

1 General

line search conditions:

$$f_{k+1}^T \leq f_k + c_1 \alpha_k \nabla f_k^T p_k, c_1, \alpha_k \in (0, 1) \quad (1)$$

$$\nabla f_{k+1}^T p_k \geq c_2 \nabla f_k^T p_k, 0 < c_1 < c_2 < 1 \quad (2)$$

$$\text{where :} \quad (3)$$

$$f_k = f(x_k) \quad (4)$$

$$f_{k+1} = f(x_k + \alpha_k p_k) \quad (5)$$

2 Quasi Newton

2.1 BFGS

properties: $O(n^2)$, self correcting, slightly more iterations than Newton Method, linear convergence order and superlinear rate of convergence

secant equation:

$$B_{k+1}(x_{k+1} - x_k) = \nabla f_{k+1} - \nabla f_k$$

$$B_{k+1} s_k = y_k$$

$$s_k = \alpha_k p_k$$

$$y_k = \nabla f_{k+1} - \nabla f_k$$

$$B_{k+1} \succ 0$$

$$s_k^T B_{k+1} s_k = s_k^T y_k > 0$$

Proof.

$$y_k^T s_k = (\nabla f_{k+1} - \nabla f_k)^T s_k$$

$$\nabla f_{k+1}^T s_k \geq c_2 \nabla f_k^T s_k$$

$$(\nabla f_{k+1} - \nabla f_k)^T s_k \geq c_2 \nabla f_k^T s_k - \nabla f_k^T s_k$$

$$y_k^T s_k \geq (c_2 - 1) \nabla f_k^T s_k$$

$$c_2 < 1, s_k \text{ is a descent dir} \implies s_k^T y_k > 0$$

Curvature condition holds. \square

constrain B by solving:

$$\min_B \|B - B_k\|$$

$$\text{s.t. } B = B^T, B s_k = y_k$$

similarly, constrain B 's inverse, H where it satisfy secant equation:

$$H_{k+1} y_k = s_k$$

$$\min_H \|H - H_k\|$$

$$\text{s.t. } H = H^T, H y_k = s_k$$

using weighted Frobenius norm:

$$\|A\|_W := \|W^{1/2} A W^{1/2}\|_F$$

$$\|X\|_F := \left(\sum_i \sum_j (X_{ij})^2 \right)^{1/2}$$

solved weight matrix W satisfy $W s_k = y_k$ solution given by:

$$H_{k+1} = (I - \rho_k s_k y_k^T) H_k (I - \rho_k y_k s_k^T)$$

$$+ \rho_k s_k^T s_k$$

$$\rho_k = \frac{1}{y_k^T s_k}$$

W is the average Hessian \bar{G} :

$$\bar{G} = \int_0^1 \nabla^2 f(x_k + \tau \alpha_k p_k) d\tau$$

initial H_0 can be chosen approximately (eg: finite differences, I)

Algorithm 1: BFGS Algorithm

$H_0, x_0, \epsilon > 0$: inverse Hessian approx., initial point, convergence tolerance

x : solution

1 $k \leftarrow 0$

2 **while** $\|\nabla f_k\| > \epsilon$ **do**

3 $\alpha_k \leftarrow \text{LineSearch}(\cdot)$

4 $x_{k+1} \leftarrow x_k + \alpha_k p_k$

5 $s_k \leftarrow x_{k+1} - x_k$

6 $y_k \leftarrow \nabla f_{k+1} - \nabla f_k$

7 $\rho_k \leftarrow \frac{1}{y_k^T s_k}$

8 $H_{k+1} \leftarrow (I - \rho_k s_k y_k^T) H_k (I - \rho_k y_k s_k^T)$

9 $+ \rho_k s_k^T s_k$

10 $k \leftarrow k + 1$

using Sherman-Morrison-Woodbury formula to obtain Hessian update equation:

$$B_{k+1} = B_k - \frac{B_k s_k s_k^T B_k}{s_k^T B_k s_k} + \frac{y_k y_k^T}{y_k^T s_k}$$

proper line search is required so that BFGS also captures curvature information

inaccurate line search can be used to reduce computation cost

3 Trust Region Methods

idea:

- models local behaviour of the objective function (eg: 2nd order Taylor series)
- set local region to explore, then simultaneously find direction and step size to take
- region size adaptively set using results from previous iterations
- step may fail due to inadequately set region, which need to be adjusted
- superlinear convergence when approximate model Hessian is equal to true Hessian

using 2nd order Taylor series model with symmetric matrix approximating Hessian

$$\min_{p \in \mathbb{R}^n} m_k(p) = f_k + g_k^T p + \frac{1}{2} p^T B_k p, \text{ st. } \|p\| \leq \Delta_k$$

$\Delta_k :=$ trust region radius

$$g_k = \nabla f(x_k)$$

full step is $(p_k = -B_k^{-1}g_k)$ taken when $B \succ 0$ and $\|B_k^{-1}g_k\| \leq \Delta_k$

evaluate goodness of model with actual function by:

$$\rho_k = \frac{f(x_k) - f(x_k + p_k)}{m_k(0) - m_k(p_k)}$$

$$\text{action} \leftarrow \begin{cases} \text{expand trust region} & , \rho_k \approx 1 \text{ (agreement)} \\ \text{shrink trust region} & , \rho_k < 0 + \text{thresh} \\ \text{keep trust region} & , o/w \end{cases}$$

Algorithm 2: Trust Region Algorithm

```

1  $k \leftarrow 0$ 
2 while  $\|\nabla f_k\| > \epsilon$  do
3    $p_k \leftarrow$ 
      $\operatorname{argmin}_p f_k + g_k^T p + \frac{1}{2} p^T B_k p, \text{ st. } \|p\| \leq \Delta_k$ 
4    $\rho_k = \frac{f(x_k) - f(x_k + p_k)}{m_k(0) - m_k(p_k)}$ 
5   if  $\rho_k < \gamma(\frac{1}{4})$  then
6      $\Delta_{k+1} \leftarrow \alpha(\frac{1}{4})\Delta_k$ 
7   else if  $\rho_k > \beta(\frac{3}{4})$  and  $\|p_k\| = \Delta_k$  then
8      $\Delta_{k+1} \leftarrow \min(2\Delta_k, \hat{\Delta})$ 
9   else
10     $\Delta_{k+1} \leftarrow \Delta_k$ 
11   if  $\rho_k > \eta(\in [0, \frac{1}{4}])$  then
12      $x_{k+1} \leftarrow x_k + p_k$ 
13   else
14      $x_{k+1} \leftarrow x_k$ 
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minimizer of the 2nd order Taylor series satisfy the following:

$$\begin{aligned} (B + \lambda I)p^* &= -g \\ \lambda(\Delta - \|p^*\|) &= 0 \\ (B + \lambda I) &\succeq 0 \end{aligned}$$

solving 2nd order Taylor series using approx methods:

- dogleg
- 2-D subspace minimization
- conjugate gradient based

Cauchy Point:

Use 1st order approx. of model and gradient descent to get next iterate, bounded within trust region.

3.1 Dogleg method

if $B \succ 0$:

$$p^B = -B^{-1}g$$

$$p^* = p^B \text{ if } \Delta \geq \|p^B\|$$

$$p^U = -\frac{g^T g}{g^T B g} g \text{ (intermediate point along direction of steepest descent)}$$

interpolate between p^U and p^B :

$$\tilde{p}(\tau) = \begin{cases} \tau p^U & , \tau \in [0, 1] \\ p^U + (\tau - 1)(p^B - p^U) & , \tau \in [1, 2] \end{cases}$$

$B \succ 0 \implies \|\tilde{p}(\tau)\|$ increases wrt. τ , $m(\tilde{p}(\tau))$ decreases wrt. τ

if $\|p^B\| \leq \Delta$: p chosen at p^B

else p chosen at intersection of $\tilde{p}(\tau)$ and trust region boundary by solving:

$$\|p^U + (\tau - 1)(p^B - p^U)\|^2 = \Delta^2$$

$$p_k^S = \operatorname{argmin}_p f_k + g_k^T p, \|p\| \leq \Delta_k$$

$$\tau_k = \operatorname{argmin}_{\tau \geq 0} m_k(\tau p_k^S), \|\tau p_k^S\| \leq \Delta_k$$

$$p_k^S = \frac{-\Delta_k g_k}{\|g_k\|}$$

$$p_k^C = \tau_k p_k^S$$

$$p_k^C = -\tau_k \frac{g_k}{\|g_k\|}$$

$$\tau_k = \begin{cases} 1 & , g_k^T B_k g_k \leq 0 \\ \min(\frac{\|g_k\|^3}{\Delta_k g_k^T B_k g_k}, 1) & , o/w \end{cases}$$

3.2 Iterative Solution

Idea: solve subproblem $\min_{\|p\| \leq \Delta} m(p)$ by applying Newton's method to find λ that matches trust region radius. This is slightly more accurate per step compared to Dogleg. Use $(B + \lambda I)p^* = -g$ to solve $\min_{\|p\| \leq \Delta} m(p)$ for λ .

If $\lambda = 0$ and $(B + \lambda I)p^* = -g$, $\|p^*\| \leq \Delta$ and $(B + \lambda I) \succeq 0$: return λ

Else: find λ s.t. $(B + \lambda I) \succeq 0$ and $\|p(\lambda)\| = \Delta$, $p(\lambda) = -(B + \lambda I)^{-1}g$. Solve and return λ .

Solve $\|p(\lambda)\| - \Delta = 0$, $\lambda > \lambda_1$ via Newton's method (root finding). Approx. this to nearly a linear problem for easy solving:

Algorithm 3: Subproblem Algo

```

1 for  $l = 0, 1, \dots$  do
2   solve  $B + \lambda^l I = R^T R$ 
3    $R^T R p_l = -g$ 
4    $R^T q_l = p_l$ 
5    $\lambda^{l+1} \leftarrow \lambda^l + (\frac{\|p_l\|}{\|q_l\|})^2 (\frac{\|p_l\| - \Delta}{\Delta})$  check  $\lambda \geq \lambda_1$ 

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4 Conjugate Gradient

4.1 linear method

Assuming unconstrained problem with strict convex quadratic objective function:

$$\frac{1}{2}x^T Ax - b^T x, A \succ 0, A^T = A$$

$\nabla(\frac{1}{2}x^T Ax - b^T x) = Ax - b$, thus $\min_x x^T Ax - b^T x$ transformed to $Ax - b = 0$.

$x_{k+1} = x_k + \alpha_k p_k$, solve for α :

$$\frac{\partial}{\partial \alpha} ((x_k + \alpha_k p_k)^T A(x_k + \alpha_k p_k) - b^T(x_k + \alpha_k p_k)) = 0$$

$$\alpha_k = \frac{-p_k^T r_k}{p_k^T A p_k}$$

4.2 Conjugate Direction

Search directions linearly independent wrt. A.

$$(\forall i \neq j) p_i^T A p_j = 0$$

Properties:

- Residual eliminated one direction at a time, resulting in max of n iterations.
- Optimal if Hessian is diagonal, if not can try preconditioning.
- Current residual is orthogonal to all previous search directions.
- Any set of conjugate directions can be used (eg: eigenvectors, Gram-Schmidt)

Expanding subspace minimizer:

Using conjugate directions to generate sequence $\{x\}$, then:

$r_k^T p_i = 0, \forall i < k$, x_k is minimizer of $\frac{1}{2}x^T Ax - b^T x$ over $\{x | x = x_0 + \text{span}\{p_0, \dots, p_{k-1}\}\}$

Proof.

$$\tilde{x} = x_0 + \sum_i \sigma_i p_i$$

\tilde{x} minimizes over $\{x_0 + \text{span}\{p_0, \dots, p_{k-1}\}\} \iff r(\tilde{x})^T p_i = 0$

$$h(\sigma) = \phi(\tilde{x})$$

$$\phi(x) = \frac{1}{2}x^T Ax - b^T x$$

h is also strictly convex quadratic,

with unique σ^* satisfying:

$$\frac{\partial h(\sigma^*)}{\partial \sigma_i} = 0, i = [0, k-1]$$

$$\frac{\partial h(\sigma^*)}{\partial \sigma_i} = \nabla \phi(\tilde{x})^T p_i = 0, i = [0, k-1]$$

$$\nabla \phi(x) = Ax - b = r$$

$$r(\tilde{x})^T p_i = 0, i = [0, k-1]$$

□

$p_i^T r_k = 0, i = [0, k-1]$ via induction:

Proof.

base case : $x_1 = x_0 + \alpha_0 p_0$ minimizes ϕ along p_0

$$\implies r_1^T p_0 = 0$$

case: $r_{k-1}^T p_i = 0, i = [0, k-2]$:

$$r_k = r_{k-1} + \alpha_{k-1} A p_{k-1}$$

$$p_{k-1}^T r_k = p_{k-1}^T r_{k-1} + \alpha_{k-1} p_{k-1}^T A p_{k-1} = 0 \text{ (by construction)}$$

$$A\text{-conjugate} \implies p_{k-1}^T A p_{k-1}$$

case: $\forall i = [0, k-2] : p_i^T r_k = 0$

$$p_i^T r_k = p_i^T r_{k-1} + \alpha_{k-1} p_i^T A p_{k-1}$$

$$p_i^T r_{k-1} = 0 \text{ (by induction hypothesis)}$$

$$\alpha_{k-1} p_i^T A p_{k-1} = 0 \text{ (conjugacy)}$$

$$p_i^T r_k = 0, i = [0, k-1]$$

□

4.3 Conjugate Gradient Method

Idea:

- uses only previous search direction to compute current search direction
- p_k set to linear combination of $-r_k$ and p_{k-1}
- impose $p_k^T A p_{k-1} = 0$

$$\begin{aligned} p_k &= -r_k + \beta_k p_{k-1} \\ p_{k-1}^T A p_k &= -p_{k-1}^T A r_k + \beta_k p_{k-1}^T A p_{k-1} \\ 0 &= -p_{k-1}^T A r_k + \beta_k p_{k-1}^T A p_{k-1} \\ \beta &= \frac{p_{k-1}^T A r_k}{p_{k-1}^T A p_{k-1}} \\ p_0 &= -(A x_0 - b) = -r_0 \end{aligned}$$

Algorithm 4: Basic Conjugate Gradient Algorithm

```

1  $r_0 \leftarrow A x_0 - b$ 
2  $p_0 = -r_0$ 
3 for  $k = [0, ..n - 1]$  do
4   if  $r_k == 0$  then
5     return  $x_k$ 
6   else
7      $\alpha_k \leftarrow \frac{-r_k^T p_k}{p_k^T A p_k}$ 
8      $x_{k+1} \leftarrow x_k + \alpha_k p_k$ 
9      $r_{k+1} \leftarrow A x_{k+1} - b$ 
10     $b_{k+1} \leftarrow \frac{p_k^T A r_{k+1}}{p_k^T A p_k}$ 
11     $p_{k+1} \leftarrow -r_{k+1} + \beta_{k+1} p_k$ 

```

p and r_k is within krylov subspace:

$$K(r_0; k) = \text{span}\{r_0, A r_0, \dots, A^k r_0\}$$

if $r_k \neq 0$:

$$r_k^T r_i = 0, i = [0, k - 1]$$

$$\text{span}\{r_0, \dots, r_k\} = \text{span}\{r_0, A r_0, \dots, A^k r_0\}$$

$$\text{span}\{p_0, \dots, p + k\} = \text{span}\{r_0, A r_0, \dots, A^k r_0\}$$

$$p_k^T A p_i = 0, i = [0, k - 1]$$

then, $\{x_k\} \rightarrow x^*$ in at most n steps.

Simplification:

$$\begin{aligned} p_{k+1} &\leftarrow -r_{k+1} + \beta_{k+1} p_k \\ \alpha_k &\leftarrow \frac{-r_k^T p_k}{p_k^T A p_k} \\ \alpha_k &\leftarrow \frac{-r_k^T (-r_k + \beta_k p_{k-1})}{p_k^T A p_k} \\ (\forall i = [0, k - 1]) r_k^T p_i &= 0 \implies \beta_k r_k^T p_{k-1} = 0 \\ \alpha_k &\leftarrow \frac{r_k^T r_k}{p_k^T A p_k} \\ r_{k+1} &= r_k + \alpha_k A p_k \\ A p_k &= \frac{r_{k+1} - r_k}{\alpha_k} \\ \beta &= \frac{p_k^T A r_{k+1}}{p_k^T A p_k} \\ p_k^T A p_k &= p_k^T \frac{r_{k+1} - r_k}{\alpha_k} = \frac{-p_k^T r_k}{\alpha_k} \text{ (conjugacy)} \\ p_k^T A p_k &= -\frac{(-r_k + \beta_k p_{k-1})^T r_k}{\alpha_k} = \frac{r_k^T r_k}{\alpha_k} \text{ (conjugacy)} \\ p_k^T A r_{k+1} &= r_{k+1}^T A p_k \\ p_k^T A r_{k+1} &= r_{k+1}^T \frac{r_{k+1} - r_k}{\alpha_k} \\ r_k &\in \text{span}\{p_k, p_{k-1}\} \text{ and } r_{k+1}^T p_i = 0, i = [0, k] \implies \\ p_k^T A r_{k+1} &= \frac{r_{k+1}^T r_{k+1}}{\alpha_k} \\ \beta_{k+1} &\leftarrow \frac{r_{k+1}^T r_{k+1}}{r_k^T r_k} \end{aligned}$$

4.4 Nonlinear Method

Minimize general convex function or nonlinear function. Variants: FR, PR.

4.4.1 FR

Modify linear CG by:

- replace residual by gradient of objective, ∇f_k
- replace α_k computation by a linear search to find approx. minimum along search direction

Equivalent to linear CG if objective is strongly convex quadratic.

Linear search for α_k with stong Wolfe condition to ensure p_k 's are descent directions wrt. objective function.

4.4.2 PR

Replace β_{k+1} computation in FR with:

$$\beta_{k+1} \leftarrow \frac{\nabla f_{k+1}^T (\nabla f_{k+1} - \nabla f_k)}{\nabla f_k^T \nabla f_k}$$