

# 1092 演算法 Homework 2

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tags: Homework

## 題目

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### 17-2 Making binary search dynamic

只需實作insert、search及print

Binary search of a sorted array takes logarithmic search time, but the time to insert a new element is linear in the size of the array. We can improve the time for insertion by keeping several sorted arrays.

Specifically, suppose that we wish to support SEARCH and INSERT on a set of  $n$  elements. Let  $k = \lceil \lg(n + 1) \rceil$ , and let the binary representation of  $n$  be  $\langle n_{k-1}, n_{k-2}, \dots, n_0 \rangle$ . We have  $k$  sorted arrays  $A_0, A_1, \dots, A_{k-1}$ , where for  $i = 0, 1, \dots, k-1$ , the length of array  $A_i$  is  $2^i$ . Each array is either full or empty, depending on whether  $n_i = 1$  or  $n_i = 0$ , respectively. The total number of elements held in all  $k$  arrays is therefore  $\sum_{i=0}^{k-1} n_i 2^i = n$ . Although each individual array is sorted, elements in different arrays bear no particular relationship to each other.

a. Describe how to perform the SEARCH operation for this data structure. Analyze its worst-case running time.

The  $A_{k-1}$  isn't have relationship for each other.

Search each array the worst-case time is  $k = \lceil \lg(n + 1) \rceil \rightarrow O(\lg(n))$ , and search each sorted array elements using binary search the worst-case time is  $O(\lg(n))$ , So the worst-case time of SEARCH operation is  $O(\lg(n)) * O(\lg(n)) = O(\lg^2(n))$

b. Describe how to perform the INSERT operation. Analyze its worst-case and amortized running times.

The worst-case time is  $O(n)$ . If now  $A_{k-1}$  need to merge to  $A_k$ , then merge two sorted array need  $O(m)$ , merge  $\sum_{i=0}^k A_i = O(km) = O(n)$ .

The andamortized time is  $O(\lg(n))$ . Merge array need  $2^k$ , Insert m times

$$T = \sum_{i=0}^{k-1} \frac{m}{2^i} 2^{i+1} \leq O(mk) \leq O(m \lg(n)) \text{ time, then each m insertion need}$$

$$\frac{O(m \lg(n))}{O(m)} = O(\lg(n)).$$