1. Q 4.11(a,c,e) textbook (a) Minimize $||Ax - b||_{\infty}$

$$\min \max_{i} |(Ax - b)_{i}|$$
 let $t = \max_{i} |(Ax - b)_{i}|$
$$\min_{x,t} t, subject \ to :$$

$$|(Ax - b)_{i}| \le t, \forall i$$

LP Formulation:

$$x \in R^{n}, t \in R$$

$$\min_{x,t} t, subject \ to.:$$

$$A_{i,:}x - t \leq b_{i}, \forall i$$

$$-A_{i,:}x - t \leq -b_{i}, \forall i$$

(c) Minimize $||Ax - b||_1, ||x||_{\infty} \le 1$

$$\begin{aligned} \min_{x} \sum_{i} |(Ax - b)_{i}|, subject\ to: \\ |x_{i}| &\leq 1, \forall i \\ let\ t_{i} &= |(Ax - b)_{i}| \\ \min_{x,t} \sum_{i} t_{i}, subject\ to: \\ |(Ax - b)_{i}| &\leq t_{i}, \forall i \\ |x_{i}| &\leq 1, \forall i \end{aligned}$$

LP Formulation:

$$x \in R^{n}, t \in R^{n}$$

$$\min_{x,t} 1^{T}t, subject \ to:$$

$$A_{i,:}x - t_{i} \leq b_{i}$$

$$-A_{i,:}x - t_{i} \leq -b_{i}$$

$$x_{i} \leq 1, \forall i$$

$$-x_{i} \leq 1, \forall i$$

(e) Minimize $||Ax - b||_1 + ||x||_{\infty}$

$$\begin{aligned} let \ t_i &= |Ax - b|_i \\ let \ v &= max|x_i| \\ \min_{t,v,x} \sum_i t_i + v, subject \ to: \\ |x_i| &\leq v, \forall i \\ |Ax - b|_i &\leq t_i, \forall i \end{aligned}$$

LP formulation:

$$x, t \in \mathbb{R}^n, v \in \mathbb{R}$$

$$\min_{t,v,x} \mathbf{1}^T t + v, subject \ to:$$

$$x_i - v \le 0, \forall i$$

$$-x_i - v \le 0, \forall i$$

$$A_{i,:}x - t_i \le +b_i, \forall i$$

$$-A_{i,:}x - t_i \le -b_i, \forall i$$

2. Q 4.16 textbook

$$x(t) \in \mathbb{R}^{n}, t \in \{0, ..., N\}$$

$$u(t) \in \mathbb{R}, t \in \{0, ..., N\}$$

$$x(t+1) = Ax(t) + bu(t), t \in \{0, ..., N\}$$

$$given :$$

$$A \in \mathbb{R}^{n \times n}$$

$$b \in \mathbb{R}^{n}$$

$$x(0) = 0$$

$$problem :$$

$$X(N) = x_{des}$$

$$f(a) = \begin{cases} |a|, & |a| leq 1 \\ 2|a| - 1, & |a| > 1 \end{cases}$$

expanding x(t):

$$\begin{split} x(0) &= 0 \\ x(1) &= Ax(0) + bu(0) = bu(0) \\ x(2) &= Ax(1) + bu(1) = A(bu(0)) + bu(1) \\ x(i) &= A^{i-1}bu(0) + A^{i-2}bu(1) + \dots + A^0bu(i-1) = \sum_{j=0}^{i-1} A^jbu(i-1-j) \\ x(i) &= \left[A^{i-1}b \quad A^{i-2}b \quad \dots \ b\right] \begin{bmatrix} u(0) \\ \vdots \\ u(i-1) \end{bmatrix} \\ x_{des} &= x(N) = \left[A^{N-1}b \quad A^{N-2}b \quad \dots \ b\right] \begin{bmatrix} u(0) \\ \vdots \\ u(N-1) \end{bmatrix} \\ let \ v_t &= |f(x(t))| \\ v_t &= max\{|u(t)|, 2|u(t)|-1\} \\ v_t &\geq |u(t)| \\ v_t &\geq 2|u(t)|-1 \\ v_t &\geq u_t \\ v_t &\geq -u_t \\ v_t &\geq 2u_t - 1 \\ v_t &\geq -2u_t - 1 \end{split}$$

LP formulation:

$$\begin{aligned} v_{0,..,N-1}, u_{0,..,N-1} &\in \mathbb{R}^{N} \\ \min_{v,u} \mathbf{1}^{T} v, subject \ to: \\ &-v_{t} + u_{t} \leq 0, \forall t \in 0, ..., N-1 \\ &-v_{t} - u_{t} \leq 0, \forall t \in 0, ..., N-1 \\ &-v_{t} + 2u_{t} \leq 1, \forall t \in 0, ..., N-1 \\ &-v_{t} - 2u_{t} \leq 1, \forall t \in 0, ..., N-1 \\ &-v_{t} - 2u_{t} \leq 1, \forall t \in 0, ..., N-1 \end{aligned}$$

$$\begin{bmatrix} A^{N-1}b \quad A^{N-2}b \quad .. \quad ..b \end{bmatrix} \begin{bmatrix} u_{0} \\ .. \\ u_{N-1} \end{bmatrix} = x_{des}$$

3. Q 4.21(a) textbook

Find explicit solution for the QCQP: minimize $c^T x$, subject to:

$$x^T A x \le 1$$
$$A \in S_{++}^n, c \ne 0$$

$$\begin{split} x^TAx &= x^TA^{1/2}A^{1/2}x = (A^{1/2}x)^TA^{1/2}x = \|A^{1/2}x\|_2^2 \\ \|A^{1/2}x\|_2^2 &\leq 1 \\ let \ y &= A^{1/2}x \\ x &= A^{-1/2}y \\ c^Tx &= c^TA^{-1/2}y \\ let \ b^T &= c^TA^{-1/2} \\ \min b^Ty \ subject \ to: \\ \|y\|_2^2 &\leq 1 \\ \min b^Ty &= \frac{-b^Tb\alpha}{\|b\|}, \alpha = \max\|y\|_2 = 1 \\ y^* &= \frac{-b}{\|b\|} \\ x^* &= A^{-1/2}y^* = A^{-1/2}\frac{-b}{\|b\|} \\ &= A^{-1/2}\frac{-(c^TA^{-1/2})^T}{\|(c^TA^{-1/2})^T\|} \\ &= \frac{-A^{-1/2}A^{-1/2}c}{\|A^{-1/2}c\|} = \frac{-A^{-1}c}{(c^TA^{-1}c)^{1/2}} \\ c^Tx^* &= \frac{-c^TA^{-1}c}{(c^TA^{-1}c)^{1/2}} \\ c^Tx^* &= -(c^TA^{-1}c)^{1/2} \end{split}$$

4. Q 4.25 textbook

$$\varepsilon_i = \{P_i u + q_i : ||u||_2 \le 1\}, i = 1, ..., K + L, P_i \in S^n$$

Find a feasible hyperplane strictly separating $\varepsilon_1,..,\varepsilon_K$ from $\varepsilon_{K+1},..,\varepsilon_{K+L}$.

$$a^{T}x + b > 0, x \in \bigcup_{i=1}^{K} \varepsilon_{i}$$
$$a^{T}x + b < 0, x \in \bigcup_{i=K+1}^{K+L} \varepsilon_{i}$$

let $\epsilon > 0$, a constant for strict separation relax inequalities to:

$$a^{T}x + b \leq -\epsilon, x \in \bigcup_{i=1}^{K} \varepsilon_{i}$$

$$a^{T}x + b \geq \epsilon, x \in \bigcup_{i=K+1}^{K+L} \varepsilon_{i}$$

$$a^{T}(P_{i}u + q_{i}) + b \leq -\epsilon, \|u\|_{2} \leq 1, i \in \{1, ..., K\} \text{ becomes}$$

$$\sup_{\|u\|_{2} \leq 1} a^{T}P_{i}u + a^{T}q_{i} + b \leq -\epsilon, i \in \{1, ..., K\}$$

$$\sup_{\|u\|_{2} \leq 1} a^{T}P_{i}u = \frac{a^{T}P_{i}(a^{T}P_{i})^{T}}{\|a^{T}P_{i}\|_{2}} = \|a^{T}P_{i}\|_{2}$$

$$\|a^{T}P_{i}\|_{2} + a^{T}q_{i} + b \leq -\epsilon, i \in \{1, ..., K\}$$

$$\|a^{T}P_{i}\|_{2} \leq -a^{T}q_{i} - b - \epsilon, i \in \{1, ..., K\}$$

$$a^{T}(P_{i}u + q_{i}) + b \ge \epsilon, \|u\|_{2} \le 1, i = \{K + 1, ..., K + L\} \text{ becomes}$$

$$\inf_{\|u\|_{2} \le 1} a^{T}P_{i}u + a^{T}q_{i} + b \ge \epsilon, i = \{K + 1, ..., K + L\}$$

$$\inf_{\|u\|_{2} \le 1} a^{T}P_{i}u = \frac{a^{T}P_{i}(-a^{T}P_{i})^{T}}{\|a^{T}P_{i}\|_{2}} = -\|a^{T}P_{i}\|_{2}$$

$$-\|a^{T}P_{i}\|_{2} + a^{T}q_{i} + b \ge \epsilon, i = \{K + 1, ..., K + L\}$$

$$\|a^{T}P_{i}\|_{2} \le a^{T}q_{i} + b - \epsilon, i = \{K + 1, ..., K + L\}$$

Second Order Cone Programming formulation:

$$\begin{aligned} & \min_{a,b} & 0 \\ & \|a^T P_i\|_2 \leq -q_i^T a - b - \epsilon, i \in \{1,..,K\} \\ & \|a^T P_i\|_2 \leq q_i^T a + b - \epsilon, i \in \{K+1,..,K+L\} \\ & \text{where } \epsilon > 0 \end{aligned}$$

5. Q 4.30 textbook Express as Geometric Programming:

$$T \in [T_{min}, T_{max}], T_{min}, T_{max} > 0$$

$$r \in [r_{min}, r_{max}], r_{min}, r_{max} > 0$$

$$w \in [w_{min}, w_{max}], w_{min}, w_{max} > 0$$

$$w \leq 0.1r$$

$$T > 0, r > 0, w > 0, c_{max} > 0$$

$$\max \alpha_4 Tr^2, subject to:$$

$$\alpha_1 \frac{Tr}{w} + \alpha_2 r + \alpha_3 rw \leq c_{max}$$

$$\alpha_i > 0, \forall i$$

$$\max \alpha_4 Tr^2 = \min \frac{1}{\alpha_4 Tr^2} = \min \frac{1}{\alpha_4} T^{-1} r^{-2}$$

$$\frac{\alpha_1}{c_{max}} Tr w^{-1} + \frac{\alpha_2}{c_{max}} r + \frac{\alpha_3}{c_{max}} rw \leq 1$$

$$T \leq T_{max} \iff \frac{T}{T_{max}} \leq 1$$

$$T \geq T_{min} \iff \frac{r_{min}}{T} \leq 1$$

$$r \leq r_{max} \iff \frac{r}{r_{max}} \leq 1$$

$$v \leq w_{max} \iff \frac{w}{w_{max}} \leq 1$$

$$w \leq w_{max} \iff \frac{w}{w_{min}} \leq 1$$

$$w \leq w_{min} \iff \frac{w_{min}}{w} \leq 1$$

$$w \leq w_{min} \iff \frac{10w}{r} \leq 1$$

$$w \leq 0.1r \iff \frac{10w}{r} \leq 1$$

GP formulation:

$$\begin{split} & \min_{T,r,w} \ \frac{1}{\alpha_4} T^{-1} r^{-2}, subject \ to: \\ & \frac{T}{T_{max}} \leq 1, \frac{T_{min}}{T} \leq 1 \\ & \frac{r}{r_{max}} \leq 1, \frac{r_{min}}{r} \leq 1 \\ & \frac{w}{w_{max}} \leq 1, \frac{w_{min}}{w} \leq 1 \\ & \frac{10w}{r} \leq 1 \\ & \frac{\alpha_1}{c_{max}} Trw^{-1} + \frac{\alpha_2}{c_{max}} r + \frac{\alpha_3}{c_{max}} rw \leq 1 \end{split}$$

- 6. Q 4.43(a-b) textbook $A: R^n \to S^m, A(x) = A_0 + x_1 A_1 + ... + x_n A_n$ Let $\lambda_1(x) \geq ... \lambda_m(x)$ denote the eigenvalues of A(x). Formulate problems as SDP.
 - (a) Minimize the maximum eigenvalue SDP formulation:

$$\min_{a,x} a, \ subject \ to:$$

$$A(x) \leq_{S^m_+} aI$$

• (b) Minimize the spread of the eigenvalues

$$let \ t = \lambda_1(x) - \lambda_m(x)$$

$$\min_{t,x,a,b} t, \ subject \ to :$$

$$A(x) \leq_{S^m_+} bI$$

$$A(x) \geq_{S^m_+} aI$$

$$t - b + a = 0$$

7. Formulate $\min_{x \in \mathbb{R}^n, z \in \mathbb{R}} \sum_{m=1}^{M} \max(a_m^T x, z) + \tau ||x||_2^2$ as LP/QP/SOCP/SDP.

$$let \ t_m = max(a_m^T x, z)$$

$$\min_{x \in \mathbb{R}^n, z \in \mathbb{R}, t \in \mathbb{R}^m} 1_{m \times 1}^T t_m + \tau ||x||_2^2, \ subject \ to:$$

$$a_m^T x \le t_m, \forall m$$

$$z \le t_m, \forall m$$

Quadratic Programming formulation:

$$X \in \mathbb{R}^{m+n+1} = \begin{bmatrix} t_{1:m}, \ x_{1:n}, \ z \end{bmatrix}$$

$$\min_{X} X^{T} \begin{bmatrix} 0_{m \times (m+n+1)} & 0_{n \times 1} \\ 0_{n \times m} & \tau I_{n \times n} & 0_{n \times 1} \end{bmatrix} X + \begin{bmatrix} 1_{1 \times m} \ 0_{1 \times (n+1)} \end{bmatrix} X, \text{ subject to :}$$

$$\begin{bmatrix} -I_{m \times m} & A & 0_{m \times 1} \\ -I_{m \times m} & 0_{m \times 5} & 1_{m \times 1} \end{bmatrix} X \leq 0_{2m \times 1}, A = \begin{bmatrix} a_{1}^{T} \\ \vdots \\ a_{m}^{T} \end{bmatrix}$$

8. Portfolio Design

$$E[x^T p] = \bar{x}p$$

$$E[(x^T p - \bar{x}^T p)(x^T p - \bar{x}^T p)^T] = E[(p^T (x - \bar{x})(p^T (x - \bar{x})^T)]$$

$$= p^T E[(x - \bar{x})(x - \bar{x})^T]p$$

$$= p^T \Sigma p$$

Quadratic Programming formulation:

$$\min_{p} p^{T} \sum p, \ subject \ to:$$

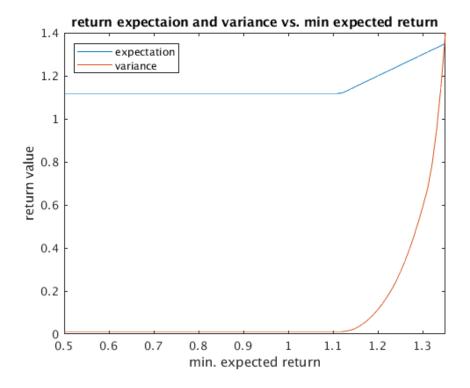
$$1^{T} p = 1$$

$$-x^{T} p \leq -r$$

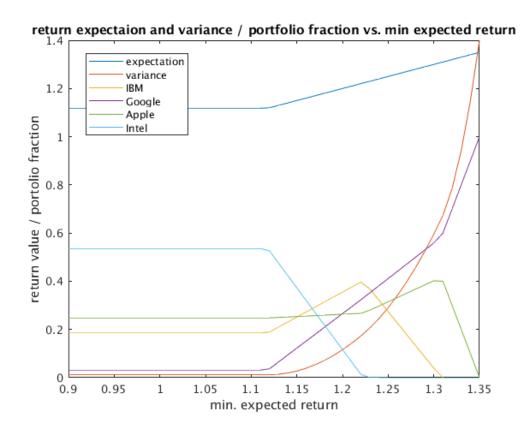
$$Diag(1_{4\times 1}) p \leq 1$$

$$-Diag(1_{4\times 1}) p \leq 0$$

Solving for a range of r yields:



Portfolio for a range of r:



Solver:

```
[rs, ret] = solve();
ps = cell2mat(ret)';
x = [1.1, 1.35, 1.25, 1.05];
H = [0.2 -0.2 -0.12 0.02;...]
      -0.2 1.4 0.02 0.0;...
      -0.12 0.02 1 -0.4;...
       0.02\ 0\ -0.4\ 0.2;
ms = ps * x;
vs = diag(ps * H * ps');
plot(rs,ms);
hold on;
plot(rs,vs);
title('return expectaion and variance vs. min expected return');
legend('expectation', 'variance');
xlabel('min. expected return');
ylabel('return value');
hold off;
%% additional plot for part b
plot(rs,ms);
hold on;
plot(rs,vs);
title('return expectaion and variance / portfolio fraction vs. min expected return'
xlabel('min. expected return');
ylabel('return value / portolio fraction');
plot(rs,ps(:,1));
plot(rs,ps(:,2));
plot(rs,ps(:,3));
plot(rs,ps(:,4));
legend('expectation','variance', 'IBM', 'Google', 'Apple', 'Intel');
hold off;
```

- 9. Optimal Control of a Unit Mass
 - a) given:

$$x(0) = 0$$

$$\dot{x}(0) = 0$$

$$f(t) = p_i, i - 1 < t \le i, i = 1, ..., 10$$

$$x(10) = 1$$

$$\dot{x}(10) = 0$$

$$minimize \sum_{i=1}^{10} p_i^2$$

$$let \ v = \dot{x}$$

$$v(i) = v(0) + \sum_{j=0}^{i} \frac{f(j)}{m} = v(0) + \sum_{j=1}^{i} p_{j}, m = 1$$

$$v(i) = \sum_{j=1}^{i} p_{j} = 1^{T} p, p \in \mathbb{R}^{i}$$

$$v(1) = p_{1}$$

$$v(2) = p_{1} + p_{2}$$
...
$$v(10) = 1^{T} p = v_{des} = 0,$$

$$x(i) = x(0) + \sum_{j=0}^{i} v(j)$$

$$x(i) = \sum_{j=1}^{i} v(j) = 1^{T} v, v \in \mathbb{R}^{i}$$

$$x(i) = \sum_{j=1}^{i} \sum_{k=1}^{j} p_{k}$$

$$x(i) = \sum_{j=1}^{i} (i - j + 1) p_{j}$$

$$x(1) = p_{1}$$

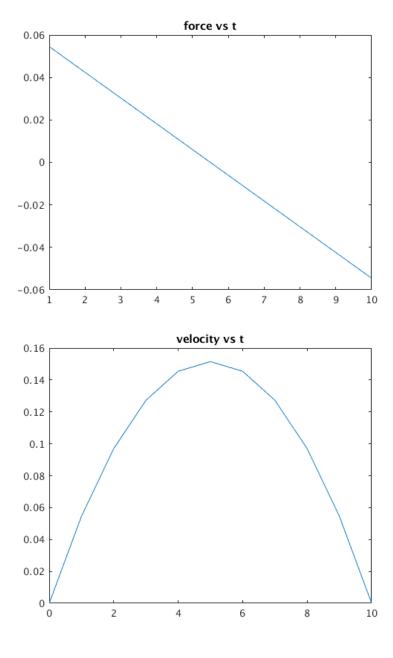
$$x(2) = 2p_{1} + 1 p_{2}$$

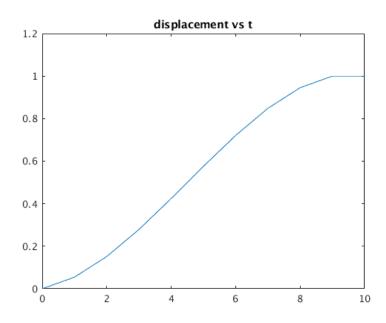
$$x(10) = 10 p_{1} + 9 p_{2} + .. + 1 p_{10}$$

QP formulation:

$$\min_{f} f^{T} I f, \ subject \ to:$$

$$\begin{bmatrix} & 1^{T} \\ 10 & 9 & \dots & 2 & 1 \end{bmatrix} f = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$$





Optimal strategy in this case is to apply a smooth symmetrical force around t = 5, so that velocity is always non-negative and displacement is always towards the destination.

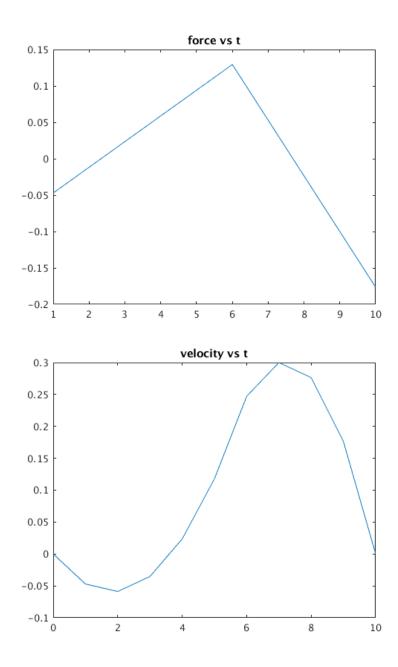
Solver:

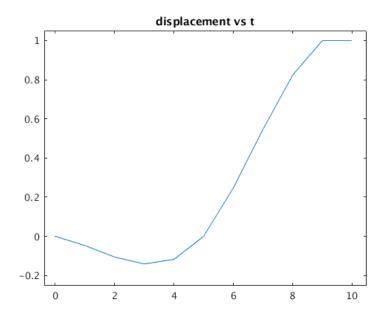
```
H = diag(ones(10,1))
f = []
A = []
b = []
Aeq = [ones(1,10); 10:-1:1]
beq = [0;1]
f = quadprog(H,f,A,b,Aeq,beq)
plot(1:10,f)
title('force vs t')
v = zeros(11,1)
for i=2:1:11
    v(i) = v(i-1) + f(i-1)
end
plot(0:1:10,v)
title('velocity vs t')
x = zeros(11,1)
for i=1:1:11
```

• b) additional constraint: x(5) = 0

$$\min_{f} f^{T} I f, \, subject \,\, to:$$

$$\begin{bmatrix} & & 1^{T} & & \\ 10 & 9 & \dots & & 2 & 1 \\ 5 & 4 & 3 & 2 & 1 & 0 & \dots & 0 \end{bmatrix} f = \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix}$$





Optimial strategy in this case is to reverse direction and then go forward again to gain enough velocity timed so that the displacement at t=5 is 0.

Solver:

```
H = diag(ones(10,1))
f = []
A = []
b = []
Aeq = [ones(1,10); 10:-1:1; 5 4 3 2 1 0 0 0 0 0]
beq = [0;1;0]
f = quadprog(H,f,A,b,Aeq,beq)
plot(1:10,f)
title('force vs t')
v = zeros(11,1)
for i=2:1:11
    v(i) = v(i-1) + f(i-1)
end
plot(0:1:10,v)
title('velocity vs t')
x = zeros(11,1)
for i=1:1:11
    if i > 1
        x(i) = x(i-1) + v(i)
    else
```

```
x(i) = v(i)
end
end
plot(0:1:10,x)
title('displacement vs t')
```