

# EXERCISE NO 8

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MTH 6312: MÉTHODES STATISTIQUES D'APPRENTISSAGE

November 4, 2015

**Exercise 1**

**Solution 1** Here the optimization problem is to find  $\beta$  that solves :

$$\begin{cases} \min_{\beta, \beta_0} |\beta| \\ \text{subject to } y_i(x_i\beta + \beta_0) \geq 1, i = 1, \dots, N \end{cases}$$

But this case is very simple and : we clearly know which vectors will be the support vectors and what will be the optimal separating plane. It will be the line defined by  $x = 7$ . The margin will be 3 and thus,  $\beta$  will be  $\frac{1}{3}$  and  $\beta_0, -\frac{7}{3}$ .

The optimization problem admits 1 solution, we could easily check that this one is a solution.

The image in annex sums up the case.

**Exercise 2**

**Solution 2** These 2 cost functions don't seem equivalent. The Gaussian process cost function is the joint likelihood and the term in  $\frac{1}{2}\mathbf{f}^\top \mathbf{K}^{-1}\mathbf{f}$  is the contribution of the prior to this joint likelihood.

Gaussian processes :

- + it is a fully statistical method, hence has an exact statistical interpretation
- + We will get exact confidence intervals for prediction.
- + we will be able to learn hyperparameters of the model (parameters of the Kernel K from the data, regularization parameter) notably by cross validation (because it is a statistical

method)

SVM :

- geometrical basis/interpretation
- No possible simple exact hyperparameters optimizations because cross validation is not relevant with SVM.
- No exact statistical interpretation

### Exercise 3

**Solution 3** Let  $\mathbf{K}$  be the matrix defined by  $\mathbf{K}_{ij} = K(x_i, x_j) = \phi(x_i)^\top \phi(x_j)$ .

$\forall (i, j) \in 1, \dots, N$ ,  $\mathbf{K}_{ij} = \phi(x_i)^\top \phi(x_j) = \phi(x_j)^\top \phi(x_i)$  if we transpose this scalar quantity.

Ant trivially,  $\forall (i, j) \in 1, \dots, N$ ,  $\mathbf{K}_{ij} = \mathbf{K}_{ji}$ , that is to say

$\mathbf{K}$  is symmetric

$$\forall i \in 1, \dots, N, \phi(x_i) = (\phi_1(x_i), \phi_2(x_i), \dots, \phi_q(x_i))^\top$$

### Exercise 4

**Solution 4** I had difficulties finding answers and solutions to this problem. I tried to find an optimization problem whose solution is the k-NN matrix and then tried to transform this

constrained optimization problem to a simple optimization problem (without constraints).

Using the same notations of the question, we can easily remark that finding  $\mathbf{M}$  the 2-NN neighbors matrix is equivalent to finding a solution to the following optimization problem :

$$\begin{cases} \min_{\mathbf{M} \in \{0,1\}^{n \times n}} \sum_{i=1}^n \sum_{j=1}^n \mathbf{M}_{ij} \mathbf{W}_{ij} \\ \text{subject to } \sum_{j=1}^n \mathbf{M}_{ij} = 2 \text{ and } \mathbf{M}_{ii} = 0 \forall i \in 1, \dots, N \end{cases}$$

In fact, here  $\sum_{j=1}^n \mathbf{M}_{ij} \mathbf{W}_{ij}$  is the diagonal element  $(i, i)$  of  $\mathbf{M}\mathbf{W}$  (as  $\mathbf{W}$  is symmetric), that is to say the total sum of the weights between the  $i^{th}$  data point to its 2-NN neighbors. This is a positive sum (as  $\mathbf{M}$  can only take values in  $\{0, 1\}$ ). Thus, minimizing all these quantities is equivalent to minimize the sum of all these quantities. This is how I got to this optimization problem for finding the exact 2-NN neighbors matrix. The constraints come from the fact that :

- (i) By line in  $\mathbf{M}$ , we must have exactly 2 1 and 2 0
- (ii) Diagonal elements must be set to 0 as a data point can't be among his own nearest neighbors.

We can rewrite this optimization problem with the matrix Trace function as :

$$\begin{cases} \min_{\mathbf{M} \in \{0,1\}^{n \times n}} \text{Tr}(\mathbf{M}\mathbf{W}) \\ \text{subject to } \sum_{j=1}^n \mathbf{M}_{ij} = 2 \text{ and } \mathbf{M}_{ii} = 0 \forall i \in 1, \dots, N \end{cases}$$

We could transform this optimization problem into an optimization problem without constraints. We should add a quantity  $Q(\mathbf{M})$  to the quantity we must optimize that satisfy the following constraints :

(ii)  $Q(\mathbf{M})$  is zero if all constraints on  $\mathbf{M}$  are respected (iii)  $Q(\mathbf{M})$  is at least the maximum value taken by  $Tr(\mathbf{M}\mathbf{W}) = \sum_{i=1}^n \sum_{j=1}^n \mathbf{M}_{ij} \mathbf{W}_{ij}$  (which is  $\sum_{i=1}^n \sum_{j=1}^n \mathbf{W}_{ij}$ , 2 \* the total sum of weights) without constraints on  $\mathbf{M}$  to avoid tradeoff effects.

This gives us the following cost function to minimize :

$$\begin{aligned} L(\mathbf{M}) &= Tr(\mathbf{M}\mathbf{W}) + \sum_{i=1}^n \sum_{j=1}^n \mathbf{W}_{ij} \times \left( \sum_{i=1}^n \left( \sum_{j=1}^n \mathbf{M}_{ij} - 2 \right)^2 + \sum_{j=1}^n \mathbf{M}_{ii} \right) \\ &= Tr(\mathbf{M}\mathbf{W}) + \sum_{i=1}^n \sum_{j=1}^n \mathbf{W}_{ij} \times \left( \sum_{i=1}^n \left( \sum_{j=1}^n \mathbf{M}_{ij} - 2 \right)^2 + Tr(\mathbf{M}) \right) \end{aligned} \quad (1)$$

- The term  $Tr(\mathbf{M})$  will force the Trace of  $\hat{\mathbf{M}}$  the solution of the problem to be 0 and  $\mathbf{M}$  to have only zeros coefficients on its diagonal.
- the term  $\sum_{i=1}^n \left( \sum_{j=1}^n \mathbf{M}_{ij} - 2 \right)^2$  will force the lines sums of  $\hat{\mathbf{M}}$  to be exactly zero. We could eventually replace in this function  $\sum_{i=1}^n \sum_{j=1}^n \mathbf{W}_{ij}$  by a constant since it is a problem constant and  $\sum_{i=1}^n \left( \sum_{j=1}^n \mathbf{M}_{ij} - 2 \right)^2$  by a trace of a more complex matrix.

The optimization problem for getting a symmetric 2-NN is the same than previously but with one more constraints on the  $\hat{\mathbf{M}}$  coefficients :

$$\left\{ \begin{array}{l} \min_{M \in \{0,1\}^{n \times n}} Tr(MW) \\ \text{subject to } \sum_{j=1}^n M_{ij} = 2, \quad M_{ii} = 0 \quad \forall i \in 1, \dots, N \\ \text{and } M_{ij} = M_{ji} \quad \forall (i, j) \in 1, \dots, N \end{array} \right.$$

Then, we get a modified cost function including that forces M to be symmetric :

$$L(M) = Tr(MW) + \sum_{i=1}^n \sum_{j=1}^n w_{ij} \times \left( \sum_{i=1}^n \left( \sum_{j=1}^n M_{ij} - 2 \right)^2 + Tr(M) + \sum_{i=1}^n \sum_{j=1}^n (M_{ij} - M_{ji})^2 \right) \quad (2)$$

These cost functions can lead to our desired solutions for any n and k (if  $k \leq n$ )...