

Scripting and Modeling with Picat

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Why Picat?

- Many complaints about Prolog
 - □ Implicit unification and non-determinism are difficult
 - Cuts and dynamic predicates are non-logical
 - Lack of constructs for programming everyday things
- No satisfactory successors
 - □ Prolog extensions are ad hoc (e.g., loops in B-Prolog)
 - Mercury requires too many declarations
 - Erlang abandons non-determinism in favor of concurrency
 - □ Oz's syntax is strange and implicit laziness is difficult
 - Curry is too close to Haskell



Features of PICAT

- Pattern-matching
 - □ Predicates and functions are defined with pattern-matching rules
- Imperative
 - □ Assignments, loops, list comprehensions
- Constraints
 - □ CP, SAT and LP/MIP
- Actors
 - □ Action rules, event-driven programming, actor-based concurrency
- Tabling
 - Memoization, dynamic programming, planning, model-checking



Outline of Talk

- A brief overview of Picat
- Scripting with Picat
- Modeling with Picat
 - □ CSP modeling
 - Dynamic programming
 - Planning
- Conclusion



Data Types

Variables – plain and attributed

- Primitive values
 - □ Integer and float
 - □ Atom

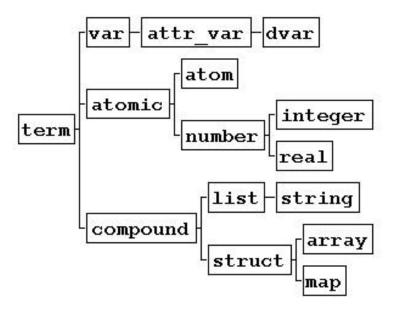
```
x1 '_' \ ab' \$%' \你好'
```

- Compound values
 - □ List

- String "abc" is the same as [a,b,c]
- □ Structure



The Type Hierarchy



Creating Structures and Lists

Structure

```
Picat> P = new_struct(point, 3)
P = point(_3b0, _3b4, _3b8)
Picat> S = $student(marry,cs,3.8)
```

List Comprehension

```
Picat> \hat{L} = [(A,I) : A \text{ in } [a,b], I \text{ in } 1..2]

L = [(a,1),(a,2),(b,1),(b,2)]
```

Range

```
Picat> L = 1..2..10 L = [1,3,5,7,9]
```

String

```
Picat> write("hello "++"world")
[h,e,l,l,o,' ',w,o,r,l,d]
```

Array

```
Picat> A = new_array(2,3)
A = \{\{3d0, 3d4, 3d8\}, \{3e0, 3e4, 3e8\}\}
```

Map

```
Picat> M = new_map([alpha= 1, beta=2])
M = (map)[alpha = 1,beta = 2]
```

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Index Notation

X[l1,...,ln]: X references a compound value

```
Picat> L = [a,b,c,d], X = L[2]
X = b

Picat> S = $student(marry,cs,3.8), GPA=S[3]
GPA = 3.8

Picat> A = {{1, 2, 3}, {4, 5, 6}}, B = A[2, 3]
B = 6
```



List Comprehension

[T: E_1 in D_1 , Cond_n, . . . , E_n in D_n , Cond_n]

```
Picat> L = [X : X in 1..5].
L = [1,2,3,4,5]

Picat> L = [(A,I): A in [a,b], I in 1..2].
L = [(a,1),(a,2),(b,1),(b,2)]

Picat> L = [X : I in 1..5] % X is local
L = [_bee8,_bef0,_bef8,_bf00,_bf08]

Picat> X=X, L = [X : I in 1..5] % X is non-local
L = [X,X,X,X,X]
```

OOP Notation

```
O.f(t1,...,tn)
Picat> Y = 13.to binary string()
                                        -- means module qualified call if O is atom
Y = ['1', '1', '0', '1']
                                        -- means f(O,t1,...,tn) otherwise.
Picat> Y = 13.to_binary_string().reverse()
Y = ['1', '0', '1', '1']
% X becomes an attributed variable
Picat> X.put(age, 35), X.put(weight, 205), A = X.age
A = 35
%X is a map
Picat> X = new_map([age=35, weight=205]), X.put(gender, male)
X = (map)([age=35, weight=205, gender=male])
Picat> S = point(1.0, 2.0), Name = S.name, Arity = S.length
Name = point
Arity = 2
Picat> I = math.pi % module qualifier
I = 3.14159
```

Predicates

Backtracking (explicit non-determinism)

```
member(X,[Y|_]) ?=> X=Y.
member(X,[_|L]) => member(X,L).

Picat> member(X,[1,2,3])
X = 1;
X = 2;
X = 3;
no
```

Control backtracking

```
Picat> once(member(X,[1,2,3]))
```

Predicate Facts

```
\begin{array}{lll} \operatorname{index}(+,-) & (-,+) & \operatorname{edge}(a,Y) ?=> Y=b. \\ \operatorname{edge}(a,b). & \operatorname{edge}(a,Y) => Y=c. \\ \operatorname{edge}(b,Y) => Y=c. \\ \operatorname{edge}(b,Y) => Y=b. \\ \operatorname{edge}(x,Y) => X=a. \\ \operatorname{edge}(x,Y) => X=a. \\ \operatorname{edge}(x,Y) => X=b. \\ \operatorname{edge}(x,Y) => X=b. \\ \operatorname{edge}(x,Y) => X=b. \\ \operatorname{edge}(x,Y) => X=c. \\ \end{array}
```

- Facts must be ground
- A call with insufficiently instantiated arguments fails
 - Picat> edge(X,Y)
 no

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Functions

- Always succeed with a return value
- Non-backtrackable

```
fib(0)=F => F=1.
fib(1)=F => F=1.
fib(N)=F,N>1 => F=fib(N-1)+fib(N-2).
```

Function facts

```
fib(0)=1.
fib(1)=1.
fib(N)=fib(N-1)+fib(N-2).
```



Assignments

- X[I1,...,In] := Exp
 Destructively update the component to Exp.
 Undo the update upon backtracking.
- Var := Exp
 The compiler changes it to Var' = Exp and replace all subsequent occurrences of Var in the body of the rule by Var'.

```
test => X = 0, X := X + 1, X := X + 2, write(X).
test => X = 0, X1 = X + 1, X2 = X1 + 2, write(X2).
```



Loops

- Types
 - □ foreach(E1 in D1, ..., En in Dn) Goal end
 - □ while (Cond) Goal end
 - □ do Goal while (Cond)
- Loops provide another way to write recurrences
- A loop forms a name scope: variables that do not occur before in the outer scope are local.
- Loops are compiled into tail-recursive predicates



Loops (Example) sum_even(L)

Using a loop

```
sum_even(L)=Sum =>
    S=0,
    foreach (X in L)
        if even(X) then S:=S+X end
    end,
    Sum=S.
```

Using a list comprehension

```
sum_even(L) = sum([X : X in L, even(X)]).
```



Tabling

- Predicates define relations where a set of facts is implicitly generated by the rules
- The process of fact generation might never end, and can contain a lot of redundancy
- Tabling memorizes calls and their answers in order to prevent infinite loops and to limit redundancy

Tabling (examples)

```
table
fib(0)=1.
fib(1)=1.
fib(N)=fib(N-1)+fib(N-2).

table(+,+,-,min)
shortest_path(X,Y,Path,W) ?=>
    Path = [(X,Y)],
    edge(X,Y,W),
shortest_path(X,Y,Path,W) =>
    Path = [(X,Z)|PathR],
    edge(X,Z,W1),
    shortest_path(Z,Y,PathR,W2),
    W = W1+W2.
```



Modules

```
module M. import M1,M2,...,Mn.
```

- The declared module name and the file name must be the same
- Files that do not begin with a module declaration are in the global module
- Atoms and structure names are global
- Picat has a global symbol table for atoms, a global symbol table for structure names, and a global symbol table for modules
- Each module has its own symbol table for the public predicates and functions



Supported Modules

- Pre-loaded and pre-imported
 - □ basic, math, io, sys
- Pre-loaded but names are not pre-imported
 - □ cp, planner, sat, os, util
- Not pre-loaded or pre-imported
 - □ Setting of PICATPATH is needed



Status of the Implementation

- Based on the B-Prolog engine
 - □ Over 20+ years of R/D
- System size
 - □ 55,000 LOC in C
 - ☐ 45,000 LOC in Picat
- Other system features
 - Debugger
 - Garbage collector
 - □ Big integers



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```
import os.

traverse(Dir), directory(Dir) =>
   List = listdir(Dir),
   printf("Inside %s%n",Dir),
   foreach(File in List)
        printf(" %s%n",File)
   end,
   foreach (File in List, File != ".", File != "..")
        FullName = Dir ++ [separator()] ++ File,
        traverse(FullName)
   end.

traverse(_Dir) => true.
```

Input Rows of Integers into an Array

```
import util.

input_data(Tri) =>
    Lines = read_file_lines("triangle.txt"),
    Tri = new_array(Lines.length),

I = 1,
    foreach(Line in Lines)
        Tri[I] = Line.split().map(to_integer).to_array(),
        I := I+1
    end.
```

foreach({I,Line} in zip(1..Lines.length,Lines))
 Tri[I] = Line.split().map(to_integer).to_array(),
end.

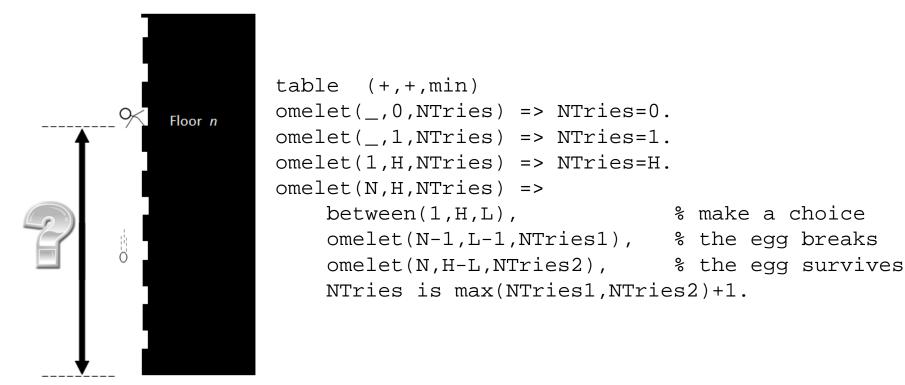


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The Omelet Problem (Or The N-Eggs Problem)



http://www.datagenetics.com/blog/july22012/

100

Maximum Path Sum (Euler Project 18 and 67)

```
table (+,+,max,nt)
path(Row,Col,Sum,Tri),Row==Tri.length =>
    Sum=Tri[Row,Col].
path(Row,Col,Sum,Tri) ?=>
    path(Row+1,Col,Sum1,Tri),
    Sum = Sum1+Tri[Row,Col].
path(Row,Col,Sum,Tri) =>
    path(Row+1,Col+1,Sum1,Tri),
    Sum = Sum1+Tri[Row,Col].
```



The planner Module

- Useful for solving planning problems
 - □ plan(State,Limit,Plan,PlanCost)
 - □ best_plan(State,Limit,Plan,PlanCost)
- Users only need to define final/1 and action/4
 - final(State) is true if State is a final state
 - □ action(State,NextState,Action,ActionCost) encodes the state transition diagram
- Uses tabling with the early termination and resourcebounded search techniques to speedup search

Ex: The Farmer's Problem

```
import planner.

go =>
    S0=[s,s,s,s],
    best_plan(S0,Plan),
    writeln(Plan).

final([n,n,n,n]) => true.

action([F,F,G,C],S1,Action,ActionCost) ?=>
    Action=farmer_wolf,
    ActionCost = 1,
    opposite(F,F1),
    S1=[F1,F1,G,C],
    not unsafe(S1).
```



Tabling is More Effective Than SAT for Planning?

- □ Nomystery (picat-lang.org/asp/nomystery.pi)
 - Picat solves all of the 30 instances
 - Clasp solves only 17 of the 30 instances
 - On solved instances, Picat is more than 100 times faster than Clasp
- □ Sokoban(picat-lang.org/asp/sokoban.pi)
 - Picat solves all of the 30 instances
 - Clasp solves only 14 of the 30 instances
- □ Ricochet Robots(picat-lang.org/asp/ricochet.pi)
 - Both Picat and Clasp solve all of the 30 instances
 - Picat is several times faster than Clasp despite that its encoding is much simpler



Constraints

- Picat can be used for constraint satisfaction and optimization problems
- Constraint Problems
 - □ Generate variables
 - Generate constraints over the variables
 - Solve the problem, finding an assignment of values to the variables that matches all the constraints
- Picat can be used as a modeling language for CP, SAT, LP/MIP
 - Loops are helpful for modeling

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SEND + MORE = MONEY

import cp.

b/A

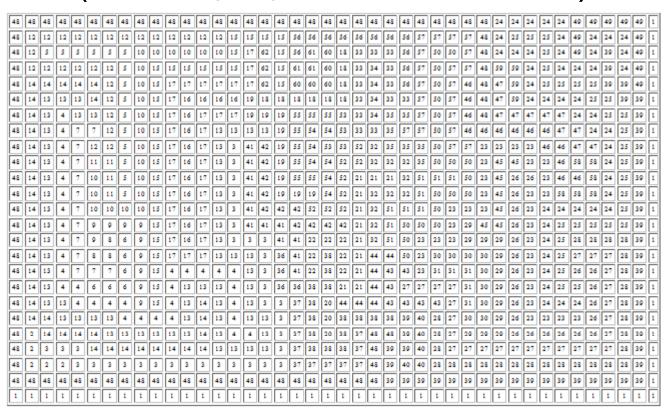
N-Queens Problem

```
import cp.

queens3(N, Q) =>
    Q = new_list(N),
    Q in 1..N,
    all_different(Q),
    all_different([$Q[I]-I : I in 1..N]),
    all_different([$Q[I]+I : I in 1..N]),
    solve([ff],Q).
```

The Number Link Problem

(picat-lang.org/asp/numberlink_b.pi)



Solved with the sat module of Picat and the Lingeling solver



Conclusion

- Picat is a logic-based multi-paradigm language
- Picat can be used as a scripting language
 - □ Future work: add modules for programming Web services
- Picat can also be used as a modeling language for DP, planning, and CSP
 - □ Future work: add modules for other solvers



Picat Vs. Haskell

<u>Commonalities</u>	<u>Differences</u>
pattern-matching	untyped vs. typed
strings, list comprehensions	strict vs. lazy
tail-recursion optimization	assignments vs. monads
Picat supports FP	multi-paradigm vs. pure FP

- Picat is more suitable to symbolic computations
 - Explicit unification
 - □ Explicit non-determinism
 - □ Tabling
 - Constraints



Picat Vs. Prolog

- Picat is more expressive
 - □ Functions, arrays, maps, loops, and list comprehensions
- Picat is more scalable because pattern-matching facilitates indexing rules
- Picat is arguably more reliable than Prolog
 - Explicit unification and non-determinism
 - Functions don't fail (at least built-in functions)
 - No cuts or dynamic predicates
 - No operator overloading
 - □ A simple static module system



Resources

- Users' Guide
 - □ http://picat-lang.org/download/picat_guide.pdf
- Hakan Kjellerstrand's Picat Page
 - □ http://www.hakank.org/picat/
- Examples
 - □ http://picat-lang.org/download/exs.pi
- Modules
 - □ http://picat-lang.org/modules.html
- Projects
 - □ http://picat-lang.org/projects.html