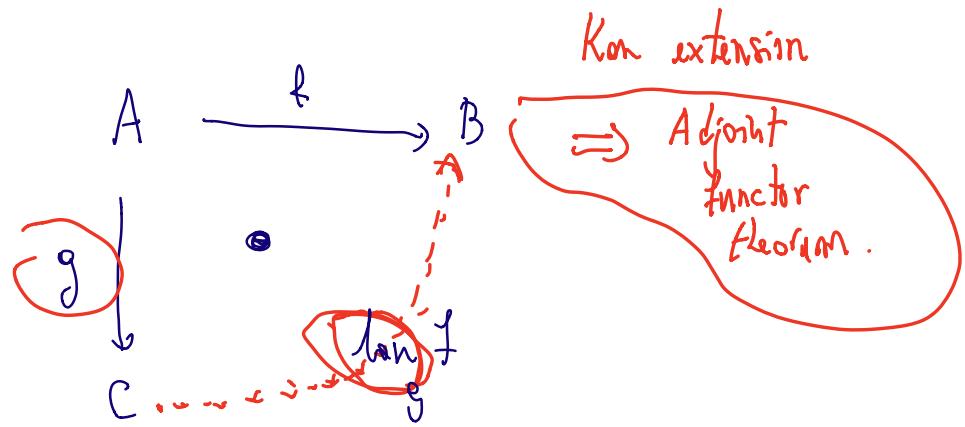


Kan extensions

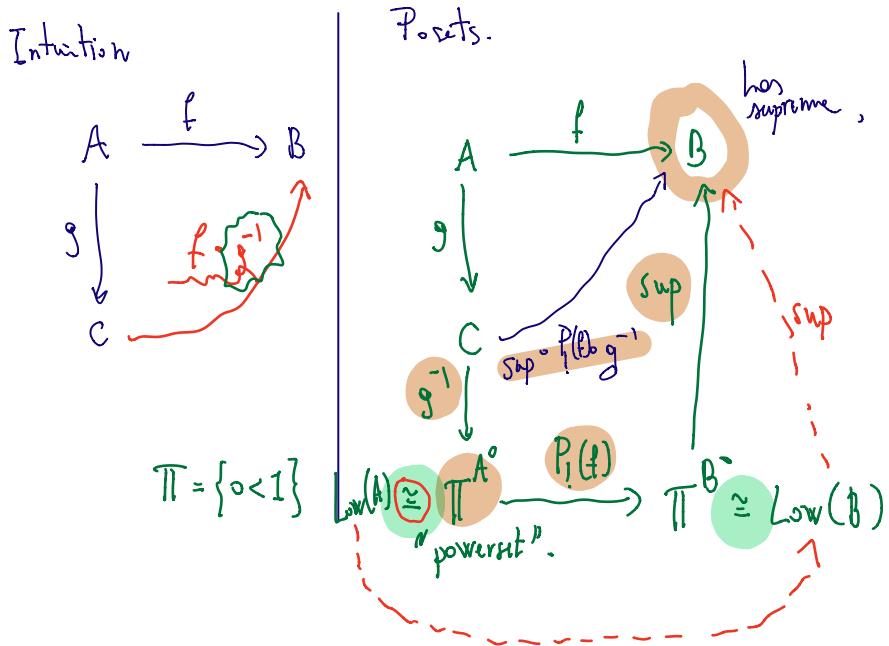


"Concrete"
(Additional assumptions)
(Concrete formula)

"Axiomatic"
(is more general)
(there is no concrete formula to compute the extension)

Historically were faithful

a very
conceptual
understanding
of the concrete
construction.



$$P_1(f)(X) = P_1(f)\left(\bigcup_{x \in X} \{x\}\right)$$

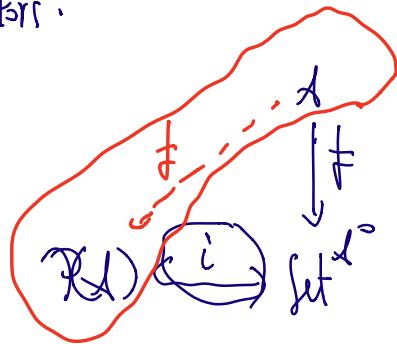
$$\bigcup_{x \in X} f(x).$$

List of stuff to do in CT.

- ① An analog of Π^{A^o} ($P(A)$) -
- ② g^{-1} (an analog of) -
- ③ $P_1(f)$
- ④ \sup (?)

(1) $\mathcal{P}A$) some notion of powerset.
(the category of small presheaves).

Def Given a locally small category A ,
 $\mathcal{P}A$) (small presheaves) is the
full subcategory of $\text{Set}^{A^{\text{op}}}$ containing
small objects of representable
functors.



Reu If A is small i is
on equivalence!

Is this a good notion of powerset?

"the Grothendieck construction"

$$\text{Loc}(P) \cong T^{P^\circ}$$

Given a small presheaf $X: \mathcal{C}^\circ \rightarrow \text{Set}$
 I want to construct a functor

$$\pi_X: \text{Elts}(X) \longrightarrow \mathcal{C}$$

Def Given a presheaf $X: \mathcal{C}^\circ \rightarrow \text{Set}$
 we define its category of elements

$\text{Elts}(X)$ $\rightarrow \text{ob} = (c, a_c)$ where c is an object of \mathcal{C} and a_c is an element of $a_c \in X(c)$.

$$X: \mathcal{C}^\circ \rightarrow \text{Set}$$

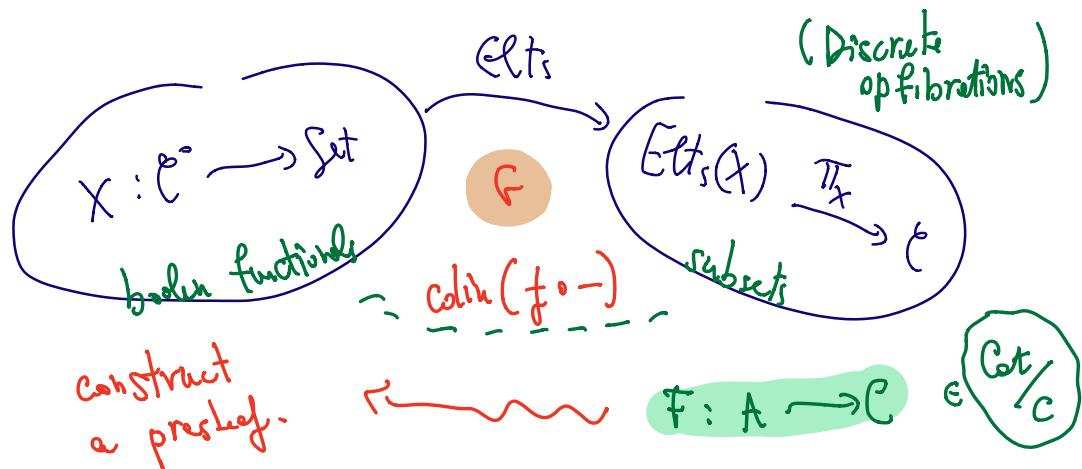
$$c \mapsto X(c)$$

mor = $(c, a_c) \rightarrow (d, a_d)$
 is a morphism in \mathcal{C}
 between $c \xrightarrow{f} d$ such that

$$X(f)(a_d) = a_c.$$

Rem $\pi_X: \text{Elt}_s(X) \longrightarrow \mathcal{C}$

$$\begin{array}{ccc} (\mathcal{C}, a_c) & \longmapsto & \mathcal{C} \\ f & \longmapsto & f_- \end{array}$$



$$A \xrightarrow{i} B \xrightarrow{\supset} \supset^o = \bigcup_{\text{sets}} \text{Set}^{(1)}.$$

$$A \xrightarrow{F} \mathcal{C} \xrightarrow{f} \text{Set}^{C^o}$$

$$F \longmapsto \text{colim}(f \circ F)$$

thus $\text{colim}(f \circ \pi_X) \cong X_-$

Back to the plan

• $P(A)$ ✓

• $T^P \cong \text{Low}(P) \checkmark$ (Groth construction).

- g^{-1} ✓
- $P_1(f)$ -
- \sup -

② g^{-1} $A \xrightarrow{g} B$
 ↓
 A° $\xleftarrow{g^{-1}}$

Def Given a functor $B(g-, -)$: $B \rightarrow \text{Set}^A$.
 $b \mapsto B(g-, b) : A^\circ \rightarrow \text{Set}$
 $a \mapsto B(g^a, b)$.

Def Safety check : facets

$$P \xrightarrow{g} Q.$$
$$\downarrow c \quad \quad \quad Q(g-, -) = "g^{-1}".$$
$$\Pi \xleftarrow{p^o}$$

$$Q(g-, q) : P^- \longrightarrow \Pi$$
$$p \mapsto \begin{cases} 1 & gp \leq q \\ 0 & \text{otherwise} \end{cases}$$
$$"g^{-1}(q)" = \{ p \in P : gp \leq q \}$$

Def A function is admissible $A \xrightarrow{f} B$

if for all b in B

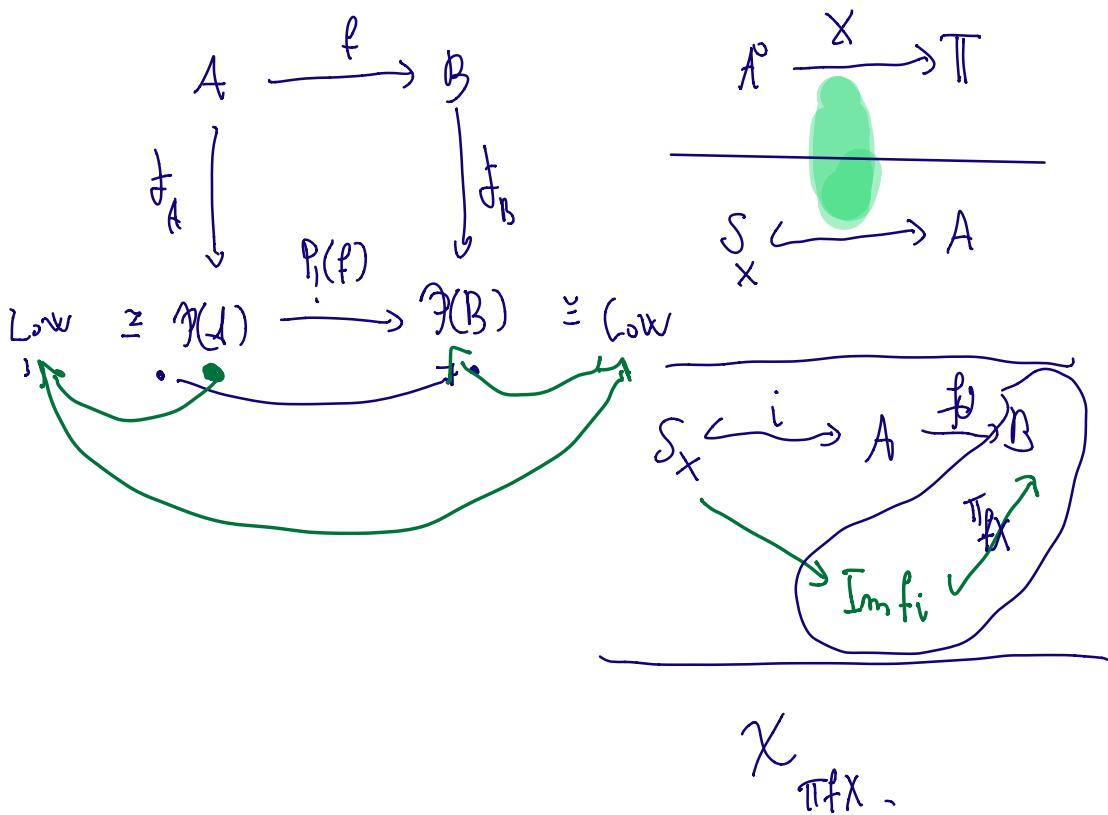
$B(g-, b)$ is a small presheaf

$$A \xrightarrow{g} B$$
$$\downarrow \quad \quad \quad B(g-, -)$$
$$P(A) \xrightarrow{p^o} \text{Set}^{B^o}$$

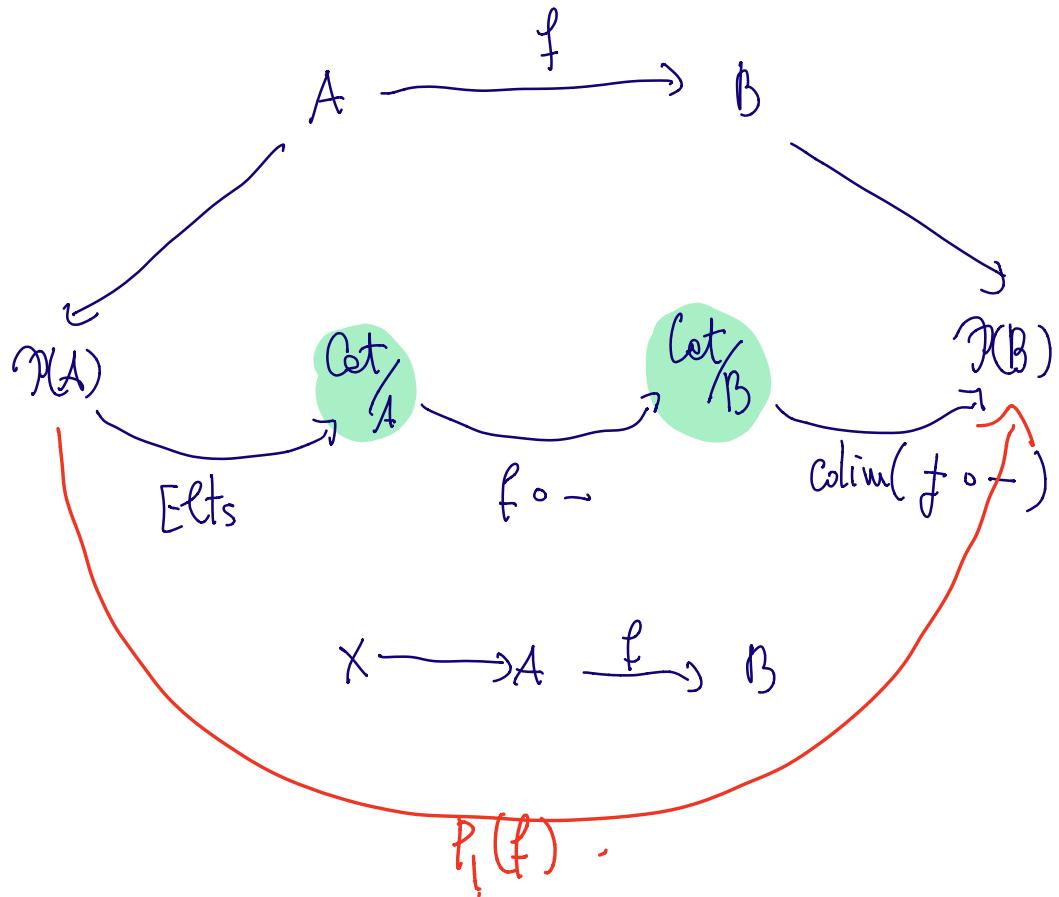
Reu functor w/ domain a small category are admissible

Reu functors with arity are admissible.

Cor Accessible functors are admissible.



We do the same!



$\boxed{P_1(f)(X) = P_1(f)\left(\bigcup_{x \in X} \{x\}\right) = \bigcup_{x \in X} f\{x\}}$

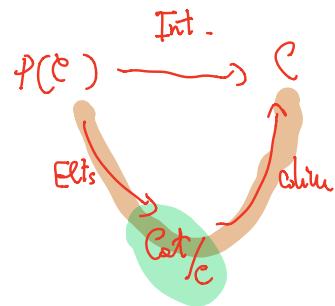
$$\begin{array}{ccc}
 A & \xrightarrow{f} & B \\
 g \downarrow & & \downarrow \text{sup} \\
 C & \xrightarrow{g'} & \mathcal{P}(B) \\
 & \xrightarrow{p_1(f)} & \mathcal{P}(B)
 \end{array}$$

thus the following are equivalent

1. B is a complete category
2. $B \xrightarrow{f} \mathcal{P}(B)$ has a left adjoint.

$$\text{Int}: \mathcal{P}(B) \longrightarrow B,$$

$1 \Rightarrow 2$



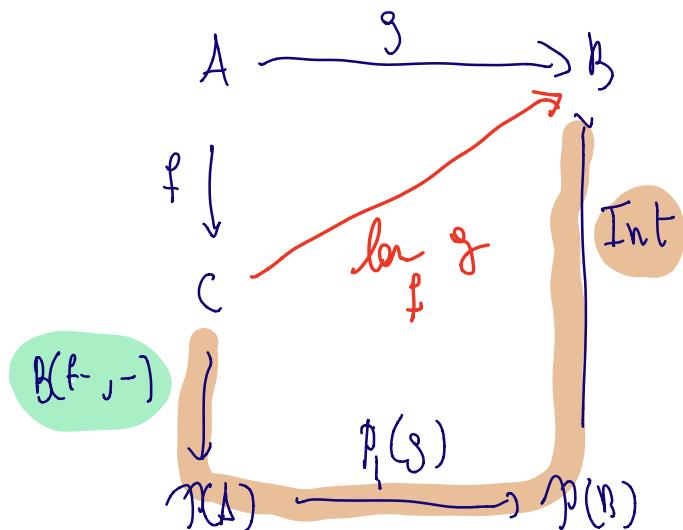
Lemma If $X: \mathcal{C}^{\text{op}} \longrightarrow \text{Set}$ is a small presheaf then there exist a small category D and functors

$$\begin{array}{ccc}
 & D & \\
 j \swarrow & \downarrow & \searrow i \\
 \text{Elts}(X) & \xrightarrow{\pi_X} & C
 \end{array}$$

such that when π_X exist
 iff ch(i) exists
 and in that case they coincide

Very good!

then if f is an admissible functor
 and B is complete, then
 we can Kan-extend!



Right left Kan extension
Left Right Kan extension.

Rm More general presentation
of Kan extensions -

Ree $A \xrightarrow{g} C$ \vdash category B

$$\begin{array}{ccc} [A, B] & \xleftarrow{g^*} & [C, B] \\ f \circ g & \longleftarrow & f \end{array}$$

thu " lon_g " : $[A, B] \xrightarrow{g} [C, B]$

$$f \xrightarrow{\quad} \text{lon}_g f$$

is left adjoint to $f \xrightarrow{g^*}$.

Ex $\text{lon}_g (h) \underset{fg}{\approx} \text{lon}_f (\text{lon}_g (h))$ -

Prop if g is fully faithful

$$(\text{lon}_g f) \circ g \underset{g}{\approx} f -$$

AFT (Adjoint functor theorem).

Lemma

Let $f: A \rightarrow B$ be a continuous functor between complete categories. Then TFAE.

$$\begin{array}{ccc} A & \xrightarrow{\text{id}} & A \\ f \downarrow & \nearrow \text{Ran } 1_f & \\ B & & \end{array}$$

① f has a right adj. g

② $\text{Ran } 1_f$ exists

And in that case they

coincide

(Adjoints are)
lex inverses

then
(AFT)

Let $f: A \rightarrow B$ be a functor between complete categories. TFAE.

- f is admissible & continuous
- f has a right adjoint.

$1 \Rightarrow 2$ Left because Inv^{-1}
exists - f

$2 \Rightarrow 1$ it is cont because it
is a left adj.

$$\begin{array}{c} A \text{ b in } B \quad B(f-, b) \text{ is null?} \\ \swarrow \text{adjunction} \\ B(f-, b) \cong A(-, g b) - \\ || \\ \underline{f(gb)} - \end{array}$$

$$B(f-, -) - B(f, 1) - \square .$$

$$B(f-, -) : A^{\text{op}} \times B \longrightarrow \text{Set} .$$

$$B(f=, -)$$