

# What Are Computers, Really?

Clarissa Littler

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# Outline

Computers

Computation

A Little Set Theory: Or, What Computation Isn't!

What Computation is!

Constructive mathematics

Wrap-up

# What Are Computers, Really?

- ▶ Physical devices that carry out computation

# What Are Computers, Really?

- ▶ Physical devices that carry out computation
- ▶ We're done!

# Questions?

Thank you all you've been a wonderful audience

# What Is Computation, Really?

- ▶ Defined computers via computation
- ▶ What's computation?

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- ▶ What's computation?
- ▶ It's what computers do, OBVI!

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- ▶ Circular definitions are bad
- ▶ Define computation independent of physical machines
- ▶ Computation  $\subseteq$  Mathematics



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- ▶ Computation  $\subset$  Mathematics

# Why Do We Care?

- ▶ Computers have limits
- ▶ Limits are independent of hardware
- ▶ Chained by the laws of physics and mathematics

# Limits of Computation

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- ▶ No program to perfectly check for infinite loops
- ▶ No program to calculate proofs of all true theorems
- ▶ Wait, what?

# Computation: Examples

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- ▶ Database queries
- ▶ Counting
- ▶ Serving web pages
- ▶ What do these have in common?

# Computation: An Informal Definition

- ▶ Computable functions:
  - ▶ finite time
  - ▶ finite rules
  - ▶ finite data

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- ▶ Computable functions:
  - ▶ finite time
  - ▶ finite rules
  - ▶ finite data
- ▶ What about operating systems and servers?

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- ▶ Cover exact definition later
- ▶ Are all mathematically definable functions computable?
- ▶ Answer: **Not even remotely!**

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  - ▶ Set of all talks you'd rather be at *right now*

# Size of Sets

- ▶ Finite sets are sets you can count up to a finite size
- ▶ The number you reach when counting elements is the *cardinality*
- ▶ We write cardinality of a set  $A$  as  $|A|$
- ▶ The set of counting numbers itself a set:  $\mathbb{N}$
- ▶ Can you count  $\mathbb{N}$  to a finite number?

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- ▶ No! A logical contradiction otherwise.



# Countably infinite

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- ▶ The number of finite strings definable over a finite alphabet is countable
  - ▶ Gödel numberings

# Uncountably infinite

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- ▶ This is larger than you can possibly wrap your head around

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- ▶ The set of computable functions is the size of the set of programs
- ▶ The set of mathematically definable functions is larger than  $\aleph_0$
- ▶ The set of computable functions is vanishingly small compared to all functions

# What Are The Implications?

- ▶ There's a reason programming is hard
- ▶ Can't even imagine a world in which all functions are computable
  - ▶ Database queries
  - ▶ Search results
  - ▶ Cryptography
- ▶ Physics and speed of information?

# Turing's Insight

- ▶ Computer was a title not an object
- ▶ Turing watched how computers work
- ▶ Think elementary school arithmetic worksheets:
  - ▶ Problems *are* scratch paper
  - ▶ Clearly defined order in which digits are added
  - ▶ Can stop and come back to problem

# Turing Machines

- ▶ A universal model of computation!
- ▶ Simplifies scratch paper: 1d not 2d
- ▶ Always uses a finite amount of scratch paper
- ▶ Has only a finite number of rules
  - ▶ can look at input and remember a finite number of things

# The Use of Turing Machines

- ▶ Universal Turing Machine
- ▶ Textual descriptions of Turing Machines are "code"
- ▶ Ties back into cardinality argument
- ▶ Discuss difficulty of problems: decidable vs. recognizable

# Decideable

- ▶ Decideable means the problem can always be solved in finite time
- ▶ Decideable problems:
  - ▶ RegExp matching
  - ▶ Arithmetic
  - ▶ Sorting
  - ▶ Checking if a proof is valid
  - ▶ (Most) type-checking



# Recognizable

- ▶ Recognizable means well-formed inputs can be solved in finite time
  - ▶ i.e. the input has a proper solution and is written in the right format
- ▶ Recognizable problems:
  - ▶ Mostly meta-properties or meta-programming: Rice's theorem
  - ▶ Termination on a given input
  - ▶ Testing a program's behavior
  - ▶ Some type-checking (i.e. Scala)
  - ▶ C++ templates

# Why Are Some Problems Not Decidable?

- ▶ Logical paradox
- ▶ Halting problem: does a program halt on a given input?
- ▶ Variant on Russels's paradox
  - ▶ Set of all sets that don't contain themselves
  - ▶ Program that terminates only if a program doesn't terminate on itself
- ▶ Once you prove one thing isn't decidable, use it to prove others

# Untyped Lambda Calculus

- ▶ A different model of computation
- ▶ Invented by Alonzo Church as foundation of mathematics
- ▶ Works as a universal model of computation!
- ▶ Theory of function abstraction and application
- ▶ `function (x) {body}`
- ▶ Variables, anonymous functions, function calls

# Church-Turing Thesis

- ▶ Lambda calculus equivalent to Turing Machines
- ▶ Church-Turing thesis says there is no stronger model of computation
- ▶ Evidence: every universal model we know can simulate each other
- ▶ Still just a hypothesis! Not proven, but very likely

# Turing Completeness and You

- ▶ Turing complete means a PL is equivalent to Turing machines
- ▶ Most languages are
- ▶ This is why you can accidentally make infinite loops

# Typed Lambda Calculus

- ▶ Lambda calculus with *types*
- ▶ Types are stronger than you're used to
- ▶ Well-typed programs can't loop
- ▶ **Not** Turing complete

# Curry-Howard

- ▶ Curry-Howard correspondence
- ▶ Types are theorems
- ▶ A program of type  $A$  is a proof that the theorem  $A$  is true
- ▶  $A \rightarrow A$  means  $A$  implies  $A$
- ▶ In Agda  
 $(A : \text{Set}) \rightarrow (a\ b\ c : A) \rightarrow a == b \rightarrow b == c \rightarrow a == c$   
is a type that expresses transitivity of equality

# Take-aways

- ▶ Programming is a **subset** of mathematics
- ▶ Understanding the limits of computation is useful
- ▶ Types can be more than bug catchers, they're theorems we can prove