Assignment # 5

Homework

Homework problems are a preparation for the quizzes. They are *not* graded. Please use the mywpi forum to post questions you have on these problems.

• 8.1, 8.2, 8.3, 8.5, 8.8, 8.9, 8.12, 8.13, 8.14, 8.18

Project

Note: For submissions on mywpi: Please submit a single pdf file containing your results. Please submit source code as a separate file, but make sure to have it listed in the pdf as well.

- 1. 8.4 (if you use sage for part of the problem, please provide the code as well)
- 2. Implement the square and multiply algorithm using a computer language of your choice. The program should print all intermediate results. For sage and Python, you can use the following template:

```
def my_pow(b,e,m):
    """ Computes b^e mod m using the square and multiply algorithm"""
    x = pow(b,e,m) # remove this line and place your code here instead
return x
```

Compute the following exponentiations $a^e \mod p$ using your program:

- (a) a = 235973, e = 456789, p = 583903
- (b) a = 984327457683, e = 2153489582, p = 994348472629

Print the output of your program (including the intermediate steps) and turn it in together with your source code.

- 3. The goal of this problem is to implement the Diffie Hellman Key Exchange.
 - (a) Choose a pseudo random generator and initialize it with a *seed* = 12345 (e.g. set_random_seed() in sage). You can use a PRNG of your choice. Please comment on the security of your scheme given that you will use the same PRNG to generate the secret DHKE parameters A, B.
 - (b) Next, we have to generate a large prime as modulus for our cyclic group. Please use a tool of your choice to generate a prime p of exactly 1024 bits. Does this prime have to be random? Justify your answer.

(c) In order to increase the security of the scheme we use a safe prime. A safe prime p_s is of the form $p_s = 2 \cdot p + 1$, where p is also prime. The group we will use for DHKE in this problem will be multiplicative group $\mathbb{Z}_{p_s}^*$. One way of choosing the DHKE parameter g is choosing it as a random element from $\mathbb{Z}_{p_s}^*$. Explain why using a safe prime helps ensuring that finding an element g of high order g is easy. Make your DHKE program choose g at random from $\mathbb{Z}_{p_s}^*$. The safe prime used should also have 1024 bits. If you have trouble generating one yourself, please use: $p_s =$

 $136493091133649836164289363338682002944029379014077033130604363922256595 \\ 037775025761962540974136000872788770954019530710825866837156972143526820 \\ 463536654299014569407235702375016873961943932744861043443226534544334946 \\ 757053788912168896006231429590913701778250440452439749688803485407941150 \\ 333770430825062090319$

- (d) Next, generate two random parameters A and B. Write your code so that A, B can only take valid parameters.
- (e) Since the *Decisional Diffie Hellman* problem is easy to solve for the given group, please apply SHA-224 to the computed $k_A = k_B$ to get a symmetric-sized key from the scheme. Hash functions are a convenient way of turning a random group element into a random bit sequence. Please use big endian conversion to represent your integer as SHA-224 input (e.g. $hex(k_A).decode('hex')$ in Python 2 or sage).
- 4. The ElGamal encryption scheme is non-deterministic: A given message m has many valid encryptions.
 - (a) Why is the ElGamal encryption scheme non-deterministic?
 - (b) How many valid ciphertexts exist for each message m (general expression)? How many are there for the system in problem 8.13 from the book (numerical answer)?
 - (c) Consider the case that for two messages $m_1 \neq m_2$ the same session parameter $y_1 = y_2$ has been chosen for the ElGamal encryption. This kind of behavior occurs if no or a bad cryptographic PRNG is being used. Show how the message m_2 can be recovered from a known message ciphertext pair $\langle m_1, c_1 \rangle$ if the same y is used.

Good Luck and Have Fun!