Increasing Performance in Byzantine Fault-Tolerant Systems with On-Demand Replica Consistency

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> EuroSys April 11, 2011



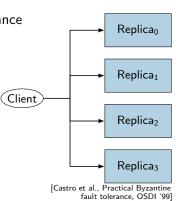


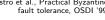


Byzantine Fault Tolerance (BFT)

- Agreement-based Byzantine fault tolerance
 - 3f + 1 replicas to tolerate f faults
 - BFT agreement protocol
 - Client-side voting
- Drawbacks
 - High resource usage
 - Performance overhead for agreement

REFIT Project Research Goal Resource-efficient BFT systems







Making BFT Systems More Resource-Efficient

REFIT Project Research Goal Resource-efficient BFT systems





Optimizing Resource Usage

- Reduced number of replicas
- Same performance
- Recent examples
 - SPARE [Distler et al., NDSS '11]
 - ZZ [Wood et al., EuroSys '11]

Optimizing **Performance**

- Default number of replicas
- Increased performance



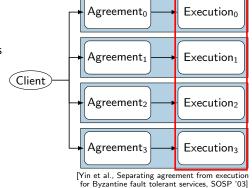


Where to Optimize?

- Agreement stage
 - BFT agreement protocol
 - Sequence of agreed requests
- Execution stage

Observations

- Service application
- Request processing



- Past optimizations have significantly reduced agreement overhead
- Non-trivial services: response times are dominated by execution stage

ODRC Approach Reducing the load on the execution stage



Basic Approach

- Traditional BFT systems
 - All 3f + 1 replicas process all requests
 - Client waits for f + 1 identical replies

Insight

In the **absence of faults**, a client only needs f + 1 replies to make progress

- ODRC
 - Each request is processed by only f + 1 replicas
 - Load distribution across replicas
 - Additional replicas process the request in case of faults



Talk Outline

ODRC

- Selective Request Execution
- On-Demand Replica Consistency
- Evaluation
- Conclusion



Talk Outline

ODRC

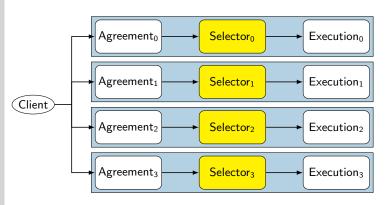
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Architecture

Selector

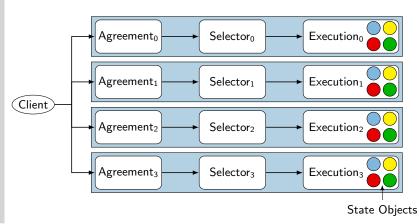
- Selects requests for execution
- Stores requests that have not been selected





Application State

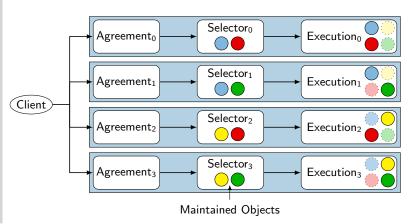
- Set of objects
 - Examples: files, directories, ...
 - Assumption: requests carry information about object access



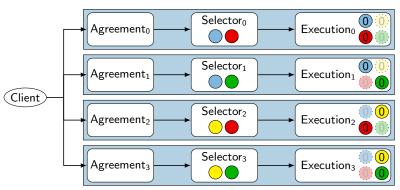


Application State Distribution

- Object distribution scheme
 - **Each** object is **maintained** on f + 1 replicas, **unmaintained** on others
 - State of unmaintained objects may be outdated

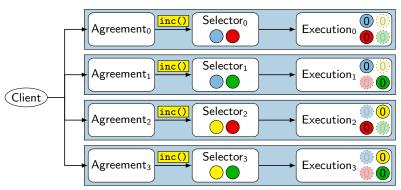






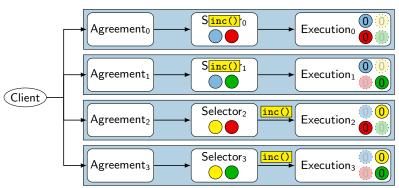
- Garbage collection of stored requests
 - Periodic object checkpoints of maintained objects
 - Stable checkpoint: f + 1 identical checkpoints





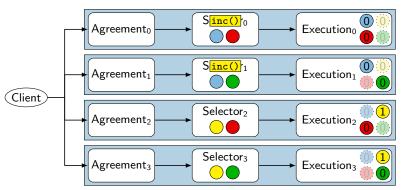
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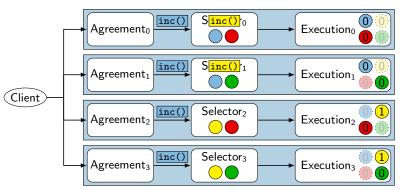
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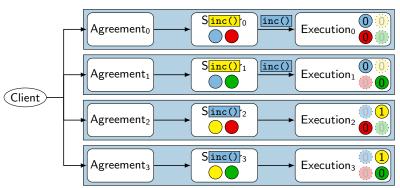
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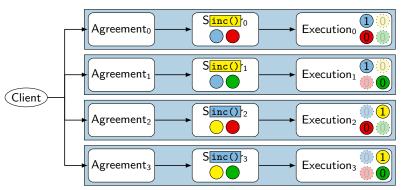
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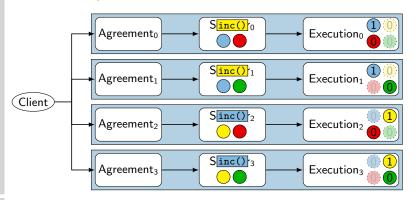
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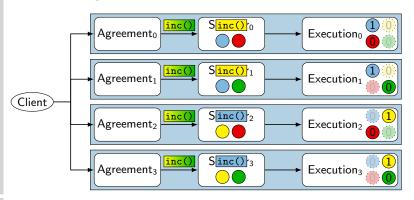


- Access of multiple objects
 - Only unmaintained objects ⇒ store request
 - At least one maintained object
 - Update unmaintained objects
 - Process request





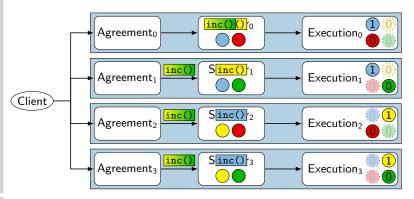
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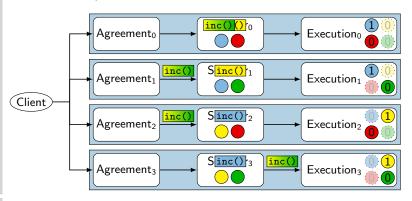
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Tobias Distler (distler@cs.fau.de)



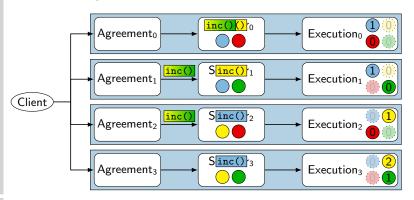


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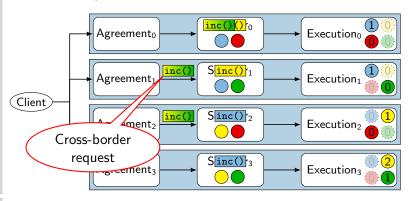


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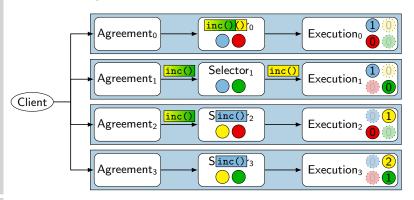


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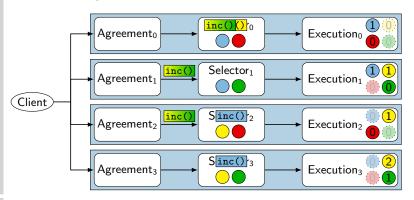


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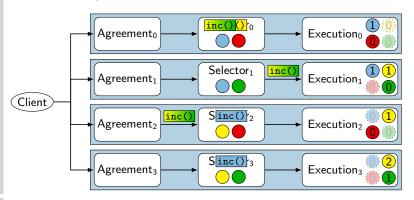


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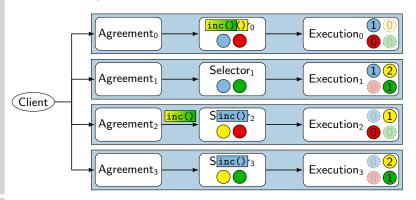


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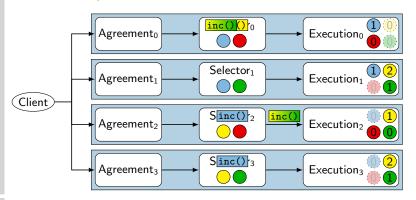


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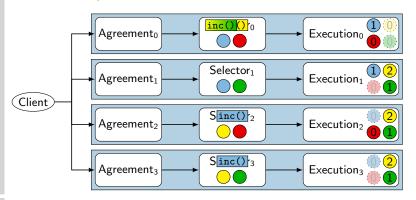


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Cross-Border Requests

- Additional consistency overhead
 - Processed by more than f+1 replicas
 - Goal: minimize number of cross-border requests





- Optimized object distribution
 - Application-centric strategies
 - Consider object dependencies

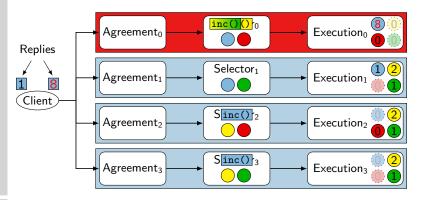




- Example: Network File System (NFS)
 - Assign files and their parent directories to the same replicas
 - Subdirectories may be assigned to different replicas

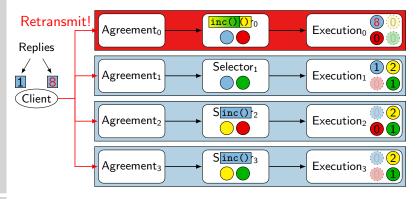


- Providing additional replies on demand
 - Standard BFT clients
 - Request retransmission after timeout
 - Additional replicas process the request



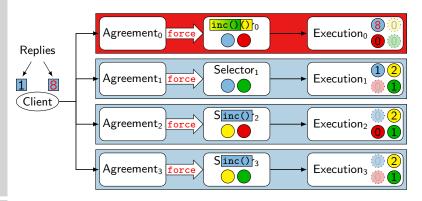


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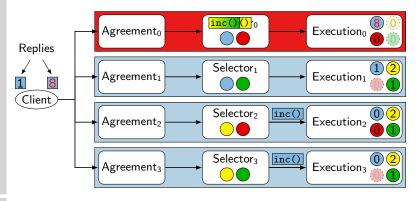


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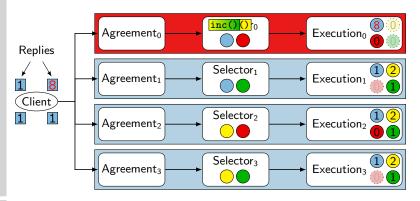


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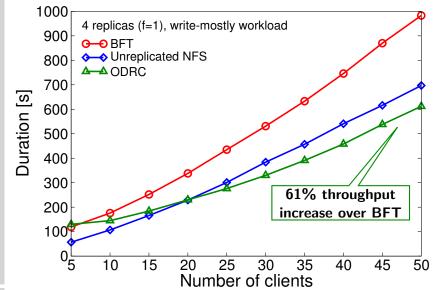
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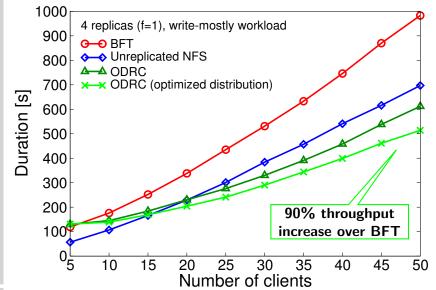


NFS Evaluation: Postmark Benchmark



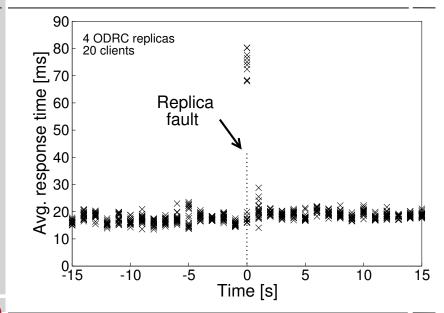


NFS Evaluation: Postmark Benchmark





NFS Evaluation: Append-Only Micro-Benchmark





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Conclusion

- Execution matters!
- Traditional BFT systems
 - All replicas process all requests
 - Consistency overhead
- ODRC
 - Selective request execution based on object access
 - On-demand replica consistency
 - Additional replies in case of faults

Thank you very much.

Questions?

