Compiler course work 1 report

Task I.1:

Aim is to extend MiniTriangle with a repeat-loop.

Should add extra token, extend AST have new print method and should let scanner and parser can identify new token.

First extend grammar:

Lexical Syntax:

Keyword repeat | until

Context-Free Syntax:

Command | repeat command until expression

Abstract Syntax:

Coding part:

Extend 'repeat', 'until' token to token.hs

```
47 | Repeat -- ^ \"repeat\"
48 | Until -- ^ \"until\"
49 | ElseIf -- ^ \"alseif\"
```

Extend command in AST.hs

```
120 | CmdRepeat {
121 | crBody :: Command, -- ^ Loop-body
122 | crUntil :: Expression, -- ^ Loop-condition
123 | cmdSrcPos :: SrcPos
```

Extend ppcommand in PPAST.hs to make parser print properly

```
ppCommand n (CmdRepeat {crBody = c, crUntil = ds, cmdSrcPos = sp}) =
indent n . showString "CmdRepeat" . spc . ppSrcPos sp . nl
. ppCommand (n+1) c
. ppExpression (n+1) ds
```

Extend mkIdOrKwd in scanner.hs to let compiler can properly identify the new token.

```
mkIdOrKwd "while" = While
mkIdOrKwd "repeat" = Repeat
mkIdOrKwd "until" = Until
```

Extend token and command in parser.y

```
74 REPEAT { (Repeat, $$) }
75 UNTIL { (Until, $$) }
```

```
REPEAT command UNTIL expression

{ CmdRepeat {crBody = $2, crUntil = $4, cmdSrcPos = $1} }
```

Task I.2

Should add extra token, extend AST have new print method and should let scanner and parser can identify new token.

First extend grammar:

Context-Free Syntax:

Expression | expression ? expression : expression

Abstract Syntax:

Expression | expression ? expression : expression ExpCond

Coding part:

Extend '?' token to token.hs since '?' is already there

```
QueMark -- ^ \"3\"
```

Exrend AST.hs:

```
153 | ExpCond {
154 | ecCond :: Expression,
155 | ecExp1 :: Expression,
156 | ecExp2 :: Expression,
157 | expSrcPos :: SrcPos
158 | }
```

Extend ppexpression to PPAST:

```
ppseq (n+1) ppexpression es

ppExpression n (ExpCond {ecCond = c, ecExp1 = e1, ecExp2 = e2, expSrcPos = sp}) =
  indent n . showString "ExpCond" . spc . ppSrcPos sp . nl
  . ppExpression (n+1) c
  . ppExpression (n+1) e1
  . ppExpression (n+1) e2
```

Scanner.hs:

```
scan l c ('?' : s) = retTkn QueMark l c (c + 1) s
```

Parser.y

```
61 '?' { (QueMark, $$) }

| expression '?' expression ':' expression
| { ExpCond {ecCond = $1, ecExp1 = $3, ecExp2 = $5, expSrcPos = srcPos $1} }
```

Task I.3:

This task should use maybe keyword to allow the condition command exist, and use the ppSeq to allow multiple Elseif command.

First extend grammar:

Lexical Syntax:

Keyword elseif

Context-Free Syntax:

Command | if expression then command

| if expression then command elseifcommands else command | if expression then command elseifcommands Elseifcommands -> Elseifcommand | elseifcommand elseifcommands Elseifcommand -> elseif expression then command Abstract Syntax: Command | if expression then command CmdIf | if expression then command elseifcommand* else command CmdIf | if expression then command elseifcommand* CmdIf Elseifcommand -> elseif expression then command CmdElIf Coding part: Token.hs: ElseIf

AST.hs:

```
96
          | CmdIf {
97
                ciCond
                         :: Expression,
                                             -- ^ Condition
98
                ciThen
                         :: Command,
                                               - ^ Then-branch
                ciElif
                         :: Maybe [Command],
99
                ciElse
                         :: Maybe Command,
                                             -- ^ Else-branch
100
                cmdSrcPos :: SrcPos
101
102
          | CmdElif {
103
104
                ceElif
                         :: Expression, — ^ Elseif-branch
105
                ceElThen :: Command,
                cmdSrcPos :: SrcPos
106
107
```

Add maybe [command] means this is sequence of command and it's conditional.

And add CmdElif is that I consider else if is a kind of command but it's following by CmdIf in parser.y do not have a command start with elseif it prevent error grammar like "elseif a then b" without if command.

PPAST.hs:

```
ppCommand n (CmdIf {ciCond = e, ciThen = c1, ciElif = cs, ciElse = c3, cmdSrcPos = sp}) =
   indent n . showString "CmdIf" . spc . ppSrcPos sp . nl
   . ppExpression (n+1) e
   . ppCommand (n+1) c1
   . maybe id (ppSeq (n+1) ppCommand) cs
   . maybe id (ppCommand (n+1)) c3

ppCommand n (CmdElif {ceElif = e, ceElThen = c1, cmdSrcPos = sp}) =
   --indent n . showString "CmdElif" . spc . ppSrcPos sp . nl
   ppExpression (n) e
   . ppCommand (n) c1
```

The ppcommand of cmdElif is a bit different to other command since I should make the command indent properly.

```
if a then a:=c elseif b then c := d
Information:
CmdIf line 1, column 1>
    ExpVar "a"
    CmdAssign <line 1, column 11>
        ExpVar "c"
    ExpVar "b"
    CmdAssign <line 1, column 30>
        ExpVar "c"
    ExpVar "d"
```

Two CmdAssign have same indent.

Scanner.hs:

```
198 mkIdOrKwd "elseif" = ElseIf
```

Simply add one token to it like task I.1

Parser.y:

```
150
      elseifcommands :: { [elseifcommand] }
151
      elseifcommands : elseifcommand { [$1] }
152
                       | elseifcommand elseifcommands { $1 : $2 }
153
154
      elseifcommand :: { Command }
155
      elseifcommand
          : ELSEIF expression THEN command
156
157
              {CmdElif {ceElif = $2, ceElThen = $4, cmdSrcPos = $1} }
158
```

should add extra command to it just follow the extended grammar.

Task I.4

This part should focus on extend the scanner, make scanner can identify the new grammar.

First extend grammar:

```
Lexical Syntax:
```

CharacterLiteral -> Graphic | EscapeCharacter

Graphic -> non-contril character except ' and \

```
EscapeCharecter -> \ (n \mid r \mid t \mid \ )
```

Coding part:

Token.hs:

AST.hs:

PPAST.hs:

```
83 Indent n . SnowString "ExpLitint". Spc . SnowS v . nt
84 ppExpression n (ExpLitChar {elcVal = v}) =
85 indent n . showString "ExpLitChar". spc . shows v . nl
```

Scanner.hs:

```
scan l c (x : s) | isDigit x = scanLitInt l c x s

isAlpha x = scanIdOrKwd l c x s

isOpChr x = scanOperator l c x s

isOpChr x = scanOperator l c x s

ix == '\'' = scanLitChar l c s

otherwise = do
```

When reach a 'start scanlitchar

```
132 ——when scanner read a ' means it can run into a lit char scan mode.

133 scanLitChar l c (x : b : xs) | x == '\\' = scanEscChar l (c+1) (b:xs)
```

When reach a \ start scanescchar

Return literal character with source position.

print error message properly.

When reach escape character return litchar, using a helper function to EscChar for change the single escape character to full escape character etc. $n \rightarrow n$

Parser.y:

```
TO LITCHAR { (LitChar {}, _)}

{ ExpLitInt {etivat = tspLivat $1, expSrcPos = tspSrcPos $1} }

LITCHAR

LITCHAR

{ ExpLitChar {elcVal = tspLCVal $1, expSrcPos = tspSrcPos $1} }
```