# Introduction to Model Checking

(Preview of Core Module)

Clemens Grabmayer

Emilio Tuosto

https://clegra.github.io

https://cs.gssi.it/emilio.tuosto/

Department of Computer Science



GRAN SASSO SCIENCE INSTITUTE



SCHOOL OF ADVANCED STUDIES Scuola Universitaria Superiore

December 2, 2024

otivation mc-schema example counterex-trace logical properties topics book module organizatio

# Model Checking

... is an effective automatable technique:

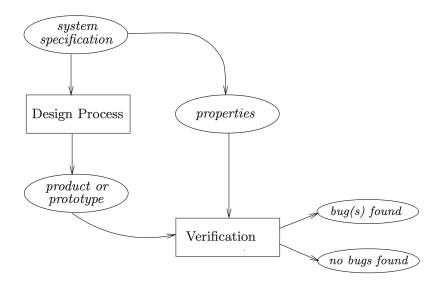
- to expose potential software design errors;
- that, given a finite-state model of a system and a formal property, systematically checks whether this property holds for that model.

#### Strengths:

- widely applicable (embedded systems, software engineering, hardware design)
- supports partial verification (of system parts)
- provides diagnostic information for debugging
- has sound mathematical underpinning (logic and process theory)

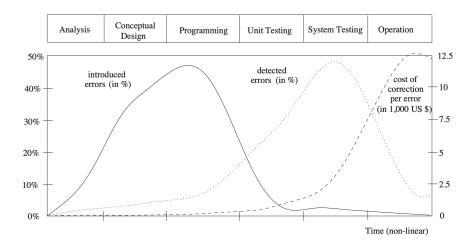
tivation mc-schema example counterex-trace logical properties topics book module organization

#### Hard-/Software Verification (traditionally)



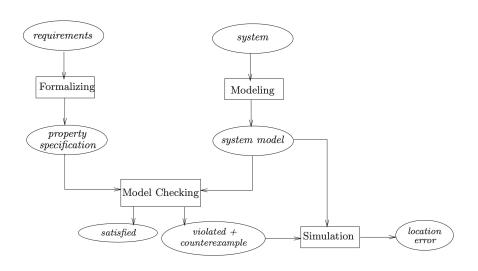
otivation mc-schema example counterex-trace logical properties topics book module organizatio

# Error introduction, detection, and repair costs



ptivation mc-schema example counterex-trace logical properties topics book module organizati

# Model checking



# Example: concurrency and non-determinism

Programs Inc, Dec, and Reset cooperate, and use a shared variable x:

```
proc Inc
                           proc Dec
                                                     proc Reset
 while true
                            while true
                                                      while true
   do
                              ob
                                                        ob
    if \times < 200
                               if x > 0
                                                         if x = 200
      then x := x + 1
                                then x := x - 1
                                                          then x := 0
    fi
                               fi
                                                         fi
   od
                              od
                                                        od
```

#### Example: concurrency and non-determinism

Programs Inc, Dec, and Reset cooperate, and use a shared variable x:

```
proc Inc
                              proc Dec
                                                            proc Reset
 while true
                                while true
                                                              while true
   dΩ
                                  dΩ
                                                                ob
     if \times < 200
                                    if x > 0
                                                                 if \times = 200
                                    then \times := \times - 1
       then x := x + 1
                                                                  then \mathbf{x} := 0
     fi
                                    fi
                                                                 fi
   od
                                  od
                                                                od
```

Question: Is  $0 \le x \le 200$  always guaranteed?

#### Modeling (by labeled transition systems)

```
\begin{array}{c} \textbf{proc Inc} \\ \textbf{while true} \\ \textbf{do} \\ \textbf{if } \times < 200 \\ \textbf{then } \times := \times + 1 \\ \textbf{fi} \\ \textbf{od} \end{array}
```

```
\begin{array}{c} \textbf{proc Dec} \\ \textbf{while true} \\ \textbf{do} \\ \textbf{if } \times > 0 \\ \textbf{then } \times := \times -1 \\ \textbf{fi} \\ \textbf{od} \end{array}
```

```
proc Reset
while true
do
if \times = 200
then \times := 0
fi
od
```

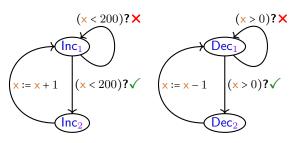
## Modeling (by labeled transition systems)

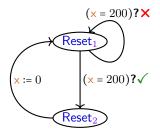
```
proc Inc
                       proc Dec
                                              proc Reset
 while true
                        while true
                                               while true
                          do
                                                do
  do
  if \times < 200 1: if \times > 0
                                          1: if \times = 200
 then x := x + 1 2: then x := x - 1
                                             2: then x := 0
    fi
                           fi
                                                  fi
  od
                          od
                                                od
```

# Modeling (by labeled transition systems)

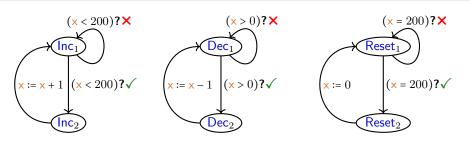
proc Inc

```
proc Dec
    while true
    do
1:         if x > 0
2:         then x := x - 1
         fi
         od
```





Labeled transition systems (LTSs)



$$Inc_1 \parallel Dec_1 \parallel Reset_1 \stackrel{?}{\vDash} \square (0 \le x \land x \le 200)$$
 (Linear-TL formula)

otivation mc-schema example <mark>counterex-trace</mark> logical properties topics book module organization

## Counterexample (offending execution trace)

 $(x = 199; Inc_1 || Dec_1 || Reset_1)$ 

otivation mc-schema example <mark>counterex-trace</mark> logical properties topics book module organization

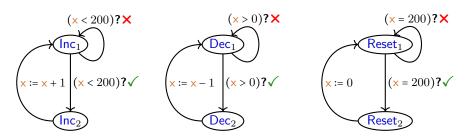
## Counterexample (offending execution trace)

 $(x = 199; | lnc_1 | | lnc_1 | | Reset_1)$ 

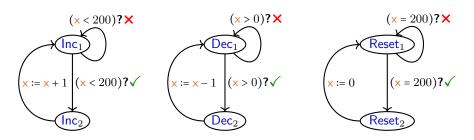
$$\begin{array}{c|c} \left( \left\langle x = 199 \; ; \; \mathsf{Inc}_1 \parallel \mathsf{Dec}_1 \parallel \mathsf{Reset}_1 \right\rangle \right) \\ & \qquad \qquad \downarrow (\mathsf{x} < 200) ? \checkmark \\ \hline \left( \left\langle x = 199 \; ; \; \mathsf{Inc}_2 \parallel \mathsf{Dec}_1 \parallel \mathsf{Reset}_1 \right\rangle \right) \end{array}$$

$$\begin{array}{c|c} & \left( x = 199 \; ; \; \operatorname{Inc_1} \parallel \operatorname{Dec_1} \parallel \operatorname{Reset_1} \right) \\ & & \downarrow (\mathsf{x} < 200) ? \checkmark \\ \hline & \left( x = 199 \; ; \; \operatorname{Inc_2} \parallel \operatorname{Dec_1} \parallel \operatorname{Reset_1} \right) \\ & & \downarrow \mathsf{x} := \mathsf{x} + 1 \\ \hline & \left( x = 200 \; ; \; \operatorname{Inc_1} \parallel \operatorname{Dec_1} \parallel \operatorname{Reset_1} \right) \\ & & \downarrow (\mathsf{x} > 0) ? \checkmark \\ \hline & \left( x = 200 \; ; \; \operatorname{Inc_1} \parallel \operatorname{Dec_2} \parallel \operatorname{Reset_1} \right) \\ & & \downarrow (\mathsf{x} = 200) ? \checkmark \\ \hline & \left( x = 200 \; ; \; \operatorname{Inc_1} \parallel \operatorname{Dec_2} \parallel \operatorname{Reset_2} \right) \\ \hline \end{array}$$

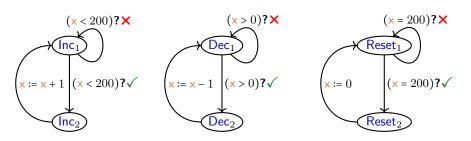
$$\begin{array}{c|c} & \left( x = 199 \; ; \; \operatorname{Inc_1} \parallel \operatorname{Dec_1} \parallel \operatorname{Reset_1} \right) \\ & & \downarrow (\mathsf{x} < 200) ? \checkmark \\ \hline & \left( x = 199 \; ; \; \operatorname{Inc_2} \parallel \operatorname{Dec_1} \parallel \operatorname{Reset_1} \right) \\ & & \downarrow \mathsf{x} := \mathsf{x} + 1 \\ \hline & \left( x = 200 \; ; \; \operatorname{Inc_1} \parallel \operatorname{Dec_1} \parallel \operatorname{Reset_1} \right) \\ & & \downarrow (\mathsf{x} > 0) ? \checkmark \\ \hline & \left( x = 200 \; ; \; \operatorname{Inc_1} \parallel \operatorname{Dec_2} \parallel \operatorname{Reset_1} \right) \\ & & \downarrow (\mathsf{x} = 200) ? \checkmark \\ \hline & \left( x = 200 \; ; \; \operatorname{Inc_1} \parallel \operatorname{Dec_2} \parallel \operatorname{Reset_2} \right) \\ \hline \end{array}$$



$$Inc_1 \parallel Dec_1 \parallel Reset_1 \not\models \Box (0 \le x \land x \le 200)$$
 (Linear-TL formula)



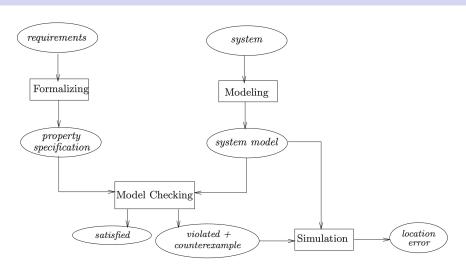
$$\begin{aligned} & \mathsf{Inc}_1 \parallel \mathsf{Dec}_1 \parallel \mathsf{Reset}_1 \quad \not\models \quad \Box \big( 0 \leq x \, \land \, x \leq 200 \big) & \quad \big( \mathsf{Linear}\text{-}\mathsf{TL} \; \mathsf{formula} \big) \\ & \mathsf{Inc}_1 \parallel \mathsf{Dec}_1 \parallel \mathsf{Reset}_1 \; \models \; \diamondsuit (x < 0) & \quad \big( \mathsf{LTL} \; \mathsf{formula} \big) \end{aligned}$$



```
\begin{aligned} & \operatorname{Inc}_1 \parallel \operatorname{Dec}_1 \parallel \operatorname{Reset}_1 \;\; \not\models \;\; \Box (0 \leq x \, \land \, x \leq 200) & \text{(Linear-TL formula)} \\ & \operatorname{Inc}_1 \parallel \operatorname{Dec}_1 \parallel \operatorname{Reset}_1 \;\; \models \;\; \diamondsuit (x < 0) & \text{(LTL formula)} \\ & \operatorname{Inc}_1 \parallel \operatorname{Dec}_1 \parallel \operatorname{Reset}_1 \;\; \not\models \;\; \forall \Box (0 \leq x \, \land \, x \leq 200) & \text{(Computation-Tree-L formula)} \\ & \operatorname{Inc}_1 \parallel \operatorname{Dec}_1 \parallel \operatorname{Reset}_1 \;\; \models \;\; \exists \Box (0 \leq x \, \land \, x \leq 200) & \text{(CTL formula)} \\ & \operatorname{Inc}_1 \parallel \operatorname{Dec}_1 \parallel \operatorname{Reset}_1 \;\; \models \;\; \forall \Box \exists \diamondsuit (x < 0) & \text{(CTL formula)} \end{aligned}
```

otivation mc-schema example counterex-trace logical properties topics book module organization

# Model checking



Any [such] verification is only as good as the model of the system.

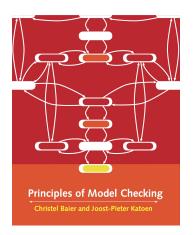
tivation mc-schema example counterex-trace logical properties <mark>topics</mark> book module organizatio

# Topics of the module

- modeling systems by labeled transition systems (LTSs)
- fairness
- Linear Temporal Logic (LTL)
  - model checking formulas
    - express properties by Büchi automata
    - model check LTSs and properties via product automata
- Computation Tree Logic (CTL)
- partial model checking
  - partially known systems (state properties/states/transitions)
- ▶ analysing system behavior with the mCRL2 model-checker toolbox

tivation mc-schema example counterex-trace logical properties topics <mark>book</mark> module organization

#### Book



pdf available:

https://is.ifmo.ru/books/\_principles\_of\_model\_checking.pdf

tivation mc-schema example counterex-trace logical properties topics book <mark>module organizati</mark>c

# Organization

Lectures (Clemens 5/Emilio 2)

- presentations on blackboard
- notes after the lecture (notes 2024/25 available)
- February (first/second week)

tivation mc-schema example counterex-trace logical properties topics book module organization

# Organization

#### Lectures (Clemens 5/Emilio 2)

- presentations on blackboard
- notes after the lecture (notes 2024/25 available)
- February (first/second week)

#### Exam

- options:
  - ▶ small verification project (of an algorithm, e.g. in mCRL2)
  - presentation about a paper
  - written exam?

otivation mc-schema example counterex-trace logical properties topics book <mark>module organizat</mark>i

# Organization

#### Lectures (Clemens 5/Emilio 2)

- presentations on blackboard
- notes after the lecture (notes 2024/25 available)
- February (first/second week)

#### Exam

- options:
  - ▶ small verification project (of an algorithm, e.g. in mCRL2)
  - presentation about a paper
  - written exam?

Thank you – we are looking forward to the course!