

RTX Software Design Report

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Chapter 1

Introduction

The purpose of this report is to outline the design of the RTX written by the group members, Clement Hoang, David Su, Peter Li, and Cole Vander Veen, as part of the SE350 course at the University of Waterloo. The OS is designed for a Keil MCB1700 Cortex-M3 board, with a LPC1768 microcontroller.

It is aimed to provide documentation for the operating system, in order to facilitate the use and understanding for anyone interested in programming for the OS. As such, this report outlines the global variables used in the OS, and then moves on to describing the kernel API in a modular and chronological way, from when we implemented it. Finally, the report closes with some analysis on the OS, and challenges that the group faced for the duration of the lab.

Chapter 2

Design Description

2.1 Global Variables and Data Structures

- `memQueue`: A data structure that models the free physical memory in the OS, by splitting the heap into blocks of equal size. It is represented by a `MemQueue` data structure, which is a linked list of `MemBlock` nodes of size `BLOCK_SIZE`. It is used by the kernel API when releasing and requesting memory, by popping a block when it is used by a process, and pushing it back in when it is released.
 - `MemBlock`: To expand, the `MemBlock` is a C-struct that holds a pointer to the next `MemBlock` in the queue. It also has reserved space in the front in case the block needs to hold an envelope.
- `gp_pcb`s: A pointer to an array of PCB structs. It holds the state of all the process control blocks that are in the OS, and is interacted with by functions that change and read PCB states. For example, setting the process priority or getting the process priority uses `gp_pcb`s to access the priority of a specific PCB.
 - PCB: a model of a process and its state. The PCB contains the following fields:
 - * `mp_sp`: stack pointer of the process
 - * `m_pid`: ID of the process
 - * `m_priority`: priority of the process
 - * `m_state`: state of the process
 - * `nextPCB`: pointer to the next PCB, if it is in a queue
 - * `msgHead`: beginning of the message queue
 - * `msgTail`: end of the message queue
- `gp_stack`:
- `p_end`
- `numOfBlocks`

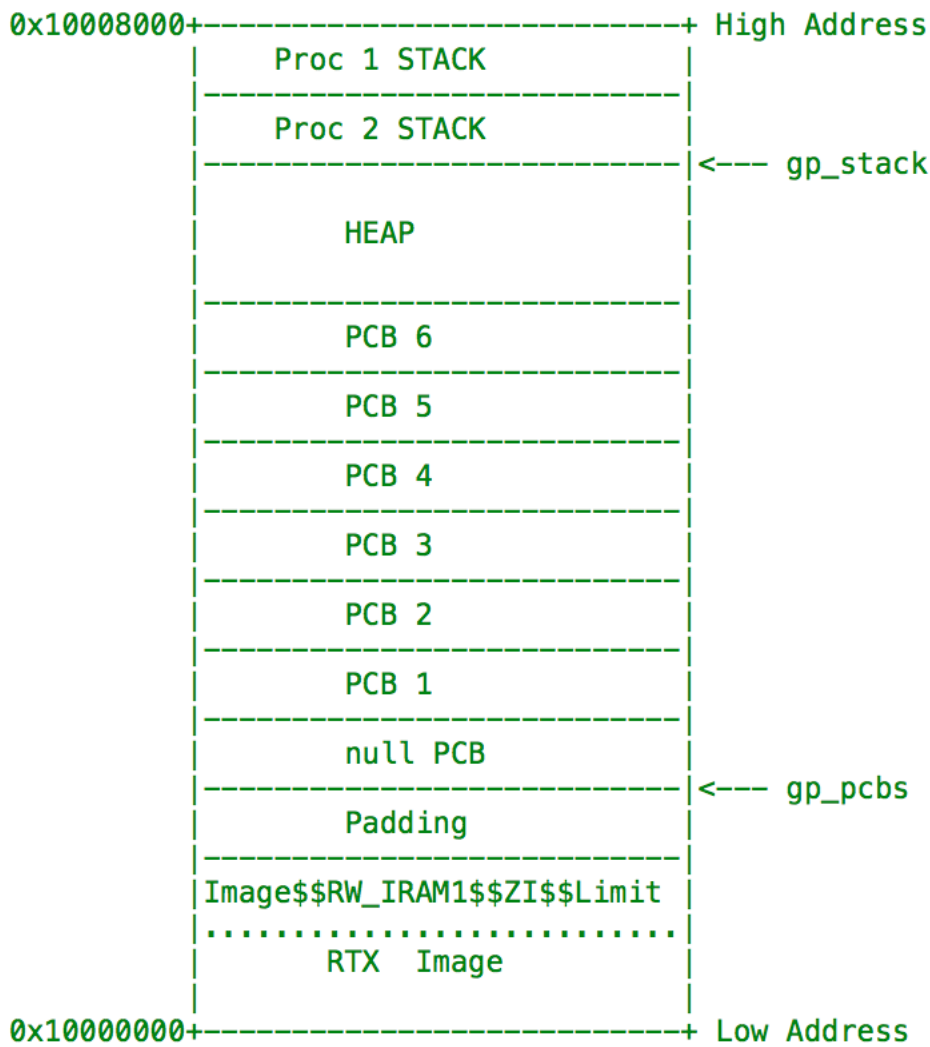


Figure 2.1: Memory Layout

2.2 Memory Management

2.2.1 Memory Structure

dsfdasfdsafdsafdsafdsaf

2.2.2 Requesting Memory Blocks

```
int k_request_memory_block(void);
```

describe input, output, effects

2.2.3 Releasing Memory Blocks

```
int k_release_memory_block(void* memory_block);
```

describe input, output, effects

Algorithm 1 k.request_memory_block

```
1: procedure REQUEST_MEMORY_BLOCK
2:   while heap is full do
3:     block the current process
4:   end while
5:   update the free space list
6:   return the address of the top of the block
7: end procedure
```

Algorithm 2 The memory release function

```
1: procedure RELEASE_MEMORY_BLOCK(*memory_block)
2:   if this block is the top block of the heap then
3:     modify heap header node (never gets overwritten)
4:   end if
5:   if there is free space immediately beneath this block then
6:     combine them by increasing this block's length
7:   else this block becomes a new block node, is added to the list
8:   end if
9:   if there is free space immediately beneath this block then
10:    combine them by increasing this block's length
11:  end if
12:  if a process is blocked on memory then
13:    unblock that process, release the processor
14:  end if
15: end procedure
```

2.3 Processor Management

2.3.1 Process Control Structures

DFASFAFD

2.3.2 Process Queues

fsadfasdfadsf

2.3.3 Process Scheduling

sdfasdfasfdasf

2.4 Process Priority Management

2.4.1 Get Process Priority

asdfadsfasf

2.4.2 Set Process Priority

dsfasdfasfsfdf

2.5 Interprocess Communication

2.5.1 Message Structure

dsfadsfadsfdasfdafs

2.5.2 Sending Messages

adsfdsafasdfasf

2.5.3 Receiving Messages

dsfafasfdasf

2.5.4 Delayed Send

sdfasfasfd

2.6 Interrupts and I-Processes

2.6.1 UART I-Process

The UART interrupt is enabled to send output to a display, and to receive input from the user. The OS handles those interrupts by registering a `UART0_IRQHandler` function that will be called whenever a UART interrupt occurs.

The file `uart_irq.c` is initialized by calling the function `uart_irq_init`. This function initializes the UART interrupts by setting the appropriate flags and choosing the correct UART port. Below is the interrupt handling pseudocode, starting with `UART0_IRQHandler`:

```
1: function UART0_IRQHANDLER
2:   push registers onto stack
3:   call c_UART0_IRQHandler_wrapper()
4:   pop registers off stack
5: end function
1: function c_UART0_IRQHANDLER_WRAPPER
2:   call c_UART0_IRQHandler()
3:   if there is another ready process with higher priority than the current process then
4:     call k_release_processor()
5:   end if
6: end function
1: function c_UART0_IRQHANDLER
2:   if receive data available then
3:     g_char_in = newly received char
4:     if g_char_in == null character then
5:       return
6:     end if
7:     echoMsg = get memory block (non blocking)
8:     if echoMsg is not null then
9:       echoMsg->mtype = ECHO
10:      if g_char_in == '\r' then
11:        echoMsg->mtext = "\n\r\0"
12:      else
13:        echoMsg->mtext = g_char_in + '\0'
14:      end if
15:      send echoMsg (no preemption) to KCD
16:    end if
17:    if cur_msg is null then
18:      cur_msg = get memory block (non blocking)
19:      if cur_msg is null (no more memory) then
20:        return
21:      end if
22:      msg_str_index = 0
23:    end if
24:    if g_char_in == '\r' or message about to overflow memory block then
25:      cur_msg->mtext += '\0'
26:      cur_msg->mtype = DEFAULT
```

```

27:         send cur_msg (no preemption) to KCD
28:         reset cur_msg, msg_str_index
29:     else
30:         cur_msg->mtext += g_char_in
31:     end if
32: else if transmit interrupt enabled then
33:     if *gp_buffer != null character then
34:         send *gp_buffer
35:         gp_buffer = next location in circular buffer
36:     else
37:         disable transmit interrupt
38:         send null character
39:         reset gp_buffer and g_buffer_end
40:     end if
41: end if
42: end function

```

If there is an incoming character, immediately forward it to the KCD (assuming there is free memory) to echo back to the user. Also, add that character to a buffer. Once a newline is encountered, send the entire buffer to the KCD to decode

If there are still messages in the transmit buffer, send the next character. If the character is a null character, disable the transmit interrupt as the transmission is finished.

Also, there is an `enable_UART_transmit()` function that sets the appropriate flags to enable the interrupt for outputting characters. This function is called by the CRT to begin outputting characters.

2.6.2 Timer I-Process

sdfasfdafd

2.7 System Processes

2.7.1 Null Process

sdfdasfadsf

2.7.2 CRT Process

sdfdsfafaf

2.8 User Processes

2.8.1 Wall Clock Process

sdfasdfafadf

2.8.2 Set Priority Process

dsfasdfasdfadsf

2.8.3 Stress Test Processes

dfdasdfasdfads

2.9 Initialization

dasfasfasfd

2.10 Testing

dfadsfasdf

2.11 Major Design Changes

dsfdafadsf

Chapter 3

Lessons Learned

3.1 Source Control and Code Management

sdfdsafsadf

Chapter 4

Team Dynamics and Individual Responsibilities

4.1 adsfadsf

dfasfasdf

Chapter 5

Timing Analysis