

CALIFORNIA STATE UNIVERSITY, NORTHRIDGE

PROTEIN FOLDING: PLANAR CONFIGURATION SPACES OF DISC  
ARRANGEMENTS AND HINGED POLYGONS

A thesis submitted in partial fulfillment of the requirements for the degree of  
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by

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## Table of Contents

<b>Signature page</b> . . . . .	<b>ii</b>
<b>Abstract</b> . . . . .	<b>iv</b>
 <b>Chapter 1</b>	
<b>Realizability Problems for Weighted Trees</b> . . . . .	<b>1</b>
1.1 Unit Disk Touching Graph Recognition Problem (UDTGRP) . . . . .	1
1.2 Perturbed Root with Unit Disk Leaves Touching Graph Recognition Problem (PRUDTGRP)	3
1.3 On the Decidability of Problem ?? . . . . .	5

ABSTRACT

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## Chapter 1

### Realizability Problems for Weighted Trees

We begin this chapter with the preliminary concepts to solving Problems ?? and ??, and then prove related Theorem ?. Recall the Unordered and Ordered Realizability Problem for a contact graph and corresponding theorem which state:

**Problem ??** Given a planar graph with positive weighted vertices, is it a contact graph of some disk arrangement where the radii equal the vertex weights?

**Problem ??** Given a planar graph with positive weighted vertices and a combinatorial embedding, is it a contact graph of some disk arrangement where the radii equal the vertex weights and the counter-clockwise order of neighbors of each disk is specified by the combinatorial embedding?

**Theorem ??** It is NP-Hard to decide whether a given tree (resp., plane tree) with positive vertex weights is the contact graph (resp., contact graph) of a disk arrangements with specified radii.

The preliminary concepts are the Hausdorff distance, the Unit Disk Touching Graph Recognition Problem, and the Perturbed Root with Unit Disk Leaves Touching Graph Recognition Problem.

#### 1.1 Unit Disk Touching Graph Recognition Problem (UDTGRP)

In Figure 1.1, we have a set of unit radius disks (circles) arranged in a manner that outlines regular, concentric hexagons.

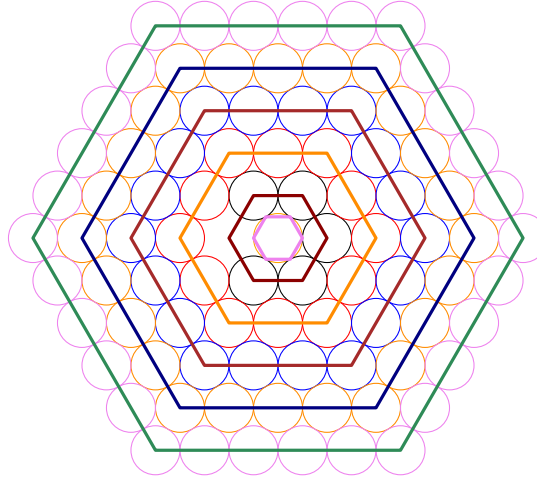


Figure 1.1: A contact graph that resembles the shape of concentric hexagons.

**Problem 1** (Approximating Polygonal Shapes with Contact Graphs). For every  $\varepsilon > 0$  and polygon  $P$ , there exists a contact graph  $G = (V, E)$  such that the Hausdorff distance  $d(P, G) < \varepsilon$

Recall problems (??) and (??): given a positive weighted tree,  $T$ , is  $T$  the (ordered) contact graph of some disk arrangement where the radii are equal to the vertex weights. For now, we'll focus on a particular family of this problem space where the weighted trees can be realized as a *snowflake*. For  $i \in \mathbb{N}$ , the construction of the snowflake tree,  $S_i$ , is as follows:

- Let  $v_0$  be a vertex that has six paths attached to it:  $p_1, p_2, \dots, p_6$ . Each path has  $i$  vertices.
- For every other path  $p_1, p_3$ , and  $p_5$ :
  - Each vertex on that path has two paths attached, one path on each side of  $p_k$ .
  - The number of vertices that lie on the path attached to the  $j^{\text{th}}$  vertex of  $p_k$  is  $i - j$ .

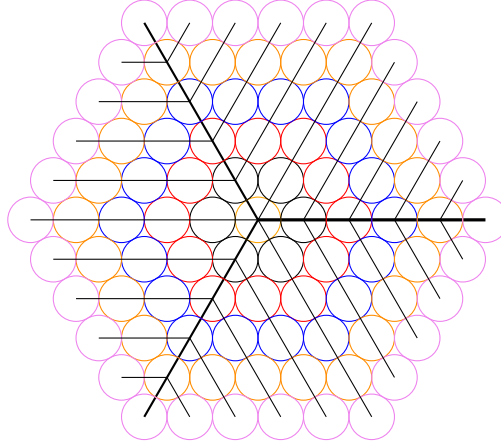


Figure 1.2: The same contact graph as in figure 1.1 overlaid with the a perfectly weighted snowflake tree.

A *perfectly weighted snowflake tree* is a snowflake tree with all vertices having weight  $\frac{1}{2}$ . A *perturbed snowflake tree* is a snowflake tree with all vertices having weight of 1 with the exception of  $v_0$ ; in a perturbed snowflake tree,  $v_0$  will have a weight of  $\frac{1}{2} + \gamma$ . For our analysis, all realizations of any snowflake, perfect or perturbed, shall have  $v_0$  fixed at origin. This is said to be the canonical position under Hausdorff distance of the snowflake tree.

**Perfectly Weighted Snowflake Tree.** Consider the graph of the triangular lattice with unit distant edges:

$$\begin{aligned}
 V &= \left\{ a \cdot (1, 0) + b \cdot \left( \frac{1}{2}, \frac{\sqrt{3}}{2} \right) : a, b \in \mathbb{Z} \right\} \\
 E &= \{ \{u, v\} : \|u - v\| = 1 \text{ and } u, v \in V \}
 \end{aligned}$$

The following graph,  $G = (V, E)$  is said to be the *unit distance graph* of the triangular lattice. We can show that no two distinct edges of this graph are non-crossing. First suppose that there were two distinct edges that crossed,  $\{u_1, v_1\}$  and  $\{u_2, v_2\}$ . With respect to  $u_1$ , there are 6 possible edges corresponding to it, with each edge  $\frac{\pi}{3}$  radians away from the next. Neither edge crosses another; and so we have a contradiction that there are no edge crossings with  $\{u_1, v_1\}$ .

The perfectly weighted snowflake tree that is a subgraph over the *unit distance graph*,  $G = (V, E)$ , of the triangular lattice. To show this, for any  $S_i$ , fix  $v_0 = 0 \cdot (1, 0) + 0 \cdot \left( \frac{1}{2}, \frac{\sqrt{3}}{2} \right) = (0, 0) \in V$  at origin. Next consider the six paths attached from origin. Fix each consecutive path  $\frac{\pi}{3}$  radians away from the next such that the following points lie on the corresponding paths:  $(1, 0) \in p_1$ ,  $\left( \frac{1}{2}, \frac{\sqrt{2}}{3} \right) \in p_2$ ,  $\left( -\frac{1}{2}, \frac{\sqrt{3}}{2} \right) \in p_3$ ,  $(-1, 0) \in p_4$ ,  $\left( -\frac{1}{2}, -\frac{\sqrt{3}}{2} \right) \in p_5$ ,  $\left( \frac{1}{2}, -\frac{\sqrt{3}}{2} \right) \in p_6$ . For  $S_i$ , there are  $i$  vertices on each path.

We define the six paths from origin as follows:

$$\begin{aligned}
p_1 &= \{ a \cdot (1, 0) = \vec{v} \mid a \in \mathbb{R}^+ \} \\
p_2 &= \left\{ a \cdot \left( \frac{1}{2}, \frac{\sqrt{3}}{2} \right) = \vec{v} \mid a \in \mathbb{R}^+ \right\} \\
p_3 &= \left\{ -a \cdot (1, 0) + a \cdot \left( \frac{1}{2}, \frac{\sqrt{3}}{2} \right) = a \left( -\frac{1}{2}, \frac{\sqrt{3}}{2} \right) = \vec{v} \mid a \in \mathbb{R}^+ \right\} \\
p_4 &= \{ a \cdot (-1, 0) = \vec{v} \mid a \in \mathbb{R}^+ \} \\
p_5 &= \left\{ a \cdot \left( -\frac{1}{2}, -\frac{\sqrt{3}}{2} \right) = \vec{v} \mid a \in \mathbb{R}^+ \right\} \\
p_6 &= \left\{ a \cdot (1, 0) - a \cdot \left( \frac{1}{2}, \frac{\sqrt{3}}{2} \right) = a \cdot \left( \frac{1}{2}, -\frac{\sqrt{3}}{2} \right) \mid a \in \mathbb{R}^+ \right\}
\end{aligned}$$

For  $S_i$  there exists  $i$  vertices on each path. We shall denote the  $i^{\text{th}}$  vertex on the  $j^{\text{th}}$  path as  $v_{j,i}$ . For each path defined above, the paths are defined as a set of vectors,  $\vec{v} = a \cdot \vec{p}$  for some  $a \in \mathbb{R}^+$  and  $\vec{p} \in \mathbb{R}^2$ . By setting  $a = 1, 2, \dots, i$ , we obtain points that are contained in  $V$ . For  $j = 1, 3, 5$  and  $l = 1, \dots, i$ , there exists two paths attached to each vertex  $v_{j,l}$ . For  $S_i$ , each path attached to the  $k^{\text{th}}$  vertex of  $p_j$ , there are  $i - k$  vertices. We will need to show that each of the  $i - k$  vertices on each corresponding path are also in  $V$ .

The triangular lattice is symmetric under rotation about  $v_0$  by  $\frac{\pi}{3}$  radians. For each vertex  $v_{1,l}$  for  $l = 1, 2, \dots, i - k$ , we place two paths from it; the first path  $\frac{\pi}{3}$  above  $p_1$  at  $v_{1,l}$  and  $\frac{-\pi}{3}$  below  $p_1$  at  $v_{1,l}$  and call these paths  $p_{1,l}^+$  and  $p_{1,l}^-$  respectively. With respect to  $v_{1,l}$ , one unit along  $p_{1,l}^+$  is a point on the triangular lattice and similarly so on  $p_{1,l}^-$ . Continuing the walk along these paths, unit distance-by-unit distance, we obtain the next point corresponding point on the the triangular lattice up to  $i - k$  distance away from  $v_{1,l}$ . This shows that each of the  $i - k$  vertices on  $p_{1,l}^-$  and  $p_{1,l}^+$  are in  $V$ . By rotating all of the paths along  $p_1$  by  $\frac{2\pi}{3}$  and  $\frac{4\pi}{3}$ , we obtain the the paths along  $p_3$  and  $p_5$  respectively, completing the construction.

## 1.2 Perturbed Root with Unit Disk Leaves Touching Graph Recognition Problem (PRUDTGRP)

The perturbed snowflake follows the construction of the perfect snowflake with the exception of  $v_0$  having weight  $\frac{1}{2} + \gamma$  where  $\gamma > 0$ . A perturbed snowflake realization has some distinct qualities from perfect snowflake realizations. The angular relationships between adjacent vertices may vary; the distance between adjacent and neighboring vertices may vary as well. Note that we regard snowflakes with unit weight as a weight of  $\frac{1}{2}$ .

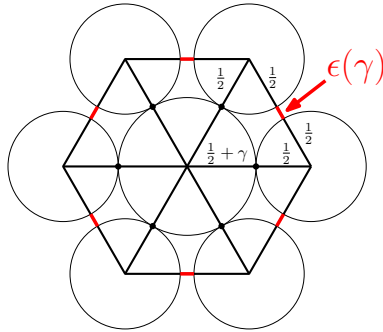


Figure 1.3: A canonical disk arrangement from a perturbed snowflake with 6 unit disks around a central disk with radius  $\frac{1}{2} + \gamma$ .

In Figure 1.3, we have a realization of disk arrangement from a perturbed snowflake. In a disk arrangement of a perfect snowflake, the disks around the central disk contact the adjacent disks. The disk arrangement

from the perturbed snowflake does not have this quality. Figure 1.3 shows a gap  $\varepsilon(\gamma)$  between adjacent disks around the central disk. This gap is formed from the perturbed weight  $\frac{1}{2} + \gamma$  of the central disk.

$$\varepsilon(\gamma) = 2\gamma + \gamma^2 \quad (1.1)$$

As the perturbed snowflake grows outer layers, we can begin to define parts of the snowflake and the corresponding disk arrangement. Let the arms extending from the center of a snowflake be *dendrites* and the arms extending off of arms be *metadendrites*. In a perturbed snowflake, the dendrites and metadendrites have some freedom to about the plane. our

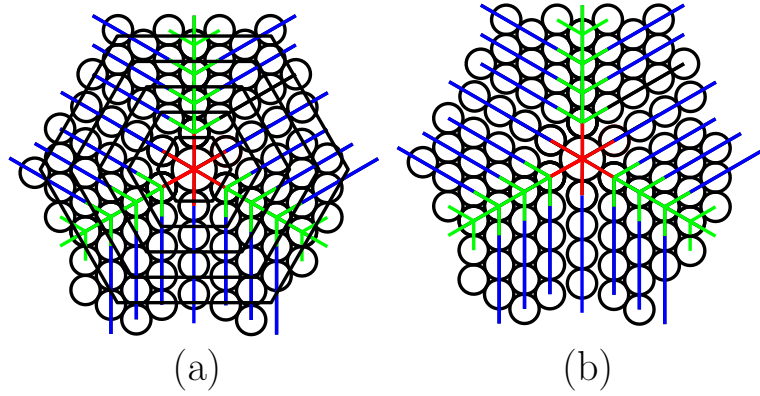


Figure 1.4

In Figure 1.4, we show an overlay of a realization of a perturbed snowflake, a corresponding disk arrangement, and concentric hexagons about the  $v_0$ .

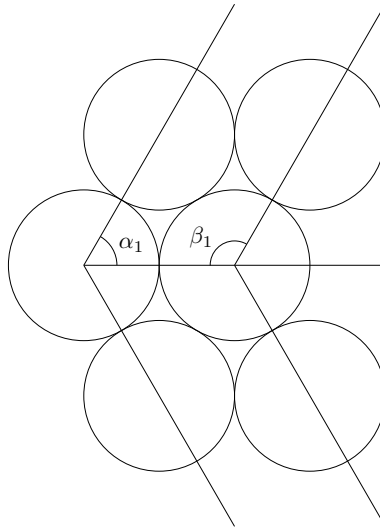


Figure 1.5: ?A?S

In Figure 1.5, we have a perturbed spine



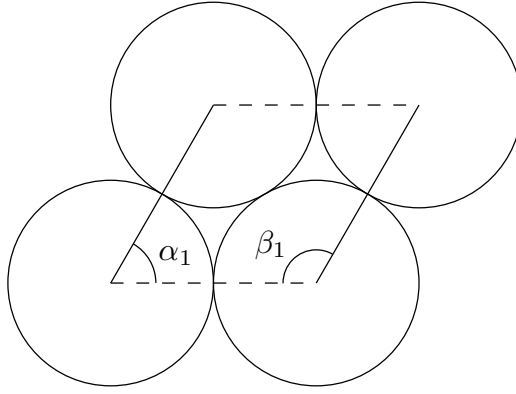


Figure 1.6: ?A?S

In Figure 1.6, we have four an arrangement of disks on the snowflake, off the spine and away from the central disk. We call this a vertebrae.

**Lemma 1.** *Given any realization of a perturbed snowflake of 7 weighted vertices, with the central vertex  $v_0$  weighted  $\frac{1}{2} + \gamma$  and the others weighted  $\frac{1}{2}$ , the total additional distance between all vertices is  $6\epsilon(\gamma)$  compared to a perfect snowflake of 7 unit weight vertices.*

*Proof.* Consider a canonical disk arrangement of a perturbed snowflake of 7 weighted vertices (see Figure 1.3). The side length of the sides formed between the center of the central disk and two adjacent disks around the central disk is  $1 + \gamma$ . Let the distance between the two adjacent disks be  $1 + \epsilon(\gamma)$ . There are a total of  $6\epsilon(\gamma)$  between adjacent centers of disks. The total perimeter of the hexagon formed about the centers of the disks in contact with the central disk is  $6 + 6\epsilon(\gamma)$ . Note that 1) the total perimeter of the hexagon formed on a perfect snowflake of 7 weighted vertices is 6 and 2) the canonical disk arrangement can be transformed to any other disk arrangement corresponding to the perturbed snowflake of 7 weighted vertices by pushing the the ring of disks around the central disk together such that all adjacent disks are in contact with each other with the exception of the disks at the end.  $\square$

### 1.3 On the Decidability of Problem ??

*Proof.* Consider a  $k \times (\sqrt{3}k)$  rectangle section of a triangular lattice, and place disks of radius 1 at each grid point as in Fig. ?????. The contact graph of these disks contains 2-cycles. Consider the spanning tree  $T$  of the contact graph indicated in Fig. ?????. The tree  $T$  decomposes into paths of collinear edges:  $T$  contains two paths along the two main diagonals, each containing  $2k - 1$  vertices; all other paths have an endpoint on a main diagonal. We now modify the disk arrangement to ensure that its contact graph is  $T$ . The disks along the main diagonal do not change. We reduce the radii of all other disks by a factor of  $1 - k^{-3}$  (as a result, they lose contact with other disks), and then successively translate them parallel in the direction of the shortest path in  $T$  to the main diagonal until the contact with the adjacent disk is reestablished. The Hausdorff distance between the union of these disks and the initial  $k \times (\sqrt{3}k)$  rectangle is clearly less than 1. However, the contact tree  $T$  with these radii no longer has a unique realization (small perturbations are possible). To show stability, we argue by induction on the hop distance from the central disk. There are  $O(i)$  disks at  $i$  hops from the central disk, most one which have radius  $(1 - k^{-3})\frac{1}{2}$ . Since all radii are 1 or  $(1 - k^{-3})\frac{1}{2}$ , the six neighbors of the central disk can differ from the regular hexagon by at most  $O(k)$ . Similarly, the disks at  $i$  hops from the center be off from the triangular grid pattern by  $O(i2^{k-3})$ , for  $i = 1, 2, \dots, k$ .

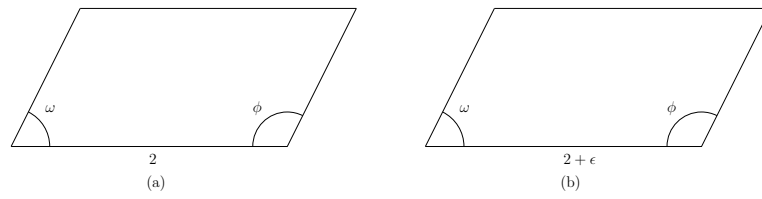


Figure 1.7

□