

CALIFORNIA STATE UNIVERSITY, NORTHRIDGE

PROTEIN FOLDING: PLANAR CONFIGURATION SPACES OF DISC  
ARRANGEMENTS AND HINGED POLYGONS

A thesis submitted in partial fulfillment of the requirements for the degree of  
Master of Science in Applied Mathematics

by

Clinton Bowen

August 2014

The thesis of Clinton Bowen is approved:

---

Dr. Silvia Fernandez

---

Date

---

Dr. John Dye

---

Date

---

Dr. Csaba Tóth, Chair

---

Date

California State University, Northridge

## Table of Contents

<b>Signature page</b> . . . . .	<b>ii</b>
<b>Abstract</b> . . . . .	<b>iv</b>
<b>Chapter 1</b>	
<b>Realizability Problems for Weighted Trees</b> . . . . .	<b>1</b>
1.1 Hausdorff Distance . . . . .	1
1.2 Weighted Trees $T_k$ . . . . .	2
1.2.1 Perturbed Weighted Trees $T_\varepsilon$ . . . . .	4

ABSTRACT

PROTEIN FOLDING: PLANAR CONFIGURATION SPACES OF DISC ARRANGEMENTS AND

HINGED POLYGONS

By

Clinton Bowen

Master of Science in Applied Mathematics

## Chapter 1

### Realizability Problems for Weighted Trees

In this chapter our goal is to prove Theorem ?? : It is NP-Hard to decide whether a given tree with positive vertex weights is the contact graph of a disk arrangements with specified radii. This chapter's approach to proving Theorem ?? introduces an ordered weighted tree  $T$  and perturbed ordered weight tree  $T_\epsilon$ , the Hausdorff distance, and then prove a lemma which shows that hexagons can be approximated by an ordered disk contact graph corresponding to the weighted tree  $T_\epsilon$ .

#### 1.1 Hausdorff Distance

Let  $A$  and  $B$  be sets in the plane. The *directed Hausdorff distance* is:

$$d(A, B) = \sup_{a \in A} \inf_{b \in B} \|a - b\| \quad (1.1)$$

$d(A, B)$  finds the furthest point  $a \in A$  from any point in  $B$ . *Hausdorff distance* is

$$D(A, B) = \max \{d(A, B), d(B, A)\} \quad (1.2)$$

In Figure 1.1, we have two sets  $X$  and  $Y$  and illustrate  $d(X, Y)$  and  $d(Y, X)$ . From this, it is possible to calculate the Hausdorff distance between  $X$  and  $Y$ .

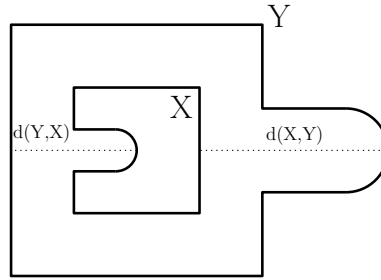


Figure 1.1: An illustrative example of  $d(X, Y)$  and  $d(Y, X)$  where  $X$  is the inner curve, and  $Y$  is the outer curve.

**$\epsilon$ -approximation** The weighted graph,  $G$ , is an  $\epsilon$ -approximation of a polygon  $P$  if the Hausdorff distance between every realization of  $G$  as a contact graph of disks and a congruent copy of  $P$  is at most epsilon. A weighted graph  $G$  is said to be a  $O(f(x))$ -approximation of a polygon  $P$  if there is a positive constant  $M$  such that for all sufficiently large values of  $x$  the Hausdorff distance between every realization such realization of  $G$  as a contact graph of disks and a congruent copy of  $P$  is at  $M \cdot |f(x)|$ . A weighted graph  $G$  is said to be a *stable* if it has the property that for every two such realizations of  $G$ , the distance between the centers of the corresponding disks is at most  $\epsilon$  after a suitable rigid transformation.

Suppose we have a unit disk  $U$  and we have a grid overlayed on the disk with side length  $\delta$ . Let  $S_1(\delta)$  be the the union of squares formed by the grid found completely in the interior of the disk  $U$ . Let  $S_2(\delta)$  be the union of squares formed by the grid with some point of the boundary of the square contained in the interior of disk  $U$ . The Hausdorff distance of  $U$  and  $S_1(\delta)$  is at most  $H(S_1(\delta), U) = \sqrt{2}\delta$ . Similarly, the Hausdorff distance of  $U$  and  $S_2(\delta)$  is at most  $H(S_2(\delta), U) = \sqrt{2}\delta$ . Thus for any  $\epsilon > 0$  choose a  $\delta$  such that  $\sqrt{2}\delta \leq \epsilon$ . Similiarily, the Hausdorff distance of  $U$  and  $S_2(\delta)$  is at most  $H(S_2(\delta), U) = \sqrt{2}\delta$ .

**Lemma 1.** For every  $\epsilon > 0$ , there exists an ordered weighted tree  $T_\epsilon$  and regular hexagon such that every realization of  $T_\epsilon$  as an ordered disk contact graph where the radii of the disks equal the vertex weights,

approximates the hexagon such that the Hausdorff distance is less than  $\varepsilon$ .

**Lemma 1 can be generalized in three different ways: (1) if all weights are equal, order does not matter, (2) if ..., and (3) ....**

## 1.2 Weighted Trees $T_k$

In this section we describe a particular family of unit weight trees and corresponding contact graphs disk arrangements called *snowflakes*. Note that we regard snowflakes with unit weight as a weight of  $r$ . For  $i \in \mathbb{N}$ , the construction of the snowflake tree,  $T_i$ , is as follows:

- Let  $v_0$  be a dvertex that has six paths attached to it:  $p_1, p_2, \dots, p_6$ . Each path has  $i$  vertices.
- For every other path  $p_1, p_3$ , and  $p_5$ :
  - Each vertex on that path has two paths attached, one path on each side of  $p_k$ .
  - The number of vertices that lie on a path attached to the  $j^{\text{th}}$  vertex of  $p_k$  is  $i - j$ .

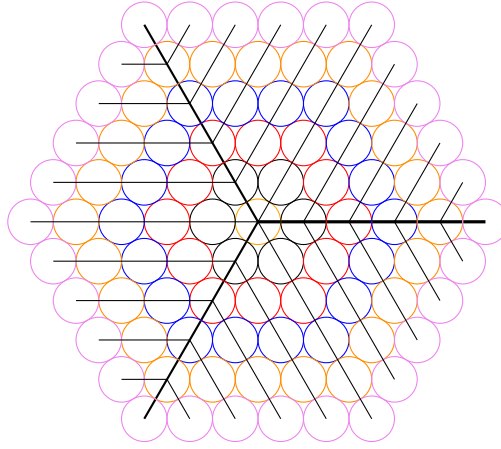


Figure 1.2: The same contact graph as in figure ?? overlaid with the a perfectly weighted snowflake tree.

A *perfectly weighted snowflake tree* is a snowflake tree with all vertices having weight  $r$ . A *perturbed snowflake tree* is a snowflake tree with all vertices having weight of 1 with the exception of  $v_0$ ; in a perturbed snowflake tree,  $v_0$  will have a weight of  $r + \varepsilon$ . For our analysis, all realizations of any snowflake, perfect or perturbed, shall have  $v_0$  fixed at origin.

**Perfectly Weighted Snowflake Tree.** Consider the graph of the triangular lattice with unit distant edges:

$$\begin{aligned} V &= \left\{ a \cdot (1, 0) + b \cdot \left( \frac{1}{2}, \frac{\sqrt{3}}{2} \right) : a, b \in \mathbb{Z} \right\} \\ E &= \{ \{u, v\} : \|u - v\| = 1 \text{ and } u, v \in V \} \end{aligned}$$

The following graph,  $G = (V, E)$  is said to be the *unit distance graph* of the triangular lattice. We can show that no two distinct edges of this graph are non-crossing. First suppose that there were two distinct edges that crossed,  $\{u_1, v_1\}$  and  $\{u_2, v_2\}$ . With respect to  $u_1$ , there are 6 possible edges corresponding to it, with each edge  $\frac{\pi}{3}$  radians away from the next. Neither edge crosses another; and so we have a contradiction that there are no edge crossings with  $\{u_1, v_1\}$ .

The perfectly weighted snowflake tree that is a subgraph over the *unit distance graph*,  $G = (V, E)$ , of the triangular lattice. To show this, for any  $S_i$ , fix  $v_0 = 0 \cdot (1, 0) + 0 \cdot \left(\frac{1}{2}, \frac{\sqrt{3}}{2}\right) = (0, 0) \in V$  at origin. Next consider the six paths attached from origin. Fix each consecutive path  $\frac{\pi}{3}$  radians away from the next such that the following points like on the corresponding paths:  $(1, 0) \in p_1, \left(\frac{1}{2}, \frac{\sqrt{3}}{2}\right) \in p_2, \left(-\frac{1}{2}, \frac{\sqrt{3}}{2}\right) \in p_3, (-1, 0) \in p_4, \left(-\frac{1}{2}, -\frac{\sqrt{3}}{2}\right) \in p_5, \left(\frac{1}{2}, -\frac{\sqrt{3}}{2}\right) \in p_6$ . For  $S_i$ , there are  $i$  vertices on each path.

We define the six paths from origin as follows:

$$\begin{aligned} p_1 &= \{a \cdot (1, 0) = \vec{v} \mid a \in \mathbb{R}^+\} \\ p_2 &= \left\{a \cdot \left(\frac{1}{2}, \frac{\sqrt{3}}{2}\right) = \vec{v} \mid a \in \mathbb{R}^+\right\} \\ p_3 &= \left\{-a \cdot (1, 0) + a \cdot \left(\frac{1}{2}, \frac{\sqrt{3}}{2}\right) = a \cdot \left(-\frac{1}{2}, \frac{\sqrt{3}}{2}\right) = \vec{v} \mid a \in \mathbb{R}^+\right\} \\ p_4 &= \{a \cdot (-1, 0) = \vec{v} \mid a \in \mathbb{R}^+\} \\ p_5 &= \left\{a \cdot \left(-\frac{1}{2}, -\frac{\sqrt{3}}{2}\right) = \vec{v} \mid a \in \mathbb{R}^+\right\} \\ p_6 &= \left\{a \cdot (1, 0) - a \cdot \left(\frac{1}{2}, \frac{\sqrt{3}}{2}\right) = a \cdot \left(\frac{1}{2}, -\frac{\sqrt{3}}{2}\right) \mid a \in \mathbb{R}^+\right\} \end{aligned}$$

For  $S_i$  there exists  $i$  vertices on each path. We shall denote the  $i^{\text{th}}$  vertex on the  $j^{\text{th}}$  path as  $v_{j,i}$ . For each path defined above, the paths are defined as a set of vectors,  $\vec{v} = a \cdot \vec{p}$  for some  $a \in \mathbb{R}^+$  and  $\vec{p} \in \mathbb{R}^2$ . By setting  $a = 1, 2, \dots, i$ , we obtain points that are contained in  $V$ . For  $j = 1, 3, 5$  and  $l = 3b \leq i$  where  $b \in \mathbb{N}$ , there exists two paths attached to each vertex  $v_{j,l}$ . We borrow the term *petiole* from botany to describe the two paths attached to  $v_{j,l}$ . In botany, the stalk that attaches to a stem of a plant is called a petiole; petioles usually have leaves attached to their ends. For  $S_i$ , each petiole attached to the  $k^{\text{th}}$  vertex of  $p_j$ , there are  $i - k$  vertices. For each vertex  $v$  on a petiole, which is not in the paths  $p_1, p_3$ , or  $p_5$ , there are two *leafs* on either side of the vertex; each leaf is a vertex that has an edge with  $v$ . The one exception to the two leafs rule is on the first vertex of the petiole off of  $p_1, p_3$ , or  $p_5$ . In this exception, attach one leaf to the side of the vertex that is closest to center vertex  $v_0$ .

The triangular lattice is symmetric under rotation about  $v_0$  by  $\frac{\pi}{3}$  radians. For each vertex  $v_{1,l}$  and  $l = 3b \leq i$  where  $b \in \mathbb{N}$ , we place two petioles from it; the first petiole  $\frac{\pi}{3}$  above  $p_1$  at  $v_{1,l}$  and  $\frac{-\pi}{3}$  below  $p_1$  at  $v_{1,l}$  and call these petioles  $p_{1,l}^+$  and  $p_{1,l}^-$  respectively. With respect to  $v_{1,l}$ , one unit along  $p_{1,l}^+$  is a point on the triangular lattice and similarly so on  $p_{1,l}^-$ . Continuing the walk along these paths, unit distance-by-unit distance, we obtain the next point corresponding point on the the triangular lattice up to  $i - k$  distance away from  $v_{1,l}$ . Without loss of generality, for each each vertex  $v$  of the petiole which are not in  $p_1$  have two associated leaf nodes  $v^+$  and  $v^-$ ;  $v^+$  is placed  $\frac{\pi}{3}$  and one unit above  $v$  and  $v^-$  is placed  $\frac{-\pi}{3}$  and one unit below  $v$ . Thus all leaf nodes are in the triangular lattices. This shows that each of the  $i - k$  vertices on  $p_{1,l}^-, p_{1,l}^+$ , and leafs are in  $V$ . By rotating all of the paths along  $p_1$  by  $\frac{2\pi}{3}$  and  $\frac{4\pi}{3}$ , we obtain the the paths along  $p_3$  and  $p_5$  respectively, completing the construction.

In Figure 1.3, we have a set of unit radii disks arranged in a manner that outlines the perfectly weighted snowflake description above.

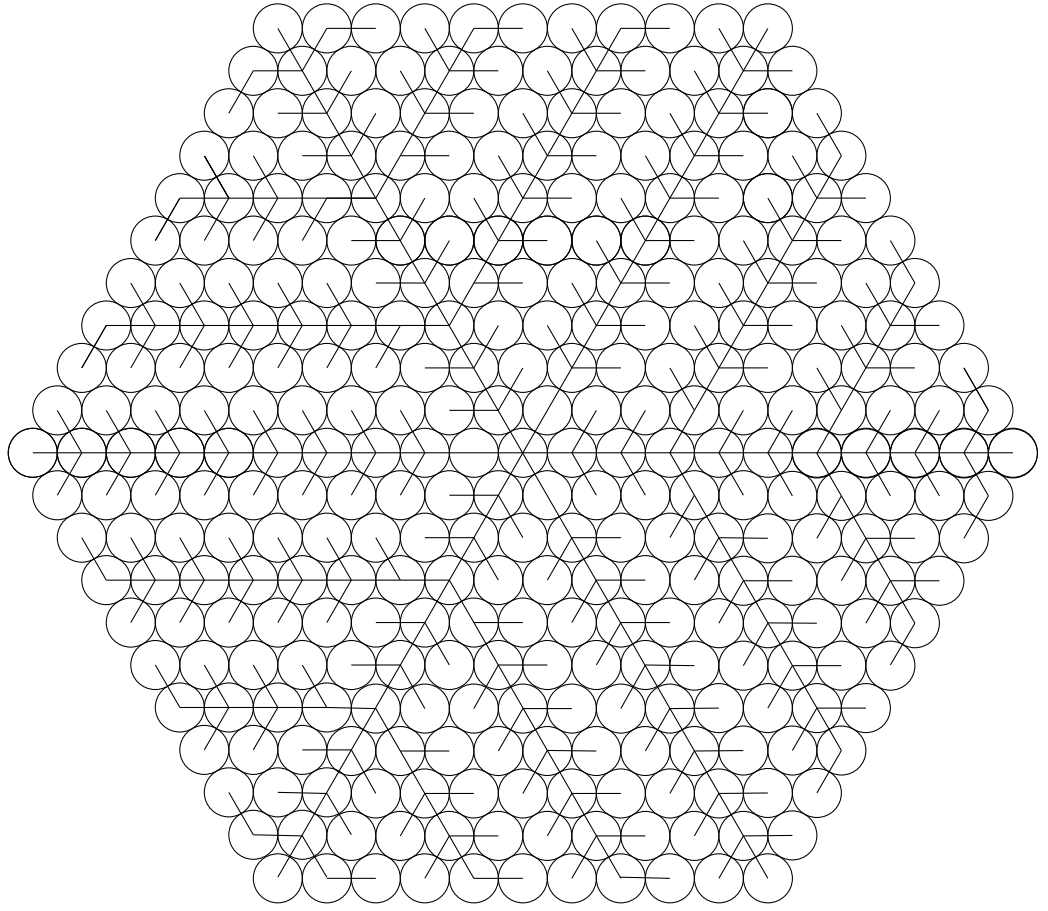


Figure 1.3: A contact graph that resembles the shape of concentric hexagons.

### 1.2.1 Perturbed Weighted Trees $T_\varepsilon$

A perturbed weighted tree  $T_\varepsilon$  is a weighted unit tree with unit weight on every vertex with the exception of the root vertex having weight  $r + \varepsilon$  where  $\varepsilon > 0$  can be realized as a disk touching graph (a disk arrangement) and  $r$  is the unit length.

The perturbed snowflake follows the construction of the perfect snowflake with the exception of  $v_0$  having weight  $r + \varepsilon$  where  $\varepsilon > 0$ . A perturbed snowflake realization has some distinct qualities from perfect snowflake realizations. The angular relationships between adjacent vertices may vary; the distance between adjacent and neighboring vertices may vary as well.

In general, the perturbation  $\varepsilon$  can modify the realization of a perfect snowflake  $S_i$  in the following ways:

#### **Modification of $S_1$ .**

Given a instance of a perturbed snowflake with  $v_0$  having weight  $r + \varepsilon$  where  $\varepsilon > 0$ , vertices neighboring  $v_0$  each have a range of placement on the plane when realized as a disk arrangement. Figure ?? shows a realization of  $S_1$  and illustrates one such example of possible gaps,  $\varepsilon$ , that could be created between adjacent disks of  $S_1$  in a perfect snowflake.

Note that (1) the adjacent disks in a perfect snowflake may or may not be adjacent in a given perturbed snowflake of  $S_1$  and (2)  $S_1 \subseteq S_i$  for any  $i \in \mathbb{N}$ . Given a snowflake in arbitrary position with  $n$  unit segments per arm, the arms of the snowflake has a maximal length of  $n$ , end to end, if in canonical position; otherwise, the arm will have an end to end length less than  $n$  (See Figure 1.4). The arm of a snowflake in arbitrary position



corresponds to a compression and shift of vertices.

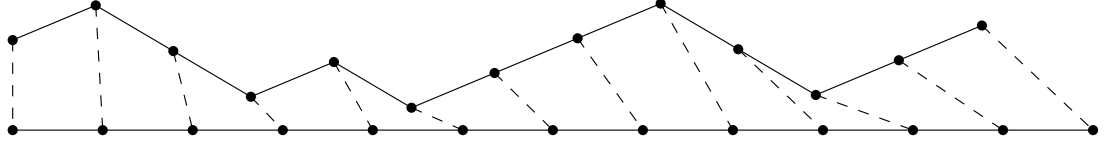


Figure 1.4: The polyline at the bottom represents a snowflake arm in canonical position. The polyline above represents a snowflake arm in non-canonical position.

We will show for any  $\varepsilon > 0$  and arbitrary position of vertices, the placement of vertices is close to canonical position. In order to show this, we show the components of a perturbed snowflake in arbitrary position are close to canonical position. The argument comprises of three parts: (1) Showing that the perturbation of  $S_1$  is small, (2) show that the displacement along the arms for all  $S_i$  for  $i \geq 1$  is small, and (3) show that the displacement along the petioles is small.

**Displacement on  $S_1$  is small.** Suppose we are given an ordered weighted tree  $T_\varepsilon$  such that the corresponding disk arrangement is a perturbed  $S_1$ . In a perfect snowflake of  $S_1$  the six disks around the central disk kiss each other. The angle formed from the center of the central disk to the centers of any two adjacent disks is  $\frac{\pi}{3}$ . The side lengths of the equilateral triangle formed by the centers of three adjacent disks, one of which is the central disk, is  $2r$ . For a perturbed  $S_1$  the central disk is weighted  $r + \varepsilon$ . This can yield a change of angular displacement from  $\frac{\pi}{3}$  to  $\frac{\pi}{3} + 2\chi$ . To find the bounds of how large or small  $\chi$  can be, we show the trigonometric relation of the half angle of the triangle corresponding to three adjacent disks (See Figure 1.5):

$$\begin{aligned}
\sin\left(\frac{\pi}{6} - \chi\right) &= \frac{1}{2r + \varepsilon} \\
\sin \frac{\pi}{6} \cos \chi &= \frac{1}{2r + \varepsilon} + \cos \frac{\pi}{6} \sin \chi \\
&\iff \\
\frac{1}{2} &\geq \frac{1}{2} \cos \chi \\
&= \frac{1}{2r + \varepsilon} + \frac{\sqrt{3}}{2} \sin \chi \\
&\geq \frac{1}{2r + \varepsilon} + \frac{\sqrt{3}}{2} \left(\chi - \frac{\chi^3}{6}\right) \\
&\iff \\
\frac{1}{2} - \frac{1}{2r + \varepsilon} &\geq \frac{\sqrt{3}}{2} \left(\chi - \frac{\chi^3}{6}\right) \quad \text{if } \chi < 1 \\
\frac{2r + \varepsilon - 2}{2(2 + \varepsilon)} &\geq \frac{5\sqrt{3}}{12} \chi \\
&\text{for } r = 1 \\
\frac{3\varepsilon}{5\sqrt{3}} = \frac{12}{5\sqrt{3}} \frac{\varepsilon}{4} &\geq \chi
\end{aligned}$$

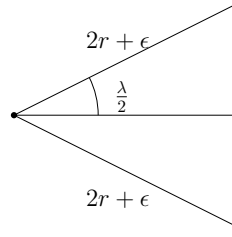


Figure 1.5: This figure depicts a triangle corresponding to the center of the central disk and two adjacent disks.

For any  $\varepsilon > 0$ , the bounds for angular displacement formed at the center of the central disk and two adjacent disks is:

$$\frac{\pi}{3} - \frac{6\varepsilon}{5\sqrt{3}} \leq \frac{\pi}{3} - 2\chi = \lambda_{\min} \leq \lambda \leq \lambda_{\max} = \frac{\pi}{3} + 2\chi \leq \frac{\pi}{3} + \frac{6\varepsilon}{5\sqrt{3}}$$

**Displacement on the arms is small.** To show that the angular displacement along the arm is small, we extend the angular argument on the perturbed  $S_1$  and by induction, show that it is small for all  $i$ .

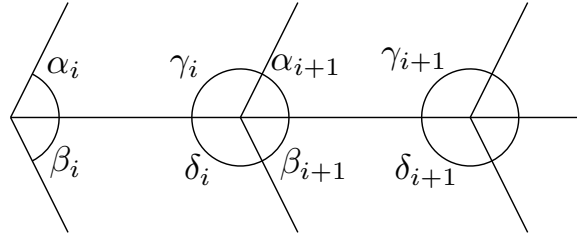


Figure 1.6: An arm depicted at the  $i^{\text{th}}$  and  $(i+1)^{\text{st}}$  vertex.

Denote the angles on the concave side of the  $i^{\text{th}}$  vertex as  $\alpha_i$  and  $\beta_i$  and the convex side of the  $(i+1)^{\text{st}}$  vertex as  $\gamma_i$  and  $\delta_i$  respectively (see Figure 1.6 for reference).

For any vertex, the sum of angles about the vertex is  $2\pi$ , e.g.:

$$\gamma_i + \delta_i + \alpha_{i+1} + \beta_{i+1} = 2\pi$$

Suppose we numbered the disks about the central disk 1 through 6. Without loss of generality, the angles  $\alpha_0$  and  $\beta_0$  correspond to the angles formed between the central angle, disks  $i$  and  $i+1$  and disks  $i+1$  and  $i+2$  respectively, for  $i = 1, 2, 3$ . The bounds for  $\alpha_0$  and  $\beta_0$  are the same as  $\lambda$  in the earlier argument, i.e.:

$$\begin{aligned} \frac{\pi}{3} - \frac{6\varepsilon}{5\sqrt{3}} &\leq \alpha_0 \leq \frac{\pi}{3} + \frac{6\varepsilon}{5\sqrt{3}} \\ \frac{\pi}{3} - \frac{6\varepsilon}{5\sqrt{3}} &\leq \beta_0 \leq \frac{\pi}{3} + \frac{6\varepsilon}{5\sqrt{3}} \end{aligned}$$

We know that  $\alpha_0 + \beta_0 \leq \frac{2\pi}{3} + \frac{12\varepsilon}{5\sqrt{3}}$ . We also know that in canonical position:

$$\begin{aligned} \pi &= \alpha_0 + \gamma_0 \\ \pi &= \beta_0 + \delta_0 \end{aligned}$$

Together, we have the following result:

$$\begin{aligned}
2\pi &= \alpha_0 + \gamma_0 + \beta_0 + \delta_0 \\
2\pi &= \alpha_0 + \gamma_0 + (2\pi - \alpha_1 - \beta_1) \\
\alpha_1 + \beta_1 &= \alpha_0 + \gamma_0 \\
&\leq \frac{2\pi}{3} + \frac{12\varepsilon}{5\sqrt{3}}
\end{aligned}$$

And so the error bounds on  $\lambda$  hold in general for  $\alpha_i$  and  $\beta_i$  for all  $i$ .

**Displacement on the petioles is small.** Note that the petioles have the same geometric structure as the arms; the exception is the number of leafs on each side of the petioles. Since we've shown that the geometric shape in arbitrary position is already close to canonical position for any  $\varepsilon > 0$ , the same argument applies here for the petioles.

We have shown the displacements of all components of the perturbed snowflake are small for any  $\varepsilon > 0$ . This shows that the structure has stability in preserving any information encoded with it.

### Proof of Lemma 1

*Proof.* Given an  $\varepsilon$ , we can construct an ordered weighted tree  $T_\varepsilon$  and hexagon such the Hausdorff distance of the ordered disk contact graph corresponding to  $T_\varepsilon$  of equal radii and the hexagon is less than  $\varepsilon$  in the following manner:

1. Let the radii of the of equally weighted disks be  $\frac{1}{kd}$  where  $d$  is the diameter of the hexagon and  $k$  is a multiple of 6, i.e.  $k = 6j$ .
2. The snowflake  $S_j$  is centered at the center of the hexagon.
3. The Hausdorff distance is at most  $\sqrt{3}\varepsilon$ .

□

### Proof of Theorem ??

*Proof.* Given an instance of a P3SAT boolean formula, we can use the snowflake reduction of the modified auxiliary construction. Using Lemma 1, we can approximate any hexagon with a tree  $T_\varepsilon$ . In the modified auxiliary construction in Chapter 3, we had four different hexagons and the skinny rhombus. We can scale the weights (radii) of the corresponding ordered weighted disk contact graph corresponding to  $T_\varepsilon$  to the rigid frame, obstacle, flag, and half sized hexagons in a modified auxiliary construction accordingly. The rhombus can be approximated by a chain of obstacle hexagons.

By approximating the polygons in the modified auxiliary construction with the snowflake, we show that Theorem ?? is a corollary by applying Lemma 1. □

## Bibliography