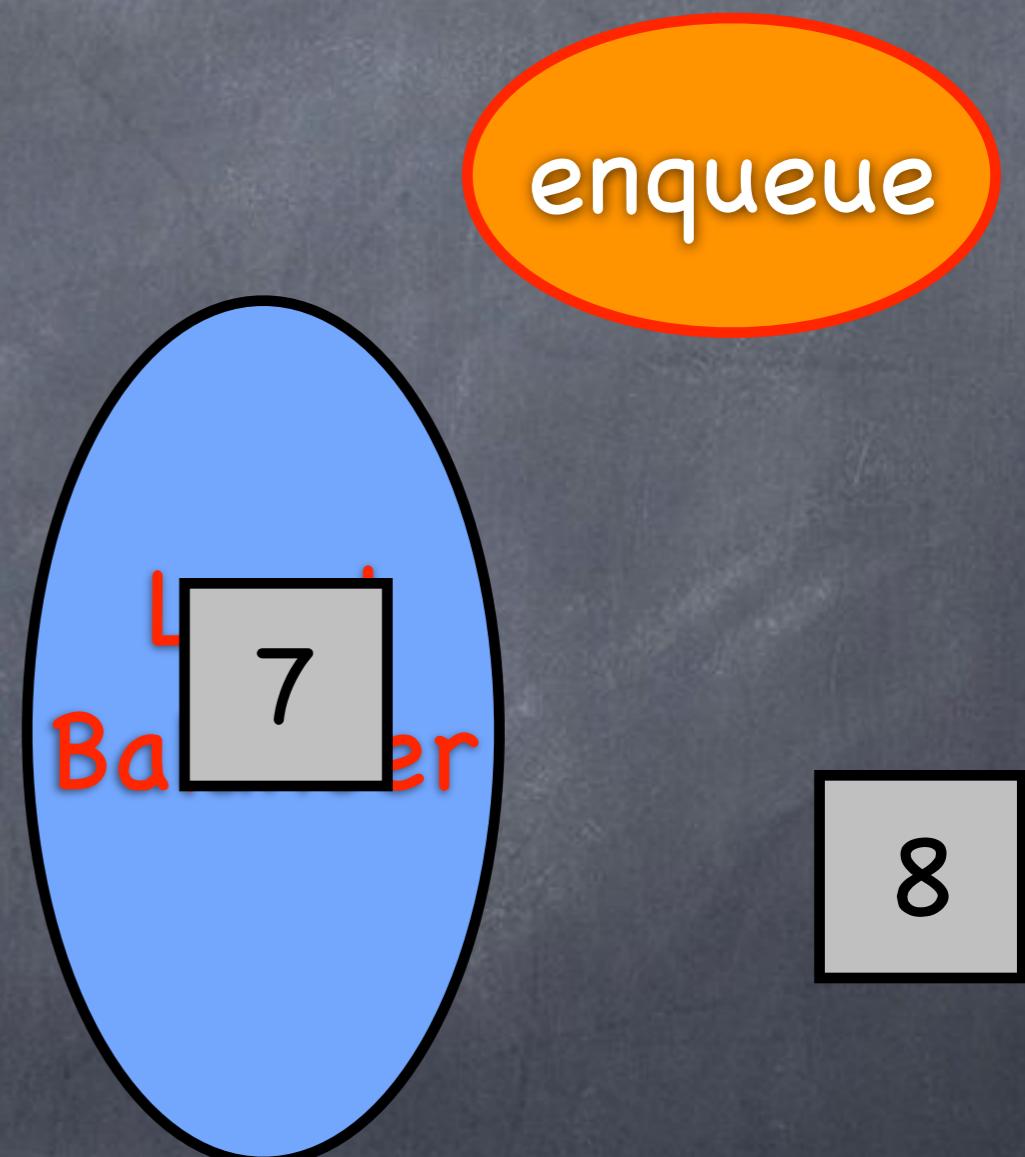
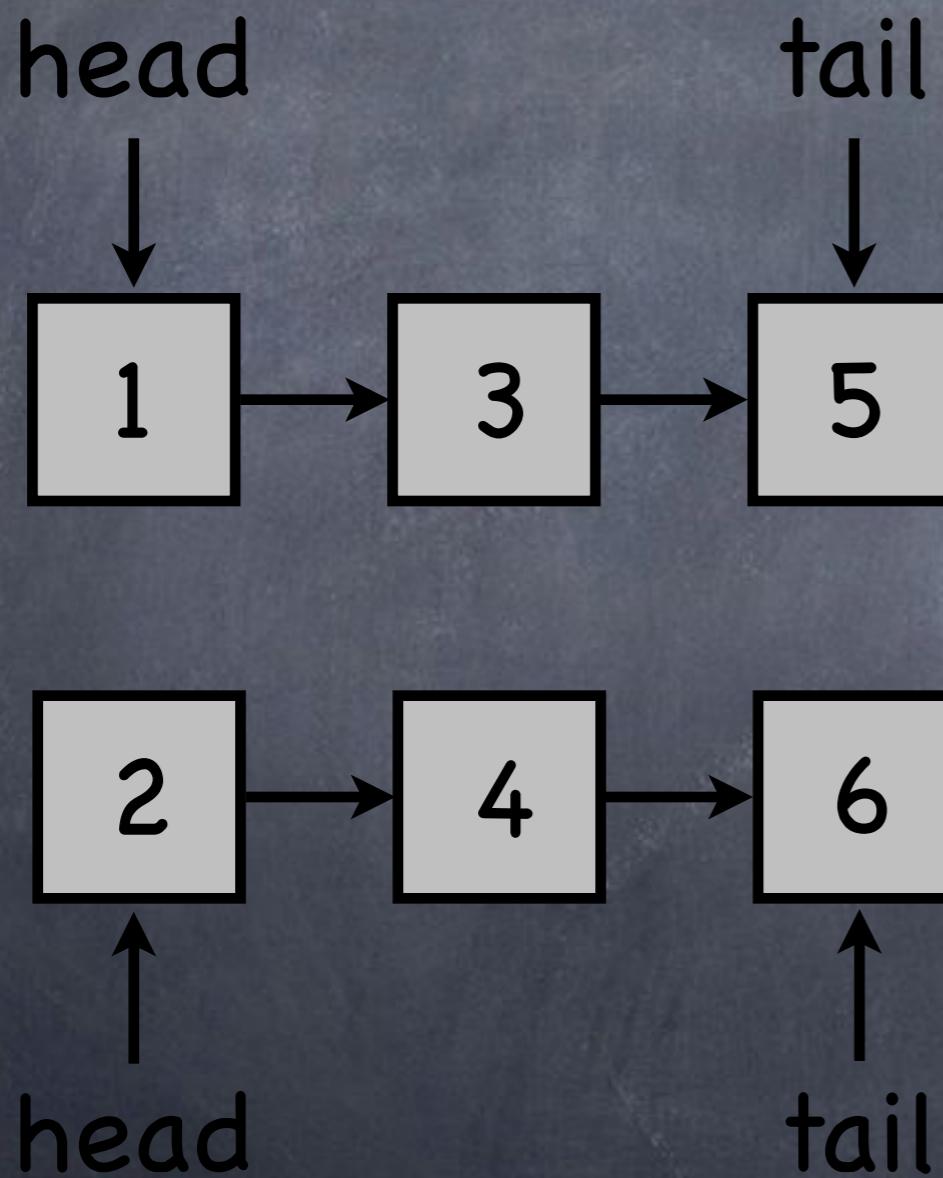
# Up to $k$ Parallel Enqueues and $k$ Parallel Dequeues



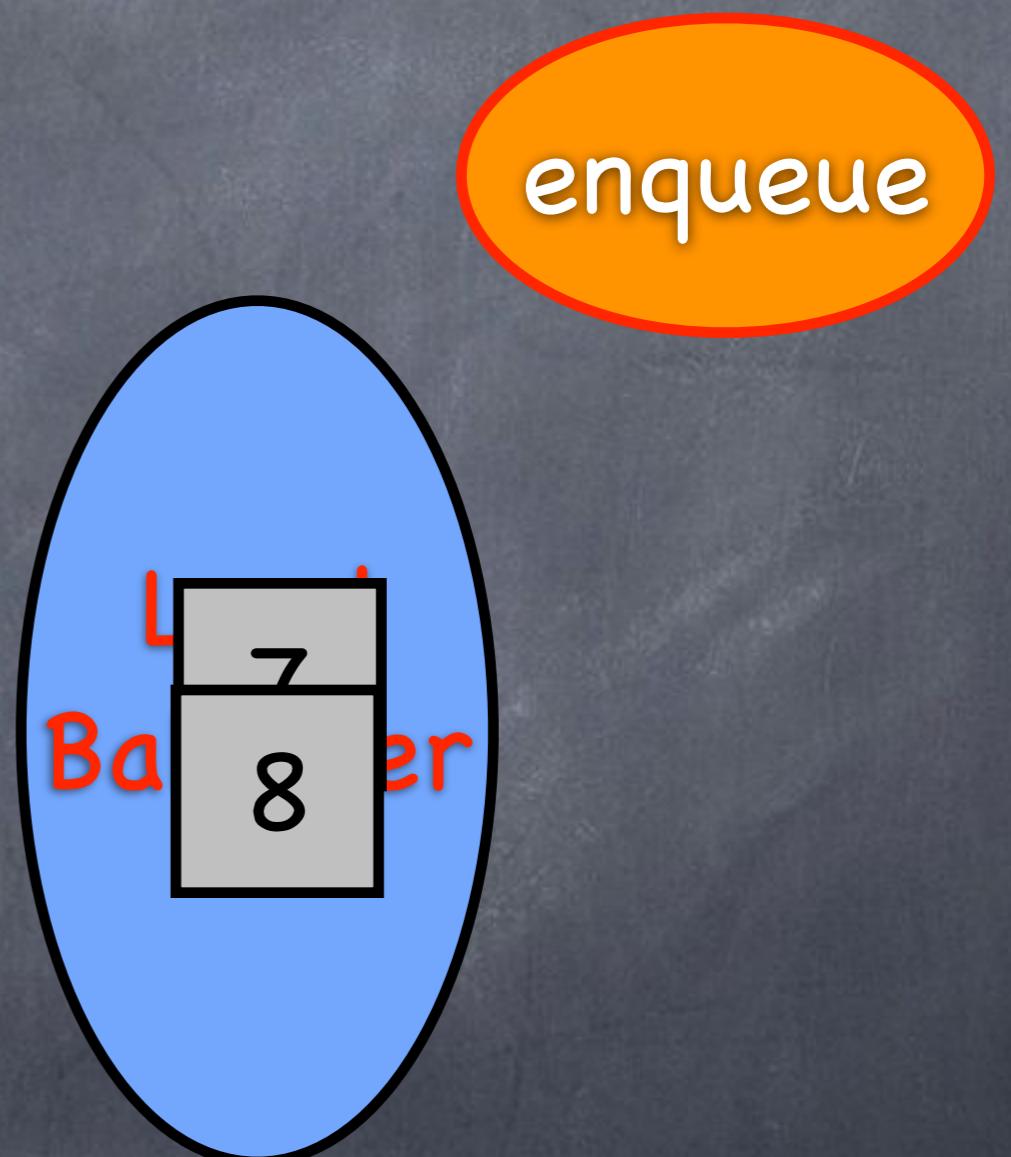
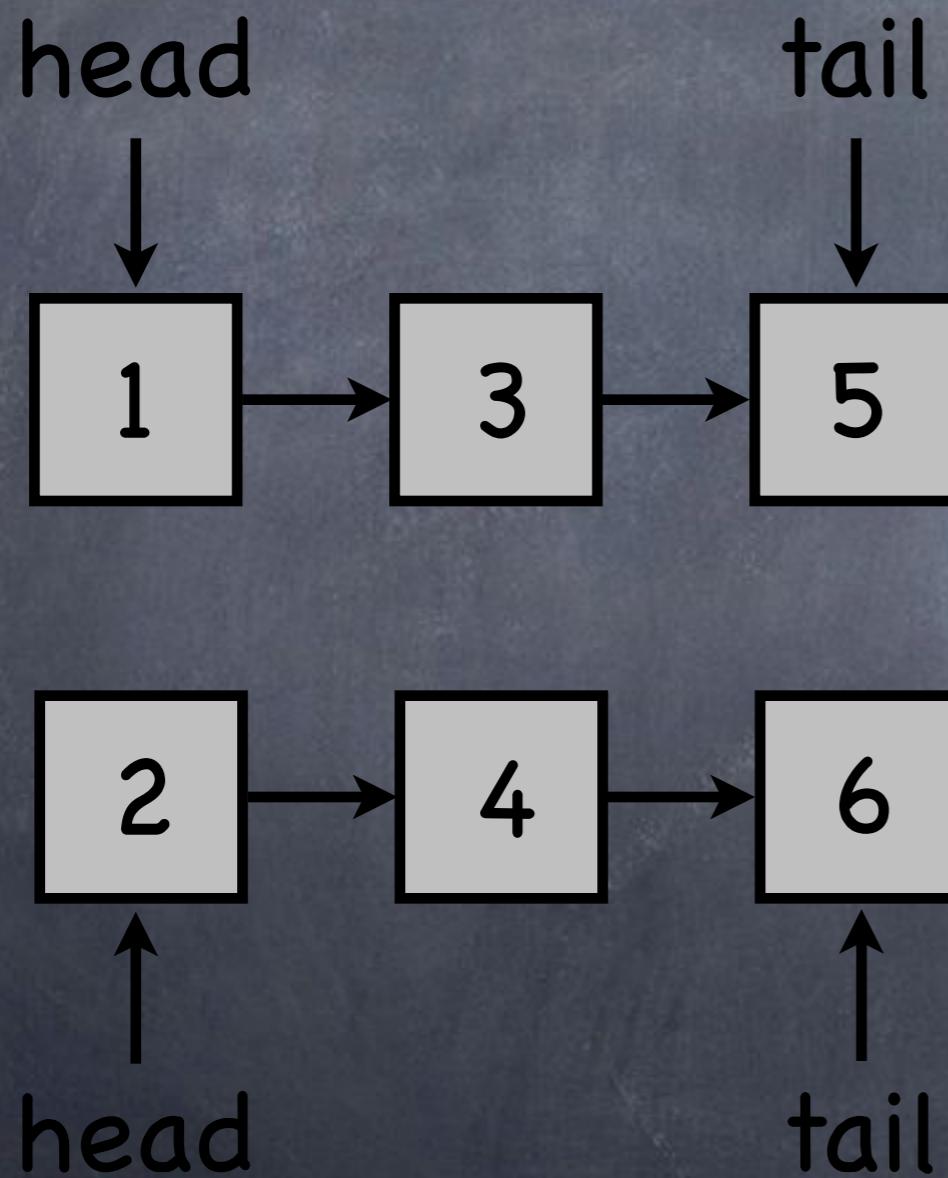
# Load Balancing



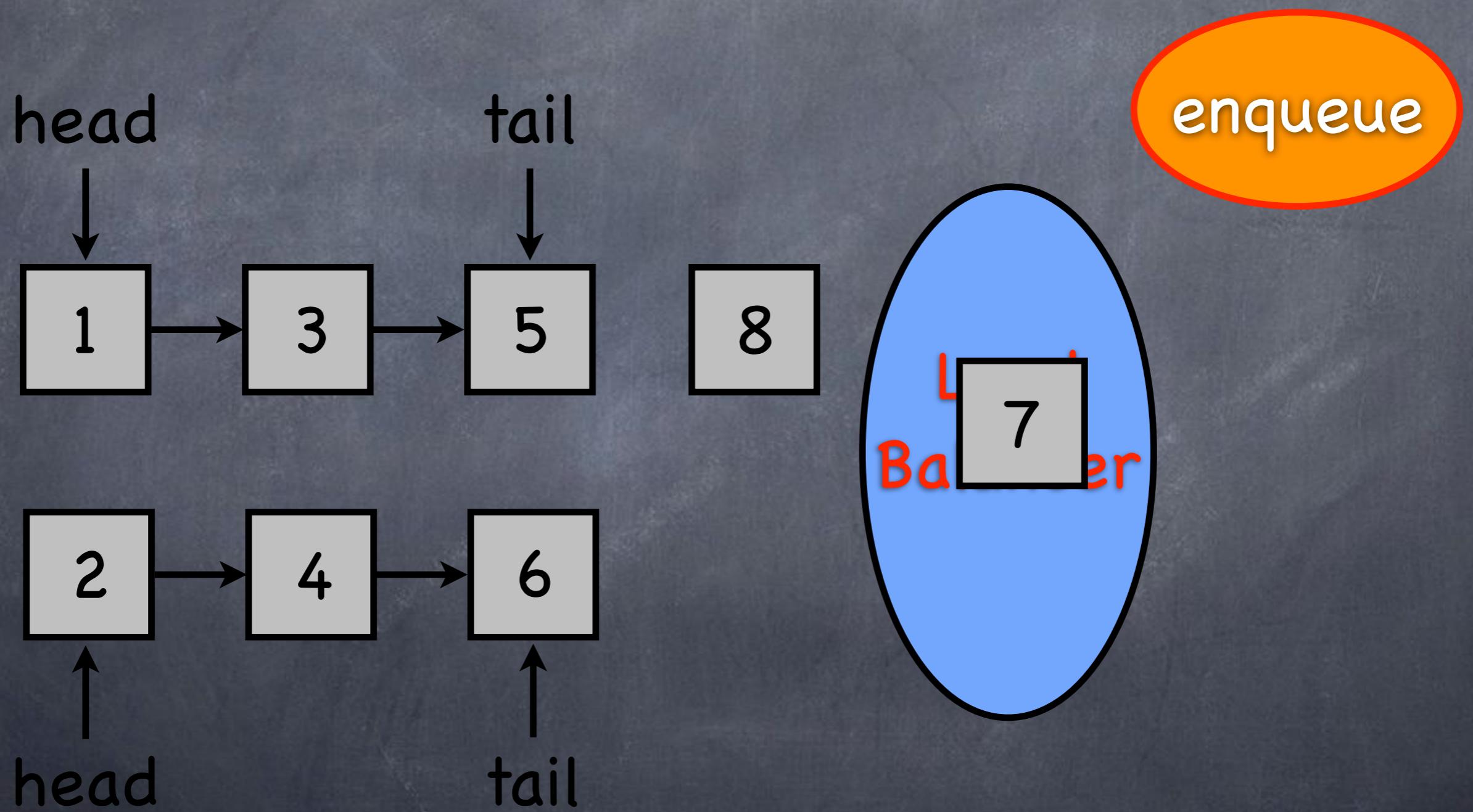
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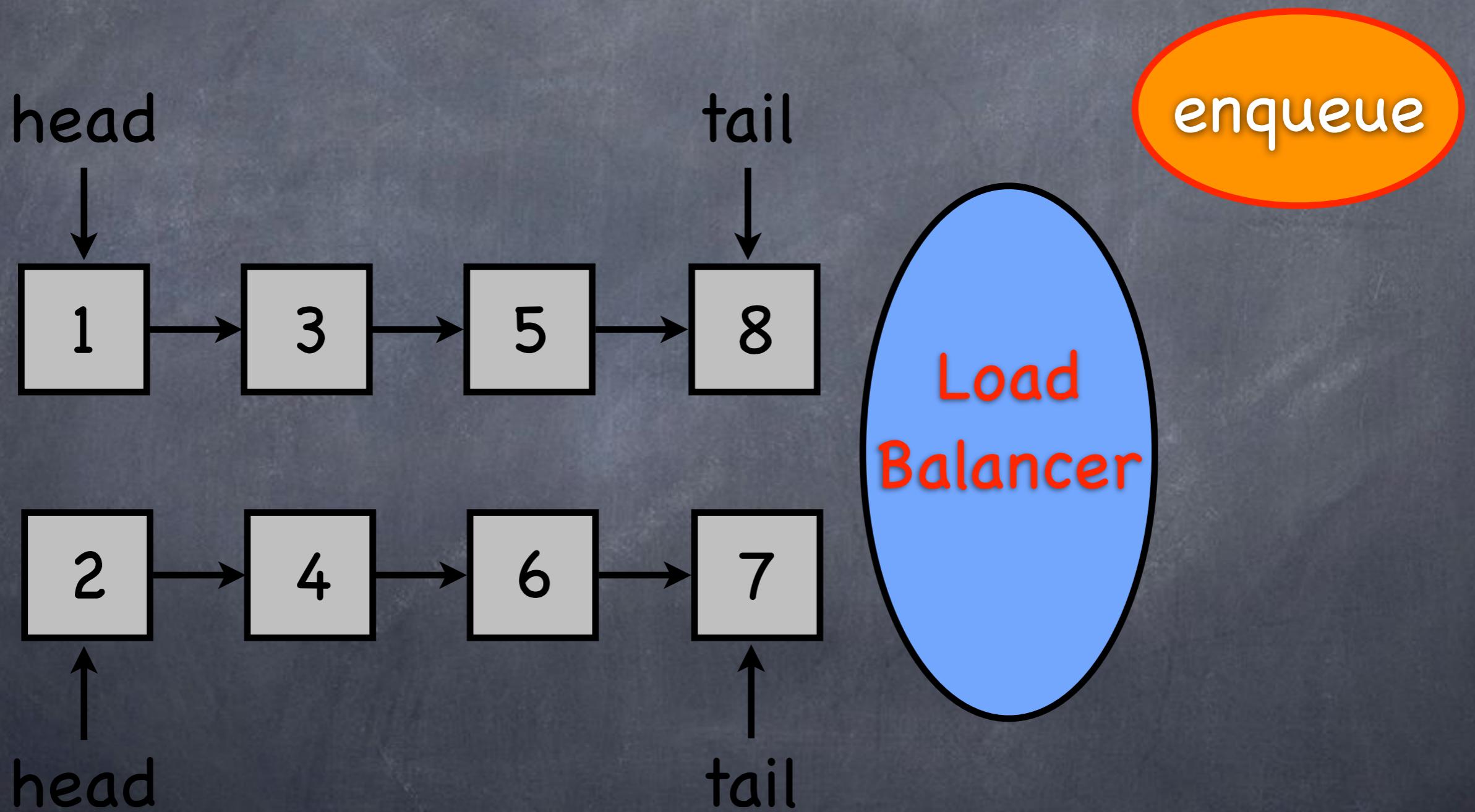
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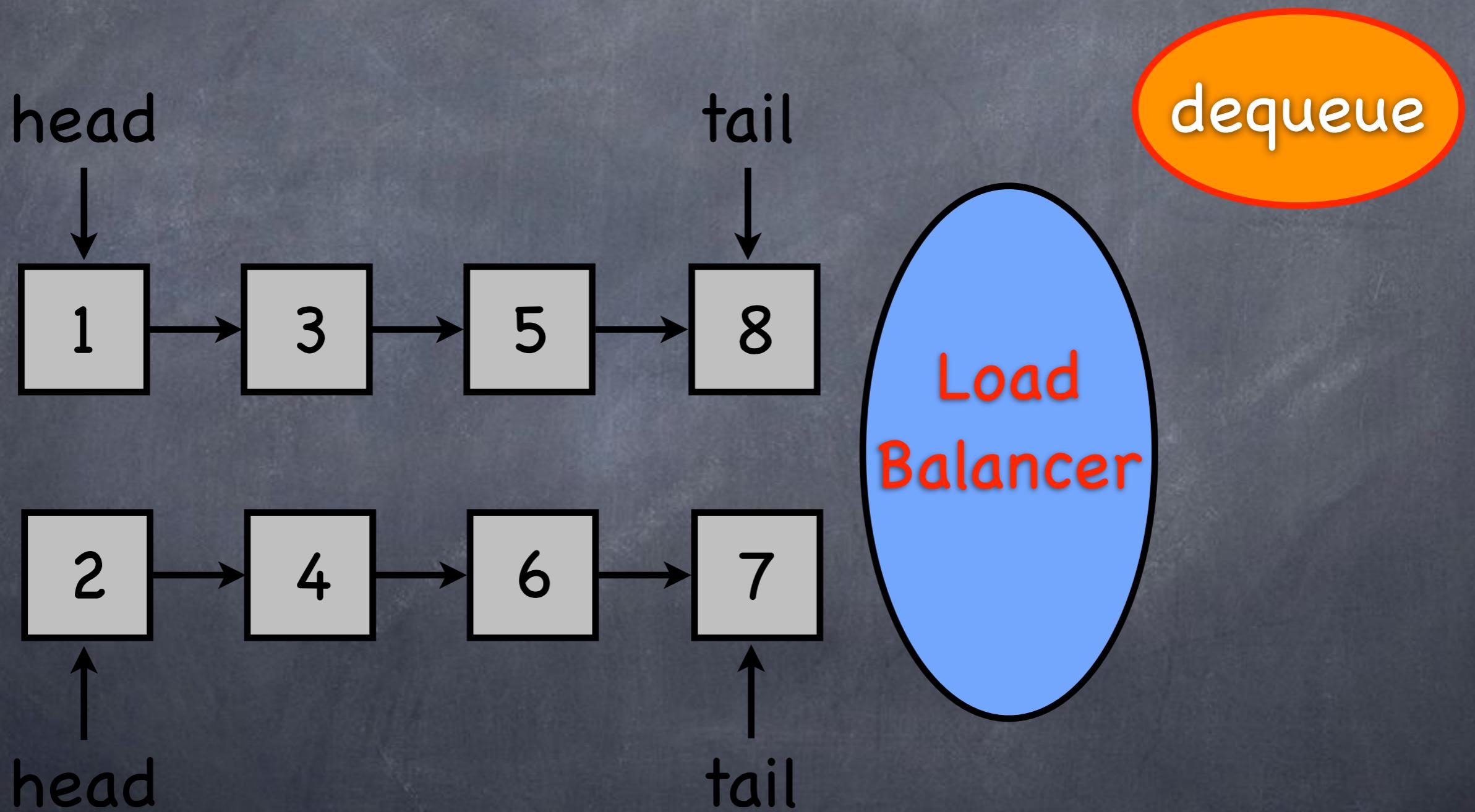
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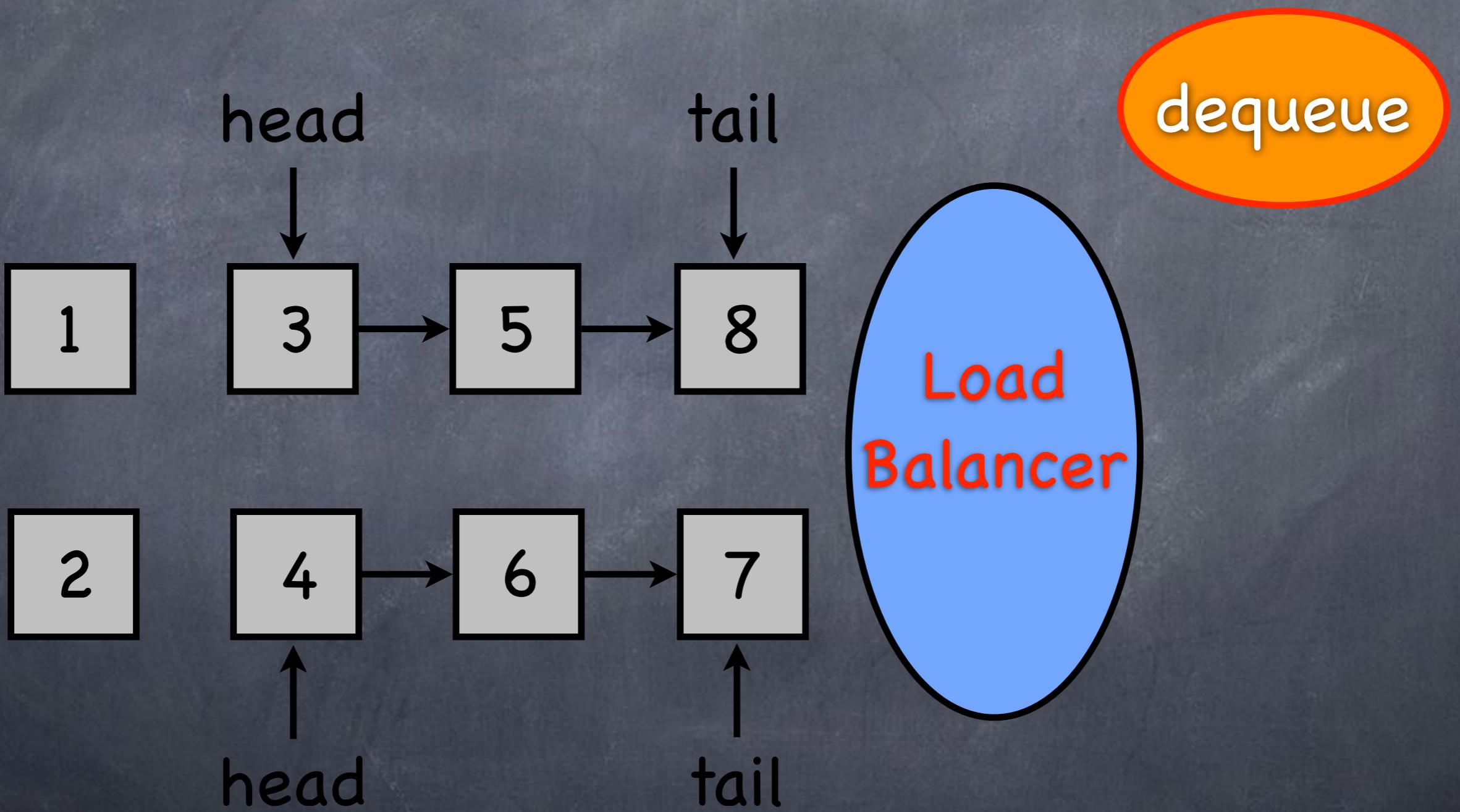
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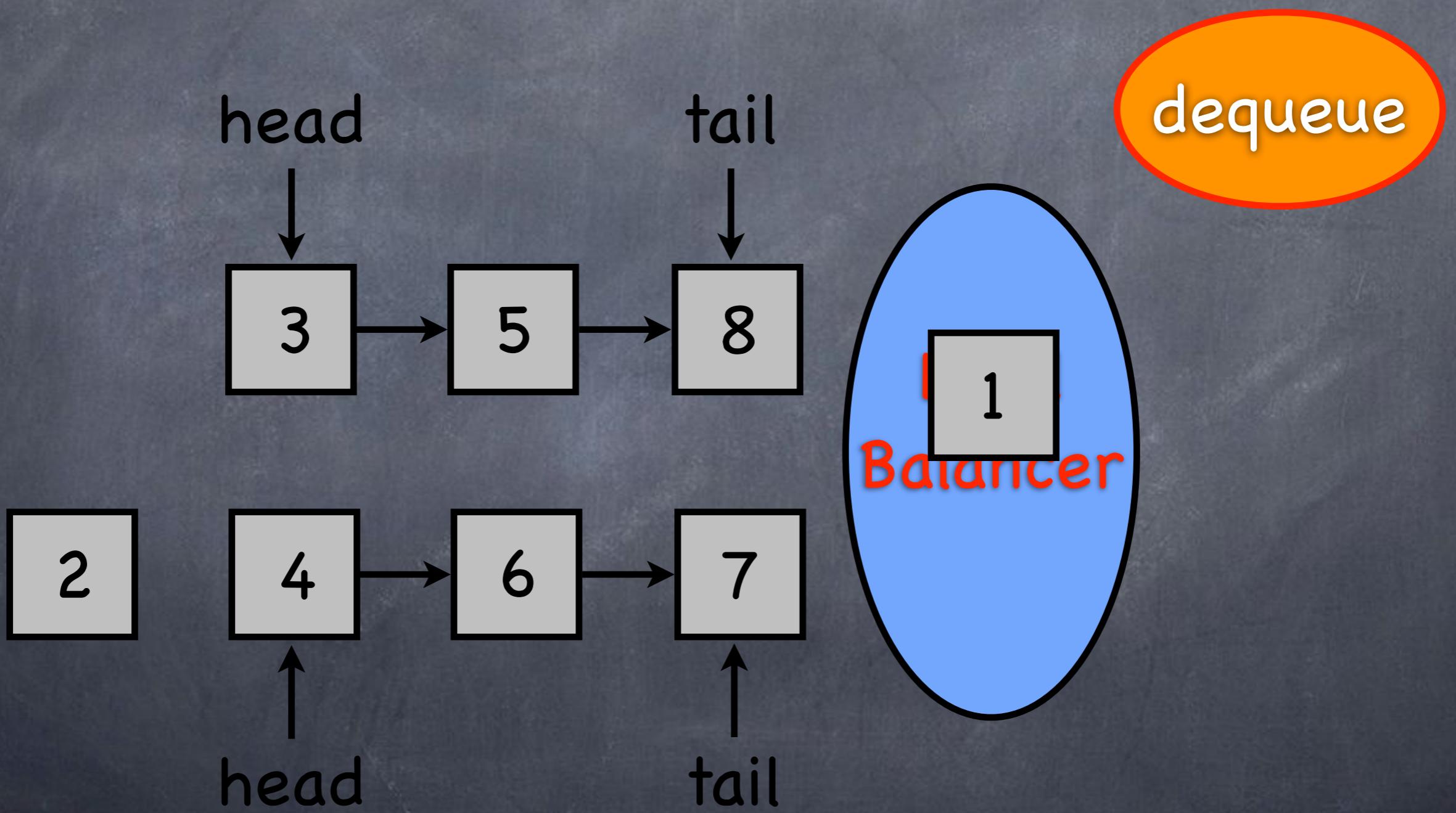
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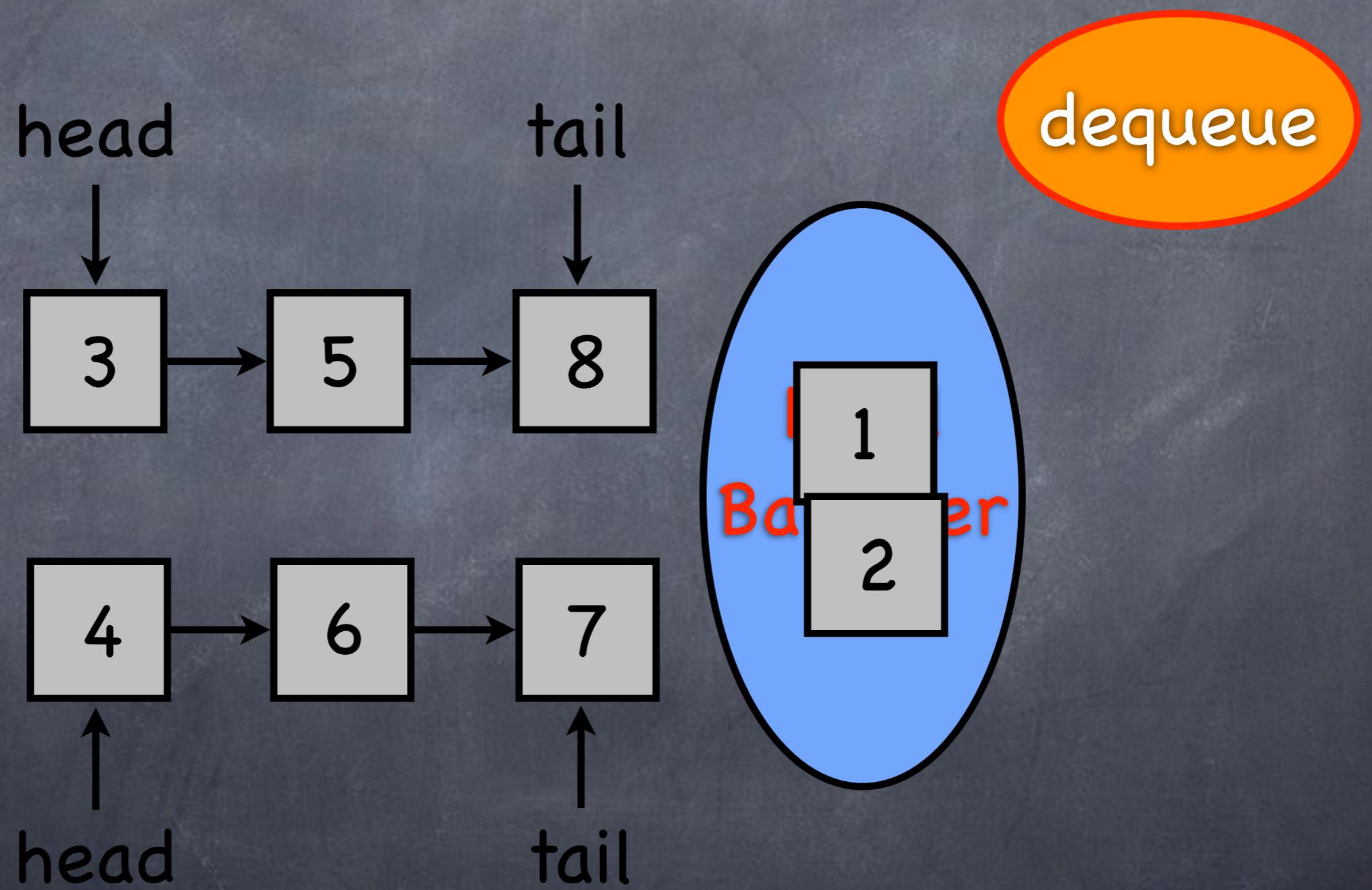
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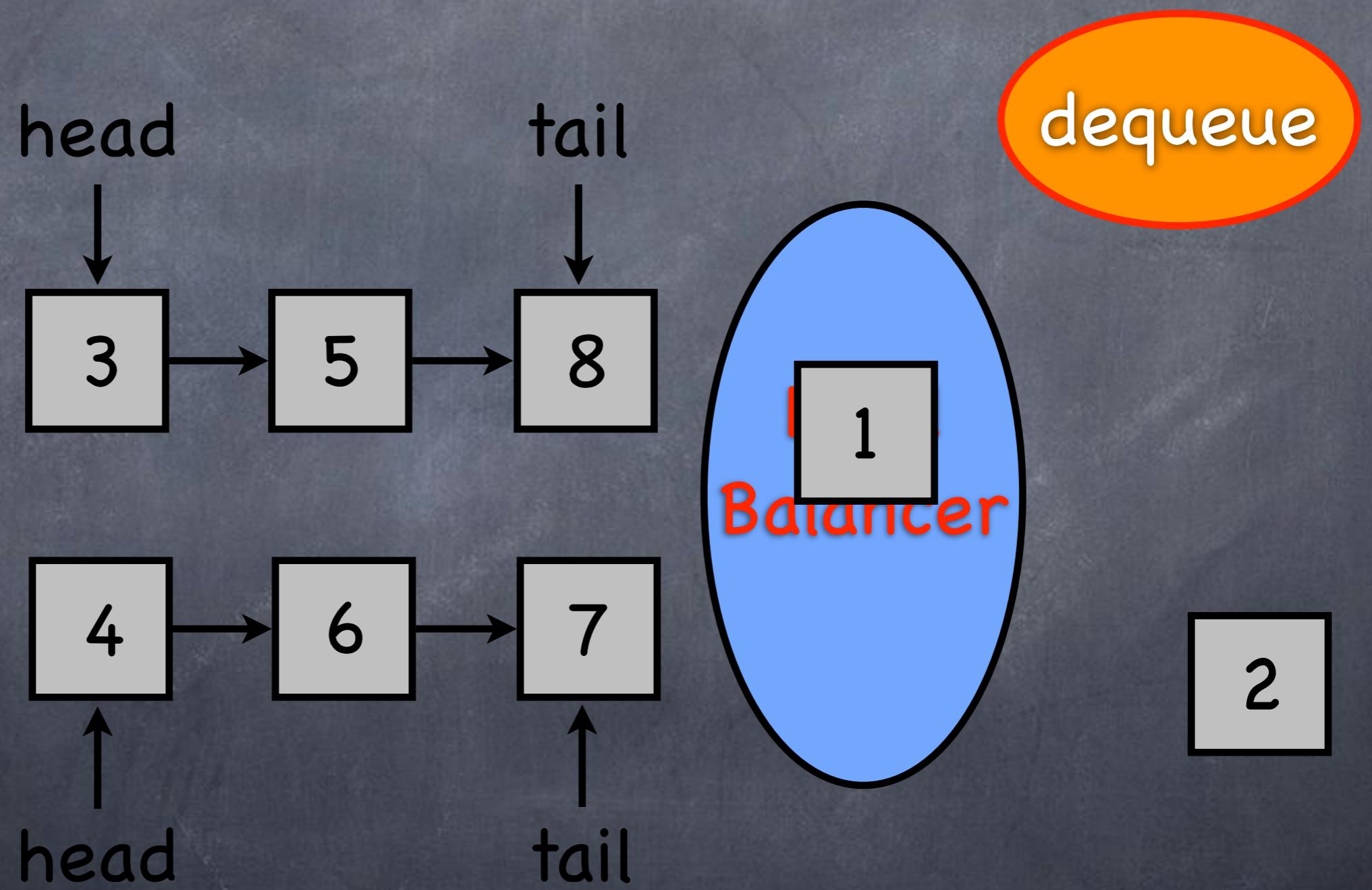
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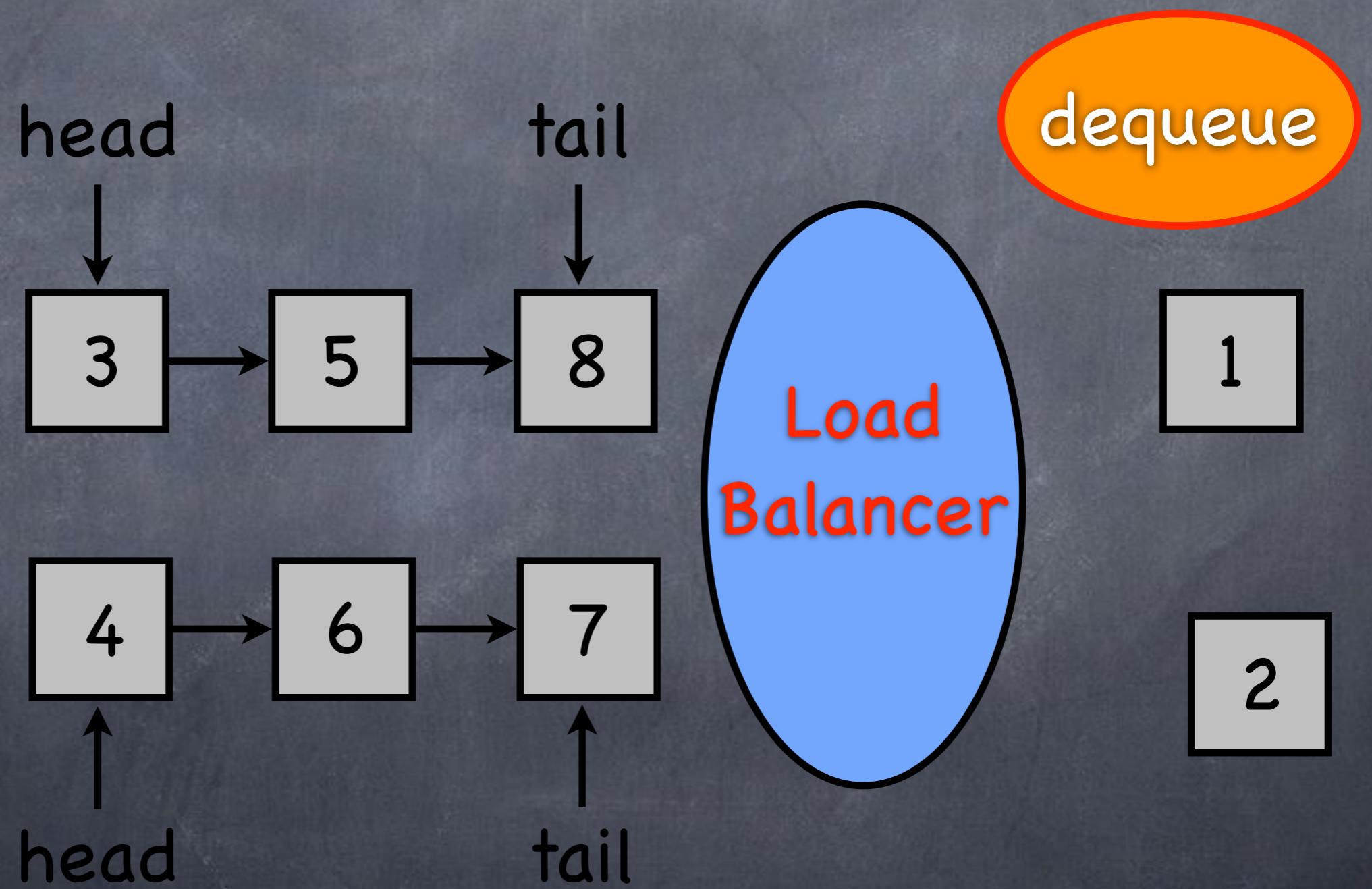
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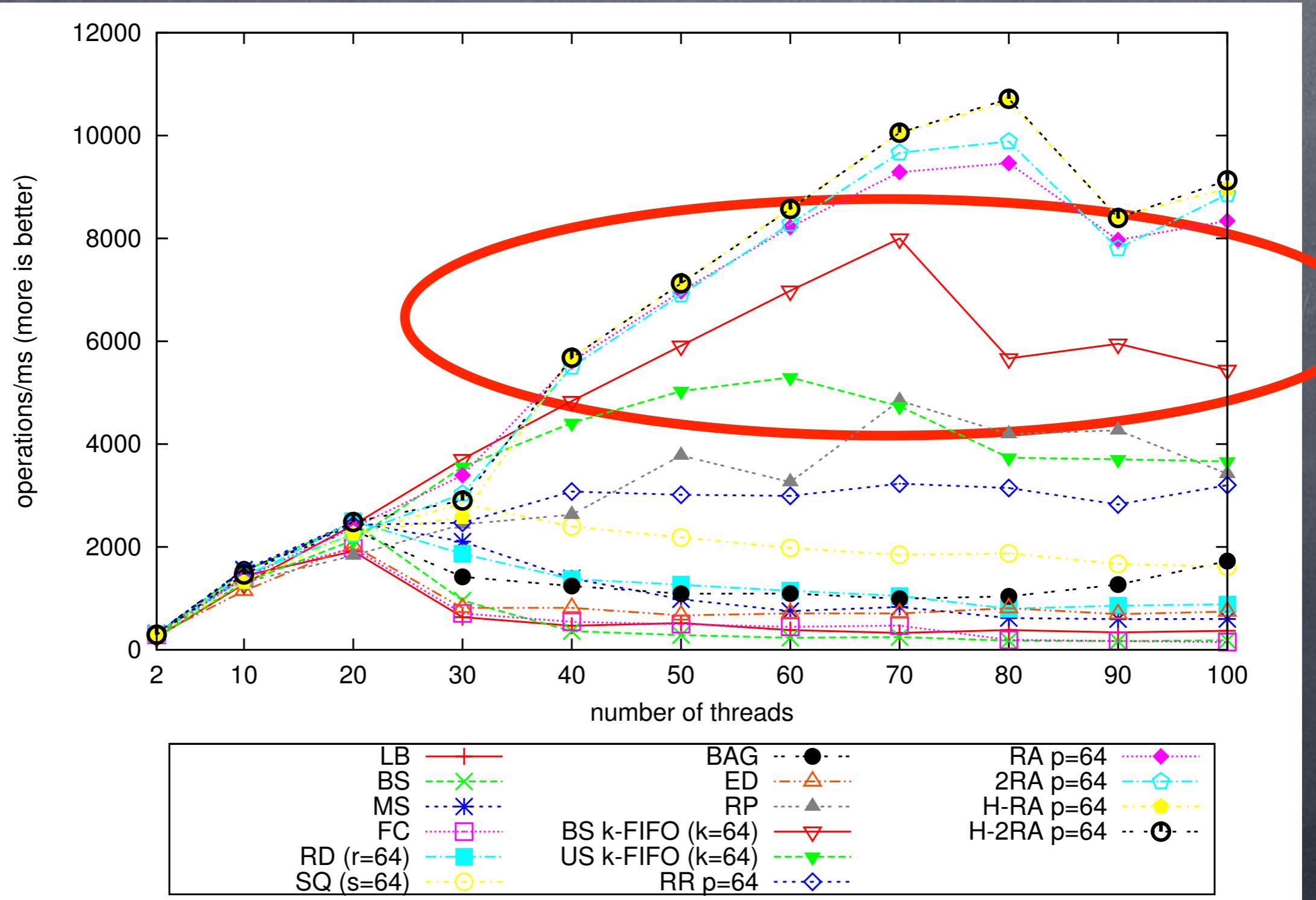


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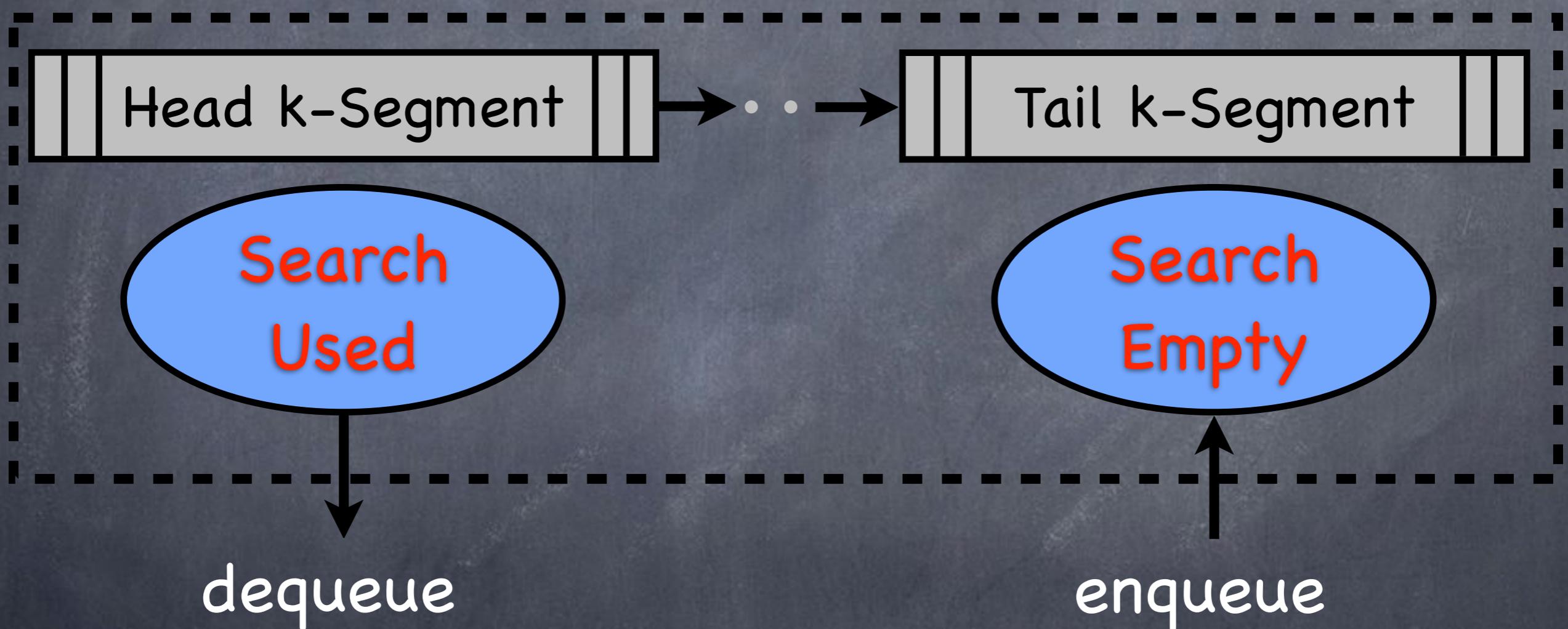
Emptiness  
Check?

# Segmented Queues



# Segmented Queues

[Afek,Korland,Yanovsky'10],[\_,Lippautz,Payer'12]



Emptiness  
Check?

```

1 bool enqueue(item):
2     while true:
3         tail_old = get_tail();
4         head_old = get_head();
5         item_old, index = find_empty_slot(tail_old, k, TESTS);
6         if tail_old == get_tail():
7             if item_old.value == EMPTY:
8                 item_new = atomic_value(item, item_old.counter + 1);
9                 if CAS(&tail_old[index], item_old, item_new):
10                     if committed(tail_old, item_new, index):
11                         return true;
12             else:
13                 if queue_full(head_old, tail_old):
14                     if segment_not_empty(head_old, k) && head == get_head():
15                         return false;
16                     advance_head(head_old, k);
17                     advance_tail(tail_old, k);
18
19 bool committed(tail_old, item_new, index):
20     if tail_old[index] != item_new:
21         return true;
22     head_current = get_head();
23     tail_current = get_tail();
24     item_empty = atomic_value(EMPTY, item_new.counter + 1);
25     if in_queue_after_head(tail_old, tail_current, head_current):
26         return true;
27     else if not_in_queue(tail_old, tail_current, head_current):
28         if !CAS(&tail_old[index], item_new, item_empty):
29             return true;
30     else: //in queue at head
31         head_new = atomic_value(head_current.value, head_current.counter + 1);
32         if CAS(&head, head_current, head_new):
33             return true;
34         if !CAS(&tail_old[index], item_new, item_empty):
35             return true;
36     return false;

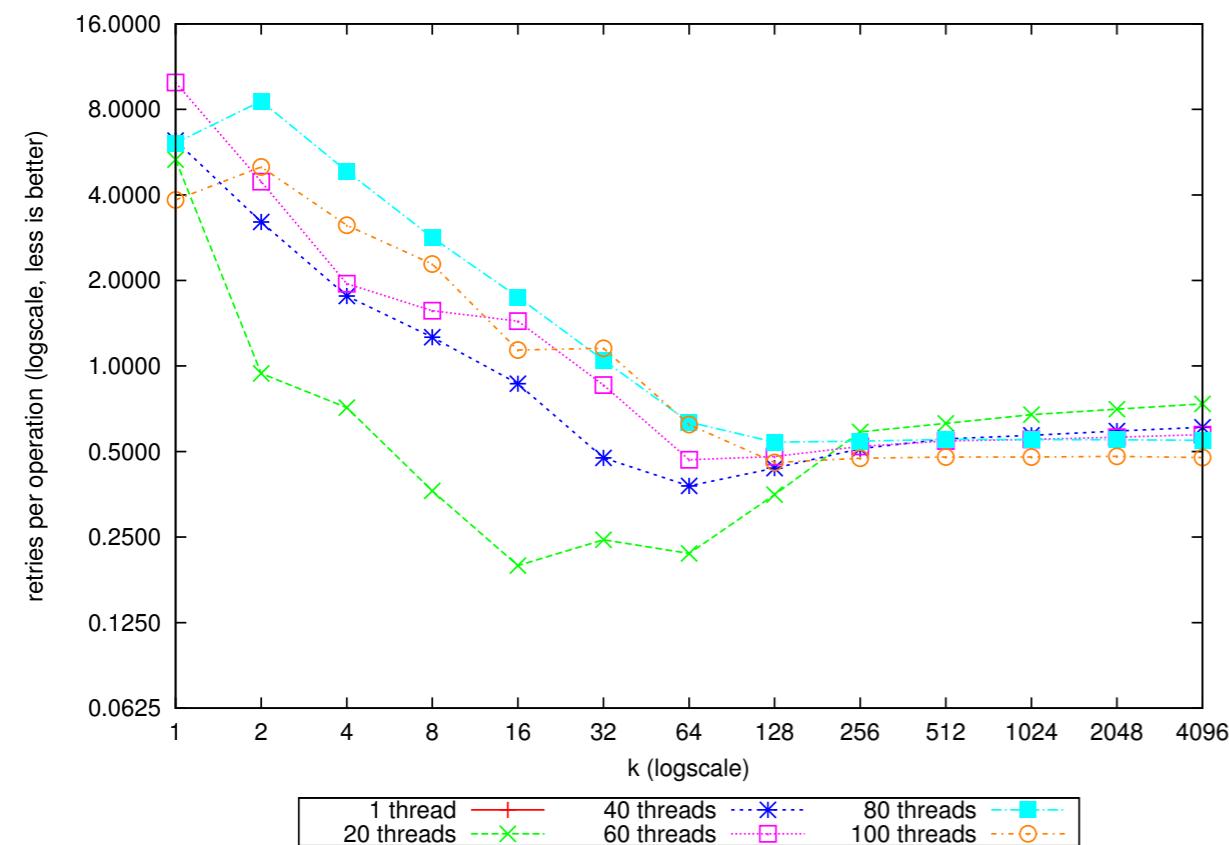
```

enqueue

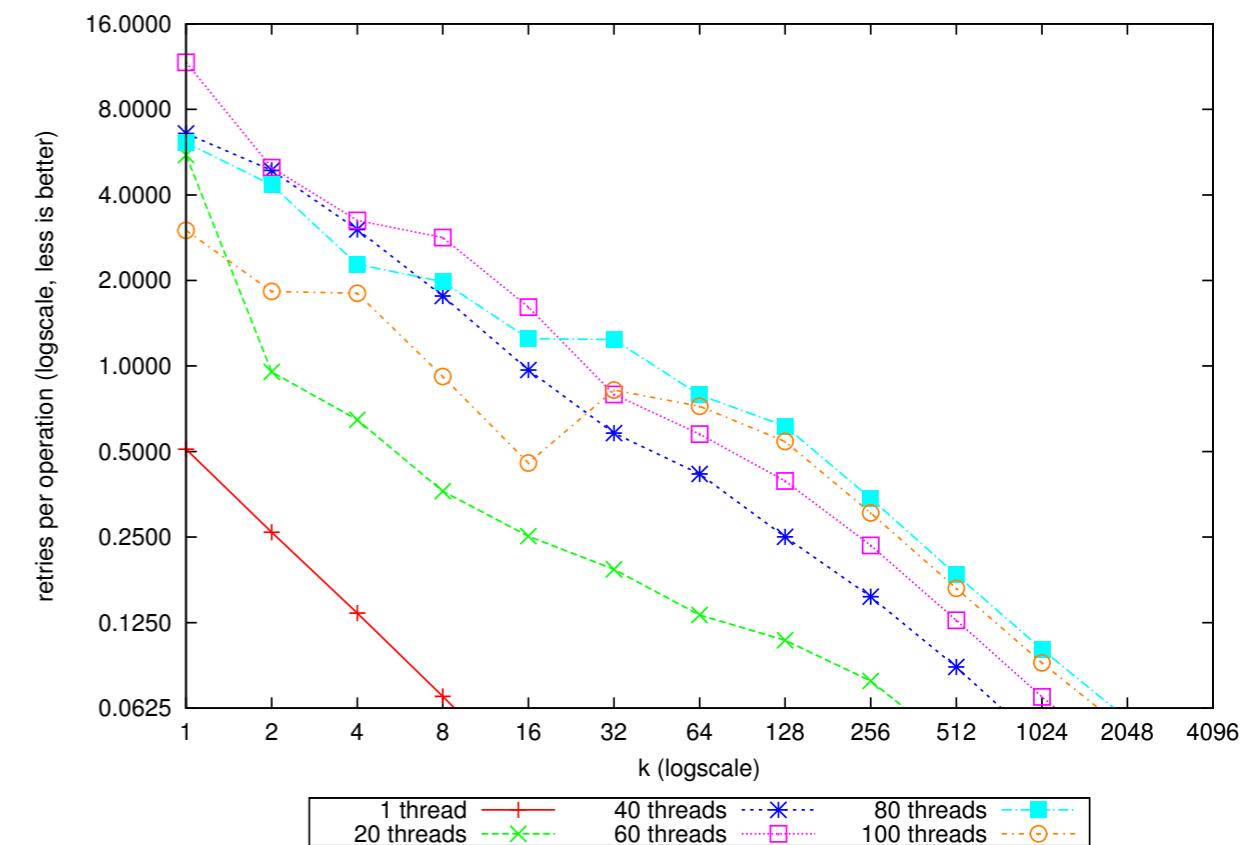
# dequeue

```
38 item dequeue():
39     while true:
40         tail_old = get_tail();
41         head_old = get_head();
42         item_old, index = find_item(head_old, k);
43         if head_old == head:
44             if item_old.value != EMPTY:
45                 if head_old.value == tail_old.value:
46                     advance_tail(tail_old, k);
47                     item_empty = atomic_value(EMPTY, item_old.counter + 1);
48                     if CAS(&head_old[index], item_old, item_empty):
49                         return item_old.value;
50             else:
51                 if head_old.value == tail_old.value && tail_old == get_tail():
52                     return null;
53                     advance_head(head_old, k);
```

# CAS Retries/Operation in $k$

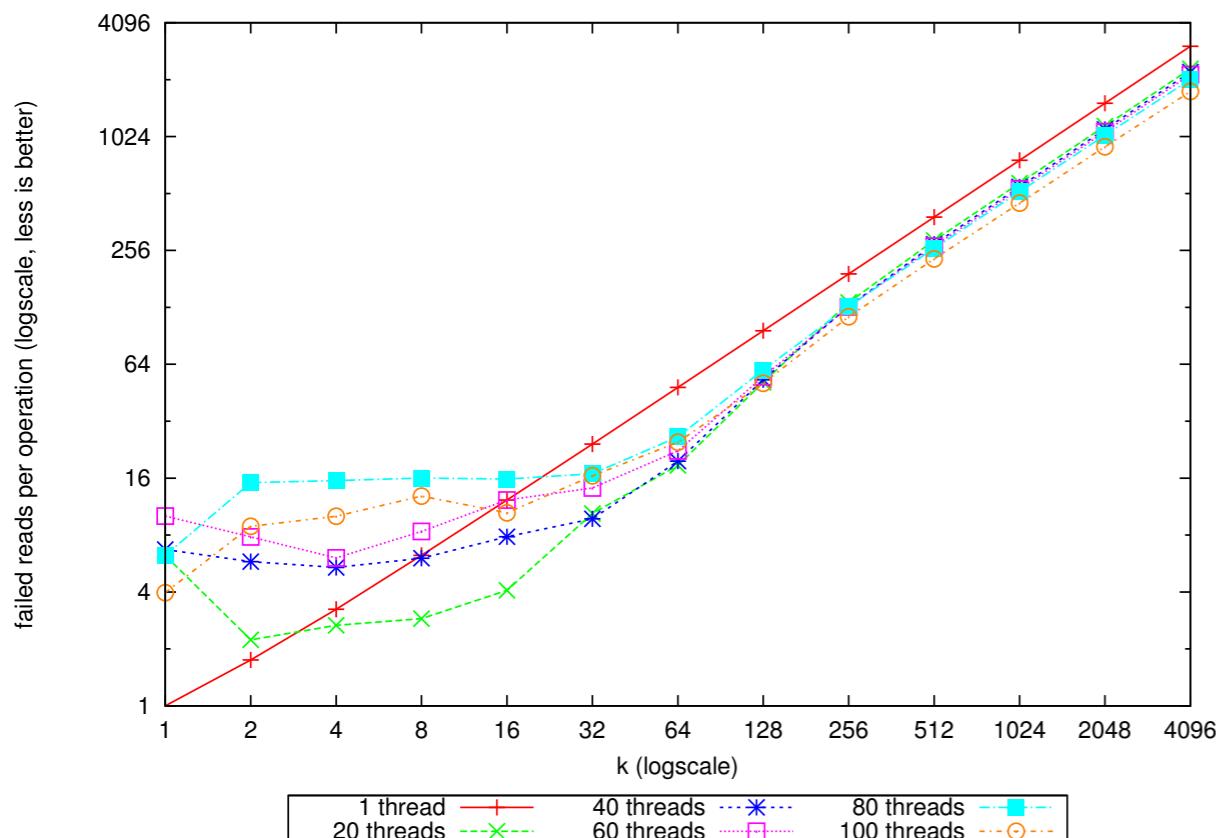


(a) BS number of retries per operation of very high contention producer-consumer benchmark ( $c = 1000, i = 0$ )

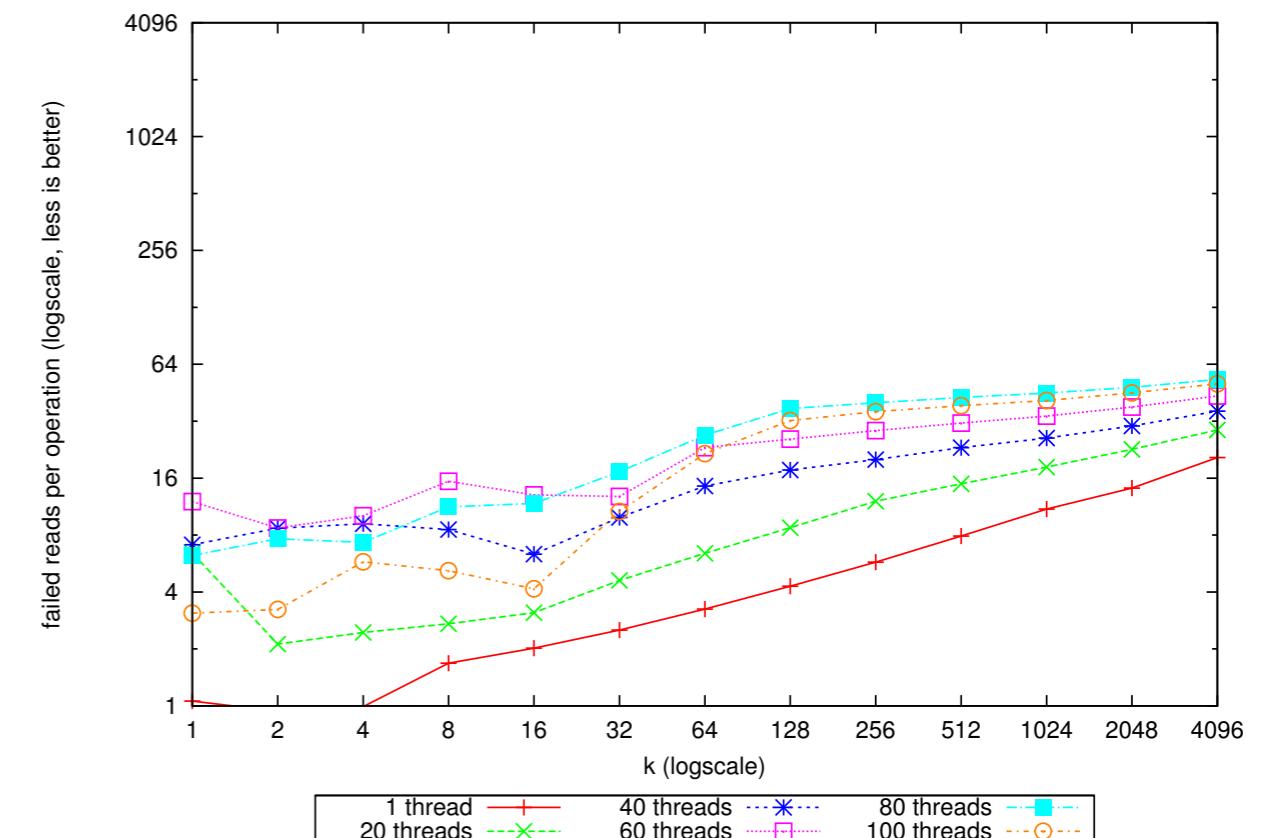


(b) BS number of retries per operation of very high contention producer-consumer benchmark ( $c = 1000, i = 5000$ )

# Failed Searches/Operation in $\kappa$

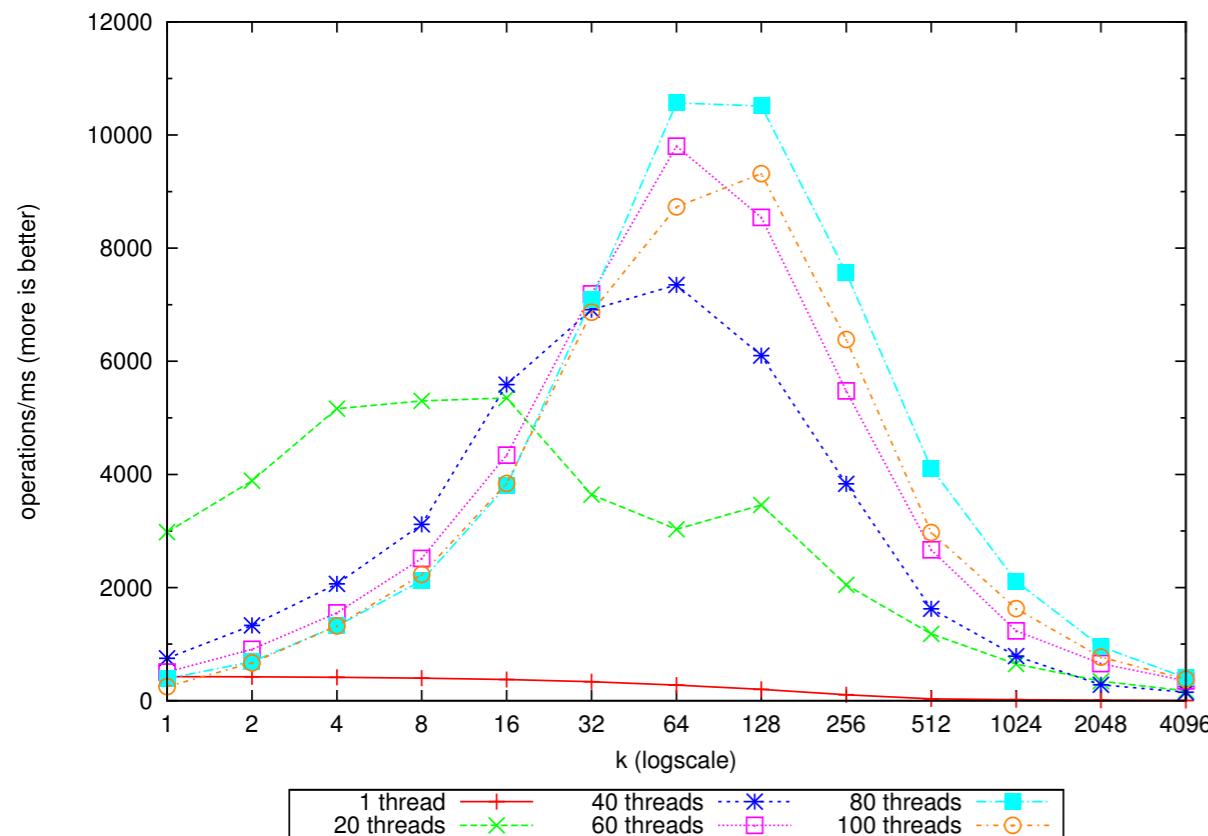


(c) BS number of failed reads per operation of very high contention producer-consumer benchmark ( $c = 1000, i = 0$ )

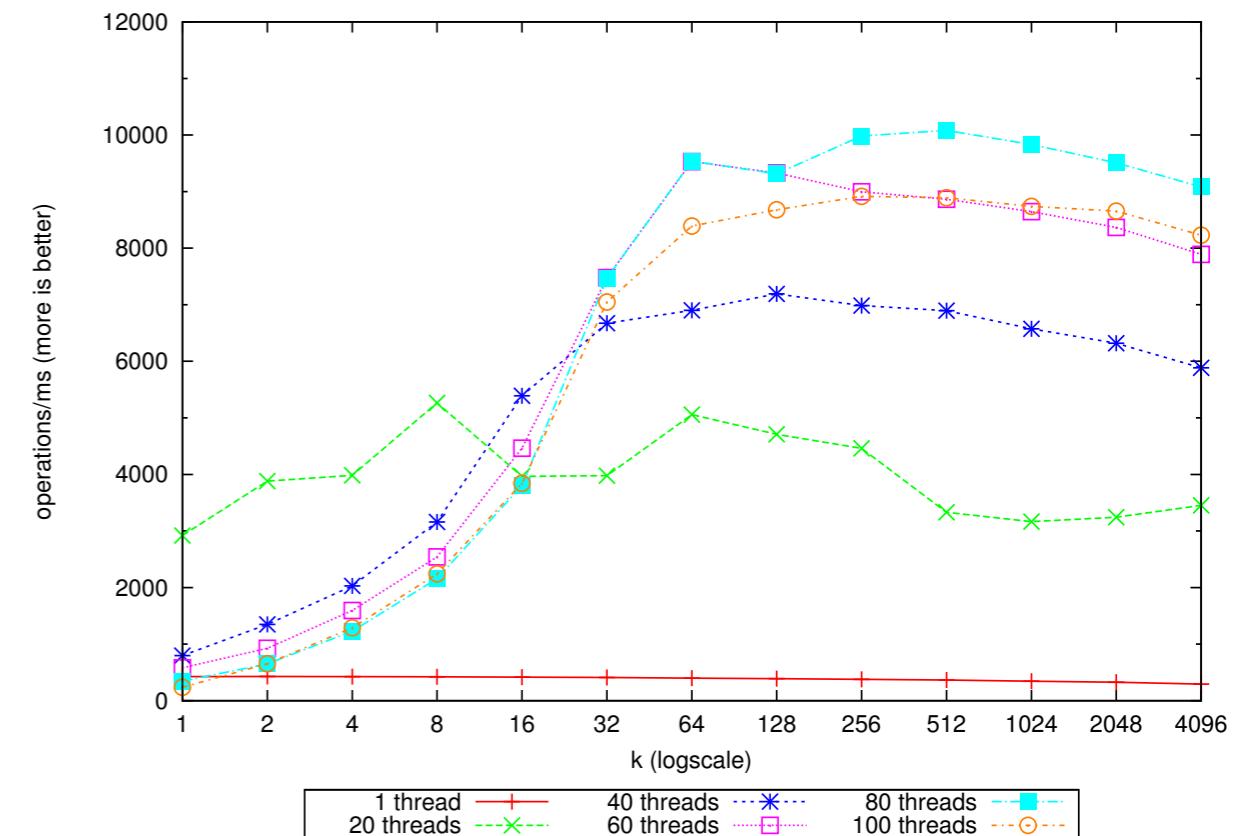


(d) BS number of failed reads per operation of very high contention producer-consumer benchmark ( $c = 1000, i = 5000$ )

# Performance in $k$

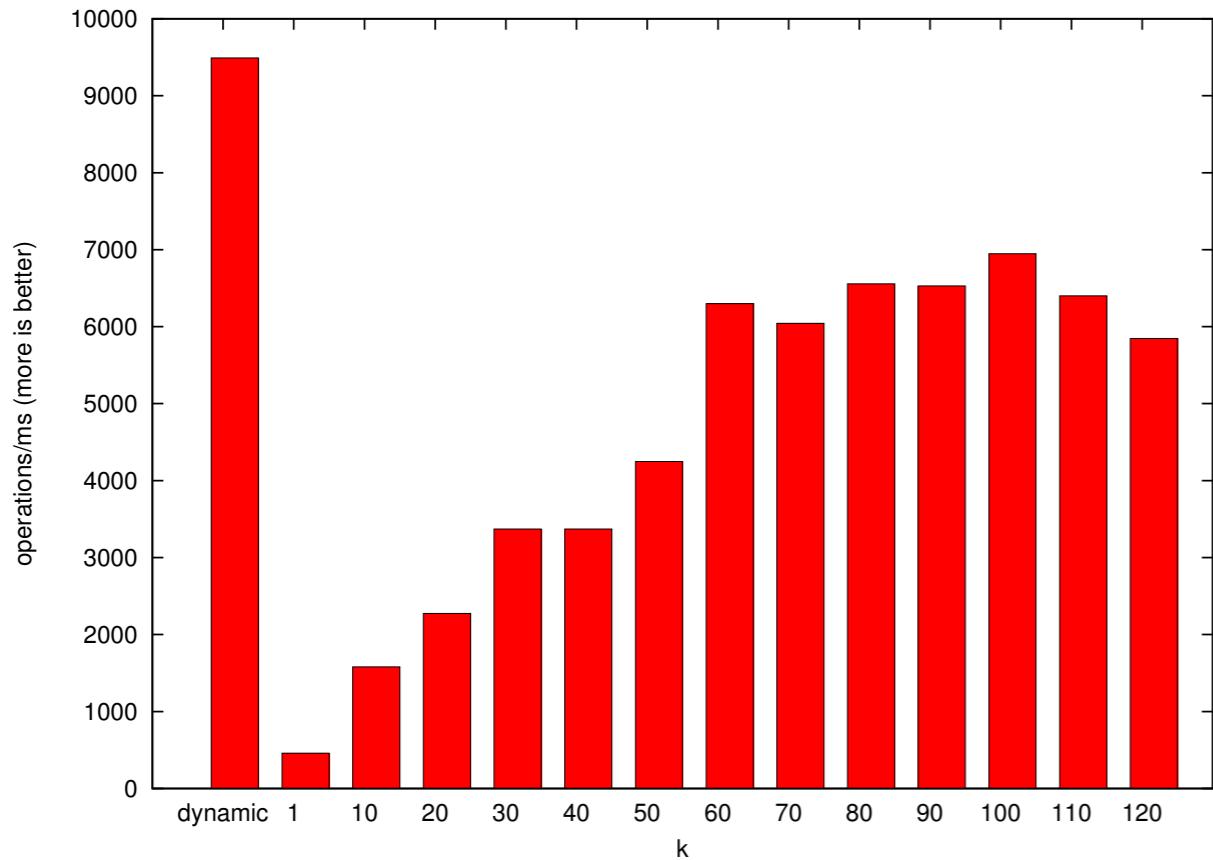


(e) BS performance of very high contention producer-consumer benchmark ( $c = 1000, i = 0$ )

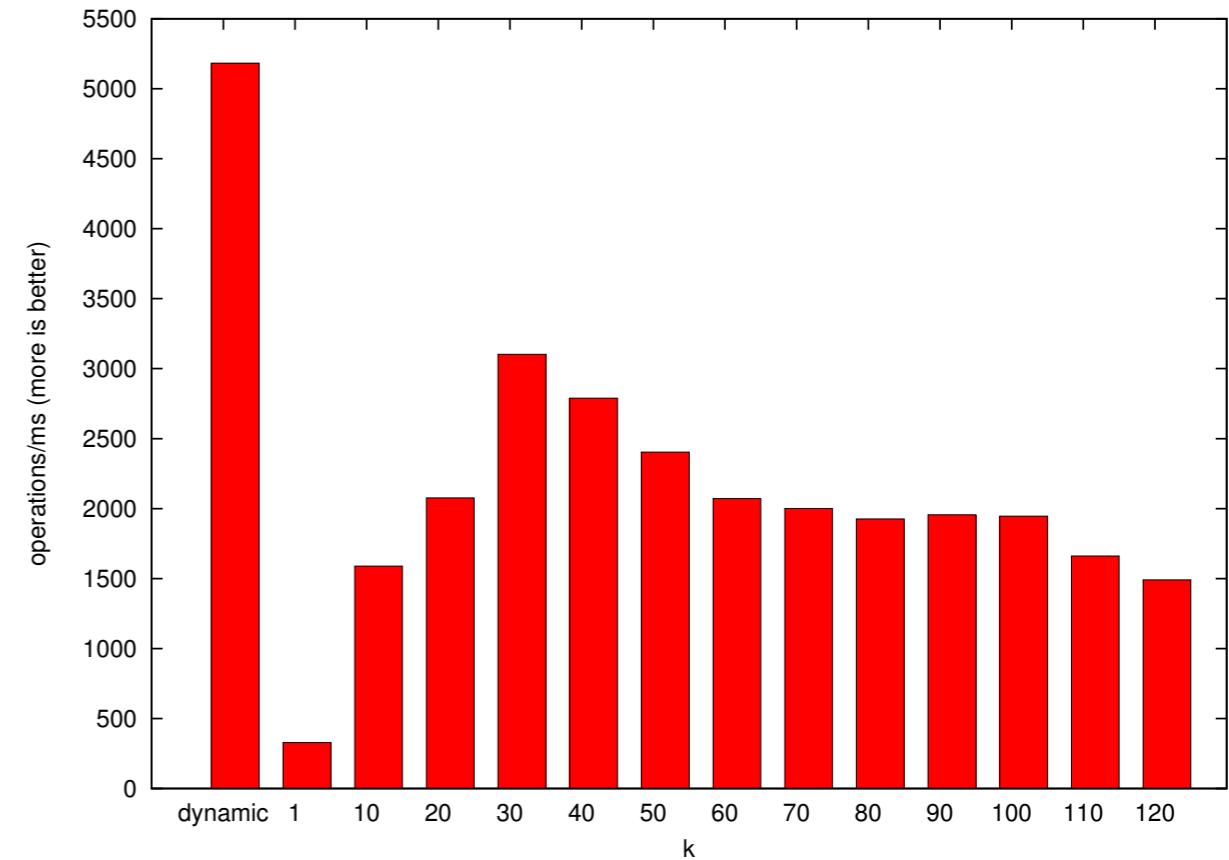


(f) BS performance of very high contention producer-consumer benchmark ( $c = 1000, i = 5000$ )

# Dynamic $k$



(a) BS  $k$ -FIFO queue



(b) US  $k$ -FIFO queue

**Fig. 4.** Variable-load producer-consumer microbenchmarks with an increasing number of static  $k$  versus a dynamically controlled  $k$  on a 40-core (2 hyperthreads per core) server

# Concurrent $k$ -FIFO Queue

- with a  $k$ -FIFO queue elements may be returned **out-of-FIFO order up to  $k$**

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- starvation-free** for finite  $k$

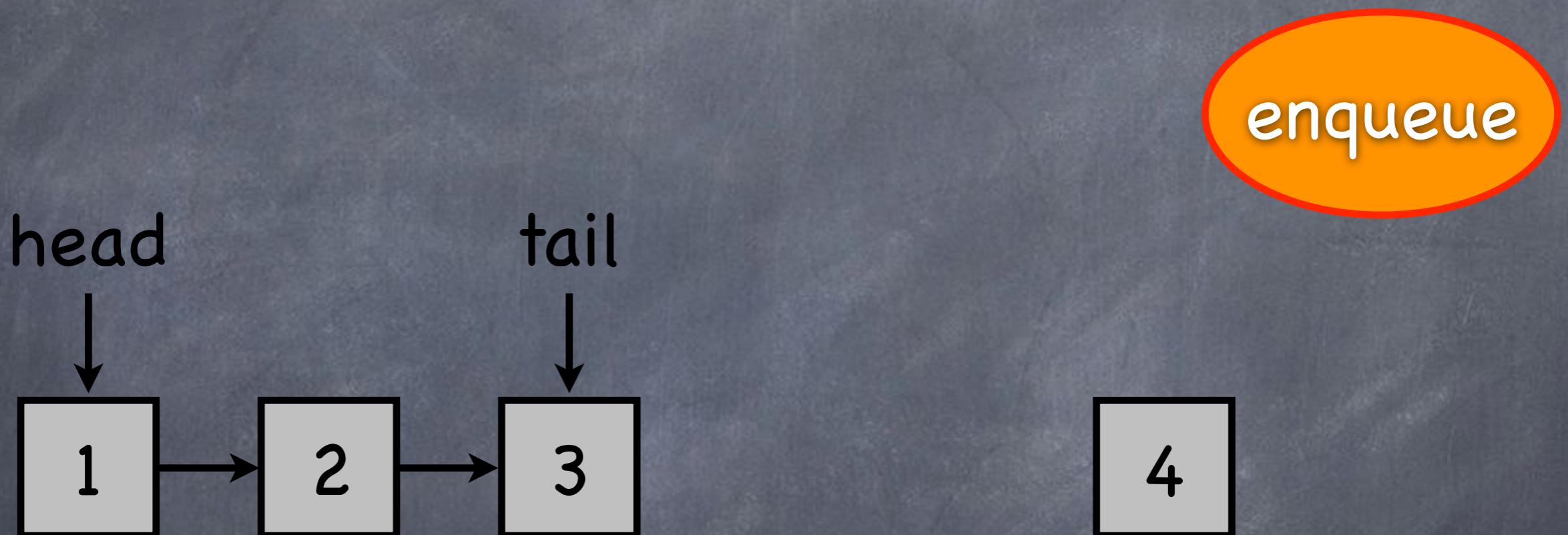
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- **starvation-free** for finite  $k$
- **0-FIFO queue = regular FIFO queue**

# Concurrent $k$ -FIFO Queue

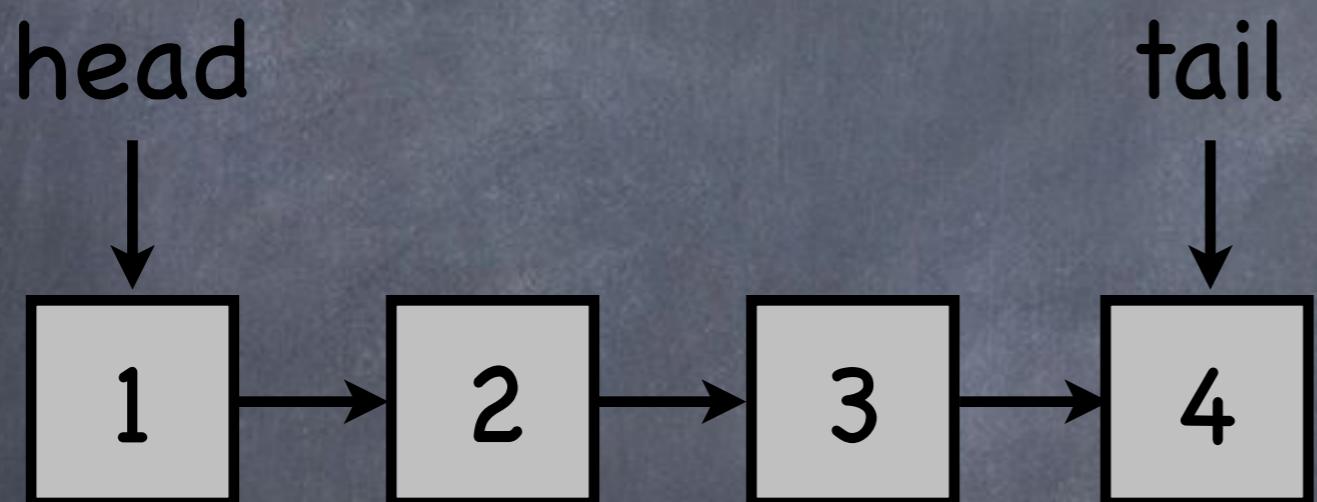
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- **starvation-free** for finite  $k$
- $0$ -FIFO queue = regular FIFO queue
- bigger  $k \rightarrow$  better performance, scalability?

# Concurrent 2-FIFO Queue (k=2)

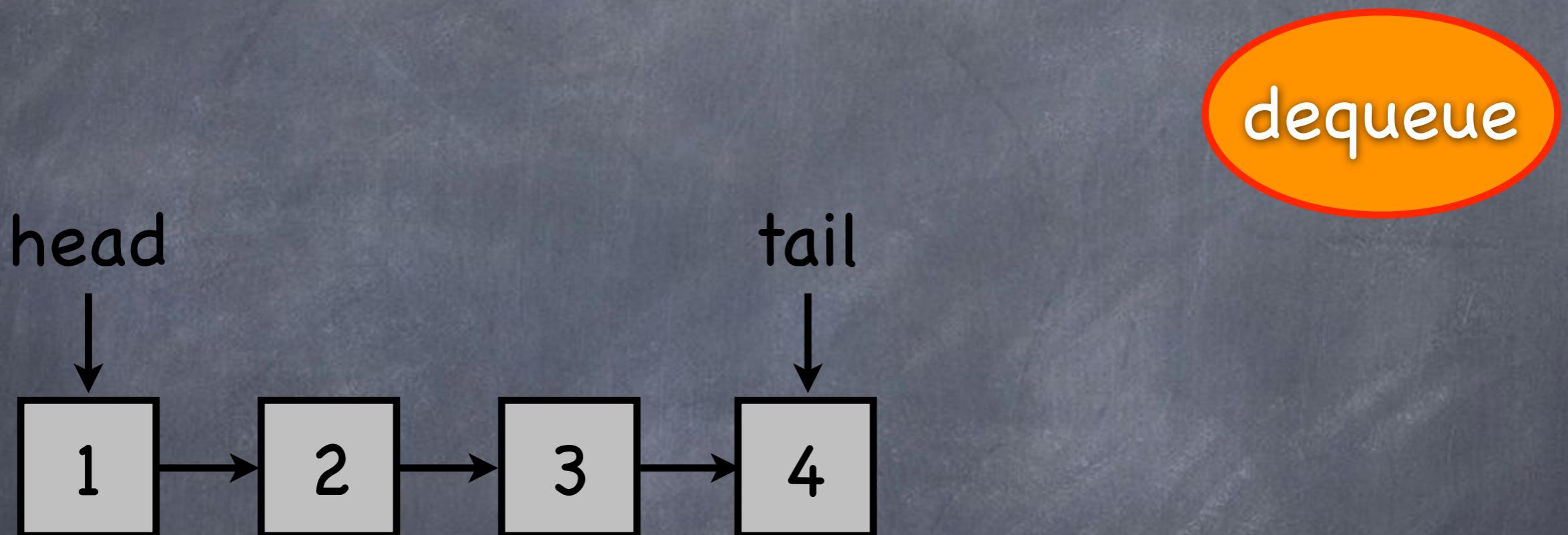


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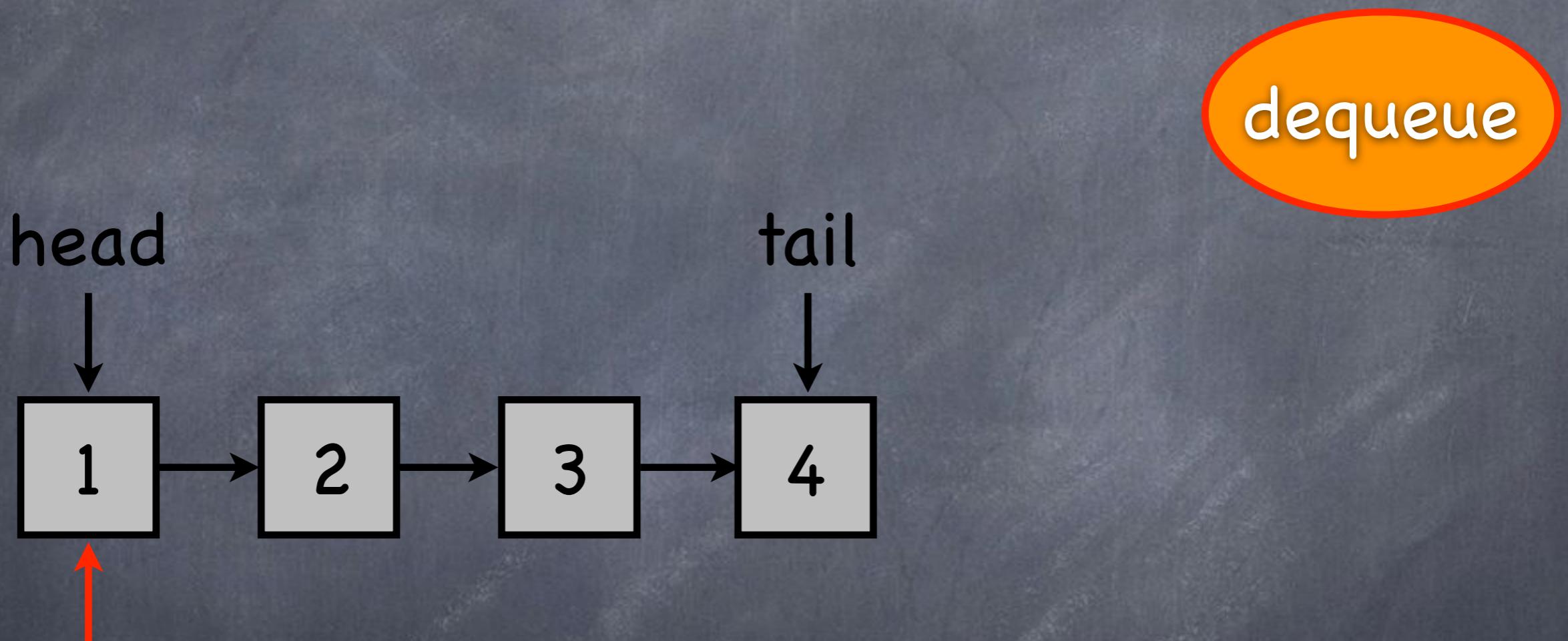
enqueue



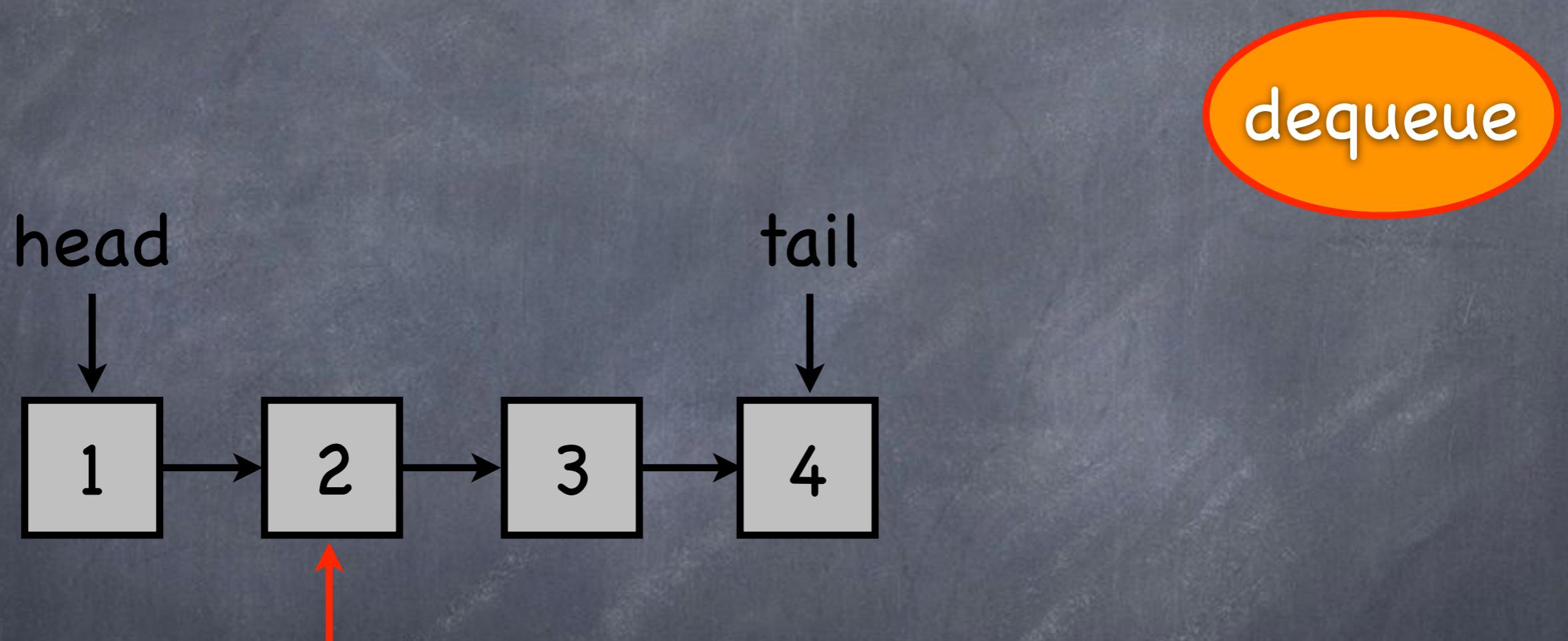
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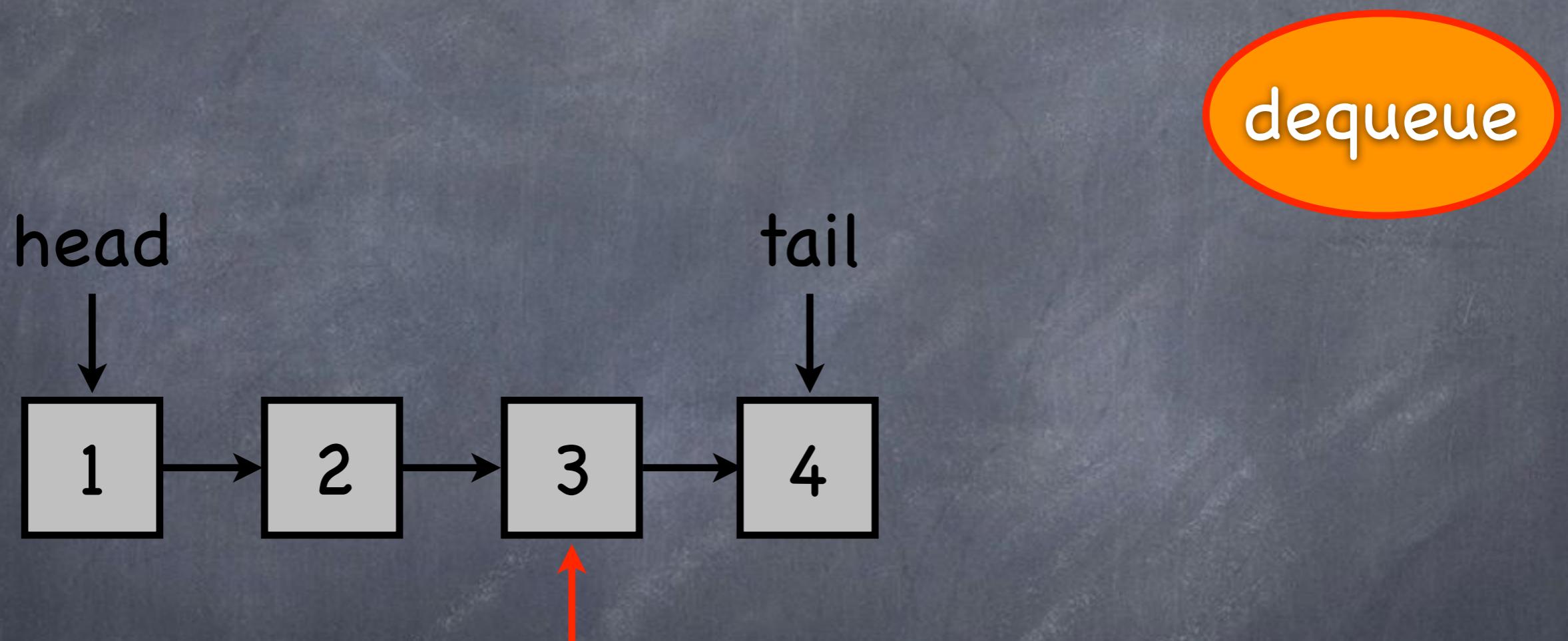
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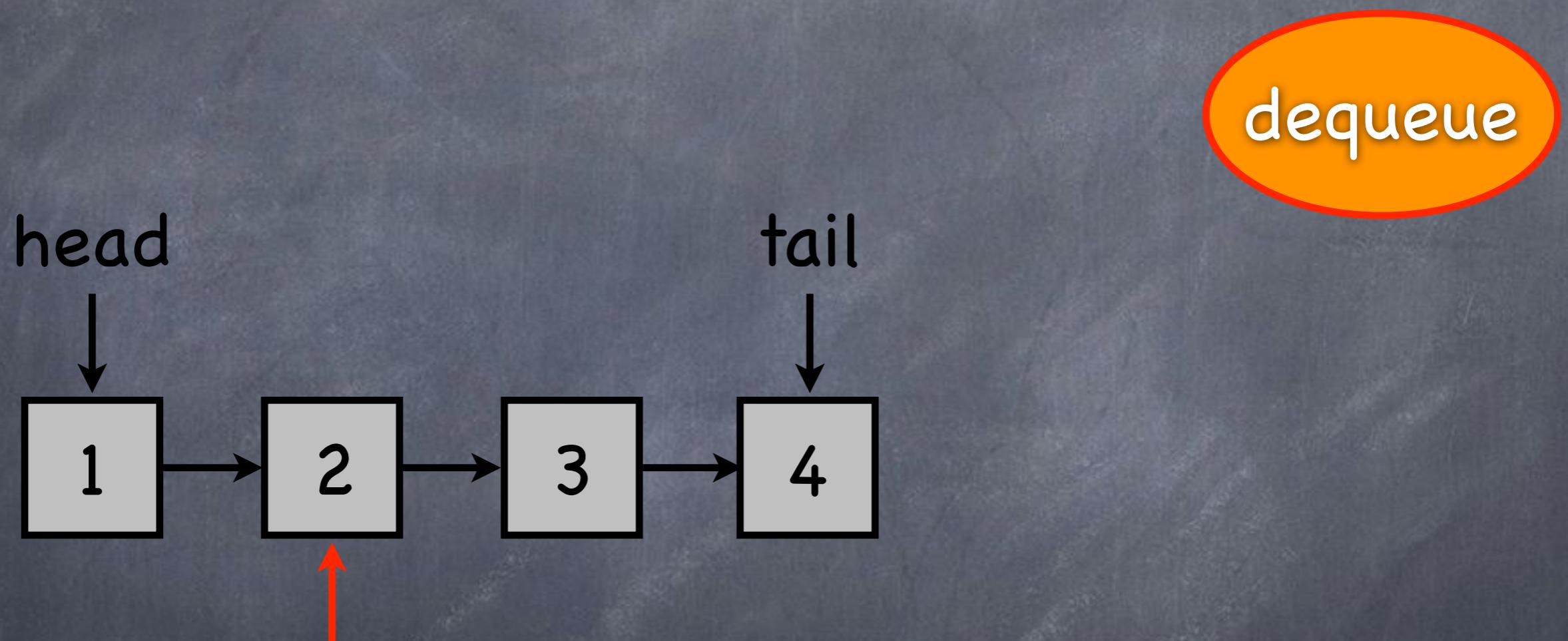
# Concurrent 2-FIFO Queue (k=2)



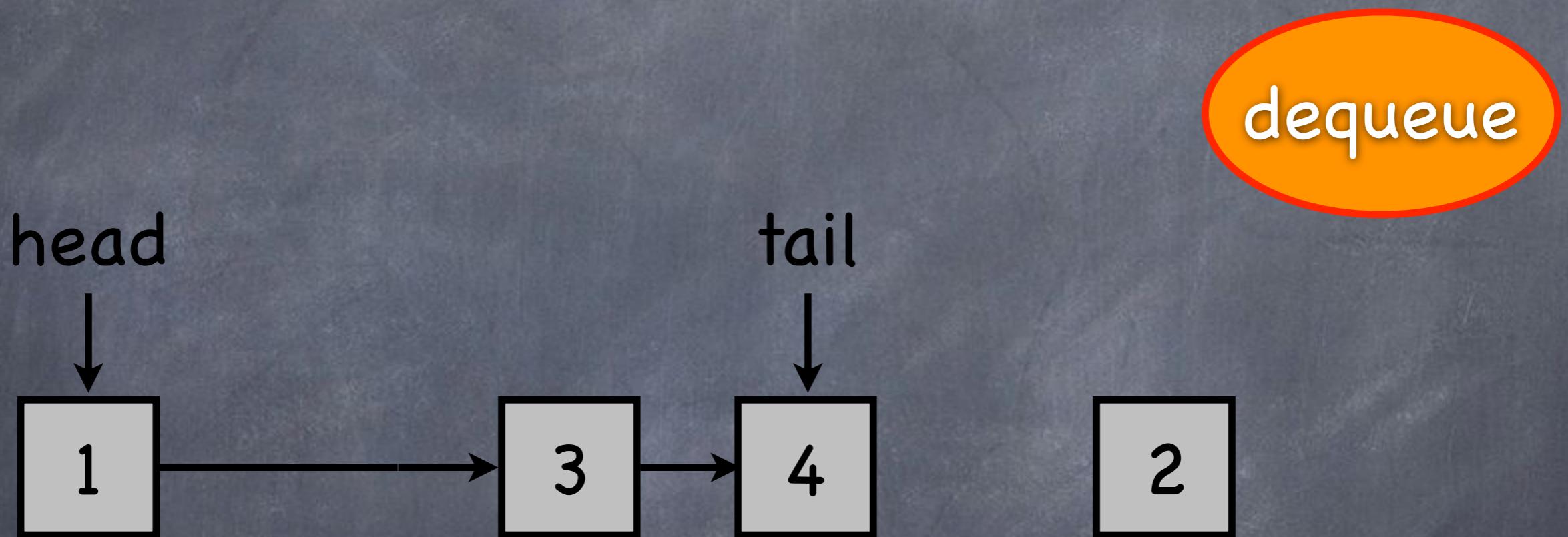
# Concurrent 2-FIFO Queue (k=2)



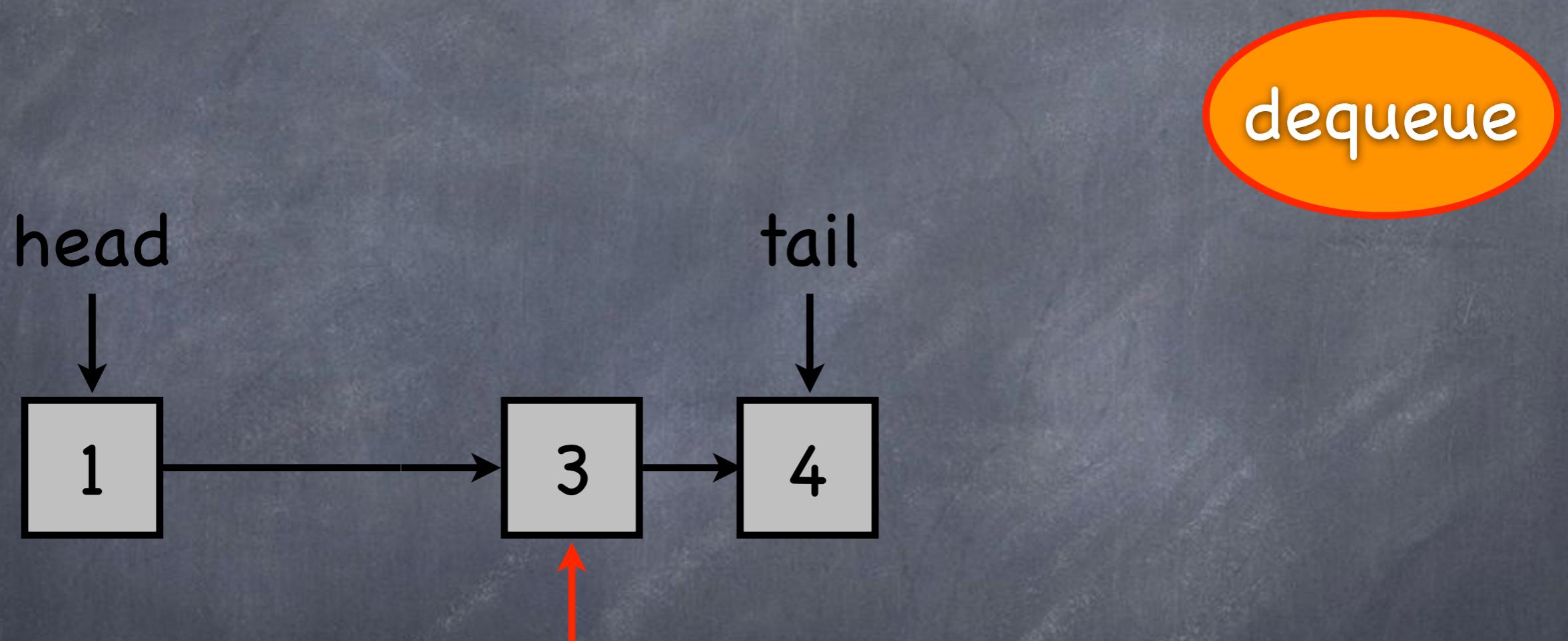
# Concurrent 2-FIFO Queue (k=2)



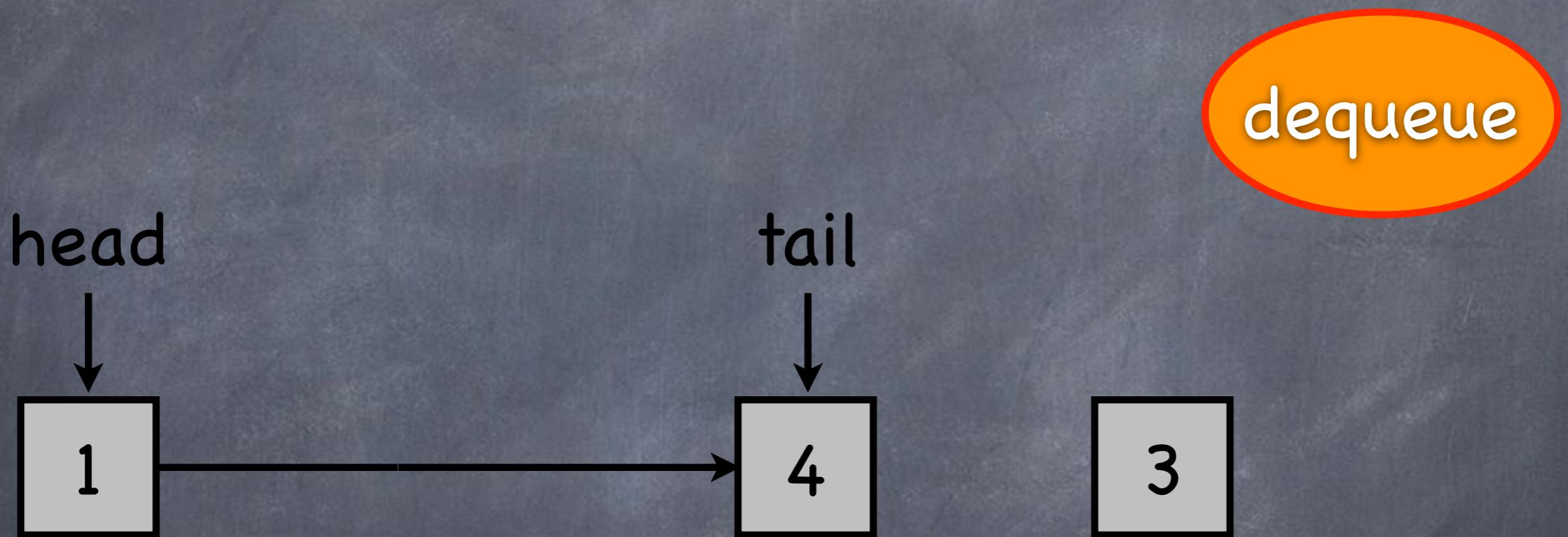
# Concurrent 2-FIFO Queue (k=2)



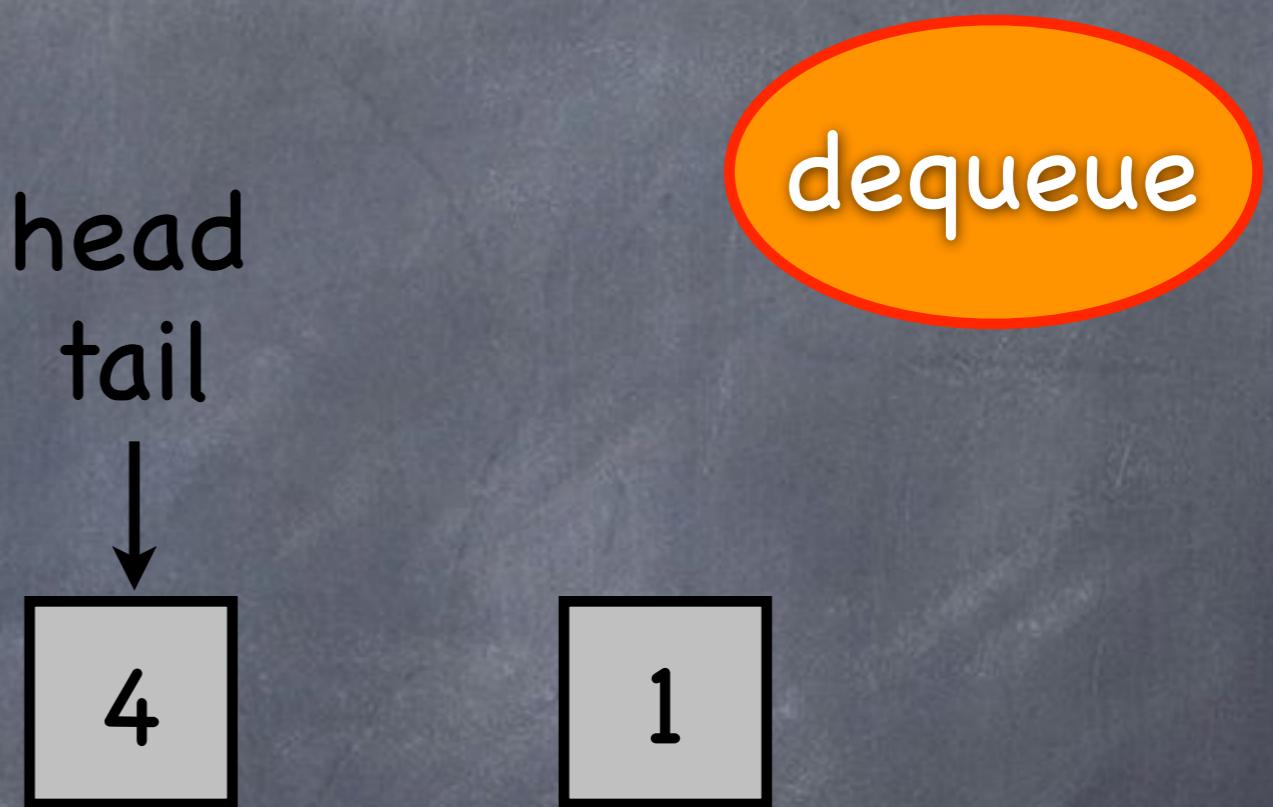
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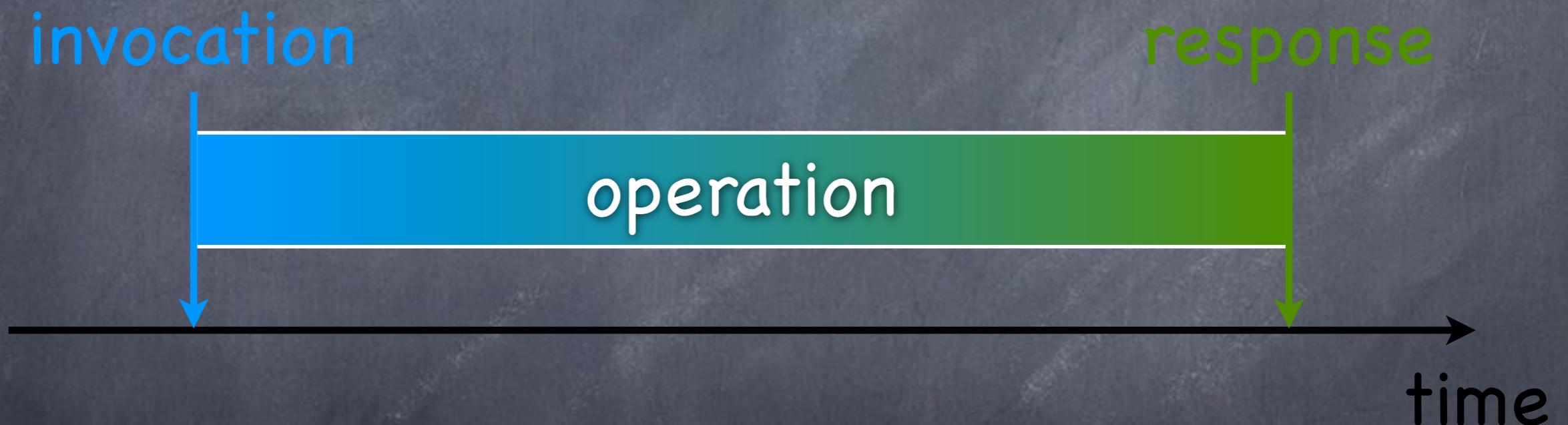


We call  $k$   
the worst-case semantical  
deviation (WCSD) of  
a  $k$ -FIFO queue from  
a regular FIFO queue

The actual semantical deviation (ASD) is the semantical deviation of a **k**-FIFO queue when applied to a given workload

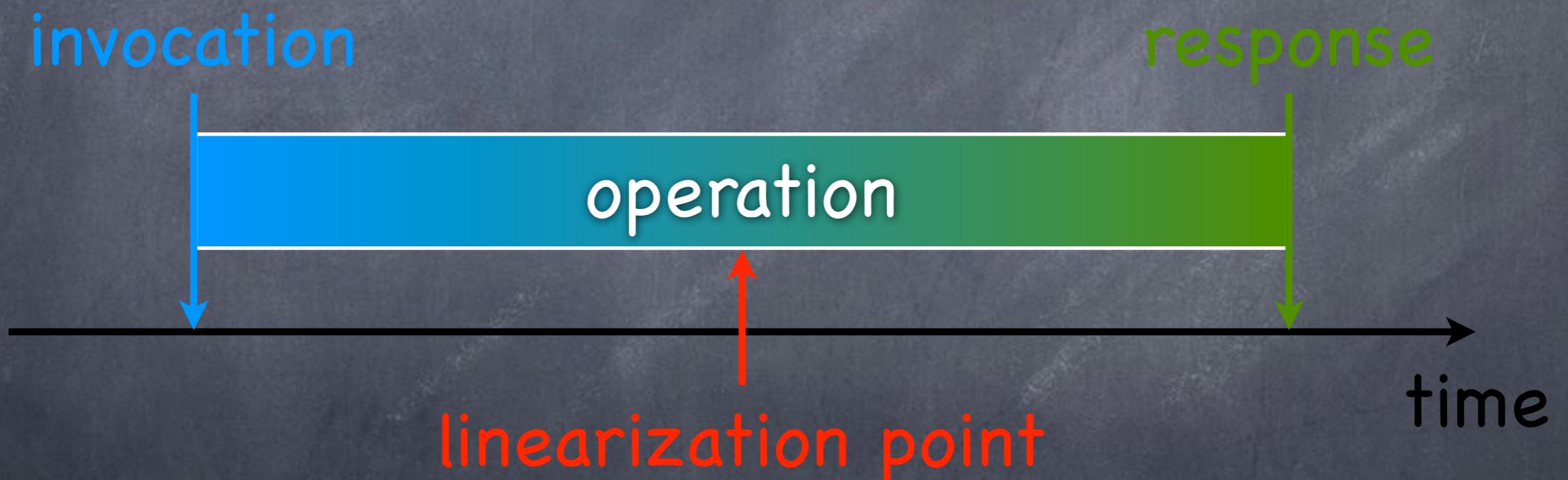
# Execution History

Sequence of Time-Stamped Invocation and Response Events

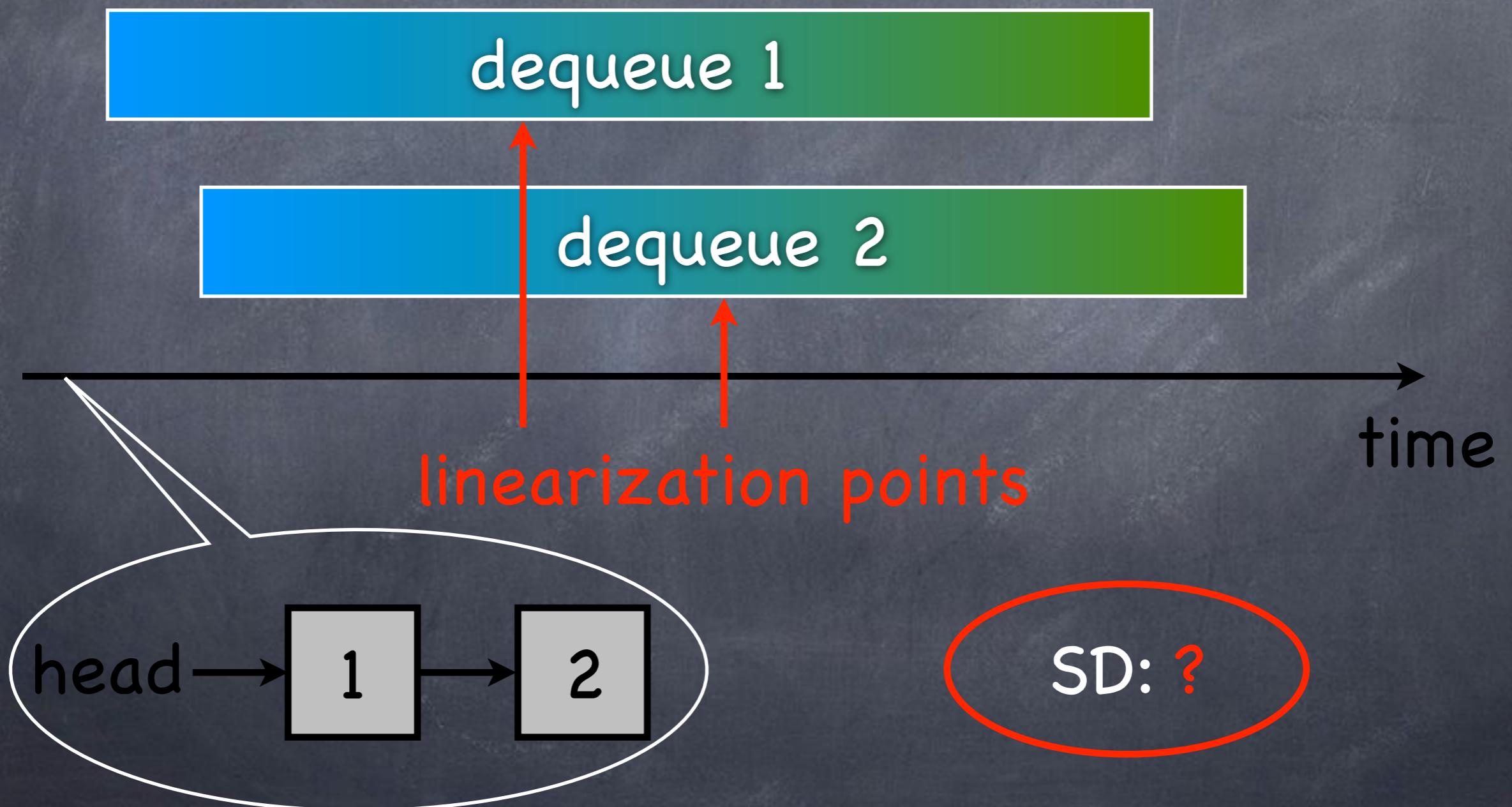


# Execution History

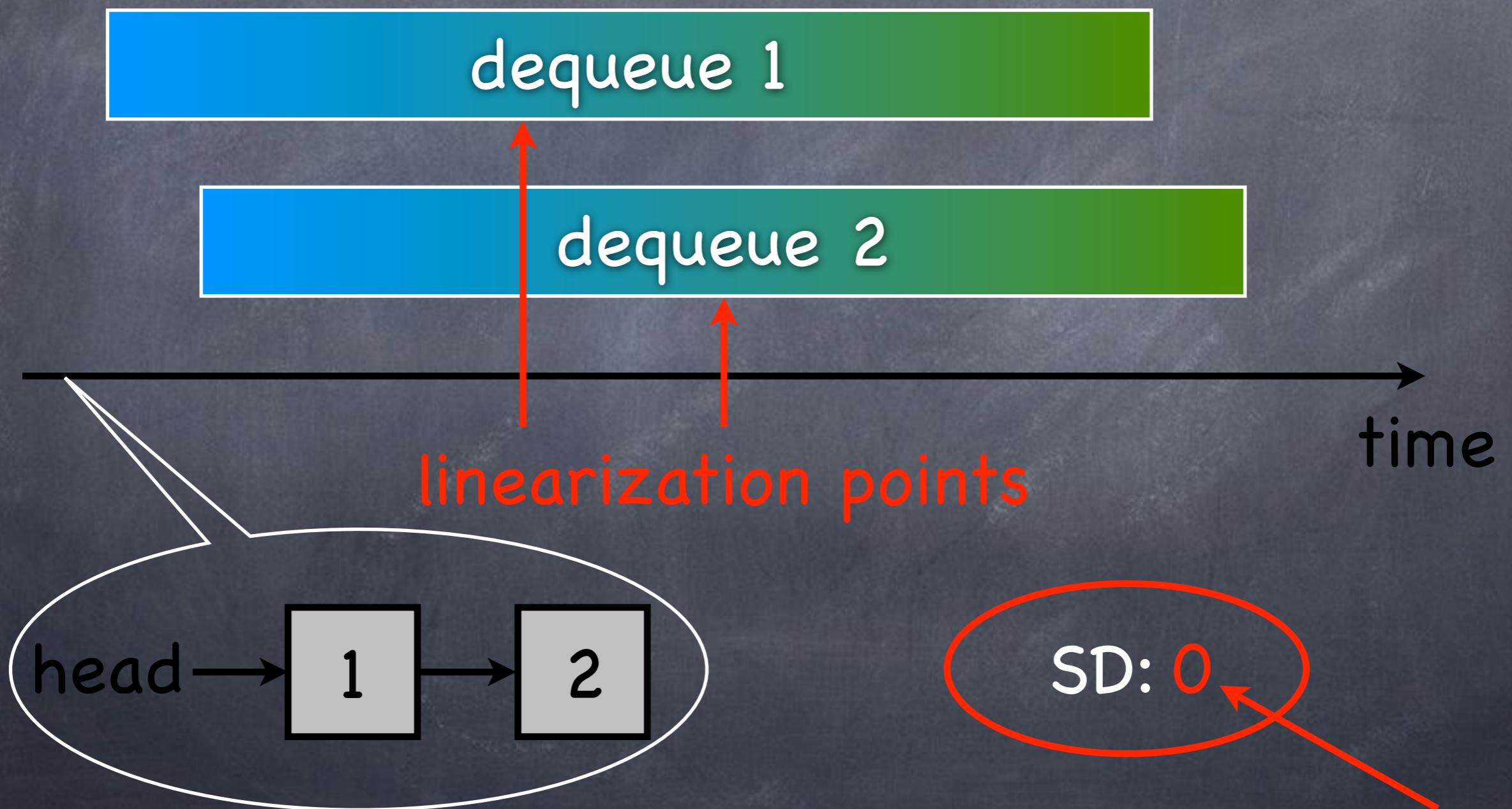
Sequence of Time-Stamped Invocation and Response Events



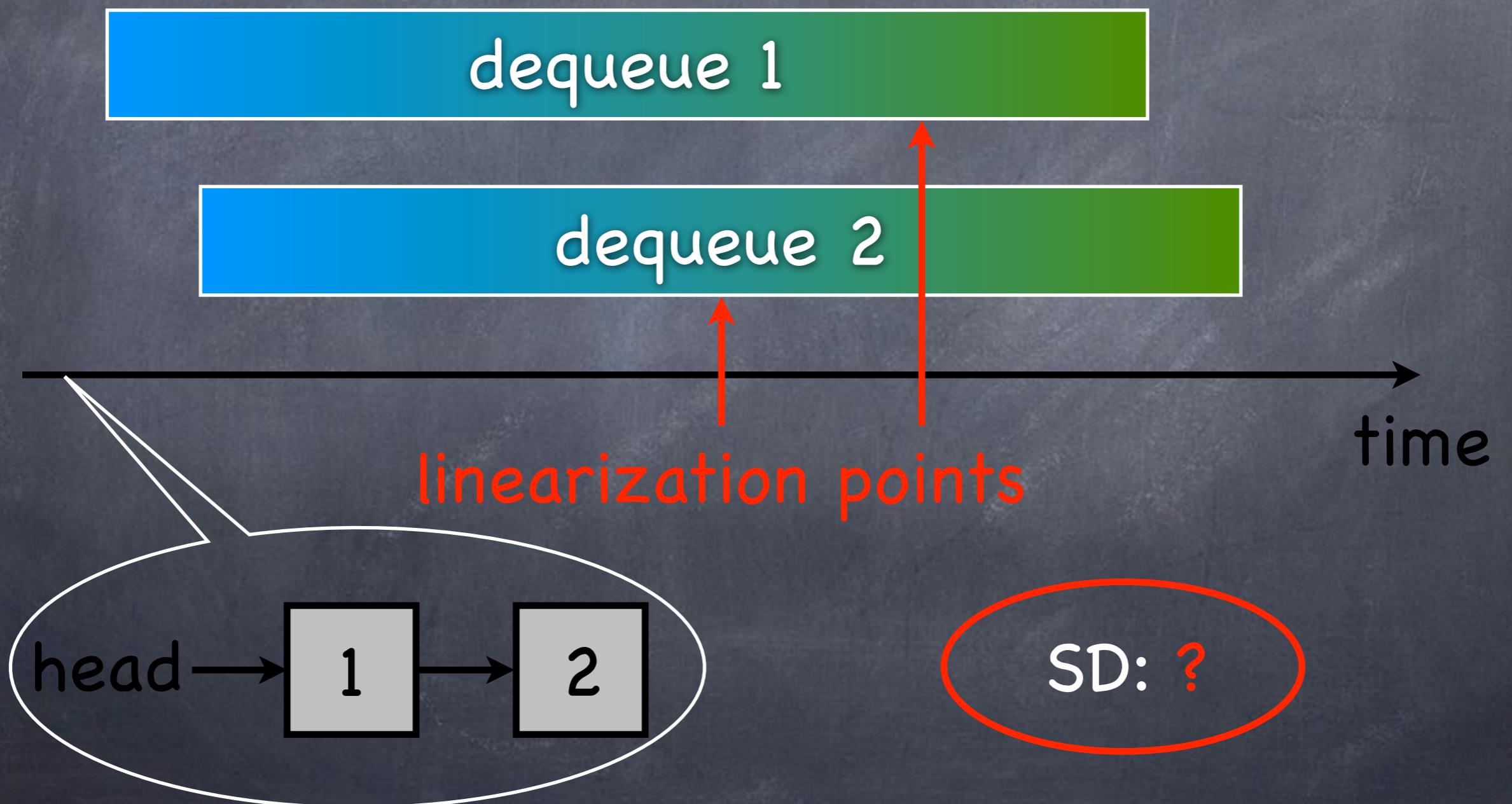
# Measuring Semantical Deviation (SD)



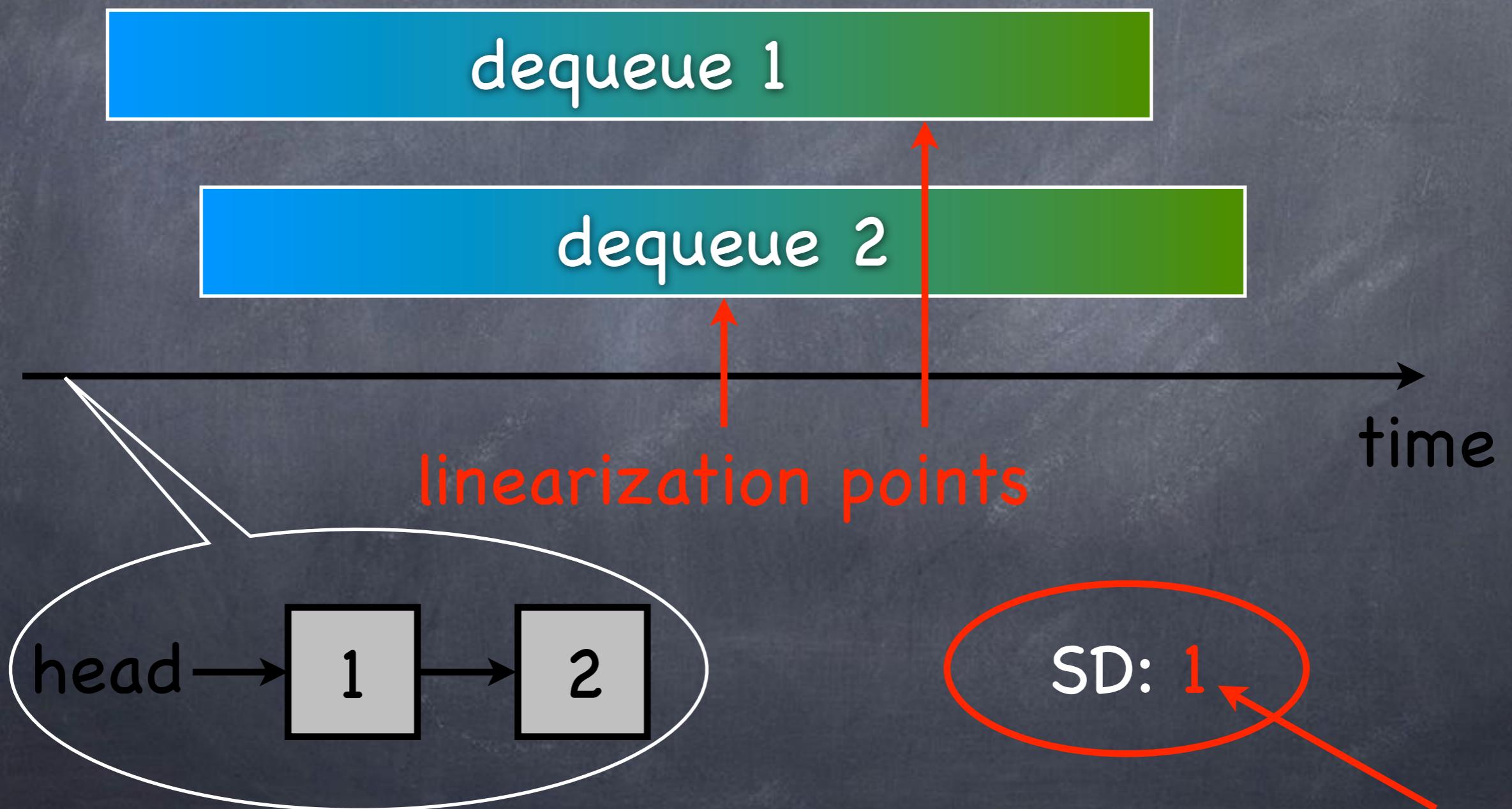
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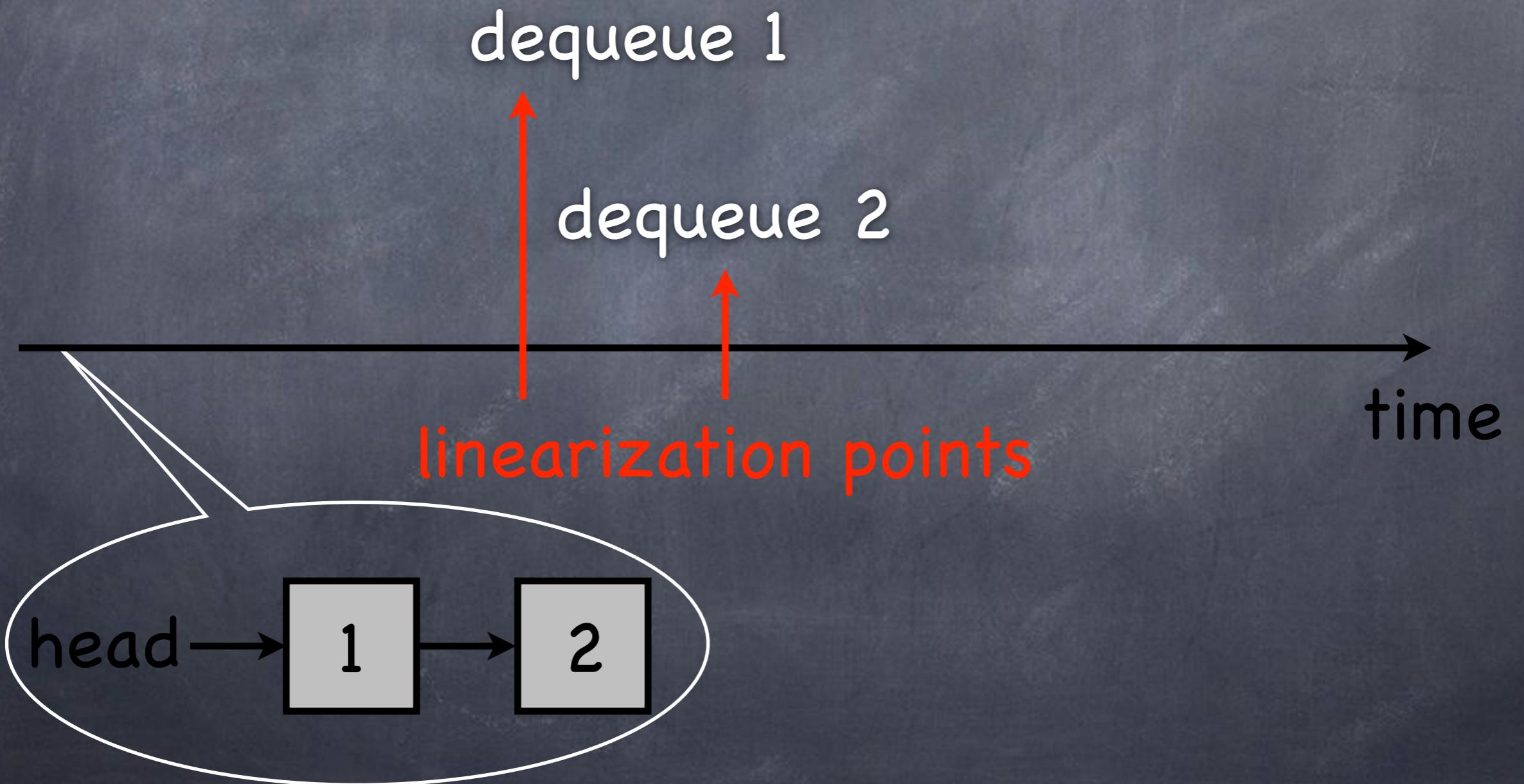


# Measuring Semantical Deviation (SD)



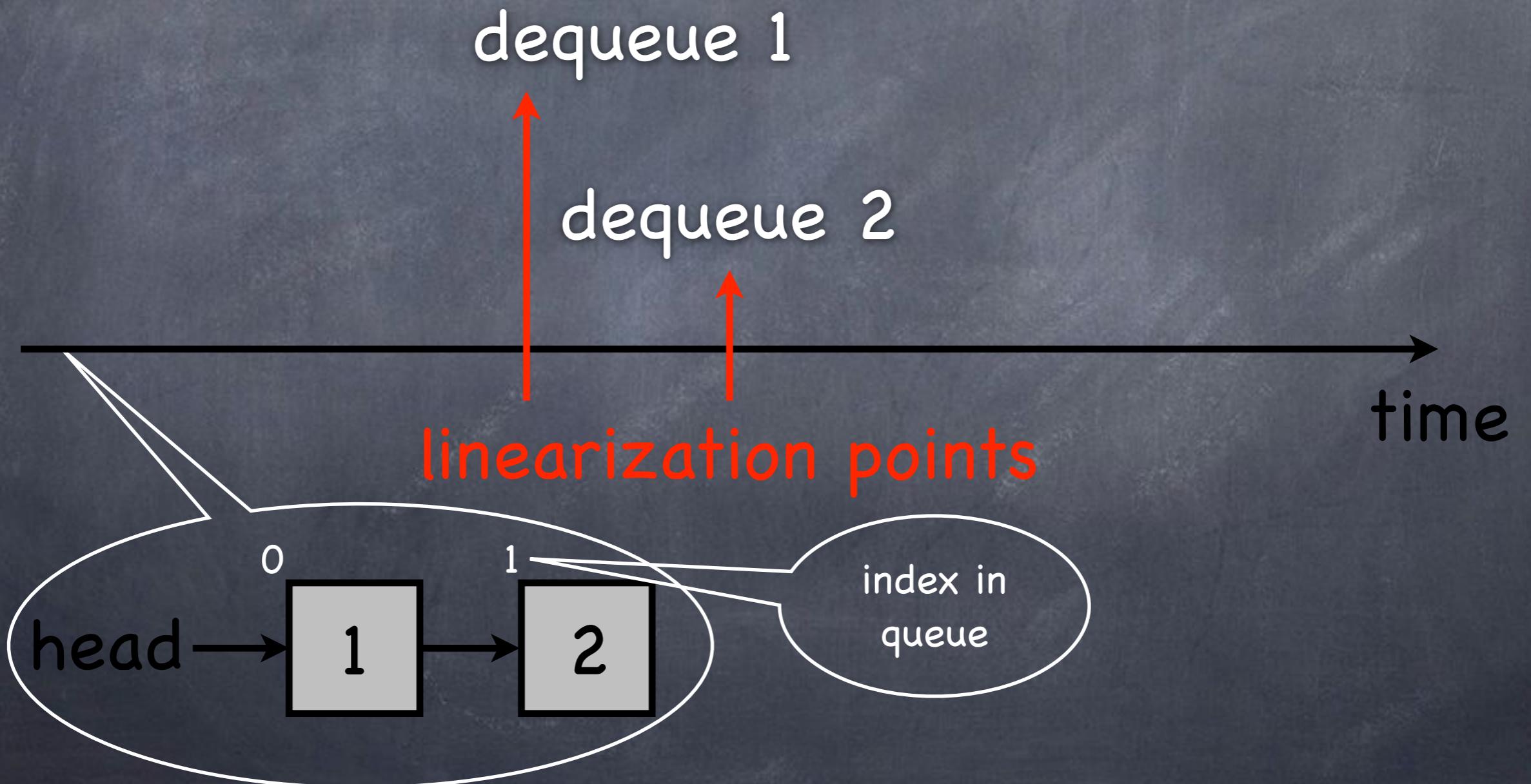
# Sequential History

Sequence of Operations (Linearization Points)



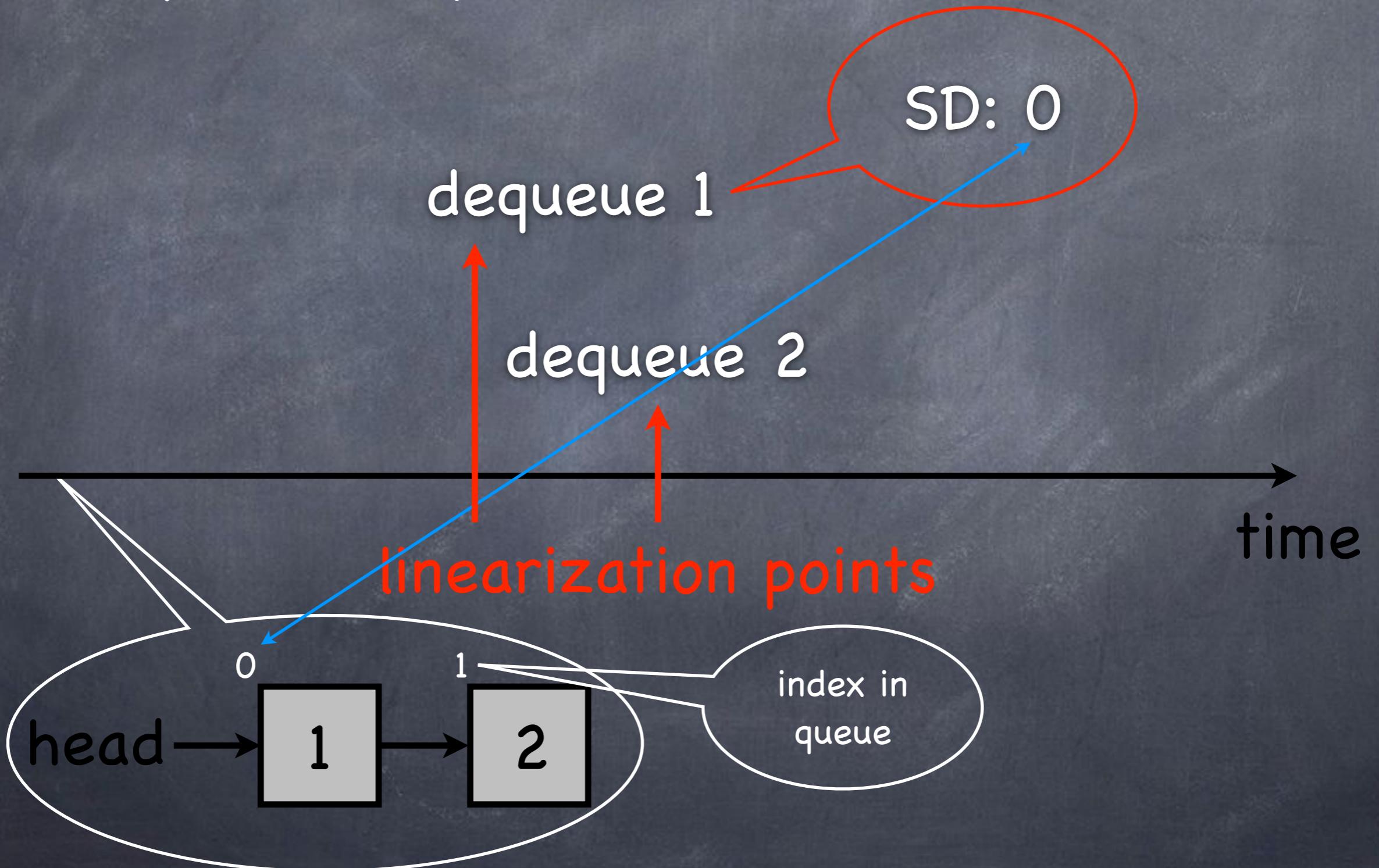
# Sequential History

Sequence of Operations (Linearization Points)



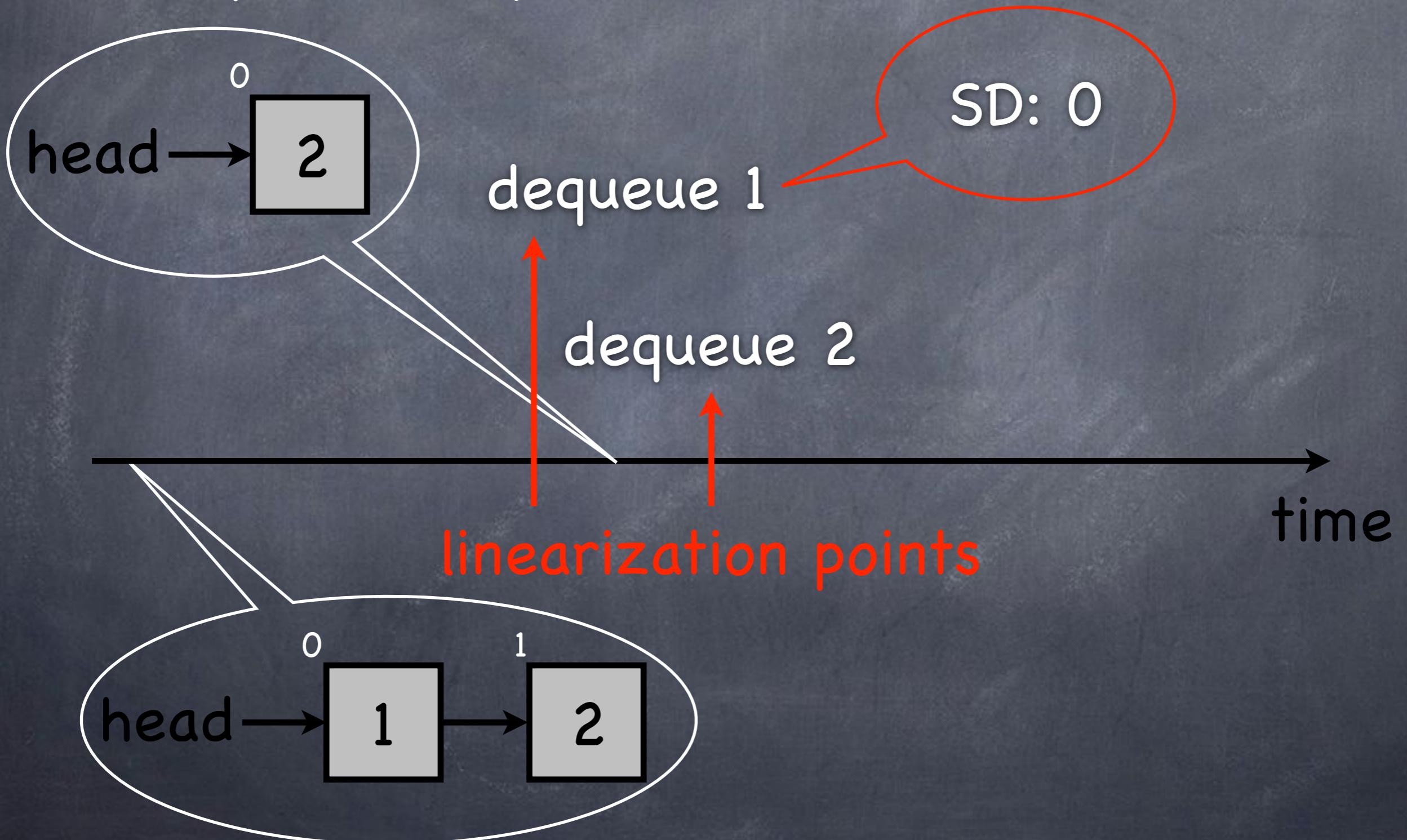
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Sequence of Operations (Linearization Points)



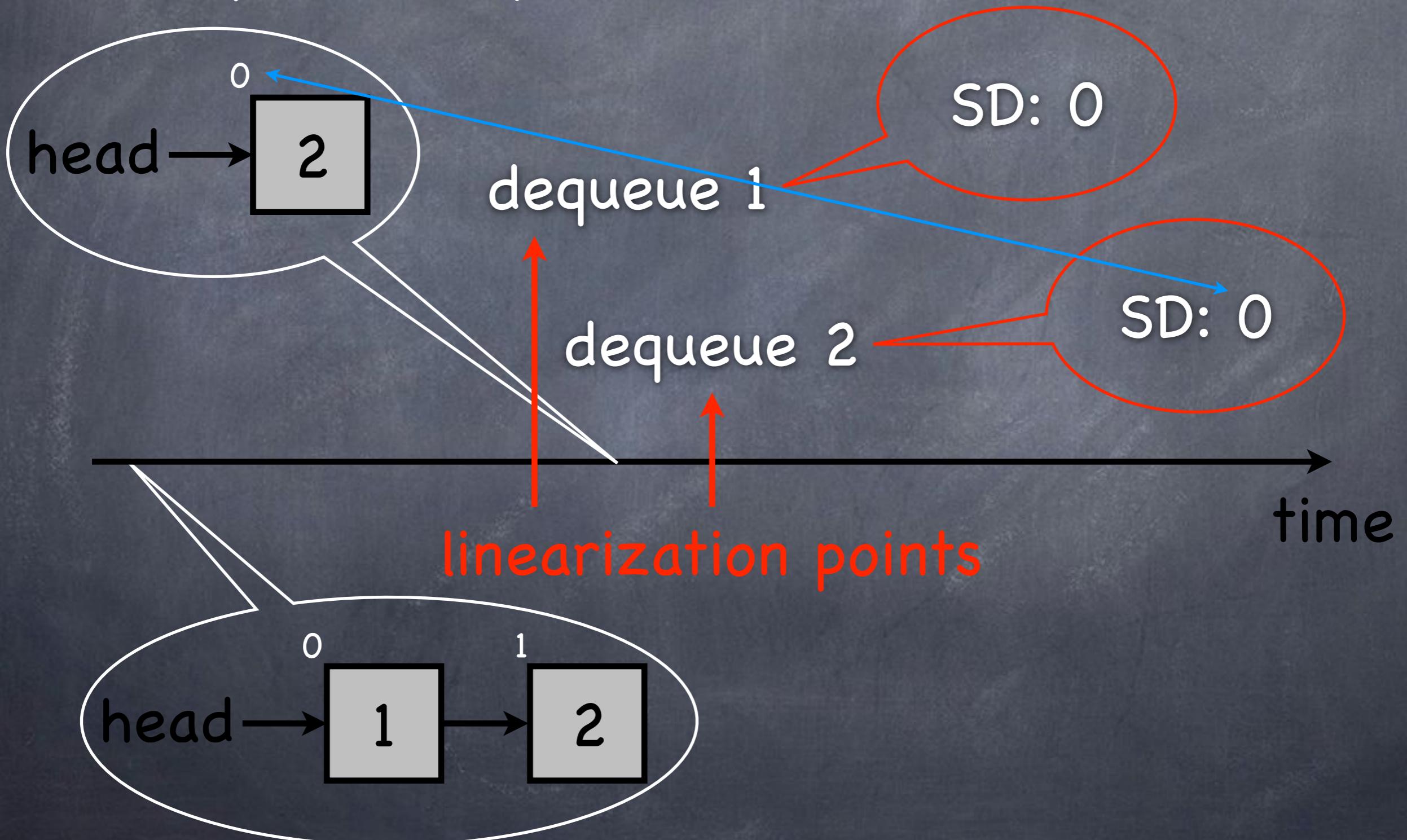
# Sequential History

Sequence of Operations (Linearization Points)



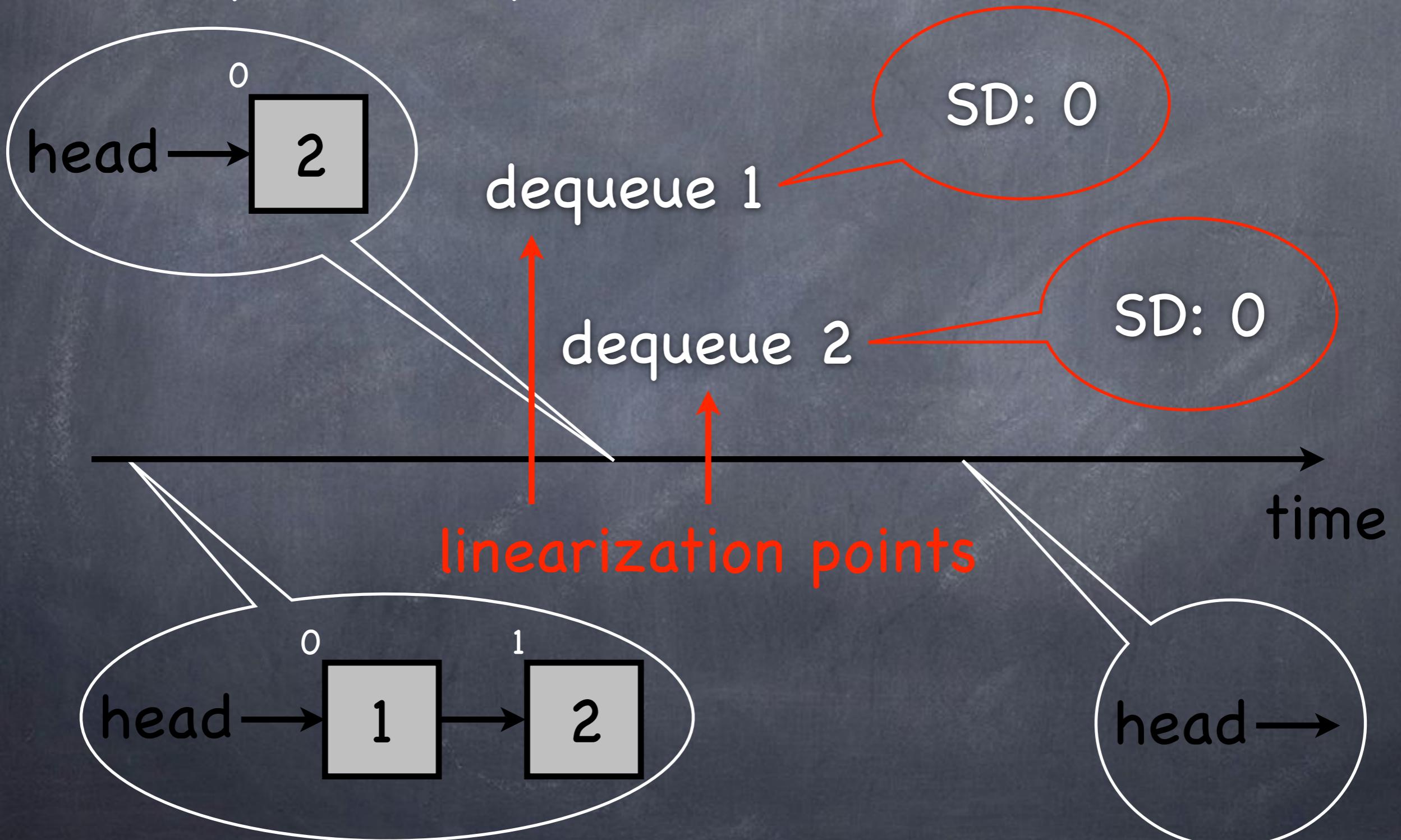
# Sequential History

Sequence of Operations (Linearization Points)



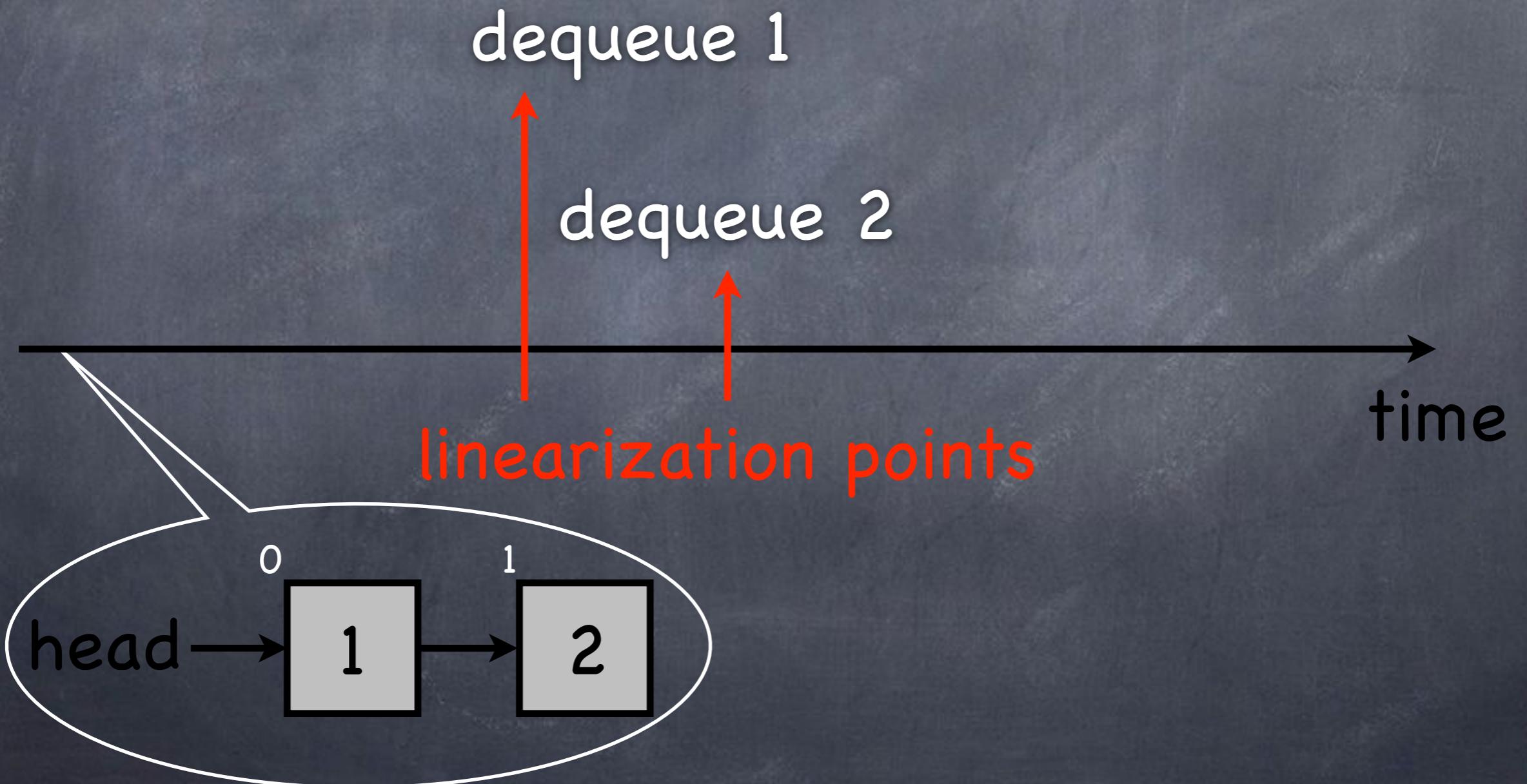
# Sequential History

Sequence of Operations (Linearization Points)



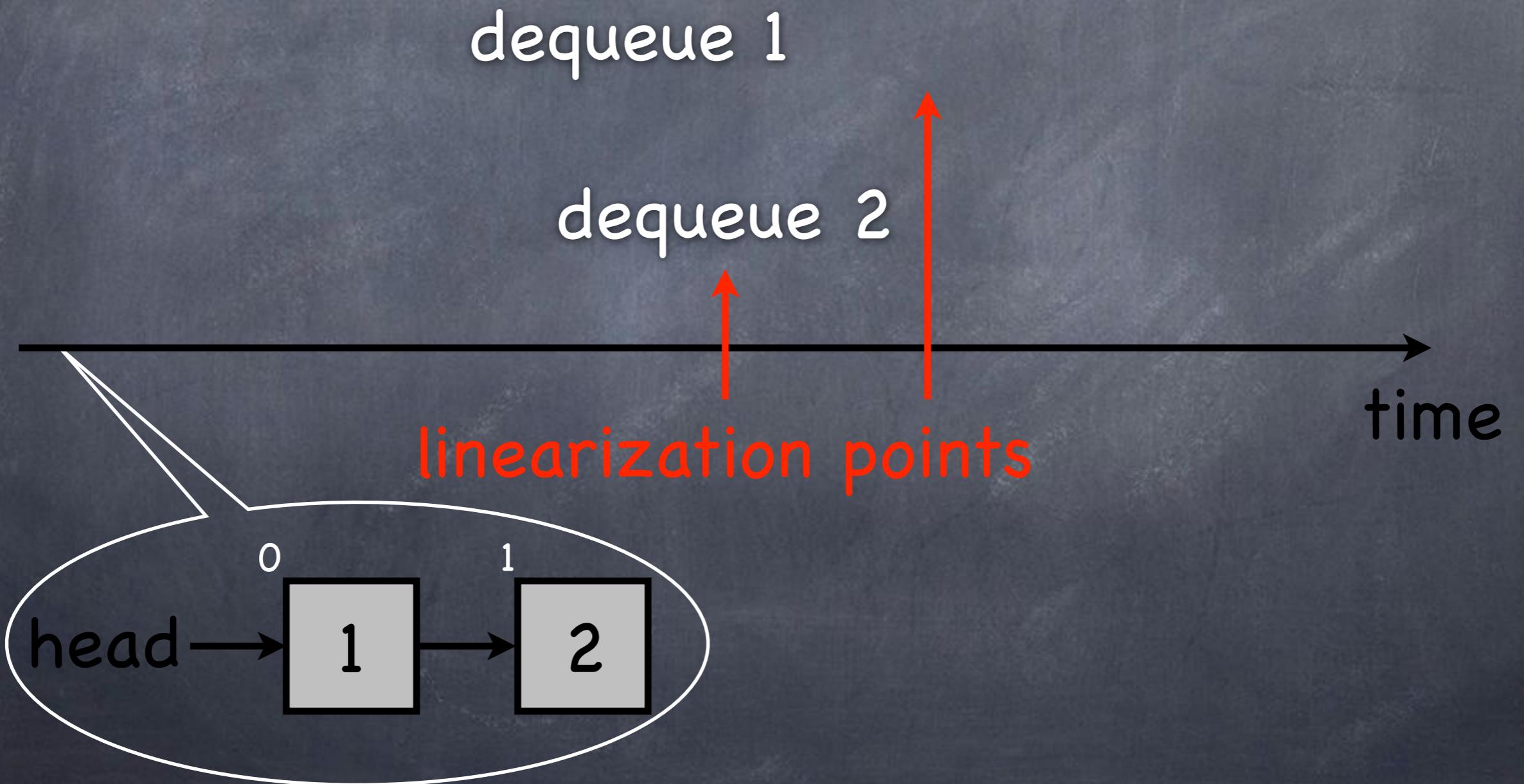
# Sequential History II

Sequence of Operations (Linearization Points)



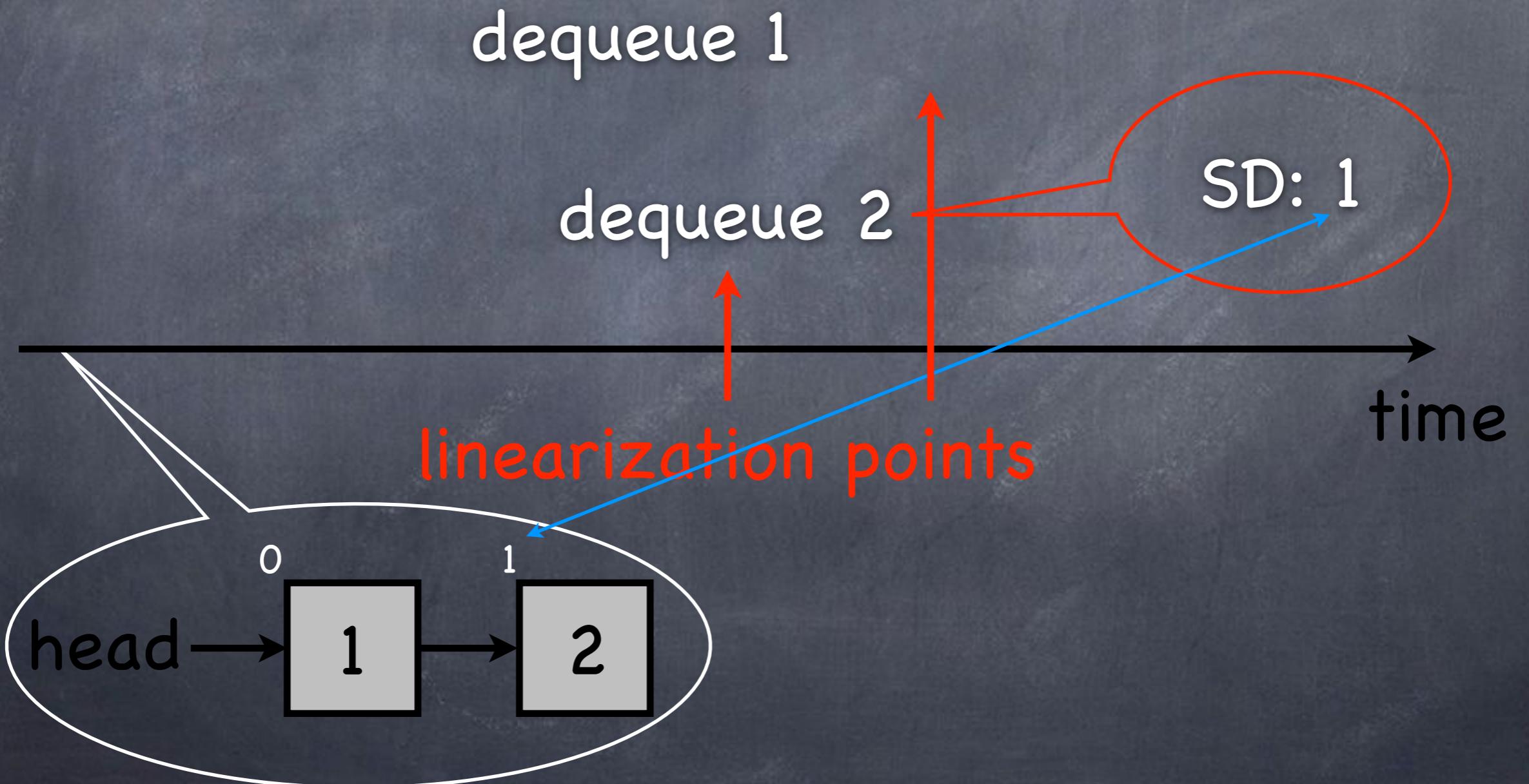
# Sequential History II

Sequence of Operations (Linearization Points)



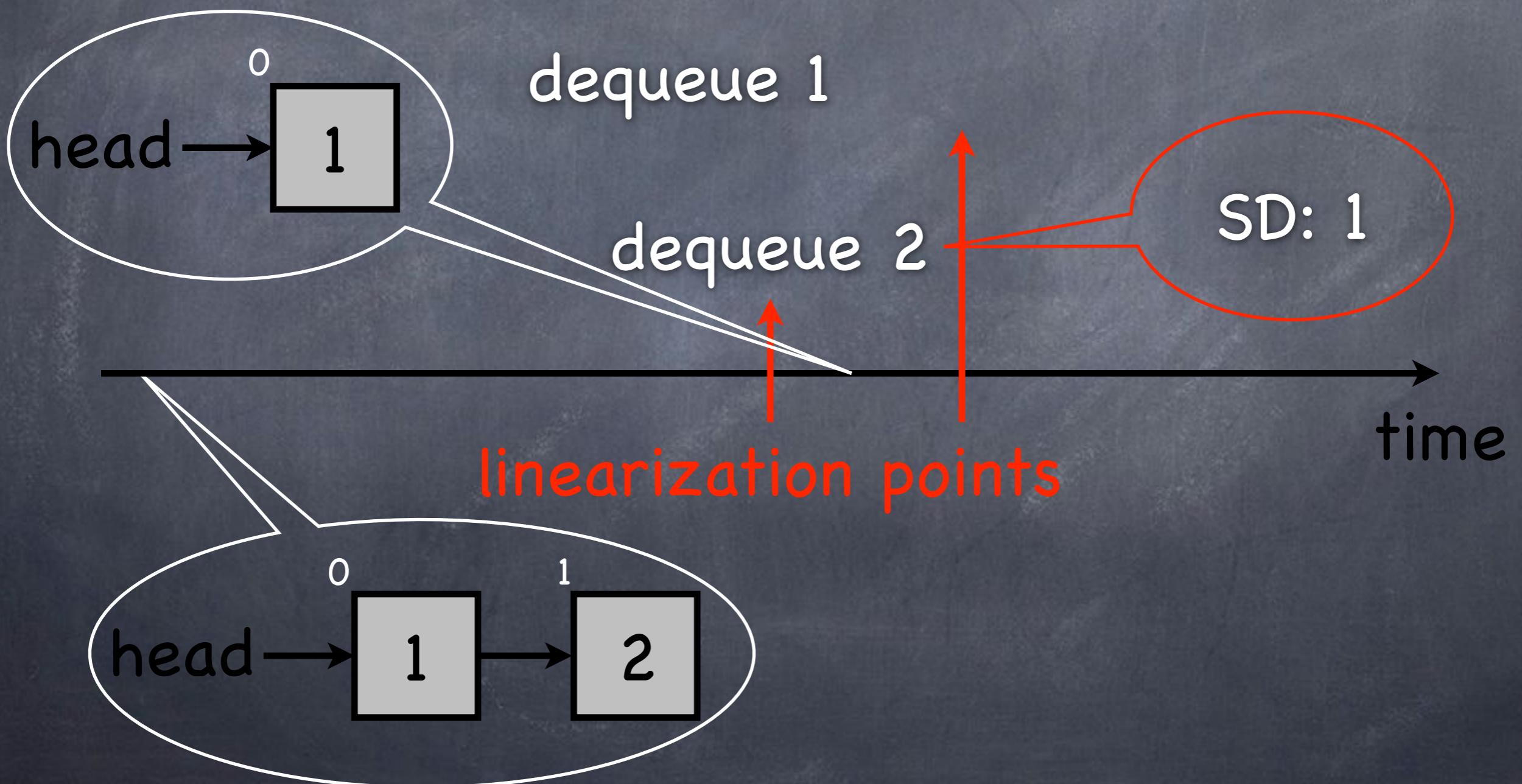
# Sequential History II

Sequence of Operations (Linearization Points)



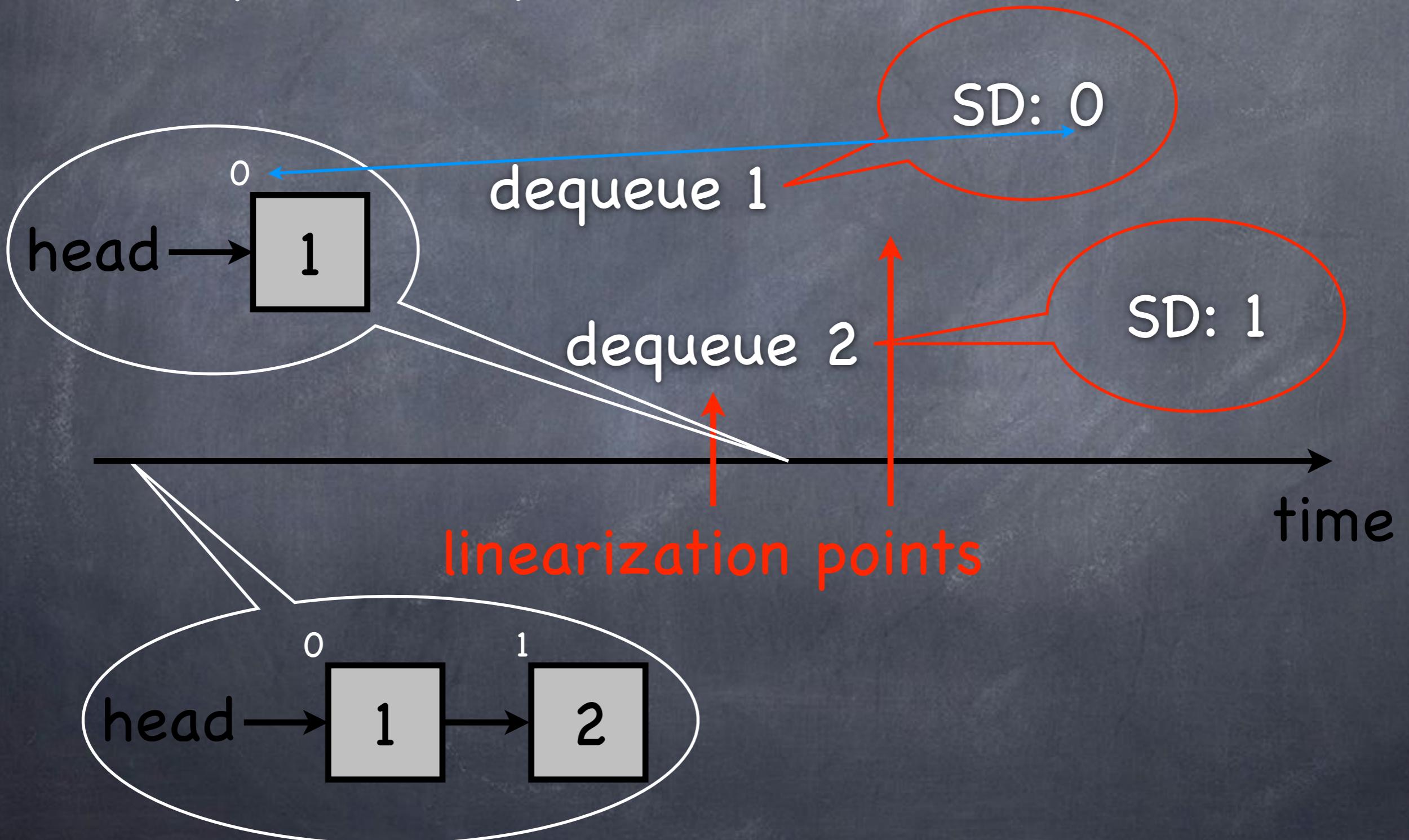
# Sequential History II

Sequence of Operations (Linearization Points)



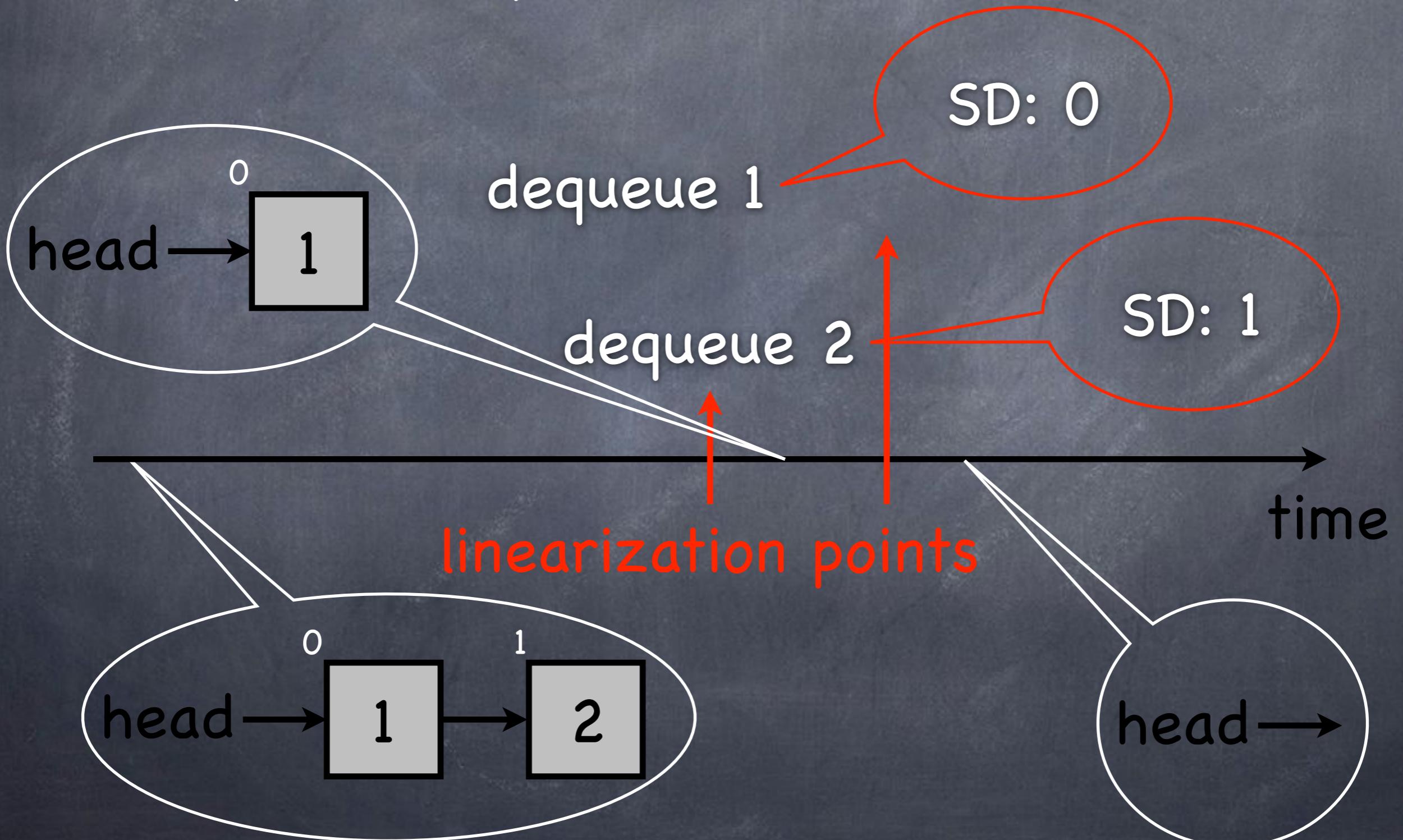
# Sequential History II

Sequence of Operations (Linearization Points)

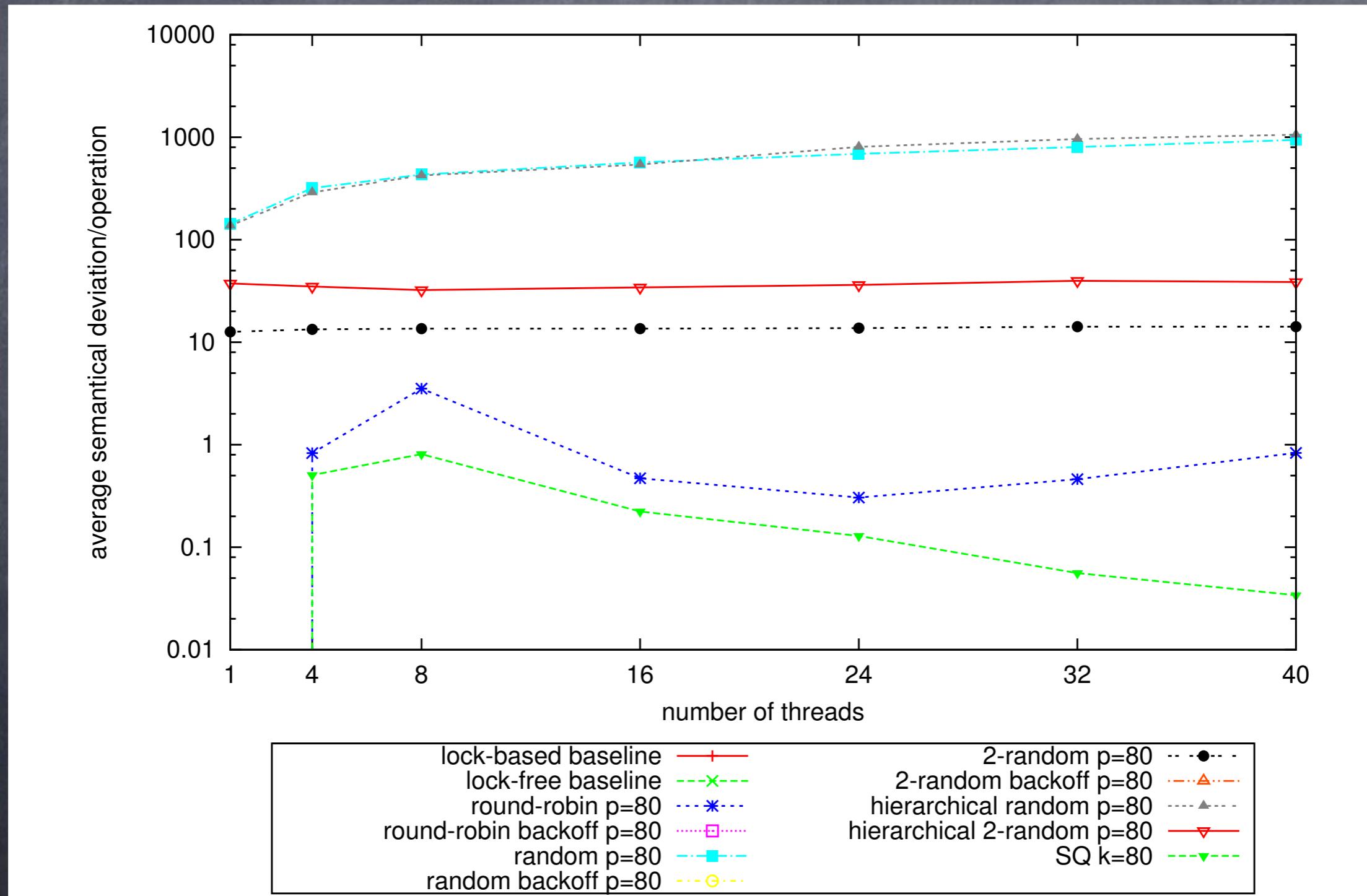


# Sequential History II

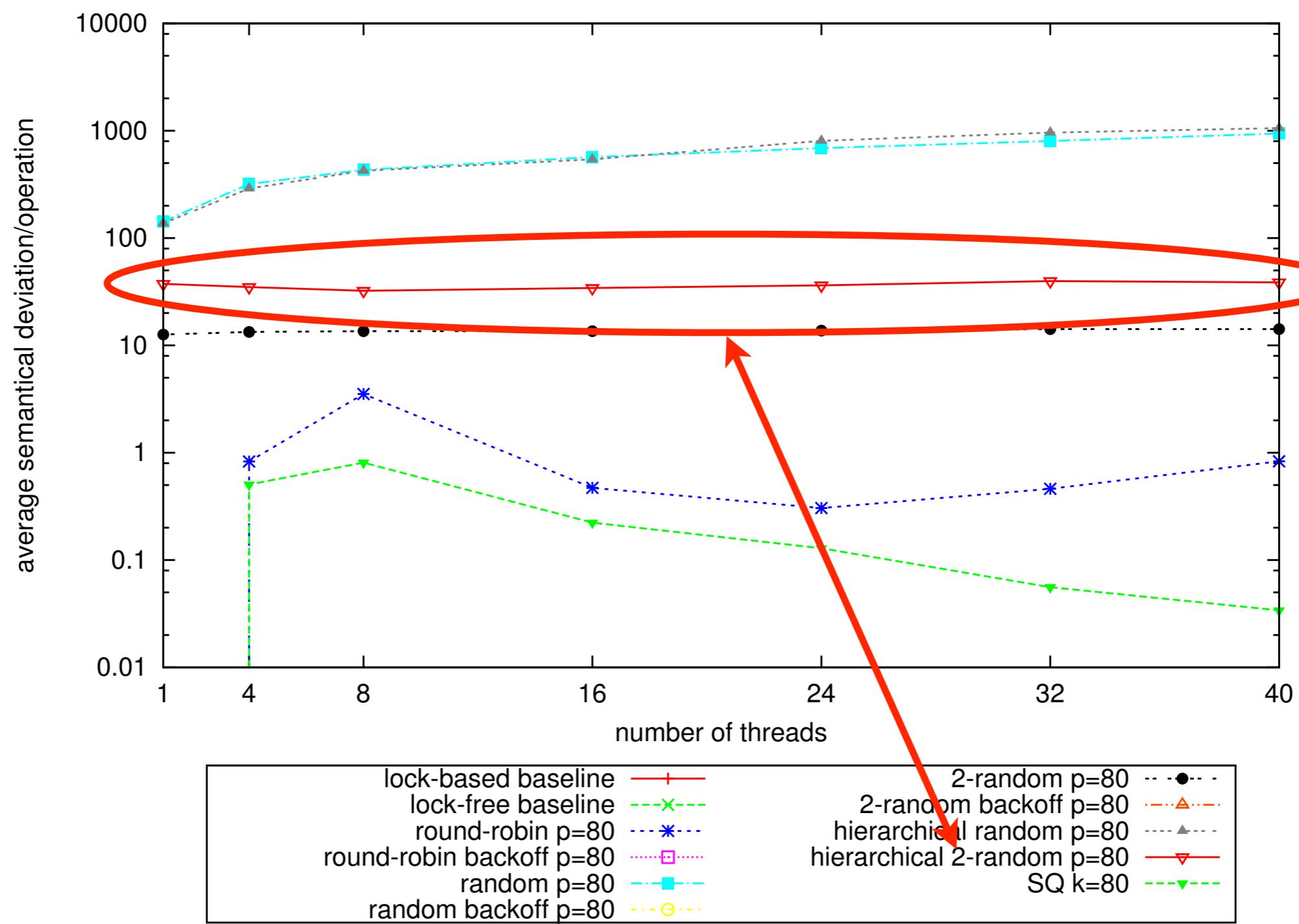
Sequence of Operations (Linearization Points)



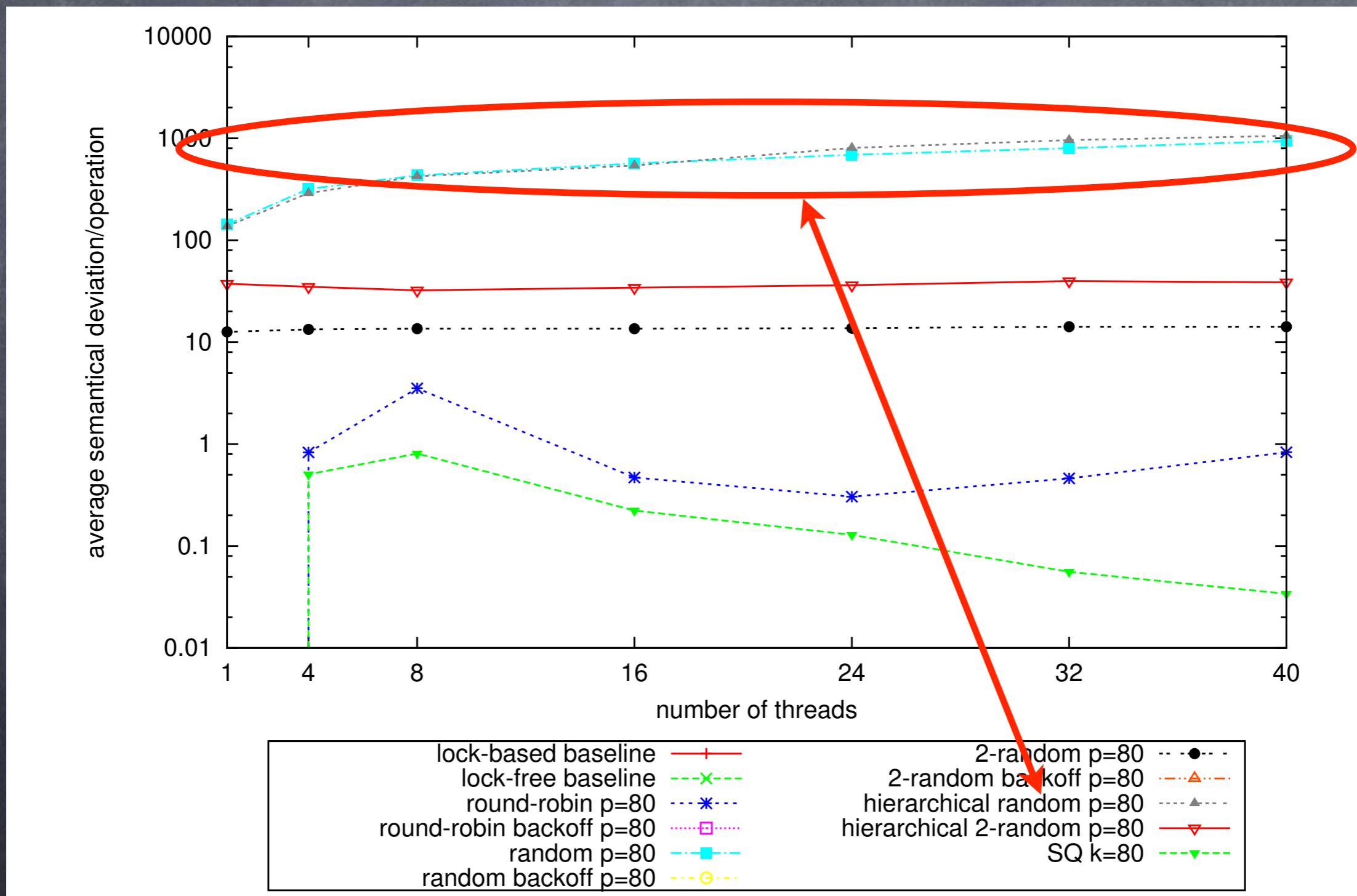
# Average Semantical Deviation



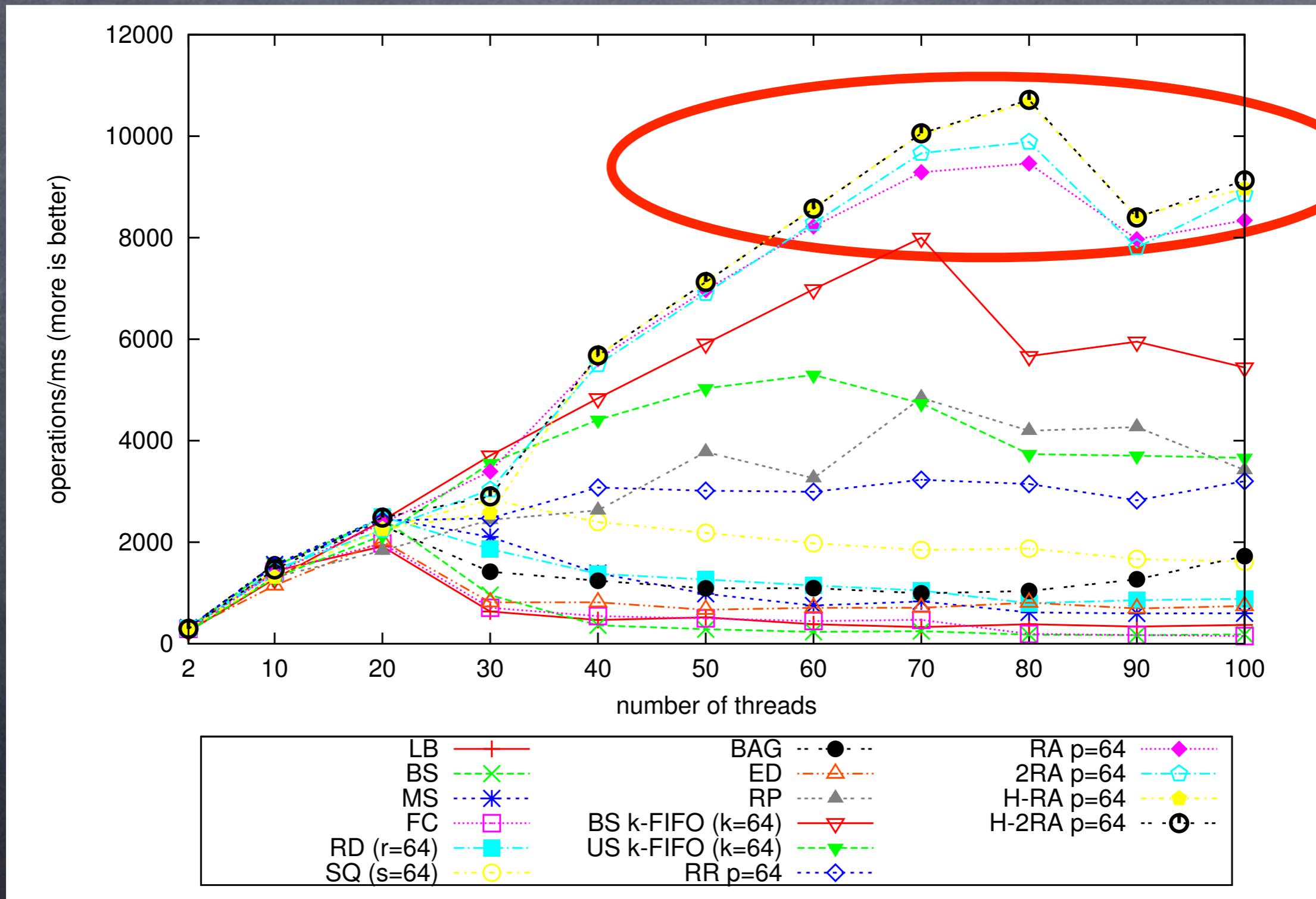
Here  $k$  may be around 40  
on average: best tradeoff



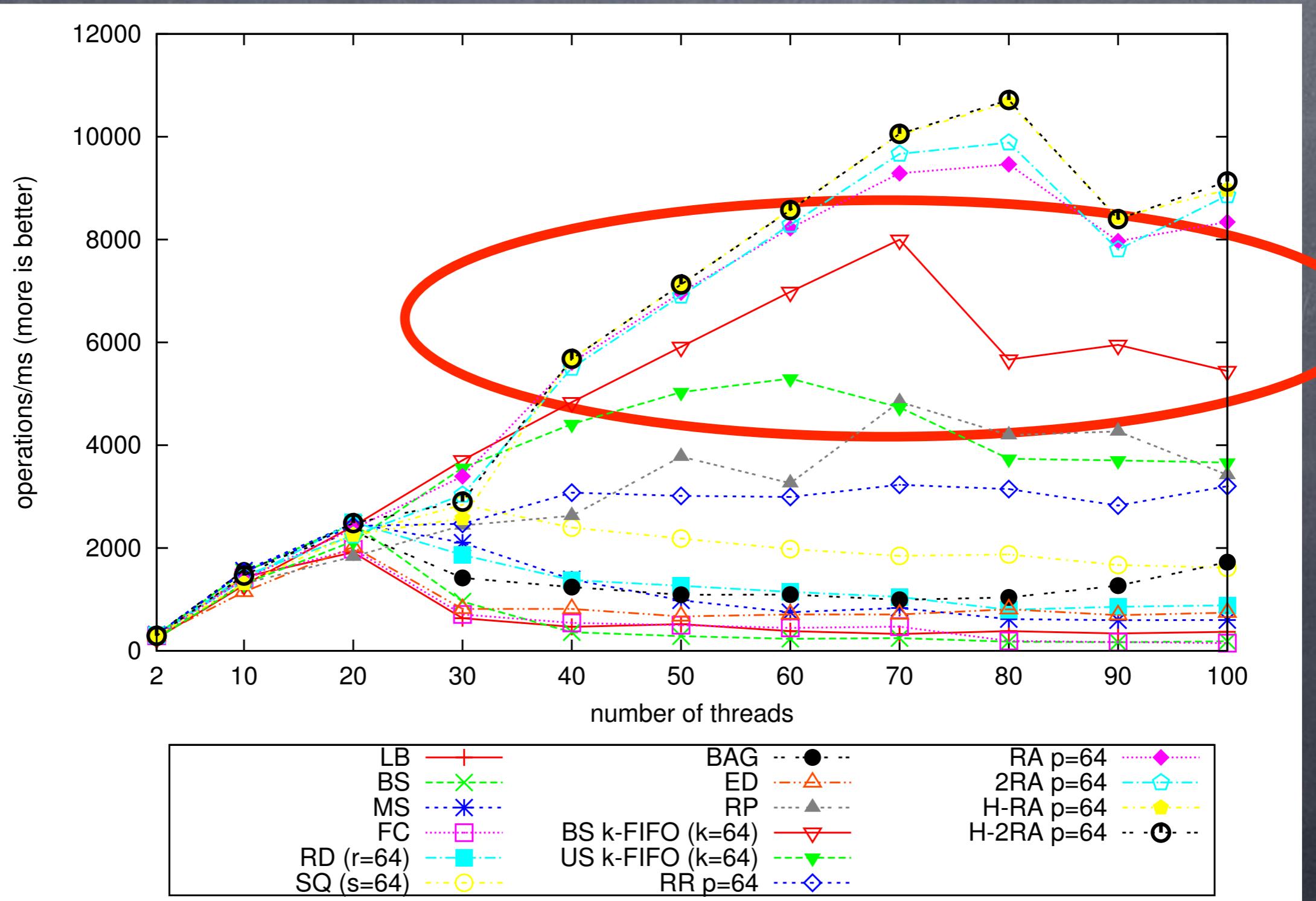
Whereas here  $\kappa$  is one order  
of magnitude bigger w/o gain



# Random vs. d-Random



# Segmented Queues



Back to Correctness?

Questions?

