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CSC460 Project 4

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**i) A final ER diagram is shown on the last page.**

**ii) Logical Database Design**

After converting out ER diagram to tables and doing the normalization, we have the following schema:

**Appointment**: {apptID, apptTime, apptSuccessful, apptType, customerID, employeeID},

**Customer** {customerID, firstName, lastName, address, dateOfBirth},

**Employee** {employeeID, firstName, lastName, jobID, deptID},

**Job** {jobID, title, salary},

**Department**{deptID, name, fee},

**StateID**{stateIDNo, issueDate, expireDate, customerID, deptID},

**Permit**{permitNo, issueDate, expireDate, class, customerID, deptID},

**License**{stateIDNo, issueDate, expireDate, class, customerID, deptID},

**Registration**{stateIDNo, issueDate, expireDate, customerID, deptID},

where customerID, employeeID are foreign keys in Appointment, deptID and jobID are foreign keys in Employee, and customerID and deptID are foreign keys in StateID, Permit, License and Registration.

**iii) Normalization Analysis**

**Appointment (from draft):** {apptID, apptTime, apptSuccessful, apptType, fee, deptID, customerID, employeeID}

Functional dependencies:

{apptID} —> {apptTime, apptSuccessful, apptType, fee, deptID, customerID, employeeID}

{deptID} —> {fee}

{employeeID, fee} —> {deptID}

ApptID is the primary key, so apptID functionally determines everything else. If two relations agree on their deptID’s they must also agree on their fees, because fees belong to the department and departments only have a single service that costs a fixed amount. Also, because employees only belong to a single department, and the four fees for the departments happen to be unique, employeeID and fee both functionally determine departmentID. That fee functionally determines deptID is somewhat coincidental, because if one service changed their fee to be the same as another, the FD wouldn’t be true anymore.

Closures:

{apptID}+ = {apptID, apptTime, apptSuccessful, apptType, fee, deptID, customerID, employeeID}

{apptTime}+ = {apptTime}

{apptSuccessful}+ = {apptSuccessful}

{apptType}+ = {apptType}

{fee}+ = {fee, deptID}

{deptID}+ = {deptID, fee}

{customerID}+ = {customerID}

{employeeID}+ ={employeeID, deptID, fee}

After closure I found that emplyoeeID also functionally determines fee.

Data Normalization

None of the attributes are set-valued so this is in 1NF.

The non-prime attributes here are apptTime, apptSuccessful, apptType, fee, deptID, customerID, and employeeID, of those deptID and fee are only partially dependent on apptID, deptID should be removed and fee should go into the department table, that information can be found using employeeID:

Appointment: {apptID, apptTime, apptSuccessful, apptType, customerID, employeeID}

{apptID}+ = {apptID, apptTime, apptSuccessful, apptType, fee, deptID, customerID, employeeID}

{apptTime}+ = {apptTime}

{apptSuccessful}+ = {apptSuccessful}

{apptType}+ = {apptType}

{customerID}+ = {customerID}

{employeeID}+ ={employeeID}

Now all the FDs aside from apptID are trivial or the superkey apptID, so this is in BCNF.

**Customer** {customerID, firstName, lastName, address, dateOfBirth}

Functional dependencies:

{customerID} —> {firstName, lastName, address, dateOfBirth}

CustomerID is the primary key, so it functionally determines everything else. Here, there are no other functional dependencies because people can have the same first but not last name, and vice versa, multiple different people may live at the same address, and not everyone in the database named ‘John’ would live at the same address. If you had George Foreman in your database you could even have people with the same first and last names and address in the database but not the same DOB.

Closures:

{customerID}+ = {customerID, firstName, lastName, address, dateOfBirth}

{firstName}+ = {firstName}

{lastName}+ = {lastName}

{address}+ = {address}

{dateOfBirth}+ = {dateOfBirth}

Data Normalization

None of the attributes are set-valued so this is in 1NF. The only non-trivial attribute in Customer is the primary key, which is a superkey. That means this is in BCNF and therefore in 2NF as well.

**Employee** {employeeID, firstName, lastName, jobID, deptID}

Functional dependencies:

{employeeID} —> { irstName, lastName, jobID, deptID}

EmployeeID is the primary key, so it functionally determines everything else. Here, there are no other functional dependencies because people can have the same first but not last name, and vice versa, multiple different people may have the same job or work in the same department, and not everyone in the database named ‘John’ would do the same job or be in the same department. If you had George Foreman in your database you could even have people with the same first and last names and jobID in the database but not the same department, or vice versa.

Closures:

{employeeID}+ = {firstName, lastName, jobID, deptID }

{firstName}+ = {firstName}

{lastName}+ = {lastName}

{jobID }+ = {jobID}

{deptID }+ = {deptID}

Data Normalization

None of the attributes are set-valued so this is in 1NF. The only non-trivial attribute in Employee is the primary key, which is a superkey. That means this is in BCNF and therefore in 2NF as well.

**Job** {jobID, title, salary}

Functional dependencies:

{jobID} —> { title, salary}

JobID is the primary key, so it functionally determines everything else. Here, there are no other functional dependencies because two jobs can have the same salary. As the rubric was laid out, it seemed that a job title always had the same salary, but a situation to have this not be true could be one where raises are given out but adding new salaries to the same title (desk clerk at 40,000 might be starting salary but perhaps it could go up to 80,000).

Closures:

{jobID}+ = {jobID, title, salary }

{title}+ = {title}

{lastName}+ = {lastName}

{jobID }+ = {jobID}

{deptID }+ = {deptID}

Data Normalization

None of the attributes are set-valued so this is in 1NF. The only non-trivial attribute in Employee is the primary key, which is a superkey. That means this is in BCNF and therefore in 2NF as well.

**Department**{deptID, name, fee}

Functional dependencies:

{deptID} 🡪 { name, fee }

{name} 🡪{fee}

DeptID is the primary key, so it functionally determines everything else. The same department name would also not be used twice in the DMV, so name is a candidate key in this relation.

Closures:

{deptID, name}+ = { deptID, name, fee }

Data Normalization

None of the attributes are set-valued so this is in 1NF. The two non-trivial attributes in Department are both superkeys because both deptID and name are candidate keys. That means this is in BCNF and therefore in 2NF as well.

**StateID**{stateIDNo, issueDate, expireDate, customerID, deptID}

**Permit**{permitNo, issueDate, expireDate, class, customerID, deptID}

**License**{stateIDNo, issueDate, expireDate, class, customerID, deptID}

**Registration**{stateIDNo, issueDate, expireDate, customerID, deptID}

**iv) Query Description**