Problem 1

Algorithm 1: Divide and Conquer Finding Min and Max

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\begin{array}{l} \mbox{DCMinMax(arr[], 1, r, min, max)} \{ \\ & \mbox{if (l == r)} \{ \\ & \mbox{min = A[l], max = A[l]} \} \\ & \mbox{else if (r-l = 1)} \{ \\ & \mbox{if (A[l] == A[r])} \{ \\ & \mbox{min = Al[l], max = A[r]} \} \\ & \mbox{else} \{ \\ & \mbox{min = A[r], max = A[l]} \} \\ & \mbox{else} \{ \\ & \mbox{DCMinMax(A[0,(l+r)/2],min,max) //Give half the array DCMinMax(A[(l+r)/2+l,r],min1,min2) if (min2 < min) \\ & \mbox{min = min2} \\ & \mbox{if (max2 > max) } \\ & \mbox{max = max2} \\ \} \end{array}
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Problem 2

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Recurrence for Pan's algorithm: M(n) = 143,640 M(n/70)
Using Master Theorem: \theta(n^p) where p = log_{70}143640 = 2.795
Recurrence for Strassens algorithm s = log_27 = 2.807
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... The recurrence for Pan's algorithm is slightly lower than Strassen's

Problem 3

See sorts.cpp

Problem 4

Prove
$$\sum_{h=1}^{i=0} 2(h-i)2^i = 2(n-\log_2(n+1))$$
 where $n = 2^{k+1} - 1$
Substitute $n \ 2(2^{h+1} - 1 - \log_2(2^{h+1})) = 2(2^{h+1} - h - 2)$
 $\sum_{i=0}^{h-1} 2(h-2)2^i$
 $= \sum_{i=0}^{h-1} 2(h-2)2^i$
 $= 2(\sum_{h=1}^{i=0} h2^i - \sum_{h=1}^{i=0} h2^i - \sum_{h=1}^{i=0} i2^i$
 $= 2(h(2^h - 1) - (h-2)2^h - 2)$
 $= 2(2^{h+1} - h - 2)$
 $\therefore \sum_{h=1}^{i=0} 2(h-i)2^i = 2(n-\log_2(n+1))$