

[Problem](#)[Submissions](#)[Leaderboard](#)[Editorial](#)**Problem Setter:** [Wanbo](#)**Problem Tester:** [Cao Peng](#), [Khongor](#)

Approach

What's the requirement if one hackerX missiles can defend both (t_i, f_i) and (t_j, f_j) when $t_i \leq t_j$?

$t_j - t_i \geq |f_i - f_j|$ (because you just need $|f_i - f_j|$ units time to change the frequency from f_i to f_j , so you can defend t_j after defend t_i by this hackerX missiles if you need. So if we regard every coming missiles (t_i, f_i) as an vertex i , and if $t_j - t_i \geq |f_i - f_j|$ then we add an edge from i to j , this edge indicates that one hackerX can defend the j^{th} missile after defending the i^{th} missile.

So the solution for this challenge seems obvious, we just need to find the "Minimum Path Cover"(Can be changed to MaxFlow problem) of this graph.

But n can be as large as 100000, and there are at most $O(n^2)$ edges, even generating the graph will get TLE, we need to optimize our solution.

We haven't taken advantage of the particularity of this graph in the previous analysis, we just look this as a more general problem.

$t_j - t_i \geq |f_i - f_j| \iff$ If $f_i < f_j$, $t_j - t_i \geq f_j - f_i \implies t_i - f_i \leq t_j - f_j \implies A_i \leq A_j$ ($A_i = t_i - f_i$) If $f_i \geq f_j$, $t_j - t_i \geq f_i - f_j \implies t_i + f_i \leq t_j + f_j \implies B_i \leq B_j$ ($B_i = t_i + f_i$)

Can we remove some constraints?

1. $t_i \leq t_j$
2. If $f_i < f_j$, $A_i \leq A_j \implies A_i + 2 * f_i \leq A_j + 2 * f_j \implies B_i \leq B_j$
3. If $f_i \geq f_j$, $B_i \leq B_j \implies B_i - 2 * f_i \leq B_j - 2 * f_j \implies A_i \leq A_j$

We can combine 2) with 3) to " $A_i \leq A_j \ \&\& \ B_i \leq B_j$ " $A_i + B_i \leq A_j + B_j \implies t_i \leq t_j \implies t_i \leq t_j$ can be removed.

\implies

$\text{edge}(i \rightarrow j) \iff A_i \leq A_j \ \&\& \ B_i \leq B_j$

We need to find the minimum path cover in the graph which $\text{edge}(i \rightarrow j) \iff A_i \leq A_j \ \&\& \ B_i \leq B_j$.

According to [Dilworth's theorem](#), the minimum path cover in this graph is equal to the longest anti-chain. So we just need to choose a maximum vertexes subset from the graph, where any of two do not have an edge. 1. Sort the (A_i, B_i) by the first key value A_i , if $A_i == A_j$, then smaller B_i comes first. 2. The longest decreasing subsequence of B_i will be the answer.(This is a very classic question that can be solved by $O(n \lg n)$).

Problem Setter's code:

C++

```
#include <map>
#include <set>
#include <list>
#include <queue>
#include <deque>
#include <stack>
#include <bitset>
#include <vector>
#include <ctime>
#include <cmath>
#include <cstdio>
#include <string>
#include <cstring>
#include <cassert>
#include <numeric>
#include <iomanip>
#include <sstream>
#include <fstream>
#include <iostream>
#include <algorithm>

using namespace std;
```



```

typedef long long      LL;
typedef pair<int, int>  PII;
typedef pair<LL, LL>    PLL;
typedef vector<int>     VI;
typedef vector<LL>      VL;
typedef vector<PII>     VPII;
typedef vector<PLL>     VPLL;
#define MM(a,x) memset(a,x,sizeof(a));
#define ALL(x) (x).begin(), (x).end()
#define P(x)      cerr<<"#x<<" = "<<(x)<<"\n"
#define PP(x,i)    cerr<<"#x<<i<<" = "<<x[i]<<"\n"
#define P2(x,y)    cerr<<"#x" = "<<(x)<<", "#y" = "<<(y)<<"\n"
#define TM(a,b)    cerr<<"#a" -> "#b": "<<1e3*(b-a)/CLOCKS_PER_SEC<<"ms"\n";
#define FOR(it,v) for(decltype(v.begin()) it=v.begin();it!=v.end();it++)
#define rep(i, n) for(int i = 0; i < n; i++)
#define UN(v) sort(ALL(v)), v.resize(unique(ALL(v))-v.begin())
#define mp make_pair
#define pb push_back
#define x first
#define y second
struct _ {_() {ios_base::sync_with_stdio(0);}} _;
template<class T> void PV(T a, T b) {while(a != b)cout << *a++, cout << (a != b ? " " : "\n");}
template<class T> inline bool chmin(T &a, T b) {return a > b ? a = b, 1 : 0;}
template<class T> inline bool chmax(T &a, T b) {return a < b ? a = b, 1 : 0;}
template<class T> string toString(T x, int len = 0) {stringstream ss; ss << x; string r = ss.str(); if(r.length() < len) r =
template<class T> void convert(string x, T& r) {stringstream ss(x); ss >> r;}
template<class A, class B> ostream& operator<<(ostream &o, pair<A, B> t) {o << "(" << t.x << ", " << t.y << ")"; return o;}
const int inf = 0x3f3f3f3f;
const int mod = int(1e9) + 7;
const int N = 111111;

int t[N], f[N];
int d[N];
int n;

int LIS(VI v) {
    fill(d, d + n, 1);
    for(int i = 0; i < n; i++)
        for(int j = 0; j < i; j++)
            if(v[i] > v[j]) chmax(d[i], d[j] + 1);
    int res = 0;
    for(int i = 0; i < n; i++) chmax(res, d[i]);
    return res;
}

int LIS1(VI v) {
    MM(d, 0x3f);
    d[1] = v[0];
    for(int i = 1; i < v.size(); i++) {
        int t = upper_bound(d + 1, d + N, v[i]) - d;
        if(t != 1 && d[t - 1] >= v[i]) continue;
        d[t] = v[i];
    }
    int res = 0;
    for(int i = 1; i < N; i++) if(d[i] != inf) res = i;
    return res;
}

int main() {
    cin >> n;
    for(int i = 0; i < n; i++) cin >> t[i] >> f[i];
    for(int i = 1; i < n; i++) assert(t[i] >= t[i - 1]);
    VI v;
    VPII vp;
    for(int i = 0; i < n; i++) vp.pb(mp(t[i] - f[i], t[i] + f[i]));
    sort(ALL(vp));
    for(int i = 0; i < n; i++) v.pb(vp[i].second);
    reverse(ALL(v));
    //PV(ALL(v));
    cout << LIS1(v) << endl;
    //P2(LIS1(v), LIS(v));
    //assert(LIS1(v) == LIS(v));

```



```

    return 0;
}

```

Problem Tester's code:

C++

```

#include <cstdio>
#include <algorithm>
#include <set>

using namespace std;

pair<int,int> a[100005];
set<int> have;

int main() {
    int n;
    scanf("%d",&n);
    for (int i = 0; i < n; ++i) {
        int x,y;
        scanf("%d%d",&x,&y);
        a[i] = make_pair(x + y,x - y);
    }
    sort(a, a + n);
    for (int i = 0; i < n; ++i) {
        //printf("%d %d\n",a[i].first,a[i].second);

        set<int>::iterator t = have.lower_bound(a[i].second);

        if ((t != have.end()) && (*t == a[i].second)) {
            continue;
        }
        if (t != have.begin()) {
            have.erase(--t);
        }
        have.insert(a[i].second);
    }
    printf("%d\n",have.size());
    return 0;
}

```

```

#include <cmath>
#include <cstdio>
#include <vector>
#include <iostream>
#include <algorithm>
#include <utility>
using namespace std;

#define Pair pair<int, int>

#define MAX 100000

Pair v[MAX];
int a[MAX];
int dp[MAX + 1];

int main() {
    int n;
    scanf("%d", &n);
    for (int i = 0; i < n; i++) {
        int x, y;
        scanf("%d%d", &x, &y);
    }
}

```



```
        v[i] = make_pair(x - y, x + y);
    }
    sort(v, v + n);
    for (int i = 0; i < n; i++)
        a[i] = v[n - i - 1].second;
    int r = 1;
    dp[1] = a[0];
    for (int i = 1; i < n; i++) {
        int low = 1, high = r;
        if (a[i] <= dp[1]) {
            dp[1] = a[i];
            continue;
        }
        while (low < high) {
            int mid = (low + high + 1) / 2;
            if (dp[mid] >= a[i])
                high = mid - 1;
            else
                low = mid;
        }
        if (low == r) r++;
        dp[low + 1] = a[i];
    }
    cout << r << endl;
    return 0;
}
```

Feedback

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