BackgroundAndResults

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1 1D Poisson Multigrid Solver

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2 Mathematical Preliminaries

The 1D Poisson problem with Dirichlet boundaries seeks to find solutions u to the ordinary differential equation

$$\begin{cases}
-u''(x) = f(x) & x \in (0,1) \\
u(0) = a \\
u(1) = b
\end{cases}$$

where f is a prescribed function and a,b are the prescribed boundary data. This equation has numerous application, such as giving the steady state solution to the 1D heat equation. One of the more simplistic ways to generate a numerical approximation of the solution to this equation is through the method of finitte differences. First, we break the domain [0,1] into n sub-intervals of equal length using grid points $x_j = j \Delta x$, where $0 \le j \le N$ and $\Delta x = 1/N$. In the interior of the grid (i.e. x_j for $1 \le j \le N - 1$, we approximate the second derivative of u via the second-order central difference method

$$\frac{u(x_{j+1}) - 2u(x_j) + u(x_{j-1})}{(\Delta x)^2} \approx u''(x_j) = -f(x_j) \qquad 1 \le j \le N - 1.$$

To this end, we seek a vector v with components v_j such that $v_j \approx u(x_j)$ such that

$$\frac{-v_{j+1} + 2v_j - v_j)}{(\Delta x)^2} = f(x_j), \qquad 1 \le j \le N - 1,$$

and $v_0 = a, v_N = b$. We will write $f_j = f(x_j)$, and g = (a, b).

The central difference condition generates a system of N-1 equations with N-1 unknowns. We can write this in as a tridiagonal linear system

$$\frac{1}{(\Delta x)^2} \begin{bmatrix} 2 & -1 & & & \\ -1 & 2 & -1 & & \\ & -1 & 2 & -1 & \\ & & \ddots & \ddots & \ddots \\ & & & -1 & 2 \end{bmatrix} \begin{bmatrix} v_1 \\ \vdots \\ \vdots \\ v_{N-1} \end{bmatrix} = \begin{bmatrix} f_1 + g_0/(\Delta x)^2 \\ f_2 \\ \vdots \\ f_{N-2} \\ f_{N-1} + g_1/(\Delta x)^2 \end{bmatrix}.$$

We remark that

$$\begin{bmatrix} 2 & -1 & & & \\ -1 & 2 & -1 & & & \\ & -1 & 2 & -1 & & \\ & & \ddots & \ddots & \ddots \\ & & & -1 & 2 \end{bmatrix}$$

is a symmetric, positive-definite matrix of dimension $(N-1) \times (N-1)$.

There are two flavors of numerical methods to solve such a system Ax = b. One is to "directly" solve the system, i.e. invert the matrix. The number of flops (floating point operations to perform this inversion ranges from $\mathcal{O}(N^3)$ (using Gaussian elemination, singular value decomposition) to $\mathcal{O}(N^2)$ (if we exploit the tridiagonal structure and use e.g. banded Cholesky). For such methods, the spatial complexity ranges from $\mathcal{O}(N^2)$ to $\mathcal{O}(N^{3/2})$ for our problem. The second method is to utilize "iterative" solvers. In this case, we generate a sequence of approximate solutions (x_n) such that $x_n \to x$ and x_{n+1} is cost-effective to compute after having x_n . For standard iterations schemes like Jacobi and Gauss-Seidel, we get flop counts of $\mathcal{O}(N^2)$ and spatial complexities of $\mathcal{O}(N)$ (SSOR reduces the former to $\mathcal{O}(N^{3/2})$). So, the iterative solvers are cheaper and at least as fast (and faster without special matrix structure), although they only give an approximate solution to the linear system. They are guarantees to converge for any b and initial guess x_0 under certain conditions on the matrix, such as positive-definiteness. These are more commonly used for large N.

We will focus on the Jacobi method. For our problem, this generates the iteration scheme takes the form

$$v_j^{(n+1)} = \frac{1}{2} \left(v_{j+1}^{(n)} - v_{j-1}^{(n)} + (\Delta x)^2 f_j \right) \qquad 1 \le j \le N - 1.$$

A variant of this, called the weighted or damped Jacobi method, compute the above as an intermediate value then weights it with $v_i^{(n)}$. Explicitly, the scheme is

$$v_j^* = \frac{1}{2} \left(v_{j+1}^{(n)} - v_{j-1}^{(n)} + (\Delta x)^2 f_j \right) \qquad 1 \le j \le N - 1$$
$$v_j^{(n+1)} = (1 - \omega) v_j^{(n)} + \omega v_j^* \qquad 1 \le j \le N - 1,$$

where $\omega \in \mathbb{R}$. It can be shown that $0 < \omega \le 1$ will give convergence (one can go a little higher, but it is approaching 1 asymptotically). A key feature of this method (for ω is the aforementioned range) is that it is very effective at damping high frequency (highly oscillatory) components of the vector that it iterates on. This means that it smooths vectors out well. However, it does not deal with smooth/low frequency components well. In fact, it can be shown that no ω will do this effectively. For this reason, we instead choose ω so that it deals with high frequencies optimally, and it is provable that $\omega = 2/3$ is optimal for our problem.

3 Multigrid

Recall that the damped Jacobi method is effective at dealing with high frequencies, but it does not fare well with low frequency modes. If we take our grid and coarsen it (say we remove every other node), then a smooth wave will look sharper (i.e. it is now high frequency). So, a reasonable idea is to do some Jacobi sweeps, move what we get to a coarser grid (where our low frequencies become high frequencies), do some relaxation there (where we can effectively deal with high frequencies), then come back up to the full grid via an interpolation process. This isn't that far from what we want. We make the following key observation: if Au = f and v is an approximation for u, then Ae = r, where e = u - v is the error and r = f - Av is the residual. We call Ae = r the residual equation. The above tells us that relaxing on Au = f with an arbitrary initial guess v is equivalent to relaxing Ae = r with initial guess 0. So, we modify our prior method as follows: do some Jacobi sweeps, compute the residual, move to the residual a coarser grid, do some relaxation on the residual equation (this will give us an approimation to the error), bring our error back up to the full grid via an interpolation process, correct our original Jacobi sweep with oour refined error, then do some more relaxation on this quantity.

This is the building block of multigrid. However, one we go down to the coarse grid, why do we stop there? We could continue this process recursively until we get to a grid where we can do a cheap direct solve. This is how we get our multigrid method, which takes the following form:

- 1. We start with our problem on a given grid.
- 2. We perform a few iterations of weighted Jacobi (damping the high frequencies).
- 3. We compute the residual f Av, where v is our approximate solution from step 2.
- 4. We move the residual to a coarser grid (where the problematic low frequencies are now high frequency and can be dealt with via weighted Jacobi). We can do this transition to the coarseion grid either by restriction/projection or using averaging methods.
- 5. If the grid is coarse enough, we do a direct solve on the residual equation (which is cheap due to the size of the problem). If not, then we move back to step 2 on our current grid (now looking at the residual equation) and proceed recursively until the grid is coarse enough to do our direct solve. Then, we move onto step 6.
- 6. We move back to the refined grid (one step up) via interpolation between data points.
- 7. We correct what we got in step two with our computation in step 6.
- 8. We do a few more iterations of weighted Jacobi.

This is called a V-cycle. We can use them to generate our multigrid scheme by taking the return from one V-cycle and making it the initial guess of the next. The amazing thing about multigrid is that both the flop count and spatial complexity are $\mathcal{O}(n)$.

4 Results

The code in 1DPoissonMG.py, along with containing the core multigrid solver using the function V-cycle, possesses various other functions with utility:

1. solver_comp computes the execution times for our multigrid method and NumPy direct solver, as well as the absolute 2-norm error for an array of grid sizes. Then, it prints this data in two tables (one table for execution times, one table for errors). One can use this function to verify the numerical accuracy of multigrid while featuring a much lower execution time than the Numpy direct solver (direct solver takes 49 seconds on a 16385 grid (N=14); I have

- not pushed it further, but it would likely take at least an order of magnitude longer).
- 2. relax_test computes the absolute error for an array of relaxation parameters, then them prints a table of these errors and outputs a plot of this data. From our prior discussion, we know that $\omega = 2/3$ is theoretically optimal, but this is not so noticable for small perturbations.
- 3. mg_test computes the execution time for multigrid for an array of grid sizes and prints this data in a table. It is similar to solver_comp, but it allows one to solve much larger problems (since we are not using the direct solver at the top level). This tells us that multigrid method can run on up to a 524289 (N=19) grid in under 1 minute.

All of these functions are contained in the Multigrid() class and have precise documentation in 1DPoissonMG.py. Here is some sample input and output:

Input to solver_comp:

drive = Multigrid() # instantiate multigrid class

N_array = [2**n for n in range(6,14)] # array of grid sizes
drive.solver_comp(N_array)

Output:

+		•	+
1	Grid Size	_	NumPy Direct Solver Execution Time (s)
	65	0.008763785001065116	0.00036002200067741796
-	129	0.012847697998950025	0.0004634220022126101
-	257	0.023098724999726983	0.001266939998458838
-	513	0.04389215799892554	0.011966152000240982
-	1025	0.08147994899991318	0.01582442700237152
-	2049	0.16037838199918042	0.09368109899878618
	4097	0.3648365989974991	0.6735369830021227
1	8193	0.6315701820021786	2.6383576019979955

+-		-4.	
 	Grid Size	 	Absolute Euclidean Error
Τ-		т.	
	65	1	9.556254761269408e-14
	129	1	3.6794935403782707e-13
	257	1	1.840396983851491e-12
	513	1	7.767679988043338e-12
	1025	1	4.857682116518334e-11
-	2049	1	1.0490398063334259e-10
	4097	1	3.743260606153312e-10
	8193	1	5.097226123983849e-10

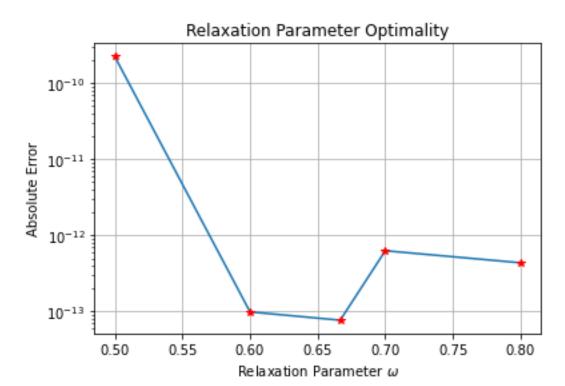
Input to relax_test:

drive = Multigrid() # instantiate multigrid class

 $omega_array = [0.5, 0.6, 2/3, 0.7, 0.8]$ # array of relaxation parameters

Output:

+			-+		+
	Grid Size	Relaxation Parameter		Error	
					_
- 1	257	0.5		2.2155681694916334e-10	
-	257	0.6		9.908908329069226e-14	l
- [257	0.666666666666666		7.713323153041377e-14	l
-	257	0.7		6.282443532386848e-13	l
١	257	0.8		4.367908411594411e-13	١
+		+	-+		+



Input to mg_test:

drive = Multigrid() # instantiate multigrid class

 $N_{array} = [2**n for n in range(6,16)] # array of grid sizes drive.mg_limit(<math>N_{array}$)

Output:

+			-+-					-+
	Grid	Size	1	Multigrid	Execution	Time	(s)	1
+			-+-					-+
-		65	1	0	.0097100470	001135	5536	1

	129	0.015764535000926116
	257	0.02671550000013667
	513	0.04875404700214858
	1025	0.08642815699931816
	2049	0.17044590799923753
	4097	0.34982147100163274
	8193	0.6816639559983741
	16385	1.3509082869968552
	32769	2.6054512939990673
+		

[]: