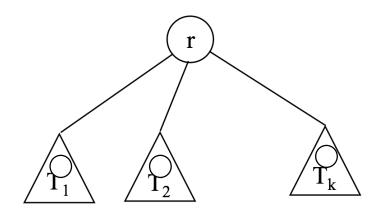
# Data Structure: Tree

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#### tree

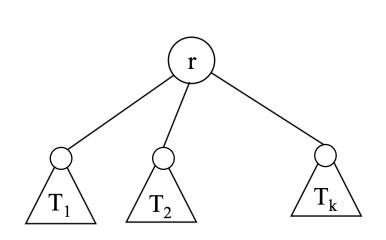
- a collection of nodes connected by edges without a cycle
- by recursive definition:
  - an empty tree or
  - a root r and subtrees T<sub>1</sub>, T<sub>2</sub>,..., T<sub>k</sub> (disjoint sets) each of whose roots are connected to r by an edge

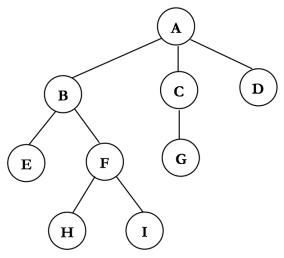


recursive definition of tree

#### tree

- Each root of  $T_1$ ,  $T_2$ ,...,  $T_k$  is a *child* of r, and r is the *parent* of each root.
- The roots of the subtrees are siblings of one another
- If there is an order among the Ti's, the tree is an ordered tree.
- The degree of a node is the number of children it has.
- The degree of a tree is the maximum degree of the nodes.
- A leaf is a node of degree 0.

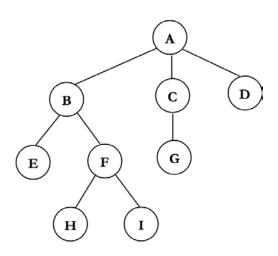




an example of tree

#### tree

- **path** between two nodes is a sequence of nodes  $n_1, n_2, ..., n_k$ , such that  $n_i$  is a parent of  $n_{i+1}$
- **length** of a path is the number of edges on the path (the path  $n_1$ ,  $n_2$ ,...  $n_k$ : length k-1)
- depth (level) of a node is the length of the (unique) path from the root to that node (root: level 0)
- height of a node is the length of the longest path from that node to a leaf (leaf: height 0)
- the height of a tree is the height of the root

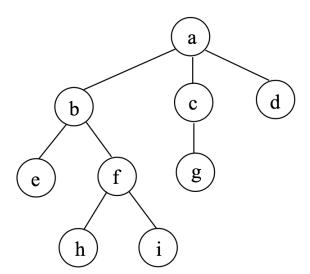


#### representation of tree

- for any node x, there exists exactly one path from the root to x?
- tree can be empty with no node?
- how many edges are in a tree with n nodes?

#### representation of tree

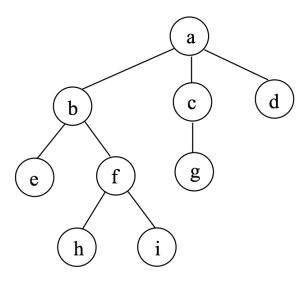
- how can we implement a tree?
  - linked list?
  - can we have pointers for the children nodes?
  - can we have fixed number of pointers to represent a tree?
    - for a tree of fixed number of degree?
    - else?



_			
data	link 1	link 2	 link n

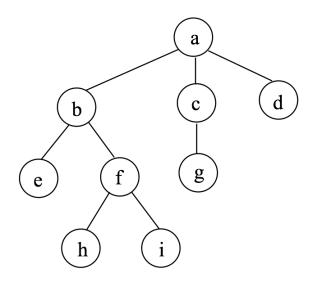
#### left child-right sibling representation

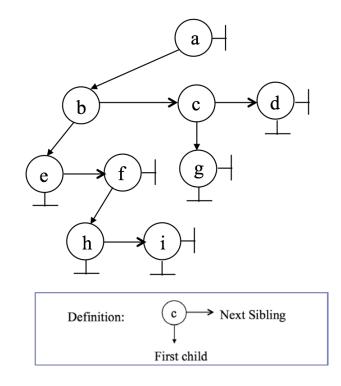
every node has at most one leftmost child and at most one closet right sibling



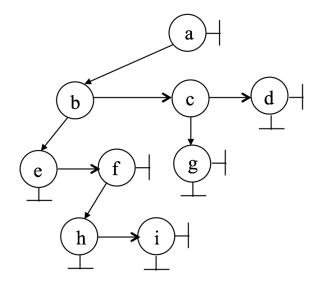
#### left child-right sibling representation

every node has at most one leftmost child and at most one closet right sibling





#### left child- right sibling representation



```
Definition: C Next Sibling
First child
```

```
struct TreeNode{
    ElementType Element;
    PtrToNode FirstChild;
    PtrToNode NextSibling;
    };
typedef struct TreeNode *PtrToNode;
```

- a finite set of nodes that is either
  - i) empty or
  - ii) a root node and two disjoint binary trees
- the tree on the left and the tree on the right are different



■ the maximum number of nodes on level i of a binary tree is 2<sup>i</sup>, i>=0

#### the proof by induction

- base: for the root at level i=0,  $2^0 = 1$
- induction hypothesis: assume that the maximum number of nodes on level i-1 > 0,  $2^{i-1}$
- induction step: on level i,

 $2 * (the maximum number of nodes on level i-1) = 2 * <math>2^{i-1} = 2^i$ 

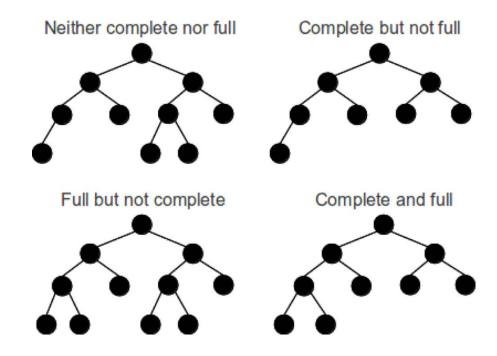
■ the maximum number of nodes in a binary tree of depth k is 2<sup>k+1</sup>-1, k>=0

■ For any nonempty binary tree T, if  $n_0$  is the number of leaf nodes, and  $n_2$  is the number of nodes of degree 2, then  $n_0 = n_2 + 1$ 

 $n = n_0 + n_1 + n_2$ ,  $n_i$  is the number of nodes with i degree  $n_i$  is the number of nodes in the tree

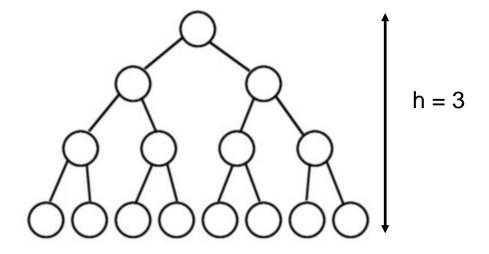
 $n = B + 1 = n_1 + 2n_2 + 1$ , B is the number of branches (edge)

- full binary tree is a binary tree in which every node has 0 or 2 children
- complete binary tree is a binary tree in which every level, except the last, is completely filled and the last level has all its nodes to the left side

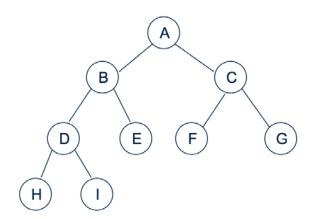


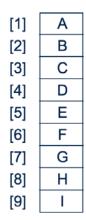
- perfect binary tree of height h is a binary tree of height h having 2<sup>h+1</sup> 1 nodes, (h >=0)
- the max number of nodes in the complete binary tree (height h) is 2 h+1 -1

$$2^{0} + 2^{1} + ... + 2^{h} = (2^{h+1} - 1)/(2 - 1) = 2^{h+1} - 1$$



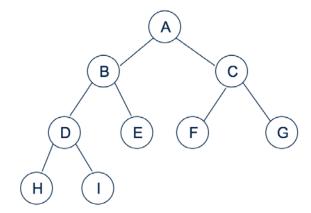
# binary tree: array representation





# binary tree: array representation

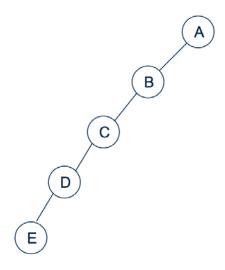
- if a complete binary tree with n nodes (i is the index) is represented sequentially,
  - leftChild(i) is at 2i for 2i <=n
  - rightChild(i) is at 2i + 1 for 2i + 1 <= n</p>
  - parent(i) is at Li/2 for i >1



[1]	Α
[2]	В
[3]	С
[4]	D
[5]	Е
[6]	F
[7]	G
[8]	Н
[9]	- 1

## binary tree: array representation

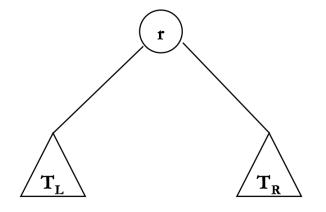
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[1]	Α		
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[5]	-		
[6]	-		
[7]	ı		
[8]	D -		
[9]			
[16]	Е		

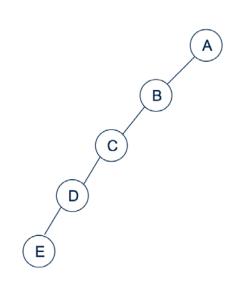
## binary tree: linked list representation

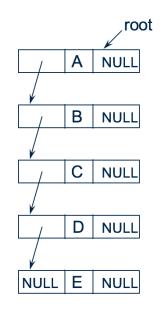
 a tree in which each node has no more than 2 children (left subtree and right subtree)

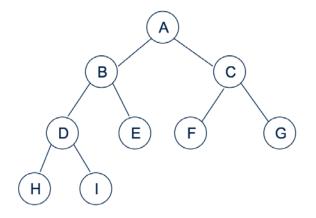


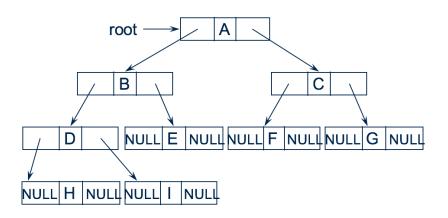
Left Element Right
--------------------

## binary tree: linked list representation



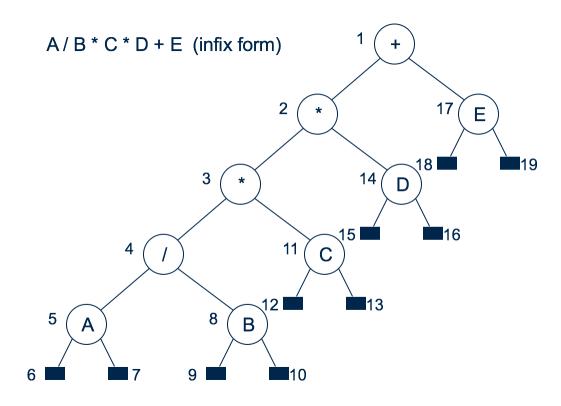






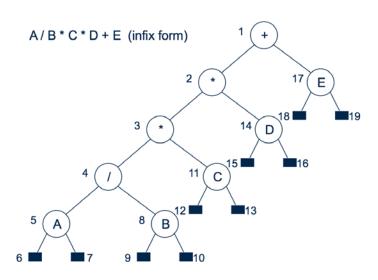
## application of binary tree

Expression Tree: intermediate representation for expressions used by the compiler



#### inorder traversal

```
void inorder(Tree ptr) {
    if(ptr) {
        inorder(ptr->left_child);
        printf("%d", ptr->data);
        inorder(ptr->right_child);
    }
}
```

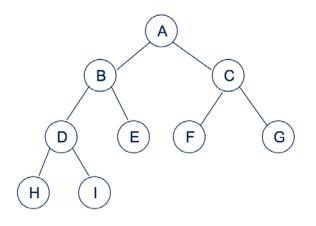


call	of	value in	action	call of	value in	action
inor	der	root		inorder	root	
	1	+		11	C	
	2	*		12	NULL	
	3	*		11	C	printf
	4	/		13	NULL	
Ι,	_5	A		2	*	printf
	6	NULL		14	D	
	5	Α	printf	15	NULL	
	7	NULL		14	D	printf
'	4	/	printf	16	NULL	
r	-8	В		1	+	printf
	9	NULL		17	$\mathbf{E}$	
	8	В	printf	18	NULL	
	10	NULL		17	$\mathbf{E}$	printf
	3	*	printf	19	NULL	

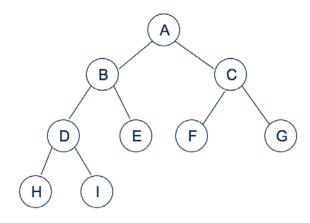
```
void preorder(Tree ptr) {
   if(ptr) {
       printf("%d", ptr->data);
       preorder(ptr->left_child);
       preorder(ptr->right_child);
void postorder(Tree ptr) {
   if(ptr) {
       postorder(ptr->left_child);
       postorder(ptr->right_child);
       printf("%d", ptr->data);
```

iterative in-order traversal using stack

```
void iterInorder (Tree node) {
           int top = -1
           Tree stack[MAX_SIZE];
           for (; ;) {
                      for (; node; node = node -> leftChild)
                                  push(node);
                      node = pop();
                                     // pop parent
                      if (!node) break;
                      printf("%d", node -> data);
                      node = node -> rightChild;
```



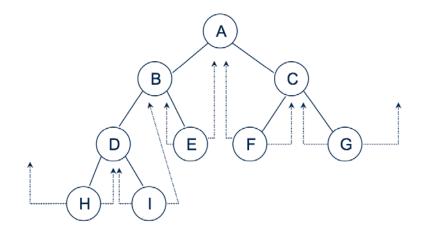
level-order traversal



level-order traversal

```
void levelOrder (Tree ptr) {
              int front = 0;
              int rear = -1;
               Tree queue[MAX];
              if (!ptr) return;
               addq(ptr);
              for (; ; ) {
                             ptr = deleteq();
                             if (ptr) {
                                            printf("%d", ptr->data);
                                            if (ptr -> leftChild)
                                                           addq(ptr -> leftChild);
                                            if (ptr -> rightChild)
                                                           addq(ptr -> rightChild);
                             else break;
```

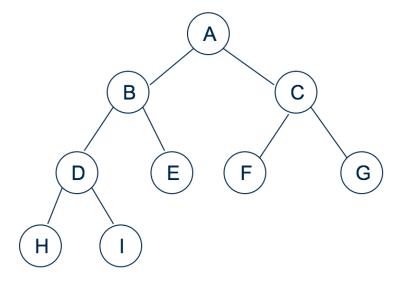
- there are n+1 null links out of 2n total links
- replace the null links by pointers, called threads to other nodes in the tree
  - if ptr -> leftChild is null, replace the null with a pointer to the node that would be visited before ptr in an in-order traversal
  - if ptr -> rightChild is null, replace the null with a pointer to the node that would be visited after ptr in an in-order traversal

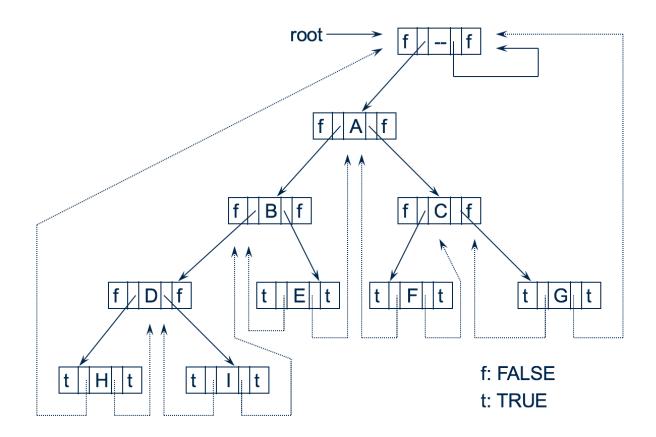


- How to distinguish actual pointers and threads?
  - →add two additional fields to the node structure
    - if ptr->left\_thread = true, ptr->left\_child contains thread
    - if ptr->left\_thread = false, ptr->left\_child contains a pointer to the left child

```
typedef struct threaded_tree *threaded_ptr;

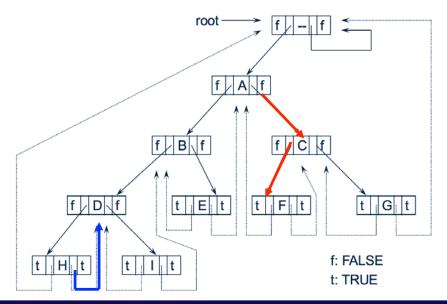
typedef struct threaded_tree {
    short int left_thread;
    threaded_ptr left_child;
    char data;
    threaded_ptr right_child;
    short int right_thread;
};
```





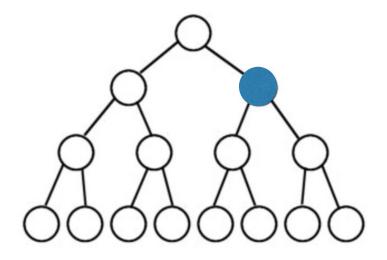
- find the in-order successor of ptr without using stack
  - if ptr -> right\_thread = TRUE, ptr -> right\_child
  - otherwise follow a path of left\_child links from the right\_child of ptr until we reach a node with left\_thread = TRUE

```
threaded_ptr insucc(threaded_ptr tree) {
    threaded_ptr temp;
    temp = tree->right_child;
    if (!tree->right_thread)
        while (!temp->left_thread)
        temp = temp->left_child;
    return temp;
}
```



- find the in-order successor of ptr without using stack
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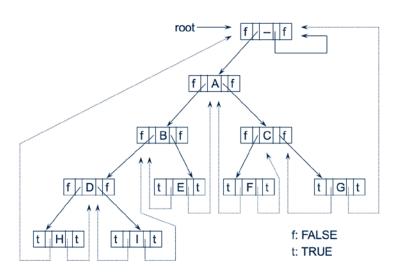
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}
```



Which node will be returned if blue node is passed into the function insucc?

- find the in-order successor of ptr without using stack
  - if ptr -> right\_thread = TRUE, ptr -> right\_child
  - otherwise follow a path of left\_child links from the right\_child of ptr until we reach a node with left\_thread = TRUE

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        temp = temp->left_child;
    return temp;
}
```



Which node will be returned if root node is passed into the function insucc?

```
void tinorder(threaded_ptr tree) {
    threaded_ptr temp = tree;
    for (;;) {
        temp = insucc(temp);
        if (temp == tree) break;
        printf("%3c", temp->data);
    }
}
```

