

國立陽明交通大學考試試卷

Exam Paper of National Yang Ming Chiao Tung University



課程名稱 Course Name PDA

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班 別 Department

日 期 Date 112/11/6

※ 考試作弊者將受記大過以上處分

※ If you cheat in exam, you will be punished

題號 NO	分數 Score
1	10
2	10
3	10
4	5
5	10
6	25
7	10+2
8	
9	
10	
11	
12	
13	
14	
15	
總 分 Total	80

1. False

10 若同樣儲存一串資料，

每一輪都要進行插入或刪除某一筆 data \rightarrow 此時使用 linklist 只要 $O(1)$ time complexity

若用 array or vector implement

則需要較高的 time complexity

因為可能會需要整個 array/vector copy 到一個新的 array/vector

$\Rightarrow O(n)$

2. No,

10 因為 TSP 屬於 NP-problem, HC 也是

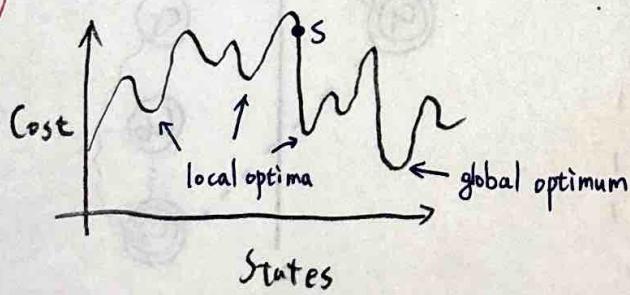
且 HC 可以 reduce 成 TSP problem

\rightarrow 因此此演算法若能在 polynomial time 內解決 TSP problem

也要能在 polynomial time 內解決 HC problem

3. 在當個 loop 的 solution 變差時，仍讓它有機會接受此 solution，

如左圖，當 s 在 local optima 時，需要有 up-hill，才有可能到達 global optimum



但也不能保證 simulated annealing 一定可以得到 global optimum，可能在得到前就因終止條件而停止。

4.

不一定

5 node a 交換的 gain : $3 - 6 = -3$

node b , : $4 - 2 = 2$

但還需要扣掉 a, b 之間的連線的 gain

\rightarrow 不確定是否能 reduce cut cost

5.

(0) balanced bipartition :

$$rW - S_{\max} \leq |A| \leq rW + S_{\max}$$

r : balanced factor, $W = |A| + |B|$, $S_{\max} = \max_i S(i)$ $\Rightarrow r=0.3, W=15, S_{\max}=4$

$$\Rightarrow 0.3 \times 15 - 4 = 0.5 \leq |A| \leq 0.3 \times 15 + 4 = 8.5, 0.7 \times 15 - 4 = 6.5 \leq |B| \leq 0.7 \times 15 + 4 = 14.5$$

$$\Rightarrow 1 \leq |A| \leq 8, 7 \leq |B| \leq 14$$

net m $\{C_1, C_2, C_5\}$ — $i=0$

net g $\{C_2, C_3, C_4\}$ ---

net k $\{C_3, C_5\}$ ~~~

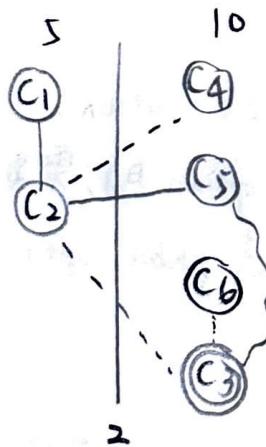
net p $\{C_3, C_6\}$

$$\text{gain}(C_1) = 0, \text{gain}(C_4) = 1$$

$$\text{gain}(C_2) = 0, \text{gain}(C_5) = 2$$

$$\text{gain}(C_3) = 2, \text{gain}(C_6) = 1$$

$i=1$ (move C_3)



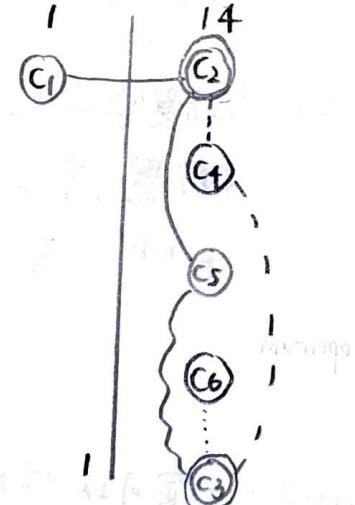
$$\text{gain}(C_2)' = 1$$

$$\text{gain}(C_4)' = 0$$

$$\text{gain}(C_5)' = 0$$

$$\text{gain}(C_6)' = -1$$

$i=2$ (move C_2)



$$\text{gain}(C_1)' = 1$$

$$\text{gain}(C_4)' = -1$$

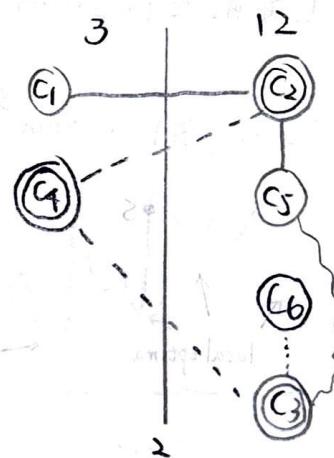
$$\text{gain}(C_5)' = -1$$

$\Rightarrow \text{gain max } C_1$
但舊不 balance $\Rightarrow C_4$

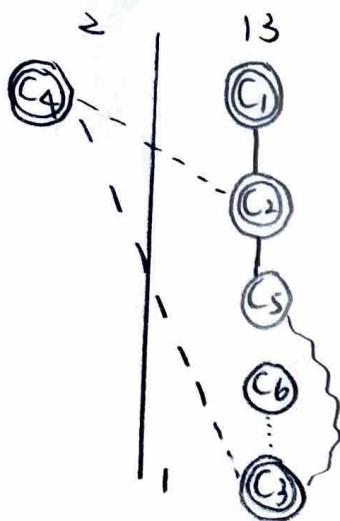
i	cell	$g(i)$	$\sum g(i)$	cutsize
0	-	-	-	4
1	C_3	2	2	2
2	C_2	1	3	1 ✓
3	C_4	-1	2	2
4	C_1	1	3	1
5	C_6	-1	2	2
6	C_5	-2	0	4

\Rightarrow 移動 C_3, C_2 *

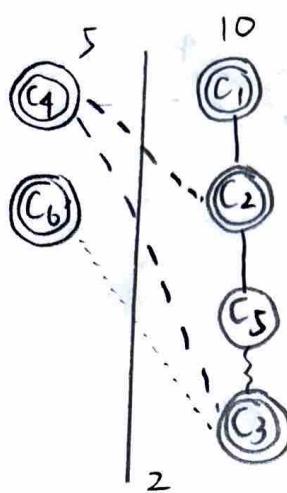
$i=3$ (move C_4)



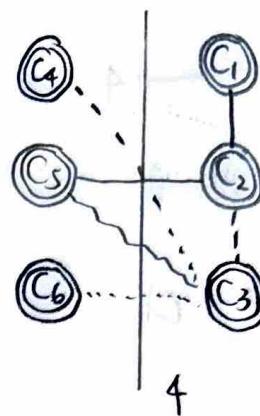
$i=4$ (move C_1)



$i=5$ (move C_6)



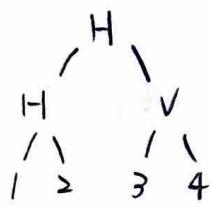
$i=6$ (move C_5)



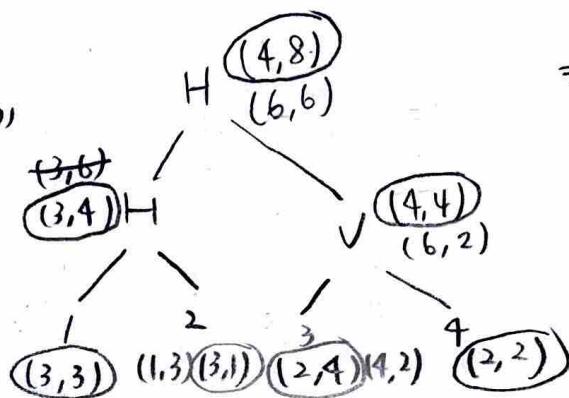
$$\text{gain}(C_5)' = -2$$

b.

~~(a)~~



(b)



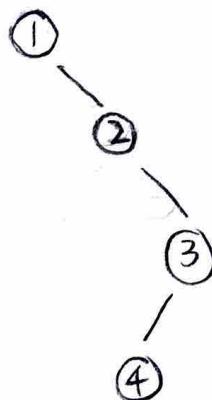
$\Rightarrow 4 \times 8 \text{ rectangle}$

$\Rightarrow \text{area} = 32$

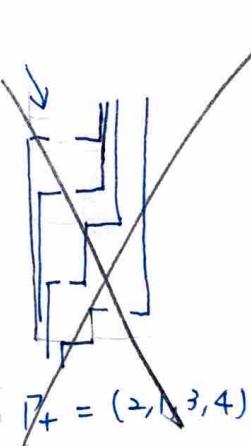
(c)



B^* tree

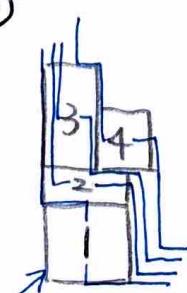
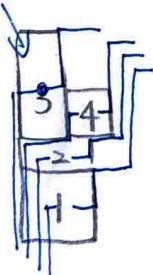


(d)



$$\Gamma_+ = (2, 1, 3, 4)$$

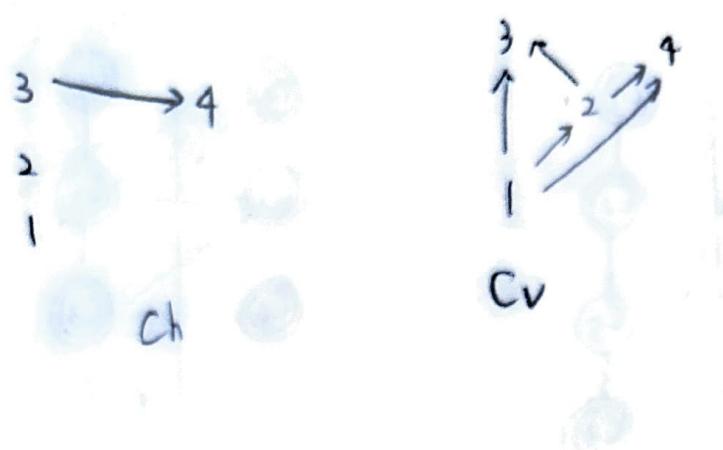
(d)



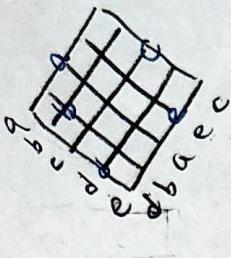
$$\Gamma_+ = \{3, 4, 2, 1\}$$

$$\Gamma_- = \{1, 2, 3, 4\}$$

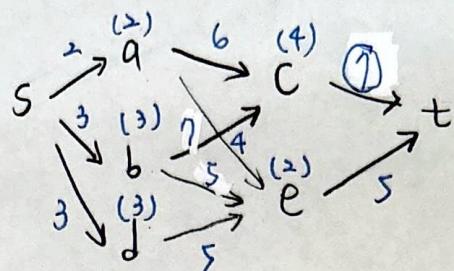
(e)



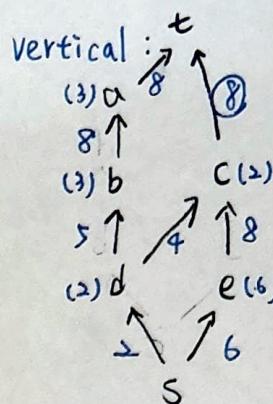
7. (a)



horizontal:



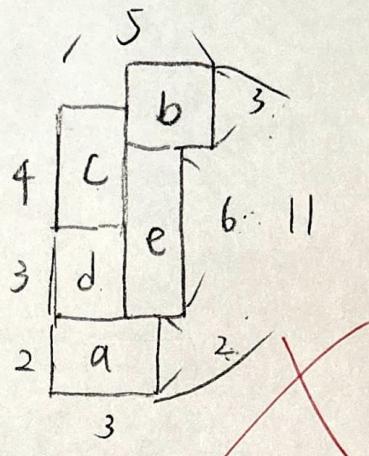
width: longest path = 7



height: longest path = 8

\Rightarrow minimum area: $7 \times 8 = 56$ *

(b)



$11 \times 5 = 55$ *