

1. Q 4.11(a,c,e) textbook
 (a) Minimize $\|Ax - b\|_\infty$

$$\begin{aligned} & \min \max_i |(Ax - b)_i| \\ \text{let } t = & \max_i |(Ax - b)_i| \\ & \min_{x,t} t, \text{ subject to :} \\ & |(Ax - b)_i| \leq t, \forall i \end{aligned}$$

LP Formulation:

$$\begin{aligned} & x \in R^n, t \in R \\ & \min_{x,t} t, \text{ subject to :} \\ & A_{i,:}x - t \leq b_i, \forall i \\ & -A_{i,:}x - t \leq -b_i, \forall i \end{aligned}$$

- (c) Minimize $\|Ax - b\|_1, \|x\|_\infty \leq 1$

$$\begin{aligned} & \min_x \sum_i |(Ax - b)_i|, \text{ subject to :} \\ & |x_i| \leq 1, \forall i \\ & \text{let } t_i = |(Ax - b)_i| \\ & \min_{x,t} \sum_i t_i, \text{ subject to :} \\ & |(Ax - b)_i| \leq t_i, \forall i \\ & |x_i| \leq 1, \forall i \end{aligned}$$

LP Formulation:

$$\begin{aligned} & x \in R^n, t \in R^n \\ & \min_{x,t} 1^T t, \text{ subject to :} \\ & A_{i,:}x - t_i \leq b_i \\ & -A_{i,:}x - t_i \leq -b_i \\ & x_i \leq 1, \forall i \\ & -x_i \leq 1, \forall i \end{aligned}$$

(e) Minimize $\|Ax - b\|_1 + \|x\|_\infty$

$$\text{let } t_i = |Ax - b|_i$$

$$\text{let } v = \max |x_i|$$

$$\min_{t,v,x} \sum_i t_i + v, \text{ subject to :}$$

$$|x_i| \leq v, \forall i$$

$$|Ax - b|_i \leq t_i, \forall i$$

LP formulation:

$$x, t \in R^n, v \in R$$

$$\min_{t,v,x} 1^T t + v, \text{ subject to :}$$

$$x_i - v \leq 0, \forall i$$

$$-x_i - v \leq 0, \forall i$$

$$A_{i,:}x - t_i \leq +b_i, \forall i$$

$$-A_{i,:}x - t_i \leq -b_i, \forall i$$

2. Q 4.16 textbook

$$\begin{aligned}x(t) &\in \mathbb{R}^n, t \in \{0, \dots, N\} \\u(t) &\in \mathbb{R}, t \in \{0, \dots, N\} \\x(t+1) &= Ax(t) + bu(t), t \in \{0, \dots, N\}\end{aligned}$$

given :

$$A \in \mathbb{R}^{n \times n}$$

$$b \in \mathbb{R}^n$$

$$x(0) = 0$$

problem :

$$\min_u \sum_{t=0}^{N-1} f(u(t)), \text{ subject to :}$$

$$X(N) = x_{des}$$

$$f(a) = \begin{cases} |a|, & |a| \leq 1 \\ 2|a| - 1, & |a| > 1 \end{cases}$$

expanding $x(t)$:

$$x(0) = 0$$

$$x(1) = Ax(0) + bu(0) = bu(0)$$

$$x(2) = Ax(1) + bu(1) = A(bu(0)) + bu(1)$$

$$x(i) = A^{i-1}bu(0) + A^{i-2}bu(1) + \dots + A^0bu(i-1) = \sum_{j=0}^{i-1} A^jbu(i-1-j)$$

$$x(i) = \begin{bmatrix} A^{i-1}b & A^{i-2}b & \dots & b \end{bmatrix} \begin{bmatrix} u(0) \\ \vdots \\ u(i-1) \end{bmatrix}$$

$$x_{des} = x(N) = \begin{bmatrix} A^{N-1}b & A^{N-2}b & \dots & b \end{bmatrix} \begin{bmatrix} u(0) \\ \vdots \\ u(N-1) \end{bmatrix}$$

$$\text{let } v_t = |f(x(t))|$$

$$v_t = \max\{|u(t)|, 2|u(t)| - 1\}$$

$$v_t \geq |u(t)|$$

$$v_t \geq 2|u(t)| - 1$$

$$v_t \geq u_t$$

$$v_t \geq -u_t$$

$$v_t \geq 2u_t - 1$$

$$v_t \geq -2u_t - 1$$

LP formulation:

$$\begin{aligned}
 & v_{0,..,N-1}, u_{0,..,N-1} \in \mathbb{R}^N \\
 & \min_{v,u} 1^T v, \text{ subject to :} \\
 & \quad -v_t + u_t \leq 0, \forall t \in 0, \dots, N-1 \\
 & \quad -v_t - u_t \leq 0, \forall t \in 0, \dots, N-1 \\
 & \quad -v_t + 2u_t \leq 1, \forall t \in 0, \dots, N-1 \\
 & \quad -v_t - 2u_t \leq 1, \forall t \in 0, \dots, N-1 \\
 & \quad [A^{N-1}b \quad A^{N-2}b \quad \dots \quad b] \begin{bmatrix} u_0 \\ \vdots \\ u_{N-1} \end{bmatrix} = x_{des}
 \end{aligned}$$

3. Q 4.21(a) textbook

Find explicit solution for the QCQP:

minimize $c^T x$, subject to:

$$x^T A x \leq 1$$

$$A \in S_{++}^n, c \neq 0$$

$$\begin{aligned}
 x^T A x &= x^T A^{1/2} A^{1/2} x = (A^{1/2} x)^T A^{1/2} x = \|A^{1/2} x\|_2^2 \\
 \|A^{1/2} x\|_2^2 &\leq 1 \\
 \text{let } y &= A^{1/2} x \\
 x &= A^{-1/2} y \\
 c^T x &= c^T A^{-1/2} y \\
 \text{let } b^T &= c^T A^{-1/2} \\
 \min b^T y &\text{ subject to :} \\
 \|y\|_2^2 &\leq 1 \\
 \min b^T y &= \frac{-b^T b}{\|b\|}, \alpha = \max \|y\|_2 = 1 \\
 y^* &= \frac{-b}{\|b\|} \\
 x^* &= A^{-1/2} y^* = A^{-1/2} \frac{-b}{\|b\|} \\
 &= A^{-1/2} \frac{-(c^T A^{-1/2})^T}{\|(c^T A^{-1/2})^T\|} \\
 &= \frac{-A^{-1/2} A^{-1/2} c}{\|A^{-1/2} c\|} = \frac{-A^{-1} c}{(c^T A^{-1} c)^{1/2}} \\
 c^T x^* &= \frac{-c^T A^{-1} c}{(c^T A^{-1} c)^{1/2}} \\
 c^T x^* &= -(c^T A^{-1} c)^{1/2}
 \end{aligned}$$

4. Q 4.25 textbook

$$\varepsilon_i = \{P_i u + q_i : \|u\|_2 \leq 1\}, i = 1, \dots, K + L, P_i \in S^n$$

Find a feasible hyperplane strictly separating $\varepsilon_1, \dots, \varepsilon_K$ from $\varepsilon_{K+1}, \dots, \varepsilon_{K+L}$.

$$a^T x + b > 0, x \in \bigcup_{i=1}^K \varepsilon_i$$

$$a^T x + b < 0, x \in \bigcup_{i=K+1}^{K+L} \varepsilon_i$$

let $\epsilon > 0$, a constant for strict separation

relax inequalities to:

$$a^T x + b \leq -\epsilon, x \in \bigcup_{i=1}^K \varepsilon_i$$

$$a^T x + b \geq \epsilon, x \in \bigcup_{i=K+1}^{K+L} \varepsilon_i$$

$$a^T(P_i u + q_i) + b \leq -\epsilon, \|u\|_2 \leq 1, i \in \{1, \dots, K\} \text{ becomes}$$

$$\sup_{\|u\|_2 \leq 1} a^T P_i u + a^T q_i + b \leq -\epsilon, i \in \{1, \dots, K\}$$

$$\sup_{\|u\|_2 \leq 1} a^T P_i u = \frac{a^T P_i (a^T P_i)^T}{\|a^T P_i\|_2} = \|a^T P_i\|_2$$

$$\|a^T P_i\|_2 + a^T q_i + b \leq -\epsilon, i \in \{1, \dots, K\}$$

$$\|a^T P_i\|_2 \leq -a^T q_i - b - \epsilon, i \in \{1, \dots, K\}$$

$$a^T(P_i u + q_i) + b \geq \epsilon, \|u\|_2 \leq 1, i = \{K + 1, \dots, K + L\} \text{ becomes}$$

$$\inf_{\|u\|_2 \leq 1} a^T P_i u + a^T q_i + b \geq \epsilon, i = \{K + 1, \dots, K + L\}$$

$$\inf_{\|u\|_2 \leq 1} a^T P_i u = \frac{a^T P_i (-a^T P_i)^T}{\|a^T P_i\|_2} = -\|a^T P_i\|_2$$

$$-\|a^T P_i\|_2 + a^T q_i + b \geq \epsilon, i = \{K + 1, \dots, K + L\}$$

$$\|a^T P_i\|_2 \leq a^T q_i + b - \epsilon, i = \{K + 1, \dots, K + L\}$$

Second Order Cone Programming formulation:

$$\min_{a, b} 0$$

$$\|a^T P_i\|_2 \leq -q_i^T a - b - \epsilon, i \in \{1, \dots, K\}$$

$$\|a^T P_i\|_2 \leq q_i^T a + b - \epsilon, i \in \{K + 1, \dots, K + L\}$$

where $\epsilon > 0$

5. Q 4.30 textbook

Express as Geometric Programming:

$$\begin{aligned}
& T \in [T_{min}, T_{max}], T_{min}, T_{max} > 0 \\
& r \in [r_{min}, r_{max}], r_{min}, r_{max} > 0 \\
& w \in [w_{min}, w_{max}], w_{min}, w_{max} > 0 \\
& w \leq 0.1r \\
& T > 0, r > 0, w > 0, c_{max} > 0 \\
& \max \alpha_4 T r^2, \text{ subject to :} \\
& \alpha_1 \frac{T r}{w} + \alpha_2 r + \alpha_3 r w \leq c_{max} \\
& \alpha_i > 0, \forall i \\
& \max \alpha_4 T r^2 = \min \frac{1}{\alpha_4 T r^2} = \min \frac{1}{\alpha_4} T^{-1} r^{-2} \\
& \frac{\alpha_1}{c_{max}} T r w^{-1} + \frac{\alpha_2}{c_{max}} r + \frac{\alpha_3}{c_{max}} r w \leq 1 \\
& T \leq T_{max} \iff \frac{T}{T_{max}} \leq 1 \\
& T \geq T_{min} \iff \frac{T_{min}}{T} \leq 1 \\
& r \leq r_{max} \iff \frac{r}{r_{max}} \leq 1 \\
& r \geq r_{min} \iff \frac{r_{min}}{r} \leq 1 \\
& w \leq w_{max} \iff \frac{w}{w_{max}} \leq 1 \\
& w \geq w_{min} \iff \frac{w_{min}}{w} \leq 1 \\
& w \leq 0.1r \iff \frac{10w}{r} \leq 1
\end{aligned}$$

GP formulation:

$$\begin{aligned}
& \min_{T, r, w} \frac{1}{\alpha_4} T^{-1} r^{-2}, \text{ subject to :} \\
& \frac{T}{T_{max}} \leq 1, \frac{T_{min}}{T} \leq 1 \\
& \frac{r}{r_{max}} \leq 1, \frac{r_{min}}{r} \leq 1 \\
& \frac{w}{w_{max}} \leq 1, \frac{w_{min}}{w} \leq 1 \\
& \frac{10w}{r} \leq 1 \\
& \frac{\alpha_1}{c_{max}} T r w^{-1} + \frac{\alpha_2}{c_{max}} r + \frac{\alpha_3}{c_{max}} r w \leq 1
\end{aligned}$$

6. Q 4.43(a-b) textbook

$$A : \mathbb{R}^n \rightarrow \mathbb{S}^m, A(x) = A_0 + x_1 A_1 + \dots + x_n A_n$$

Let $\lambda_1(x) \geq \dots \geq \lambda_m(x)$ denote the eigenvalues of $A(x)$.

Formulate problems as SDP.

- (a) Minimize the maximum eigenvalue
SDP formulation:

$$\min_{a,x} a, \text{ subject to :}$$

$$A(x) \leq_{\mathbb{S}_+^m} aI$$

- (b) Minimize the spread of the eigenvalues

$$\text{let } t = \lambda_1(x) - \lambda_m(x)$$

$$\min_{t,x,a,b} t, \text{ subject to :}$$

$$A(x) \leq_{\mathbb{S}_+^m} bI$$

$$A(x) \geq_{\mathbb{S}_+^m} aI$$

$$t - b + a = 0$$

7. Formulate $\min_{x \in \mathbb{R}^n, z \in \mathbb{R}} \sum_{m=1}^M \max(a_m^T x, z) + \tau \|x\|_2^2$ as LP/QP/SOCP/SDP.

$$\begin{aligned} & \text{let } t_m = \max(a_m^T x, z) \\ & \min_{x \in \mathbb{R}^n, z \in \mathbb{R}, t \in \mathbb{R}^m} 1_{m \times 1}^T t_m + \tau \|x\|_2^2, \text{ subject to :} \\ & a_m^T x \leq t_m, \forall m \\ & z \leq t_m, \forall m \end{aligned}$$

Quadratic Programming formulation:

$$\begin{aligned} X & \in \mathbb{R}^{m+n+1} = [t_{1:m}, x_{1:n}, z] \\ \min_X X^T & \begin{bmatrix} 0_{m \times (m+n+1)} \\ 0_{n \times m} & \tau I_{n \times n} & 0_{n \times 1} \\ 0_{m \times (m+n+1)} \end{bmatrix} X + [1_{1 \times m} \ 0_{1 \times (n+1)}] X, \text{ subject to :} \\ & \begin{bmatrix} -I_{m \times m} & A & 0_{m \times 1} \\ -I_{m \times m} & 0_{m \times 5} & 1_{m \times 1} \end{bmatrix} X \leq 0_{2m \times 1}, A = \begin{bmatrix} a_1^T \\ \vdots \\ a_m^T \end{bmatrix} \end{aligned}$$

8. Portfolio Design

$$\begin{aligned}
E[x^T p] &= \bar{x}^T p \\
E[(x^T p - \bar{x}^T p)(x^T p - \bar{x}^T p)^T] &= E[(p^T (x - \bar{x}))(p^T (x - \bar{x}))^T] \\
&= p^T E[(x - \bar{x})(x - \bar{x})^T] p \\
&= p^T \Sigma p
\end{aligned}$$

Quadratic Programming formulation:

$$\min_p p^T \Sigma p, \text{ subject to :}$$

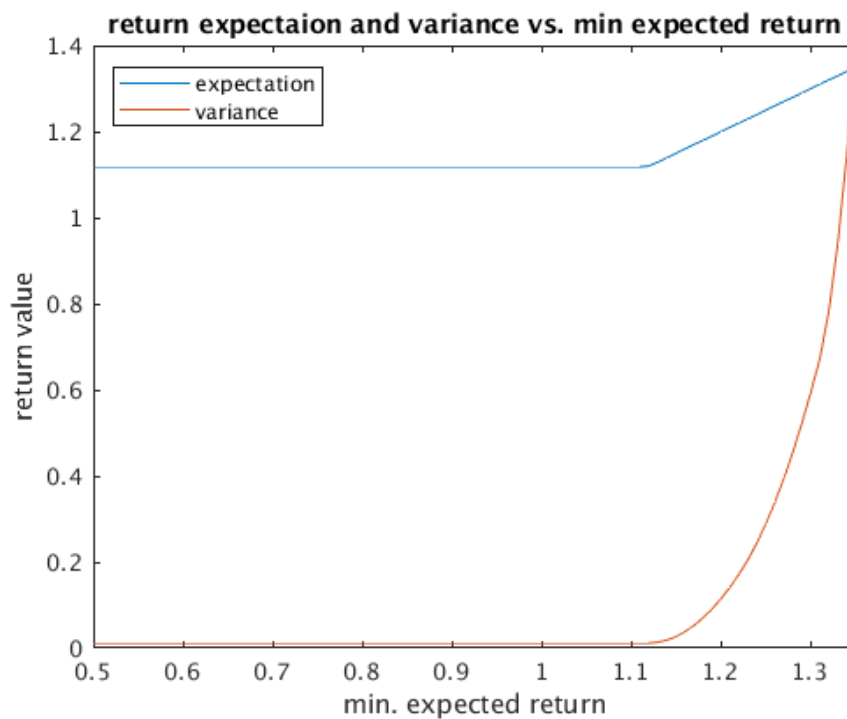
$$1^T p = 1$$

$$-x^T p \leq -r$$

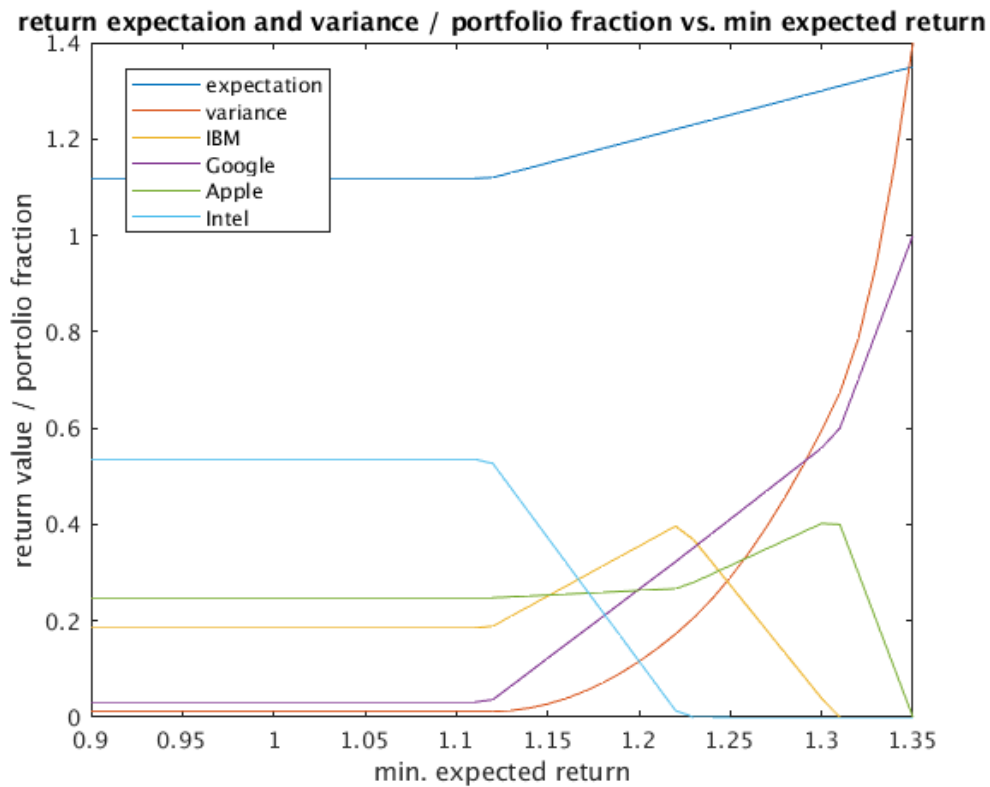
$$\text{Diag}(1_{4 \times 1}) p \leq 1$$

$$-\text{Diag}(1_{4 \times 1}) p \leq 0$$

Solving for a range of r yields:



Portfolio for a range of r :



Solver:

```
function [rs, s] = solve()
    x = [1.1, 1.35, 1.25, 1.05];
    H = [ 0.2 -0.2 -0.12 0.02;...
        -0.2 1.4 0.02 0.0;...
        -0.12 0.02 1 -0.4;...
        0.02 0 -0.4 0.2];

    Aeq = ones(1,4);
    beq = 1;
    A = -x;
    rs = 0.5:0.01:1.35
    lb = zeros(4,1);
    ub = ones(4,1);
    s = arrayfun(@(b) quadprog(H, [], A, -b, Aeq, beq, lb, ub), rs, ...
        'UniformOutput', false);
end
```

```
[rs, ret] = solve();

ps = cell2mat(ret)';
x = [1.1, 1.35, 1.25, 1.05]';
H = [ 0.2 -0.2 -0.12 0.02;...
      -0.2 1.4 0.02 0.0;...
      -0.12 0.02 1 -0.4;...
      0.02 0 -0.4 0.2];
ms = ps * x;
vs = diag(ps * H * ps');

plot(rs,ms);
hold on;
plot(rs,vs);
title('return expectaion and variance vs. min expected return');
legend('expectation','variance');
xlabel('min. expected return');
ylabel('return value');
hold off;

%% additional plot for part b
plot(rs,ms);
hold on;
plot(rs,vs);
title('return expectaion and variance / portfolio fraction vs. min expected return');
xlabel('min. expected return');
ylabel('return value / portolio fraction');
plot(rs,ps(:,1));
plot(rs,ps(:,2));
plot(rs,ps(:,3));
plot(rs,ps(:,4));
legend('expectation','variance', 'IBM', 'Google', 'Apple', 'Intel');
hold off;
```

9. Optimal Control of a Unit Mass

- a) given:

$$x(0) = 0$$

$$\dot{x}(0) = 0$$

$$f(t) = p_i, i - 1 < t \leq i, i = 1, \dots, 10$$

$$x(10) = 1$$

$$\dot{x}(10) = 0$$

$$\text{minimize } \sum_{i=1}^{10} p_i^2$$

$$\text{let } v = \dot{x}$$

$$v(i) = v(0) + \sum_{j=0}^i \frac{f(j)}{m} = v(0) + \sum_{j=1}^i p_j, m = 1$$

$$v(i) = \sum_{j=1}^i p_j = 1^T p, p \in \mathbb{R}^i$$

$$v(1) = p_1$$

$$v(2) = p_1 + p_2$$

...

$$v(10) = 1^T p = v_{des} = 0,$$

$$x(i) = x(0) + \sum_{j=0}^i v(j)$$

$$x(i) = \sum_{j=1}^i v(j) = 1^T v, v \in \mathbb{R}^i$$

$$x(i) = \sum_{j=1}^i \sum_{k=1}^j p_k$$

$$x(i) = \sum_{j=1}^i (i - j + 1) p_j$$

$$x(1) = p_1$$

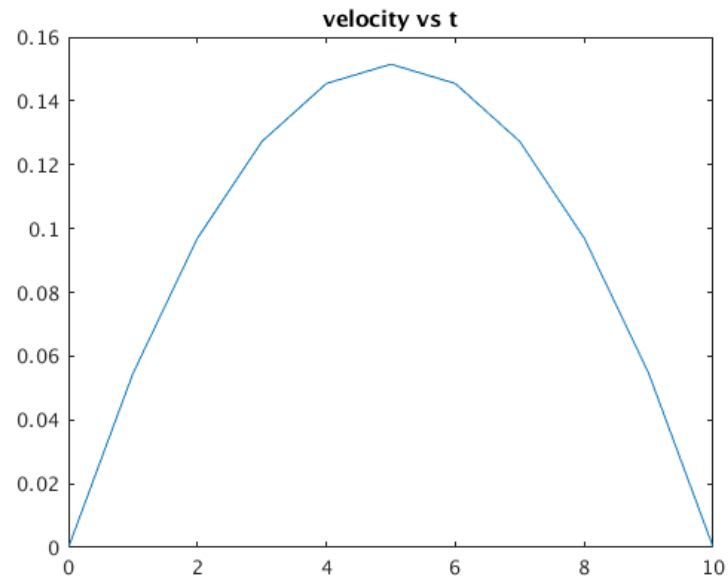
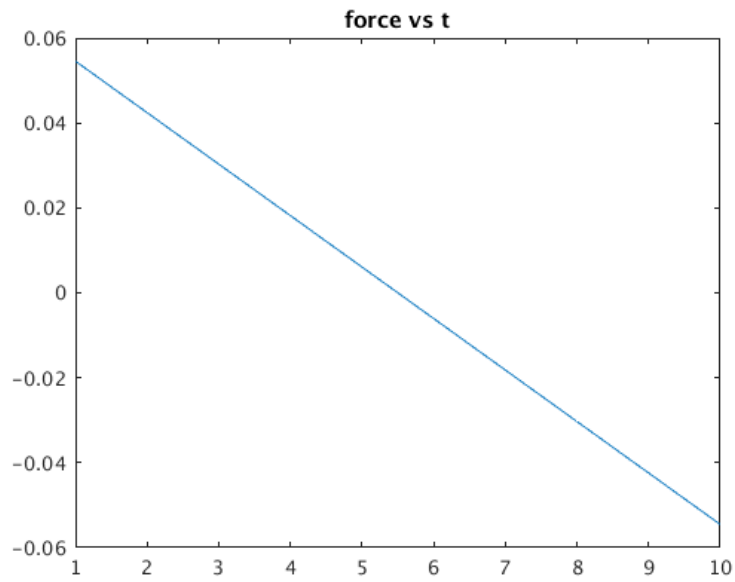
$$x(2) = 2p_1 + 1p_2$$

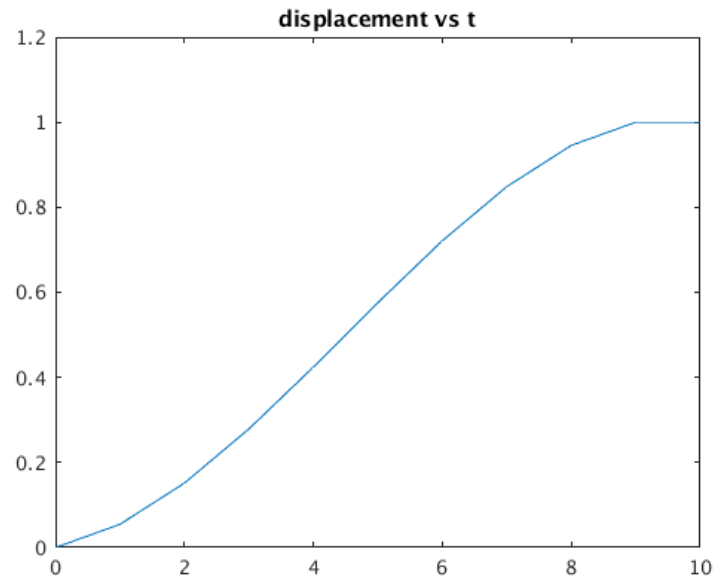
$$x(10) = 10p_1 + 9p_2 + \dots + 1p_{10}$$

QP formulation:

$$\min_f f^T I f, \text{ subject to :}$$

$$\begin{bmatrix} 10 & 9 & \dots & 2 & 1 \end{bmatrix} f = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$$





Optimal strategy in this case is to apply a smooth symmetrical force around $t = 5$, so that velocity is always non-negative and displacement is always towards the destination.

Solver:

```
H = diag(ones(10,1))
f = []
A = []
b = []
Aeq = [ones(1,10); 10:-1:1]
beq = [0;1]
f = quadprog(H,f,A,b,Aeq,beq)
```

```
plot(1:10,f)
title('force vs t')
```

```
v = zeros(11,1)
for i=2:1:11
    v(i) = v(i-1) + f(i-1)
end
```

```
plot(0:1:10,v)
title('velocity vs t')
```

```
x = zeros(11,1)
for i=1:1:11
```

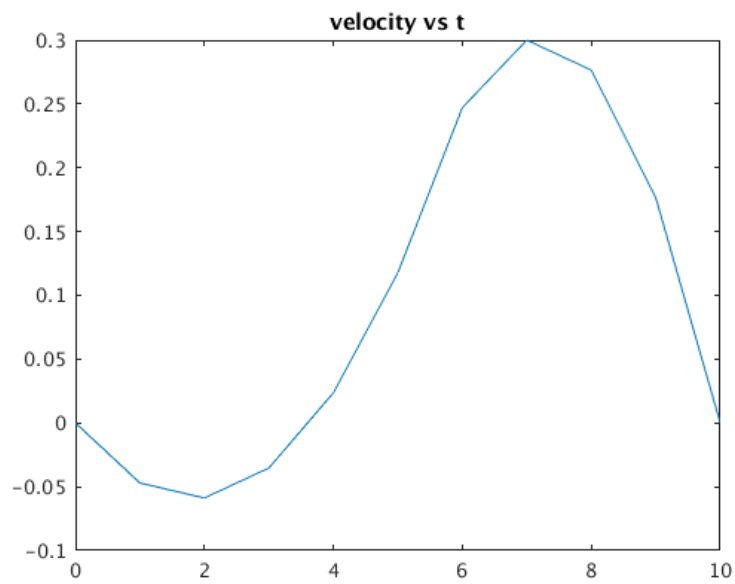
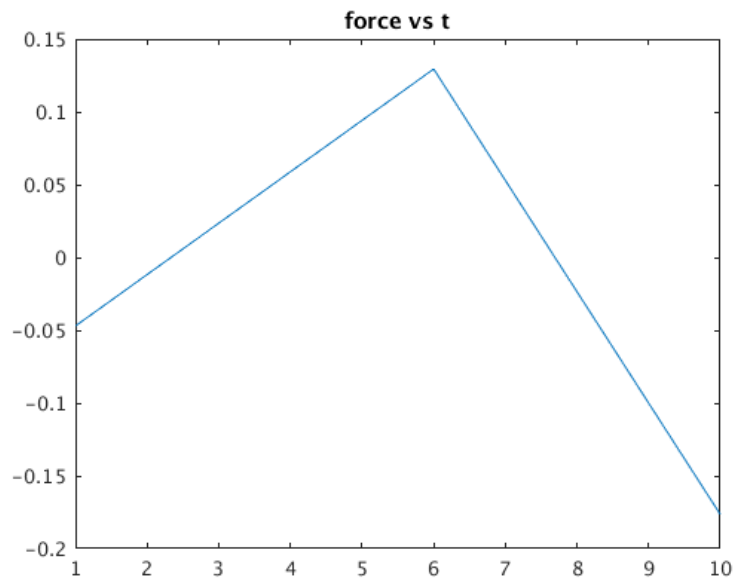
```
        if i > 1
            x(i) = x(i-1) + v(i)
        else
            x(i) = v(i)
        end
    end

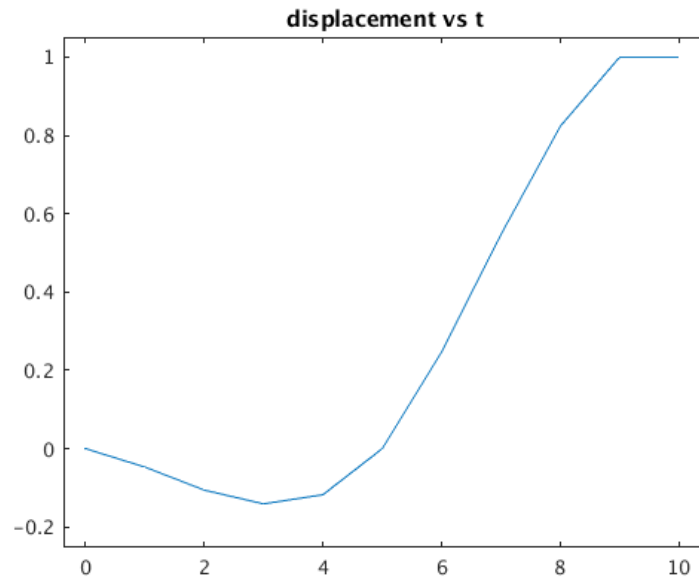
    plot(0:1:10,x)
    title('displacement vs t')
```

- b) additional constraint: $x(5) = 0$

$$\min_f f^T I f, \text{ subject to :}$$

$$\begin{bmatrix} 10 & 9 & \dots & 1^T & & & & & & \\ 5 & 4 & 3 & 2 & 1 & 0 & \dots & 0 & & \end{bmatrix} f = \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix}$$





Optimal strategy in this case is to reverse direction and then go forward again to gain enough velocity timed so that the displacement at $t=5$ is 0.

Solver:

```
H = diag(ones(10,1));
f = []; A = []; b = [];
Aeq = [ones(1,10); 10:-1:1; 5 4 3 2 1 0 0 0 0 0];
beq = [0;1;0];
f = quadprog(H,f,A,b,Aeq,beq);
plot(1:10,f)
title('force vs t')
v = zeros(11,1)
for i=2:1:11
    v(i) = v(i-1) + f(i-1)
end
plot(0:1:10,v)
title('velocity vs t')
x = zeros(11,1)
for i=1:1:11
    if i > 1
        x(i) = x(i-1) + v(i)
    else
        x(i) = v(i)
    end
end
plot(0:1:10,x)
title('displacement vs t')
```

10. Least-Square Deconvolution

- a) Find the deconvolution filter, g of size m , and the best D for the given impulse response, h , of the channel.

A convolution matrix is constructed for use in QP. Constraints are also made for the convolution value $\text{conv}(g, h)[t = D] = 1$ and filter to allow non-zero only in the appropriate indices in the system of equations. Then D is allowed to change and we collect feasible solutions for the filter and take the best one based on the lowest loss of the objective.

Assume 1-indexing as in Matlab

given: D

$$X = \begin{bmatrix} 0_{n \times 1} \\ g_{m-1} \\ \dots \\ g_0 \\ 0_{n \times 1} \end{bmatrix}$$

$t_0 = \text{size}(A, 1) - 1$ //this row index corresponds to $\text{conv}(g, h)[0]$

$$A = \begin{bmatrix} h_0 & \dots & h_{n-1} & 0_{1 \times (m+n)} \\ 0 & h_0 & \dots & h_{n-1} & 0_{1 \times (m+n-1)} \\ \dots & & & & \\ & 0_{1 \times (m+n)} & h_0 & \dots & h_{n-1} \end{bmatrix}$$

$$A_2 = A$$

$$A_2(t_0 + D, :) = 0$$

$$H = A_2^T A_2$$

$$X_i = 0, \forall i \in \{1, \dots, n\} \cup \{n + m + 1, \dots, 2n + m\}$$

$$A_{eq} = A(t_0 + D, :)$$

QP formulation for a given D :

$$\min_X X^T H X, \text{ subject to :}$$

$$A_{eq} X = 1$$

$$X_i = 0, \forall i \in \{1, \dots, n\} \cup \{n + m + 1, \dots, 2n + m\}$$

Vary D in range $[-(m + n - 1), 1]$ and solve QP to find best g using lowest $X^T H X$.

Solver:

```

n = length(h);
m = 20;

% setup convolution matrix
e = zeros(m+2*n,n);
for i=1:10
    e(:,i) = ones(m+2*n,1) * h(i);
end
A = full(spdia(e, 0:9, m+n+1, m+2*n));

% row index for conv(f,g)[t=0]
t_0 = size(A,1)-1;

gs = []; losses = []; ds = [];
% solve for different D offset from t_0 using QP
for D=-(m+n-1):1
    fprintf('d: %d\n',-D);
    A2 = [ A(1:t_0+D-1,:);
           zeros(1,m+2*n);
           A(t_0+D+1:end,:) ];
    H = A2'*A2;
    Aeq = A(t_0+D,:);
    beq = 1;
    lb = [zeros(n,1); ones(m,1) * -Inf; zeros(n,1)];
    ub = [zeros(n,1); ones(m,1) * Inf; zeros(n,1)];
    X = quadprog(H,[],[],[],Aeq,beq,lb,ub);
    if length(X) ~= 0 % guard for feasible answer
        g = flip(X(n+1:n+1+m-1)); % flip to get [g_0, ...]
        loss = X'*H*X;
        losses = [ losses; loss];
        gs = [gs; g'];
        ds = [ds; -D];
    end
end

% obtain the best answer for g
[~,idx] = sort(losses(:,1));
losses_sorted = losses(idx,:);
gs_sorted = gs(idx,:);
ds_sorted = ds(idx,:);
g_best = gs_sorted(1,:);
d_best = ds_sorted(1,:);
fprintf("D best: %d\n", d_best);

```

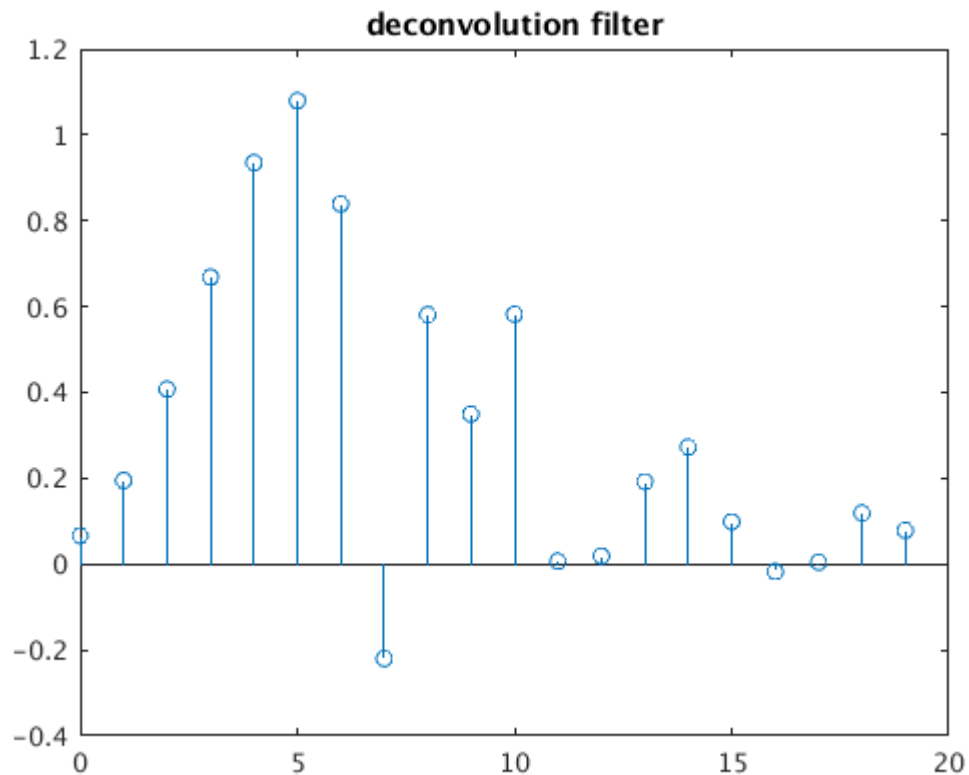
```

temp = conv(g_best,h);
assert(abs(temp(d_best+1)-1.0)<1e-15);
stem([0:length(temp)-1], temp);
title('conv(g,h)');

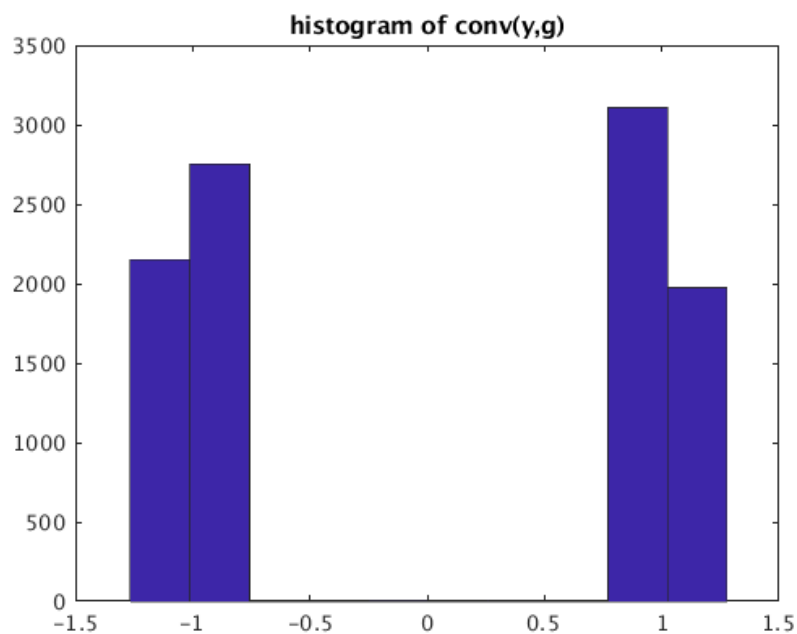
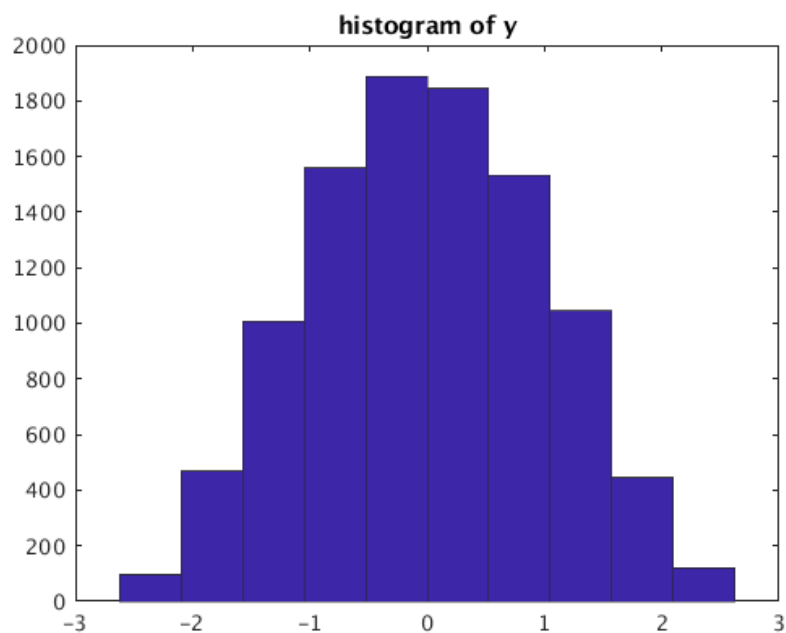
% verify on arbitrary data
samples = rand(10000,1);
output = conv(g_best, conv(h, samples));
[r,lags] = xcorr(samples,output);
[~,i]=max(r);
l = -lags(i);
assert(l==d_best);

```

Optimal parameters: $D=8$, g plotted below:



- b) Plot histogram of y (channel output) and z (filtered signal)



We have recovered the original signal. The output distribution is reasonable in that it has binary modes which is what we expect from given knowledge of the original signal input format.