## August 2, 2013

# 1 Model

The atomic unit of the factor graph is shown in Figure 1.

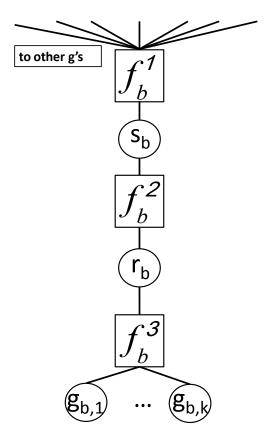


Figure 1: Factor graph

We require efficient computation of the messages.

## 2 Top-down messages

#### 2.1 $\mu_{f_b^1 \to s_b}$

The general form of the message is given by

$$\mu_{f_b^1 \to s_b}(s_b) = \sum_{\vec{q}} f_b^1(s_b, \vec{q}) \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k})$$
(1)

We divide the cases into  $s_b = 0$ , and  $s_b = 1$ .

$$\mu_{f_b^1 \to s_b}(s_b = 0) = \sum_{\vec{g}} f_b^1(s_b = 0, \vec{g}) \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k})$$
 (2)

$$= (1 - \epsilon) \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)$$
(3)

$$\mu_{f_b^1 \to s_b}(s_b = 1) = \sum_{\vec{q}} f_b^1(s_b = 1, \vec{g}) \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k})$$
(4)

$$= \epsilon \times \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)$$
 (5)

$$+ \sum_{\vec{g}} \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k}) - \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)$$

$$= \epsilon \times \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)$$
 (6)

$$+1 - \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)$$

$$= \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)(\epsilon - 1) + 1 \tag{7}$$

$$= 1 - (1 - \epsilon) \prod_{g_{b',k}} \mu_{g_{b',k} \to f_b^1}(g_{b',k} = 0)$$
 (8)

$$= 1 - \mu_{f_i^1 \to s_b}(s_b = 0) \tag{9}$$

where we have made use of the property that messages from a node/factor sum to 1

## **2.2** $\mu_{s_b \to f_b^2}$

if  $s_b$  has not been observed, this message is given by

$$\mu_{s_b \to f_b^2}(s_b) = \mu_{f_b^1 \to s_b}(s_b) \tag{10}$$

otherwise, it is a constant message that indicates its observed state.

2.3  $\mu_{f_b^2 \to r_b}$ 

$$\mu_{f_b^2 \to r_b}(r_b) = \sum_{s_b} f_b^2(s_b, r_b) \mu_{s_b \to f_b^2}(s_b)$$
(11)

$$= \sum_{s_b} P(r_b \mid s_b) \mu_{s_b \to f_b^2}(s_b)$$
 (12)

 $s_b$  only has two values, so this is easy to compute.

2.4  $\mu_{r_b \to f_b^3}$ 

$$\mu_{r_b \to f_b^3}(r_b) = \mu_{f_b^2 \to r_b}(r_b) \tag{13}$$

**2.5**  $\mu_{f_b^3 \to g_{b,k}}$ 

$$\mu_{f_b^3 \to g_{b,k}}(g_{b,k}) = \sum_{r_b} \sum_{k' \neq k} f_b^3(r_b, \vec{g}_b) \mu_{r_b \to f_b^3}(r_b) \prod_{k' \neq k} \mu_{g_{b,k} \to f_b^3}(g_{b,k'})$$
(14)

$$= \sum_{r_b} \mu_{r_b \to f_b^3}(r_b) \sum_{k' \neq k} P(\vec{g}_b \mid r_b) \prod_{k' \neq k} \mu_{g_{b,k} \to f_b^3}(g_{b,k'})$$
 (15)

$$= \sum_{r_b} \mu_{r_b \to f_b^3}(r_b) \sum_{k' \neq k} P(\vec{g}_{b,k} \mid r_b) \prod_{k' \neq k} P(\vec{g}_{b,k'} \mid r_b) \mu_{g_{b,k} \to f_b^3}(g_{b,k'})$$
(16)

$$= \sum_{r_b} \mu_{r_b \to f_b^3}(r_b) P(\vec{g}_{b,k} \mid r_b) \prod_{k' \neq k} \sum_{g_{b,k'}} P(\vec{g}_{b,k'} \mid r_b) \mu_{g_{b,k} \to f_b^3}(g_{b,k'})$$
(17)

The number of values  $r_b$  can take should be small (no more than, say 10? 20?), so this should be possible to compute.

**2.6**  $\mu_{g_{b,k}\to f_b^1}$ 

$$\mu_{g_{b,k}\to f_b^1}(g_{b,k}) = \mu_{f_b^3\to g_{b,k}}(g_{b,k}) \prod_{b'\neq b} \mu_{f_{b'}^1\to g_{b,k}}(g_{b,k})$$
(18)

# 3 Bottom-up messages

3.1  $\mu_{f_b^1 \to g_{b',k}}$ 

Recall that  $f_b^1$  connects to other bricks. We denote one such brick by b'.

Because  $f_b^1$  only cares about whether  $g_{b',k}$  points to b, we reduce the domain of  $g_{b',k}$  to 0, 1 which means "doesn't point to b" and "points to b", respectively. The messages can be "expanded" by  $g_{b',k}$  to its original domain when  $g_{b',k}$  needs to send a message to  $f_b^3$ .

The general form of this message is

$$\mu_{f_b^1 \to g_{b',k}}(g_{b,k'}) = \sum_{s_b} \sum_{g \neq g_{b',k}} f_b^1(s_b, \mathbf{g}) \mu_{s_b \to f_b^1}(s_b) \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g).$$
 (19)

We divide the cases into  $g_{b',k} = 0$ , and  $g_{b',k} = 1$ 

$$\begin{split} \mu_{f_b^1 \to g_{b',k}}(g_{b',k} = 0) &= \mu_{s_b \to f_b^1}(s_b = 0) \times (1 - \epsilon) \times \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0) \\ &+ \mu_{s_b \to f_b^1}(s_b = 1) \times \epsilon \times \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0) \\ &+ \mu_{s_b \to f_b^1}(s_b = 1) (\prod_{g \neq g_{b',k}} \sum_g \mu_{g \to f_b^1}(g) - \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0)) \\ &= \mu_{s_b \to f_b^1}(s_b = 0) \times (1 - \epsilon) \times \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0) \\ &+ \epsilon \times \mu_{s_b \to f_b^1}(s_b = 1) \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0) \\ &+ \mu_{s_b \to f_b^1}(s_b = 1) (1 - \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0)) \\ &= \mu_{s_b \to f_b^1}(s_b = 0) \times (1 - \epsilon) \times \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0) \\ &+ \mu_{s_b \to f_b^1}(s_b = 1) (1 - \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0)) \\ &+ \mu_{s_b \to f_b^1}(s_b = 1) (1 - \prod_{g \neq g_{b',k}} \mu_{g \to f_b^1}(g = 0)) \end{split}$$

where we have made use of the property that messages from a node/factor sum to 1.

$$\mu_{f_b^1 \to g_{b',k}}(g_{b,k'} = 1) = \mu_{s_b \to f_b^1}(s_b = 1) \prod_{g \neq g_{b',k}} \sum_g \mu_{g \to f_b^1}(g)$$
(23)

$$= \mu_{s_b \to f_b^1}(s_b = 1) \tag{24}$$

 $3.2 \quad \mu_{g_{b,k} \to f_b^3}$ 

$$\mu_{g_{b,k}\to f_b^3}(g_{b,k}) = \prod_{b'\in V(b)} \mu_{f_{b'}^3\to g_{b,k}}(g_{b,k})$$
(25)

wher V(b) is the after set of brick b. Note that  $\mu_{f_{b'}^3 \to g_{b,k}}(g_{b,k})$  needs to be "expanded" to the proper domain of  $g_{b,k}$ 

3.3  $\mu_{f_b^3 \to r_b}$ 

$$\mu_{f_b^3 \to r_b}(r_b) = \sum_{\vec{q}} f_b^3(\vec{q}, r_b) \prod_k \mu_{g_{b,k} \to f_b^3}(g_{b,k})$$
(26)

$$= \sum_{\vec{q}} \prod_{k} P(g_{b,k} \mid r_b) \prod_{k} \mu_{g_{b,k} \to f_b^3}(g_{b,k})$$
 (27)

$$= \prod_{k} \sum_{g_{b,k}} P(g_{b,k} \mid r_b) \mu_{g_{b,k} \to f_b^3}(g_{b,k})$$
 (28)

3.4  $\mu_{r_b \to f_b^2}$ 

$$\mu_{r_b \to f_b^2}(r_b) = \mu_{f_b^3 \to r_b}(r_b) \tag{29}$$

 $\mathbf{3.5} \quad \mu_{f_b^2 \to s_b}$ 

$$\mu_{f_b^2 \to s_b}(s_b) = \sum_{r_b} f_b^2(r_b, s_b) \mu_{r_b \to f_b^2}(r_b)$$
(30)

$$= \sum_{r_b} P(r_b \mid s_b) \mu_{r_b \to f_b^2}(r_b)$$
 (31)

These messages should be trivial to compute, since the set of values  $r_b$  is small.

3.6  $\mu_{s_b \to f_b^1}$ 

$$\mu_{s_b \to f_b^(s_b)} = \mu_{f_b^2 \to s_b}(s_b) \tag{32}$$