

# HW3

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## 1 HW 3: Turing Machines

### 1.1 Problem 1

A *Turing machine with left-reset* is similar to an ordinary Turing machine but the transition function has the form

$$\delta : Q \times \Gamma \rightarrow Q \times \Gamma \times \{R, RESET\}$$

The semantics of the *RESET* move is that the Turing machine will go all the back to the left-most side of the tape, instead of just moving a step left. Prove that Turing machines with left reset are equivalent to ordinary Turing machines. To do this you'll need to

- show that you can turn a Turing machine with left reset into an ordinary Turing machine

(the hard part)

- show that you can turn an ordinary Turing machine into a Turing machine with left reset

(the easy part)

The easiest way to do this is to simulate one type of movement with the other.

## 1.2 Problem 2

Show, by construction, that the Turing-decidable languages are closed under

1. union
2. intersection
3. complement

## 1.3 Problem 3

Give an informal description of the following languages

- $\{w \mid w \text{ contains twice as many 0s as 1s}\}$
- $\{w \mid w \text{ has matching parentheses and a length that's a power of 2}\}$

## 1.4 Problem 4

Show, by construction, that the Turing *recognizable* languages are closed under

- Kleene star
- concatenation
- reversal

Are the *recognizable* languages closed under complement? If not, can you give a counter-example?