

Mathematical Logic

homework 4

Problem 1 Can you derive the rule of contradiction from the modified contradiction?

Solution:

1. $\Gamma \quad \neg\varphi \quad \psi$ (premise)
2. $\Gamma \quad \neg\varphi \quad \neg\psi$ (premise)
3. $\Gamma \quad \neg\varphi \quad \varphi$ (modified contradiction by 1 and 2)
4. $\Gamma \quad \varphi \quad \varphi$ (assumption)
5. $\Gamma \quad \varphi$ (modified contradiction by 3 and 4)

Problem 2 Prove: (a) $\frac{\Gamma \quad \varphi}{\Gamma \quad \neg\neg\varphi}$ (b) $\frac{\Gamma \quad \neg\neg\varphi}{\Gamma \quad \varphi}$

Solution: (a)

1. $\Gamma \quad \varphi$ (premise)
2. $\Gamma \quad \neg\varphi \quad \varphi$ (assumption)
3. $\Gamma \quad \neg\varphi \quad \neg\varphi$ (assumption)
4. $\Gamma \quad \neg\varphi \quad \neg\neg\varphi$ (modified contradiction by 2 and 3)
5. $\Gamma \quad \neg\neg\varphi \quad \neg\neg\varphi$ (assumption)
6. $\Gamma \quad \neg\neg\varphi$ (case analysis by 4 and 5)

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(b)

1. $\Gamma \quad \neg\neg\varphi$ (premise)
2. $\Gamma \quad \neg\varphi \quad \neg\neg\varphi$ (assumption)
3. $\Gamma \quad \neg\varphi \quad \neg\varphi$ (assumption)
4. $\Gamma \quad \neg\varphi \quad \varphi$ (modified contradiction by 2 and 3)
5. $\Gamma \quad \varphi \quad \varphi$ (assumption)
6. $\Gamma \quad \varphi$ (case analysis by 4 and 5)

Problem 3 Is the following derivable?

$$\overline{\Gamma \quad \exists x\varphi \quad \forall x\varphi}$$

Solution: We let

$$\mathfrak{I} = \{\{0^{\mathfrak{I}}, 1^{\mathfrak{I}}\}, R^{\mathfrak{I}}, \beta\}$$

and $0^{\mathfrak{I}} \in R^{\mathfrak{I}}, 1^{\mathfrak{I}} \notin R^{\mathfrak{I}}$. if $\varphi := Rx$.

Then $\mathfrak{I} \frac{0}{v_0} \frac{1}{v_1} \models \Gamma \wedge \exists x\varphi$. $\mathfrak{I} \frac{0}{v_0} \frac{1}{v_1} \not\models \forall x\varphi$.

So $\{\Gamma \exists x\varphi\} \not\models \forall x\varphi$, it is not correct, so it is not derivable.

Problem 4 Let $S = \{R\}$ with unary relation symbol R .

$$\Phi = \{\exists xRx\} \cup \{\neg Ry | \text{for every variable } y\}$$

Prove that:

- Φ is consistent.
- There is no term $t \in T^S$, with $\Phi \vdash Rt$.

Solution: (1) We construct

$$\mathfrak{I} = \{\{0^{\mathfrak{I}}, 1^{\mathfrak{I}}\}, R^{\mathfrak{I}}, \beta\}$$

where $\beta(v_i) = 1^{\mathfrak{I}}$, and $0^{\mathfrak{I}} \in R^{\mathfrak{I}}, 1^{\mathfrak{I}} \notin R^{\mathfrak{I}}$.

(i) because $0^{\mathfrak{J}} \in R^{\mathfrak{J}}$, so $\exists x Rx$ is satisfiable.

(ii) and $\beta(v_i) = 1 \notin R^{\mathfrak{J}}$, so $\{\neg Ry \mid \text{for every variable } y\}$ is satisfiable.

Φ is satisfiable $\Rightarrow \Phi$ is consistent (**Lemma 2.5**).

(2) We construct

$$\mathfrak{J} = \{\{1^{\mathfrak{J}}\}, R^{\mathfrak{J}}, \beta\}$$

Where $\beta(v_i) = 1^{\mathfrak{J}}$, and $1^{\mathfrak{J}} \notin R^{\mathfrak{J}}$. Because $\{1\}$ is universe.

So no term $t \in T^S$, with $\Phi \vdash Rt$. It is satisfiable.

Problem 5 Let $S = \{R\}$ with unary relation symbol R . and

$$\Phi = \{Rx \vee Ry\}$$

Prove that:

- Φ is consistent.
- $\Phi \not\models Rx$, $\Phi \not\models Ry$.
- $\mathfrak{J}^{\Phi} \not\models \Phi$

Solution: (1) Because $\Phi \not\models Rx \wedge Ry$, so Φ is not inconsistent.

(2) Construct

$$\mathfrak{J} = \{\{0^{\mathfrak{J}}, 1^{\mathfrak{J}}\}, R^{\mathfrak{J}}, \beta\}$$

and $0^{\mathfrak{J}} \in R^{\mathfrak{J}}$, $1^{\mathfrak{J}} \notin R^{\mathfrak{J}}$.

Then $\mathfrak{J} \frac{0 \ 1}{x \ y} \models \Phi$, and $\mathfrak{J} \frac{0 \ 1}{x \ y} \not\models Ry$.

$\Phi \not\models Ry \Rightarrow \Phi \vdash Ry$ **is not correct**, so $\Phi \not\models Ry$.

same as $\Phi \not\models Rx$.

(3) Because $\Phi \not\models Rx$, so does not exist $\bar{t} \in R^{\mathfrak{J}^{\Phi}}$. then $\mathfrak{J}^{\Phi} \not\models Rx$ and $\mathfrak{J}^{\Phi} \not\models Ry \Rightarrow \mathfrak{J}^{\Phi} \not\models Rx \vee Ry$