$July\ 22,\ 2013$

1 Model

The atomic unit of the factor graph is shown in Figure 1.

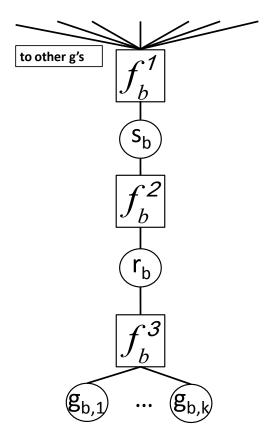


Figure 1: Factor graph

We require efficient computation of the messages.

2 Top-down messages

2.1 $\mu_{f_b^1 \to s_b}$

The general form of the message is given by

$$\mu_{f_b^1 \to s_b}(s_b) = \sum_{\vec{q}} f_b^1(s_b, \vec{q}) \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k})$$
(1)

We divide the cases into $s_b = 0$, and $s_b = 1$.

$$\mu_{f_b^1 \to s_b}(s_b = 0) = \sum_{\vec{g}} f_b^1(s_b = 0, \vec{g}) \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k})$$
 (2)

$$= (1 - \epsilon) \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)$$
(3)

$$\mu_{f_b^1 \to s_b}(s_b = 1) = \sum_{\vec{q}} f_b^1(s_b = 1, \vec{g}) \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k})$$
(4)

$$= \epsilon \times \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)$$
 (5)

$$+ \sum_{\vec{g}} \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k}) - \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)$$

$$= \epsilon \times \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)$$
 (6)

$$+1 - \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)$$

$$= \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)(\epsilon - 1) + 1 \tag{7}$$

$$= 1 - (1 - \epsilon) \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)$$
 (8)

$$= 1 - \mu_{f_i^1 \to s_b}(s_b = 0) \tag{9}$$

where we have made use of the property that messages from a node/factor sum to 1

2.2 $\mu_{s_b \to f_b^2}$

if s_b has not been observed, this message is given by

$$\mu_{s_b \to f_b^2}(s_b) = \mu_{f_b^1 \to s_b}(s_b) \tag{10}$$

otherwise, it is a constant message that indicates its observed state.

2.3 $\mu_{f_b^2 \to r_b}$

$$\mu_{f_b^2 \to r_b}(r_b) = \sum_{s_b} f_b^2(s_b, r_b) \mu_{s_b \to f_b^2}(s_b)$$
(11)

$$= \sum_{s_b} P(r_b \mid s_b) \mu_{s_b \to f_b^2}(s_b)$$
 (12)

 s_b only has two values, so this is easy to compute.

 $2.4 \quad \mu_{r_b \to f_b^3}$

$$\mu_{r_b \to f_b^3}(r_b) = \mu_{f_b^2 \to r_b}(r_b) \tag{13}$$

2.5 $\mu_{f_b^3 \to g_{b,k}}$

$$\mu_{f_b^3 \to g_{b,k}}(g_{b,k}) = \sum_{r_b} \sum_{k' \neq k} f_b^3(r_b, \vec{g}_b) \mu_{r_b \to f_b^2}(r_b) \prod_{k' \neq k} \mu_{g_{b,k} \to f_b^3}(g_{b,k'})$$
(14)

$$= \sum_{r_b} \mu_{r_b \to f_b^2}(r_b) \sum_{k' \neq k} P(\vec{g}_b \mid r_b) \prod_{k' \neq k} \mu_{g_{b,k} \to f_b^3}(g_{b,k'})$$
(15)

$$= \sum_{r_b} \mu_{r_b \to f_b^2}(r_b) \sum_{k' \neq k} P(\vec{g}_{b,k} \mid r_b) \prod_{k' \neq k} P(\vec{g}_{b,k'} \mid r_b) \mu_{g_{b,k} \to f_b^3}(g_{b,k'})$$
(16)

$$= \sum_{r_b} \mu_{r_b \to f_b^2}(r_b) P(\vec{g}_{b,k} \mid r_b) \prod_{k' \neq k} \sum_{g_{b,k'}} P(\vec{g}_{b,k'} \mid r_b) \mu_{g_{b,k} \to f_b^3}(g_{b,k'})$$
(17)

The number of values r_b can take should be small (no more than, say 10? 20?), so this should be possible to compute.

3 Bottom-up messages

$3.1 \quad \mu_{f_b^1 \to g_{b',k}}$

Recall that f_b^1 connects to other bricks. We denote one such brick by b'.

Because f_b^1 only cares about whether $g_{b',k}$ points to b, we reduce the domain of $g_{b',k}$ to 0, 1 which means "doesn't point to b" and "points to b", respectively. The messages can be "expanded" by $g_{b',k}$ to its original domain when $g_{b',k}$ needs to send a message to f_b^3 .

The general form of this message is

$$\mu_{f_b^1 \to g_{b',k}}(g_{b,k'}) = \sum_{s_b} \sum_{g \neq g_{b',k}} f_b^1(s_b, \mathbf{g}) \mu_{s_b \to f_b^1}(s_b) \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g).$$
 (18)

We divide the cases into $g_{b',k} = 0$, and $g_{b',k} = 1$

$$\begin{split} \mu_{f_b^1 \to g_{b',k}}(g_{b',k} = 0) &= & \mu_{s_b \to f_b^1}(s_b = 0) \times (1 - \epsilon) \times \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0) \\ &+ \mu_{s_b \to f_b^1}(s_b = 1) \times \epsilon \times \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0) \\ &+ \mu_{s_b \to f_b^1}(s_b = 1) (\prod_{g \neq g_{b',k}} \sum_{g} \mu_{g \to f_b^1}(g) - \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0)) \\ &= & \mu_{s_b \to f_b^1}(s_b = 0) \times (1 - \epsilon) \times \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0) \\ &+ \epsilon \times \mu_{s_b \to f_b^1}(s_b = 1) (\prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0)) \\ &= & \mu_{s_b \to f_b^1}(s_b = 1) (1 - \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0)) \\ &= & \mu_{s_b \to f_b^1}(s_b = 0) \times (1 - \epsilon) \times \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0) \\ &+ \mu_{s_b \to f_b^1}(s_b = 1) (\prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0)) \\ &= & \mu_{s_b \to f_b^1}(s_b = 0) \times (1 - \epsilon) \times \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0) \\ &+ \mu_{s_b \to f_b^1}(s_b = 0) \times (1 - \epsilon) \times \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0) \\ &+ \mu_{s_b \to f_b^1}(s_b = 0) \times (1 - \epsilon) \times \frac{\prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0)}{\mu_{g_{b',k} \to f_b^1}(g = 0)} \\ &+ \mu_{s_b \to f_b^1}(s_b = 0) \times \frac{\mu_{g \to f_b^1}(s = 0)}{\mu_{g_{b',k} \to f_b^1}(g = 0)} (\epsilon - 1) + 1)) \\ &= & \mu_{s_b \to f_b^1}(s_b = 0) \times \frac{\mu_{s_b \to f_b^1}(s_b = 0)}{\mu_{g_{b',k} \to f_b^1}(g = 0)} (\epsilon - 1) + 1)) \\ &= & \frac{\mu_{s_b \to f_b^1}(s_b = 0)^2}{\mu_{g_{b',k} \to f_b^1}(s_b = 0)}(\epsilon - 1) + \mu_{g_{b',k} \to f_b^1}(g = 0)} \\ &+ \mu_{s_b \to f_b^1}(s_b = 1)\mu_{s_b \to f_b^1}(s_b = 0)(\epsilon - 1) + \mu_{g_{b',k} \to f_b^1}(g = 0)} \\ &= & \frac{\mu_{s_b \to f_b^1}(s_b = 0)^2 + \mu_{s_b \to f_b^1}(s_b = 1)\mu_{s_b \to f_b^1}(s_b = 0)(\epsilon - 1) + \mu_{g_{b',k} \to f_b^1}(g = 0)}{\mu_{g_{b',k} \to f_b^1}(s_b = 0)(\epsilon - 1) + \mu_{g_{b',k} \to f_b^1}(g = 0)} \\ &= & \frac{\mu_{s_b \to f_b^1}(s_b = 0)^2 + \mu_{s_b \to f_b^1}(s_b = 1)\mu_{s_b \to f_b^1}(s_b = 0)(\epsilon - 1) + \mu_{g_{b',k} \to f_b^1}(g = 0)}{\mu_{g_{b',k} \to f_b^1}(g = 0)} \end{split}$$

where $\mu_{g_{b',k}\to f_b^1}^0$ means the old message sent by $g_{b',k}$, and we have made use of the property that messages from a node/factor sum to 1. Note that we can compute this message without having to explicitly examine all the other $g'_{b,k}s$; we re-used the computation from computing $\mu_{s_b\to f_b^1}(s_b=0)$.

$$\mu_{f_b^1 \to g_{b',k}}(g_{b,k'} = 1) = \mu_{s_b \to f_b^1}(s_b = 1) \prod_{g \neq g_{b',k}} \sum_g \mu_{g \to f_b^1}(g)$$
 (26)

$$= \mu_{s_b \to f_b^1}(s_b = 1) \tag{27}$$

 $3.2 \quad \mu_{g_{b,k} \to f_b^3}$

$$\mu_{g_{b,k}\to f_b^3}(g_{b,k}) = \prod_{b'\in V(b)} \mu_{f_{b'}^3\to g_{b,k}}(g_{b,k})$$
(28)

wher V(b) is the after set of brick b. Note that $\mu_{f_{b'}^3 \to g_{b,k}}(g_{b,k})$ needs to be "expanded" to the proper domain of $g_{b,k}$

3.3 $\mu_{f_b^3 \to r_b}$

$$\mu_{f_b^3 \to r_b}(r_b) = \sum_{\vec{g}} f_b^3(\vec{g}, r_b) \prod_k \mu_{g_{b,k} \to f_b^3}(g_{b,k})$$
(29)

$$= \sum_{\vec{n}} \prod_{k} P(g_{b,k} \mid r_b) \prod_{k} \mu_{g_{b,k} \to f_b^3}(g_{b,k})$$
 (30)

$$= \prod_{k} \sum_{g_{b,k}} P(g_{b,k} \mid r_b) \mu_{g_{b,k} \to f_b^3}(g_{b,k})$$
 (31)

3.4 $\mu_{r_b \to f_b^2}$

$$\mu_{r_b \to f_b^2}(r_b) = \mu_{f_b^3 \to r_b}(r_b) \tag{32}$$

 $3.5 \quad \mu_{f_b^2 \to s_b}$

$$\mu_{f_b^2 \to s_b}(s_b) = \sum_{r_b} f_b^2(r_b, s_b) \mu_{r_b \to f_b^2}(r_b)$$
(33)

$$= \sum_{r_b} P(r_b \mid s_b) \mu_{r_b \to f_b^2}(r_b)$$
 (34)

These messages should be trivial to compute, since the set of values r_b is small.

3.6 $\mu_{s_b \to f_b^1}$

$$\mu_{s_b \to f_{\iota}^{(s_b)}} = \mu_{f_b^2 \to s_b}(s_b) \tag{35}$$