

# Introduction to Model Checking

## (Preview of Core Module)

Clemens Grabmayer

<https://clegra.github.io>

Emilio Tuosto

<https://cs.gssi.it/emilio.tuosto/>

Department of Computer Science



December 2, 2024

# Model Checking

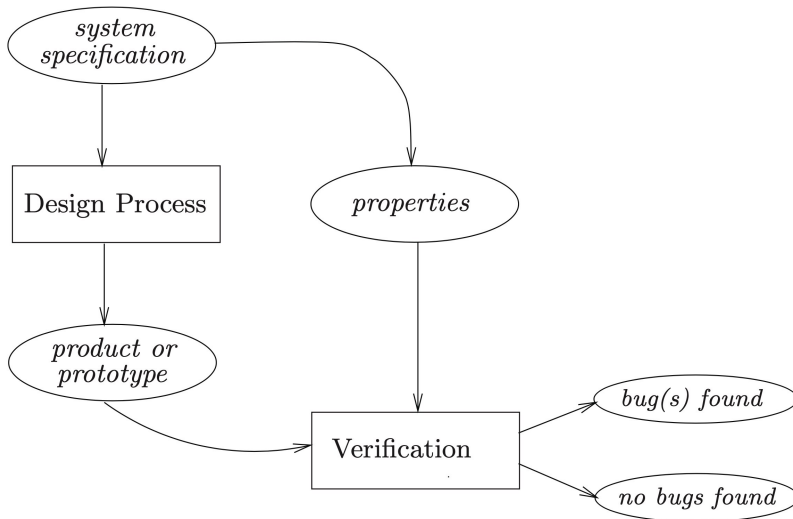
... is an effective automatable technique:

- ▶ *to expose potential software design errors;*
- ▶ *that, given a finite-state model of a system and a formal property, systematically checks whether this property holds for that model.*

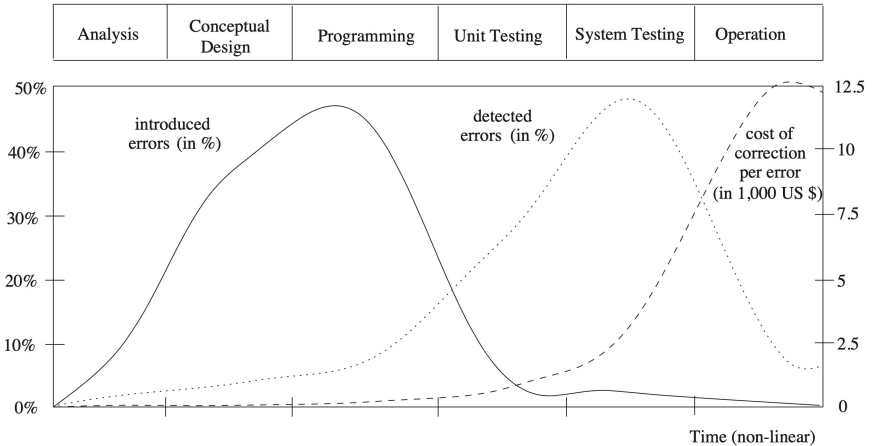
Strengths:

- ▶ widely applicable  
(embedded systems, software engineering, hardware design)
- ▶ supports **partial verification** (of system parts)
- ▶ provides **diagnostic information** for debugging
- ▶ has sound **mathematical underpinning** (logic and process theory)

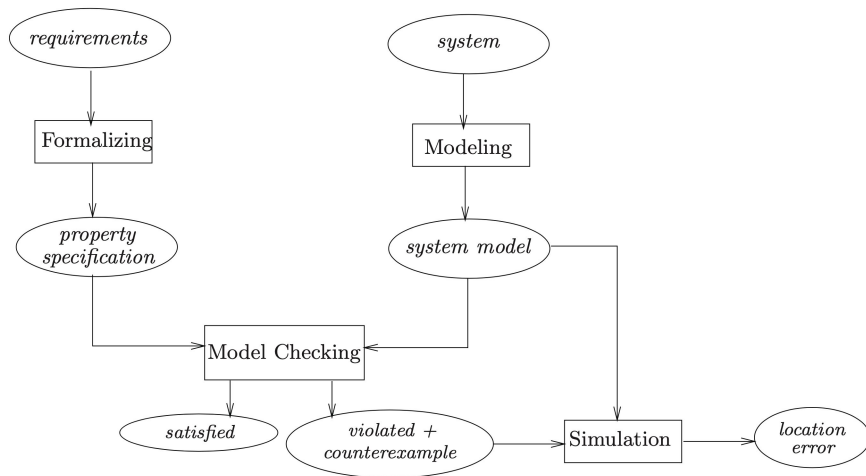
# Hard-/Software Verification (traditionally)



# Error introduction, detection, and repair costs



# Model checking



# Example: concurrency and non-determinism

Programs [Inc](#), [Dec](#), and [Reset](#) cooperate, and use a shared variable  $x$ :

```
proc Inc
  while true
  do
    if  $x < 200$ 
    then  $x := x + 1$ 
    fi
  od
```

```
proc Dec
  while true
  do
    if  $x > 0$ 
    then  $x := x - 1$ 
    fi
  od
```

```
proc Reset
  while true
  do
    if  $x = 200$ 
    then  $x := 0$ 
    fi
  od
```

# Example: concurrency and non-determinism

Programs [Inc](#), [Dec](#), and [Reset](#) cooperate, and use a shared variable  $x$ :

```
proc Inc
  while true
    do
      if  $x < 200$ 
        then  $x := x + 1$ 
      fi
    od
```

```
proc Dec
  while true
    do
      if  $x > 0$ 
        then  $x := x - 1$ 
      fi
    od
```

```
proc Reset
  while true
    do
      if  $x = 200$ 
        then  $x := 0$ 
      fi
    od
```

Question: Is  $0 \leq x \leq 200$  always guaranteed?

# Modeling (by labeled transition systems)

**proc** [Inc](#)

**while** true

**do**

**if**  $x < 200$

**then**  $x := x + 1$

**fi**

**od**

**proc** [Dec](#)

**while** true

**do**

**if**  $x > 0$

**then**  $x := x - 1$

**fi**

**od**

**proc** [Reset](#)

**while** true

**do**

**if**  $x = 200$

**then**  $x := 0$

**fi**

**od**



# Modeling (by labeled transition systems)

**proc** Inc

**while** true

**do**

1: **if**  $x < 200$

2: **then**  $x := x + 1$

**fi**

**od**

**proc** Dec

**while** true

**do**

1: **if**  $x > 0$

2: **then**  $x := x - 1$

**fi**

**od**

**proc** Reset

**while** true

**do**

1: **if**  $x = 200$

2: **then**  $x := 0$

**fi**

**od**

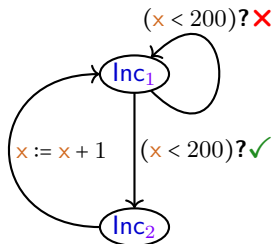
# Modeling (by labeled transition systems)

**proc Inc**

**while true**

**do**

1: **if**  $x < 200$   
 2: **then**  $x := x + 1$   
**fi**  
**od**

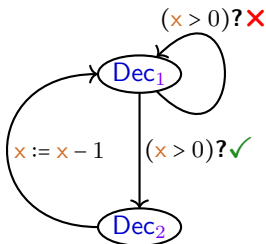


**proc Dec**

**while true**

**do**

1: **if**  $x > 0$   
 2: **then**  $x := x - 1$   
**fi**  
**od**

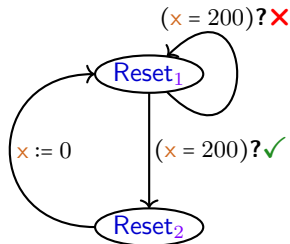


**proc Reset**

**while true**

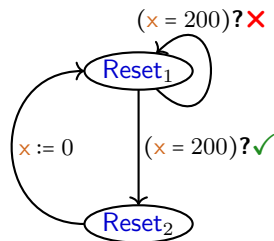
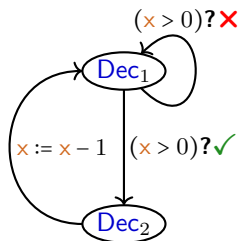
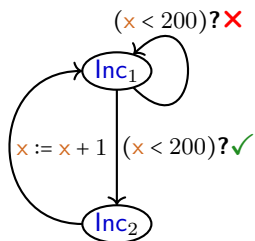
**do**

1: **if**  $x = 200$   
 2: **then**  $x := 0$   
**fi**  
**od**



Labeled transition systems (LTSs)

# Formalizing properties (in temporal logic)



$$Inc_1 \parallel Dec_1 \parallel Reset_1 \stackrel{?}{\models} \Box(0 \leq x \wedge x \leq 200) \quad (\text{Linear-TL formula})$$

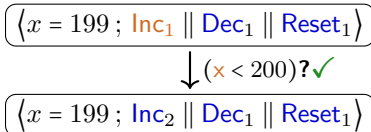
# Counterexample (offending execution trace)

$$\langle x = 199 ; \text{Inc}_1 \parallel \text{Dec}_1 \parallel \text{Reset}_1 \rangle$$

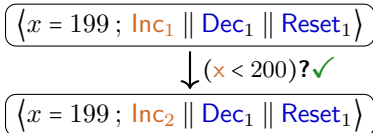
# Counterexample (offending execution trace)

$$\langle x = 199 ; \text{Inc}_1 \parallel \text{Dec}_1 \parallel \text{Reset}_1 \rangle$$

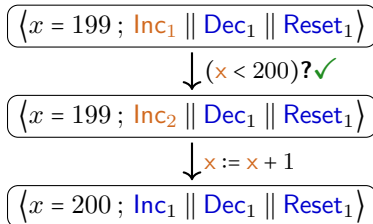
# Counterexample (offending execution trace)



# Counterexample (offending execution trace)

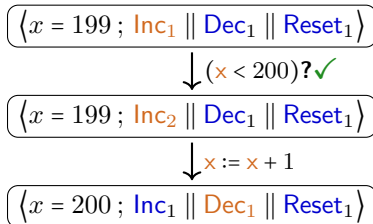


# Counterexample (offending execution trace)

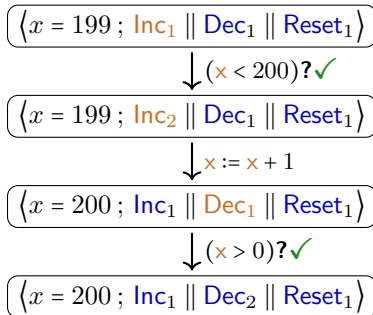




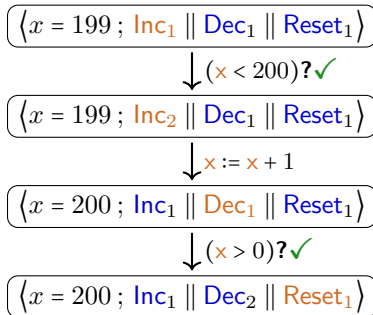
# Counterexample (offending execution trace)



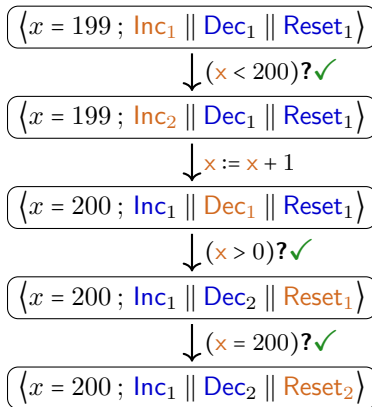
# Counterexample (offending execution trace)



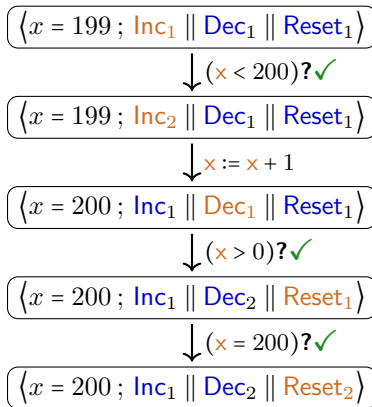
# Counterexample (offending execution trace)



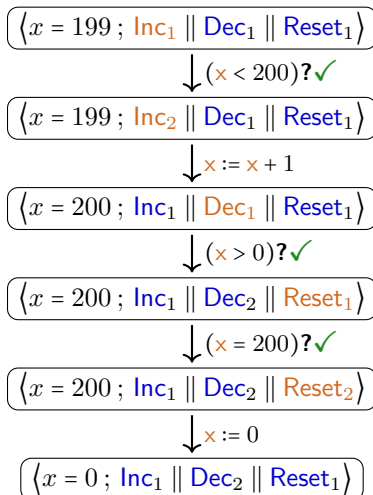
# Counterexample (offending execution trace)



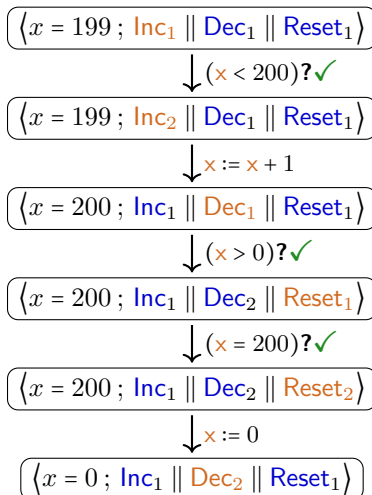
# Counterexample (offending execution trace)



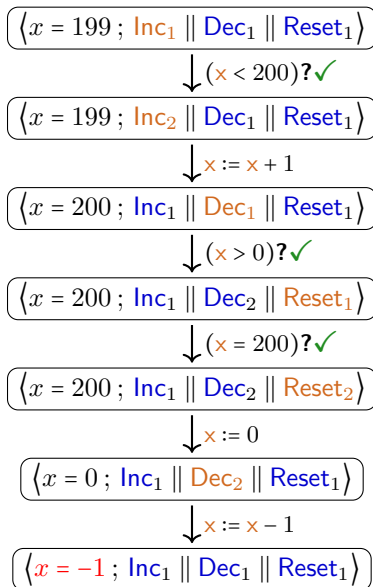
# Counterexample (offending execution trace)



# Counterexample (offending execution trace)

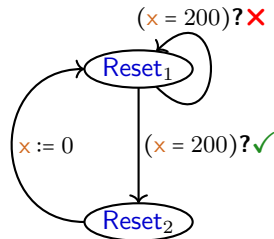
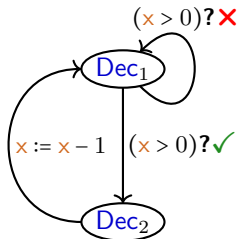
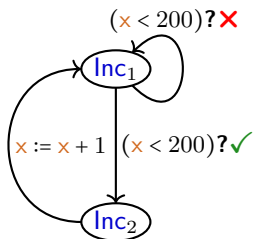


# Counterexample (offending execution trace)



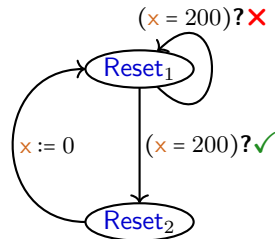
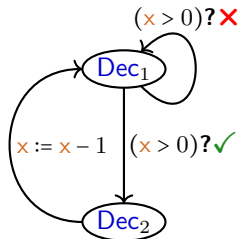
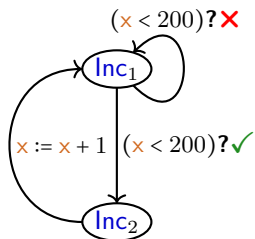


# Formalizing properties (in temporal logic)



$\text{Inc}_1 \parallel \text{Dec}_1 \parallel \text{Reset}_1 \not\models \Box(0 \leq x \wedge x \leq 200)$  (Linear-TL formula)

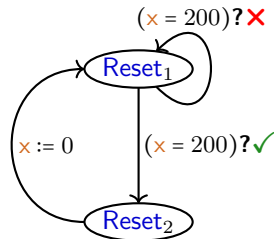
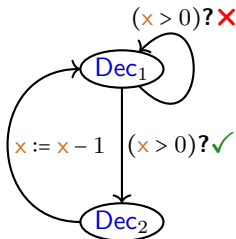
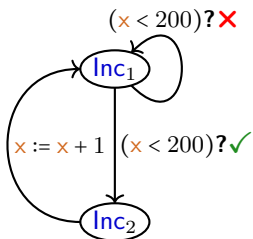
# Formalizing properties (in temporal logic)



$Inc_1 \parallel Dec_1 \parallel Reset_1 \not\models \Box(0 \leq x \wedge x \leq 200)$  (Linear-TL formula)

$Inc_1 \parallel Dec_1 \parallel Reset_1 \models \Diamond(x < 0)$  (LTL formula)

# Formalizing properties (in temporal logic)



$\text{Inc}_1 \parallel \text{Dec}_1 \parallel \text{Reset}_1 \not\models \Box(0 \leq x \wedge x \leq 200)$  (Linear-TL formula)

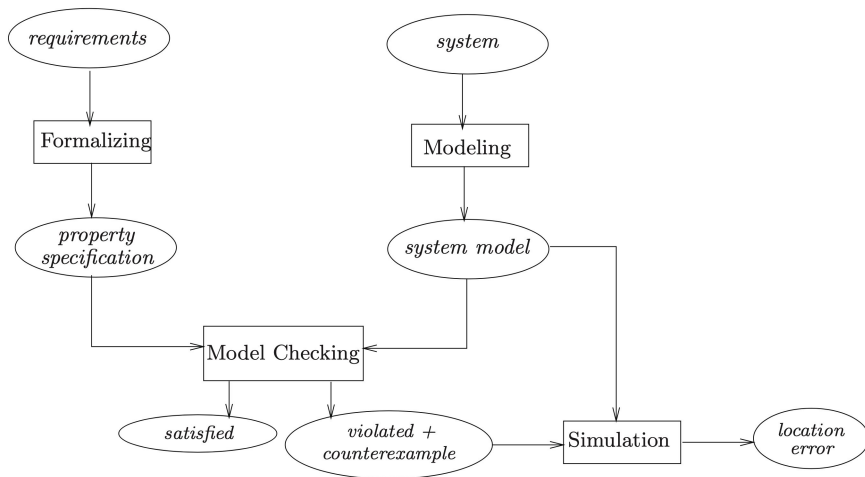
$\text{Inc}_1 \parallel \text{Dec}_1 \parallel \text{Reset}_1 \models \Diamond(x < 0)$  (LTL formula)

$\text{Inc}_1 \parallel \text{Dec}_1 \parallel \text{Reset}_1 \not\models \forall \Box(0 \leq x \wedge x \leq 200)$  (Computation-Tree-L formula)

$\text{Inc}_1 \parallel \text{Dec}_1 \parallel \text{Reset}_1 \models \exists \Box(0 \leq x \wedge x \leq 200)$  (CTL formula)

$\text{Inc}_1 \parallel \text{Dec}_1 \parallel \text{Reset}_1 \models \forall \Box \exists \Diamond(x < 0)$  (CTL formula)

# Model checking

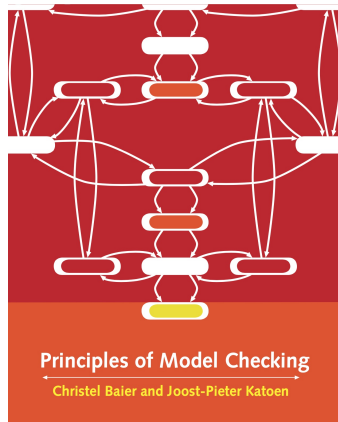


*Any [such] verification is only as good as the model of the system.*

# Topics of the module

- ▶ modeling systems by **labeled transition systems (LTSs)**
- ▶ **fairness**
- ▶ **Linear Temporal Logic (LTL)**
  - ▶ model checking formulas
    - ▶ express properties by **Büchi automata**
    - ▶ model check LTSs and properties via **product automata**
- ▶ **Computation Tree Logic (CTL)**
- ▶ **partial** model checking
  - ▶ partially known systems (state properties/states/transitions)
- ▶ analysing system behavior with the **mCRL2 model-checker toolbox**

# Book



- pdf available:

[https://is.ifmo.ru/books/\\_principles\\_of\\_model\\_checking.pdf](https://is.ifmo.ru/books/_principles_of_model_checking.pdf)

# Organization

## Lectures (Clemens 5/Emilio 2)

- ▶ presentations on blackboard
- ▶ notes after the lecture (notes 2024/25 available)
- ▶ February (first/second week)

# Organization

## Lectures (Clemens 5/Emilio 2)

- ▶ presentations on blackboard
- ▶ notes after the lecture (notes 2024/25 available)
- ▶ February (first/second week)

## Exam

- ▶ options:
  - ▶ small verification project (of an algorithm, e.g. in [mCRL2](#))
  - ▶ presentation about a paper
  - ▶ written exam?



# Organization

## Lectures (Clemens 5/Emilio 2)

- ▶ presentations on blackboard
- ▶ notes after the lecture (notes 2024/25 available)
- ▶ February (first/second week)

## Exam

- ▶ options:
  - ▶ small verification project (of an algorithm, e.g. in [mCRL2](#))
  - ▶ presentation about a paper
  - ▶ written exam?

Thank you – we are looking forward to the course!