Graph Optimization Lab session 4

Giuliana Carello

DEIB, Politecnico di Milano

2023/2024

Exercise 1: column generation for multicommodity flow problem

Consider the multicommodity flow problem. Write a .mod and a .run files to solve the multicommodity flow problem with the column generation applied to the path-based formulation. Solve the two instances available online. Use the parameters as defined in the file ex4-parameters.mod and the scheme for the column generation in the file column-generation-scheme.run.

Exercise 2: arc based formulation for multicommodity flow problem

Consider the multicommodity flow problem. Write a .mod and a .run files to solve the multicommodity flow problem with the arc-based formulation. Solve the two instances available online. Use the parameters name as defined in the file ex4-parameters.mod. (Exercise 1 Lab session 2)

(Restricted) Master problem

$$\begin{aligned} & \min & & \sum_{p \in P} c_p x_p \\ & \text{s. t.} & & \sum_{p \in P_i} x_p = d_k, \forall k \in K \\ & & \sum_{p \in P_{ij}} x_p \leq u_{ij}, \forall (i,j) \in A \\ & & x_p \geq 0, \forall p \in P \end{aligned}$$

(Restricted) Master problem

Restricted) Master problem

min
$$\sum_{p \in P} c_p x_p$$

s. t. $\sum_{p \in P_k} x_p = d_k, \forall k \in K$

$$\sum_{p \in P_{ij}} x_p \leq u_{ij}, \forall (i,j) \in A$$

$$x_p \geq 0, \forall p \in P$$

Pricing problem

$$\min \sum_{(i,j) \in A} g_{ij} z_{ij}$$

$$\sum_{(i,j) \in A} z_{ji} - \sum_{(i,j) \in A} z_{ij} = b_i, \forall i \in N$$

$$z_{ij} \in \{0,1\} \quad \forall (i,j) \in A$$

$$\min \sum_{(i,j)\in A} g_{ij}z_{ij}$$

$$\sum_{(j,i)\in A} z_{ji} - \sum_{(i,j)\in A} z_{ij} = b_i, \forall i \in \Lambda$$

$$z_{ij} \in \{0,1\} \quad \forall (i,j) \in A$$

where $g_{ij} = c_{ij} - \mu_{ij}$ and

$$b_i = \left\{ egin{array}{l} -1, i = s_k \ 1, i = t_k \ 0, ext{otherwise} \end{array}
ight.$$

Column generation scheme

- 1. Define an initial set of paths P
- 2. Solve the restricted master problem on P
- **3.** For each commodity *k*:
 - ▶ Build the pricing problem $(g_{ij} = c_{ij} \mu_{ij})$
 - Solve the pricing problem
 - ightharpoonup Compare the optimal solution of the pricing with σ_k
 - If the optimal solution of the pricing is lesser than σ_k then a new path must be added
- 4. If at least one path has been added, go to 1

```
Defining and solving more than one problem problem ProblemName:variables, objective function, constraints; to define a problem with the list of variables, objective fuction and constraints
```

problem ProblemName; option solver gurobi; solve ProblemName; to solve the problem

```
Defining the set of paths set N; param n_p; set P := 1..n_p; set Path {P} within N cross N; param or {P} within N; #Origin param dest {P} within N; # Destination param cp {P}; # Cost
```

```
Updating the set of paths let n_p := n_p + 1; let Path[n_p] := {(i,j) in A: z[i,j] = 1}; let cp[n_p] := sum {(i,j) in Path[n_p]} c[i,j];
```

- ConstraintName.dual; to access dual variable associated with ConstraintName
- ▶ param ParameterName symbolic in N;
- ▶ let := : assign a value to a parameter that can be changed
- ▶ for : for statement
- ▶ if : if statement