CALIFORNIA STATE UNIVERSITY, NORTHRIDGE

PROTEIN FOLDING: PLANAR CONFIGURATION SPACES

OF DISC ARRANGEMENTS AND HINGED POLYGONS:

PROTEIN FOLDING IN FLATLAND

A thesis submitted in partial fulfillment of the requirements
For the degree of Master of Science in Mathematics

By

Clinton Bowen

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The thesis of Clinton Bowen is approved:

Dr. John Dye | Date

Dr. Silvia Fernandez | Date

Dr. Bernardo Abrego | Date

Dr. Csaba Toth, Chair | Date

California State University, Northridge

DEDICATIONS

ACKNOWLEDGEMENTS

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ABSTRACT

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Insert Abstract here

Abstract

We look into the decidability of whether a hinged configuration locks.

1 Introduction

We look into the decidability of continuity on planar configuration space using regular, unitary hexagonal polygons. These polygons can also represent unit disk configurations [**?**]

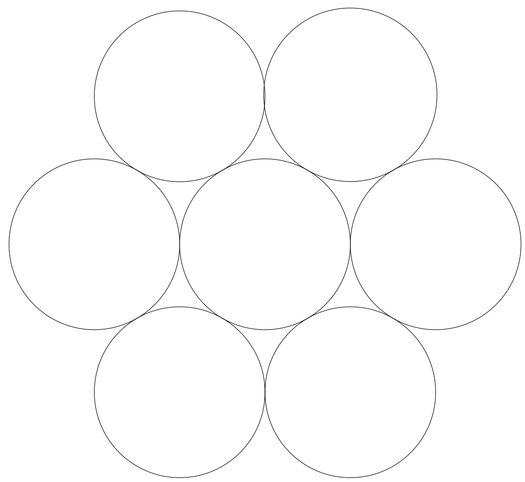

Figure 1: A locked 7 ball configuration

Motivation Protein folding, graphite, crystalline structures in metallurgy; disc packing; hexagonal configurations; Determine whether chemical structures are realizable.

Outline Section 2 covers the necessary mathematical concepts to understanding the problem. Section 3 explains the problem, Section 4 covers the results and findings about the problem. Section 5, the conclusion, offers final remarks on the problem.

2 Background

Here we review some of the necessary mathematics behind the problem. The definitions found in this chapter are those found in [**?**, **?**, **?**].

Here goes something

2.1 Circle Packing

It turns out the circle packings are an equivalent way to to represent linkages and their corresponding problems. Before we establish the relation, we will cover some fundamental concepts of circle packings. A *circle packing*, P, embedded in a plane is a set of circles with disjoint interiors $\{C_i\}_{i=1}^n$ such that for any circle $C \in \{C_i\}_{i=1}^n$, C is tangent to a different circle of $\{C_i\}_{i=1}^n$.

Any circle embedded in a plane has a given center point and radius. This information of planar embedded circle packings allows us to establish the relationship to linkages with the following construction:

(i) let the centerpoints of the circle packing be a set of vertices V;

(ii) if two circles in a circle packing are tangent, we define an edge between their centerpoints. The distance of this edge is the sum the radii of the two tangent circles.

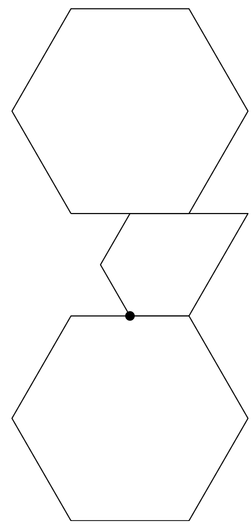

Figure 2: blah blah blah

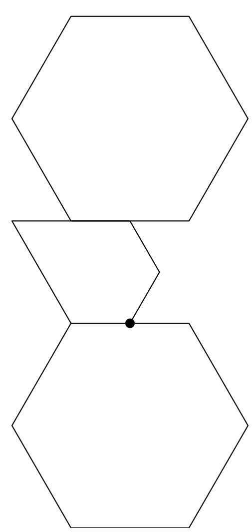

Figure 3: blah blah blah

This construction establishes a relationship between linkages and circle packings. It begs questioning as to whether every connected simple planar graph has a circle packing. The question is answered in the following theorem.

Theorem 2.1 (Circle Packing Theorem)**.** *For every connected simple planar graph G there is a circle packing in the plane whose intersection graph is (isomorphic to) G.*

A proof of Theorem 2.1 is found in chapter 7 of [**?**]. Theorem 2.1 also gives us the ability to establish an equivalent definition of configuration spaces on circle packings and allows us to pose the same realizability problems found with simple planar graphs. To narrow the focus of the types of circle packing realizabilty problems that we are interested in, we add the following restriction: all circles in a circle packing have unit diameter.

2.1.1 Realizability Problems in Unit Disk Packings

In [**?**], it was shown that unit disk graph recognition is NP-Hard.

2.2 Area Packing Problem

2.2.1 Hinged Polygons

Definition 2.1 (Polygonal Chain)**.** A polygonal chain $P = (v_0, v_1, \ldots, v_{n-1})$ is a sequence of consecutively joined segments (or edges) $e_i = v_i v_{i+1}$ of fixed lengths $l_i = |e_i|$, in a plane. [**?**]

A chain is said to be closed if $v_{n-1} = v_1$, otherwise it is said to be open. Hinged polygons have been researched for decades and related to linkage problems [**?**, **?**].

Consider the locked configuration of figure 4. We can configure the hexagons to be locked by placing hinged points as follows:

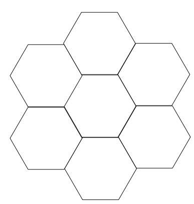

Figure 4: A locked 7 hexagonal configuration. (needs to modify picture by placing red points for hing points.)

2.2.2 Hinged Hexagons of Fixed Size

The Shapes Figure 5 is a locking shape: Figure 5 shall reside in the boundary of a lattice and have a hinge

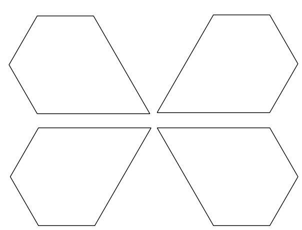

Figure 5: A locking shape in the lattice boundary's channel.

point at one vertex where the locking shape and boundary meet.

Junctions We define junctions to be the point three hexagons meet in a hexagonal lattice, e.g. Figure **??**.

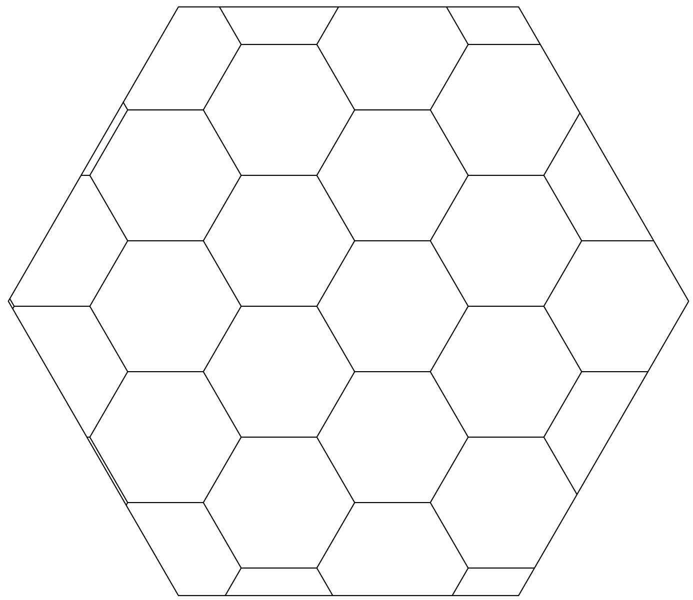

Figure 6: A portion of a hexagonal lattice.

Central Scaling

Junctions in Conjunctive Normal Form Explain the configurations we're interested in.

3 Configuration Spaces of Polygonal Chains

3.1 SAT Problems

Problem 3.1 (Satisfiability Problem). Let $\{x_i\}_{i=1}^n$ be boolean variables, and $t_i \in \{x_i\}_{i=1}^n \cup \{\bar{x}_i\}_{i=1}^n$. A *clause* is is said to be a disjuction of distinct terms:

$$t_1 \vee \cdots \vee t_{j_k} = C_k$$

Then the *satisfiability problem* is the decidability of a conjuction of a set of clauses, i.e.:

$$\wedge_{i=1}^m C_i$$

[**?**] A *3-SAT problem* is a SAT problem with all clauses having only three boolean variables.

Definition 3.1 (Planar 3-SAT Problem)**.** Given a boolean 3-SAT formula B, define the associated graph of B as follows:

$$G(B) = (\{\, v_x \,|\, v_x \text{ represents a variable in } B \,\} \cup \{\, v_C \,|\, v_C \text{ represent a clause in } B \,\}, \{\, (v_x, v_C) \,|\, x \in C \text{ or } \bar{x} \in C \,\}) \tag{1}$$

If $G(B)$ in equation (**??**) is planar, then B is said to be a *Planar 3-SAT Problem* [**?**].

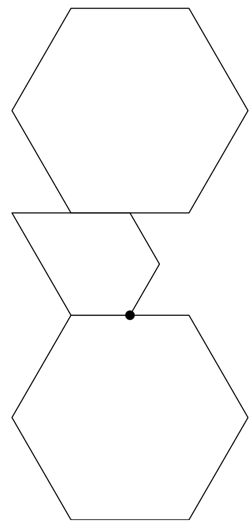

Figure 7: left switch between polygons

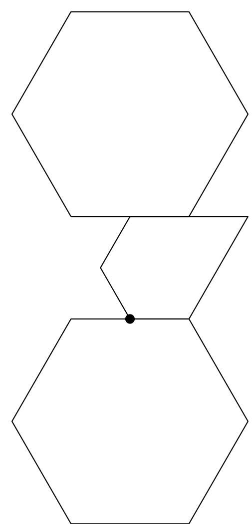

Figure 8: right switch between polygons

4 Problem

4.1 Problem Statement

text

4.2 Decidability of Problem

test

4.3 Hexagonal Locked Configuration

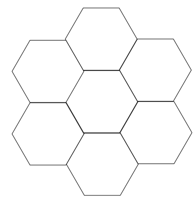

Figure 9: 7 hexagonal configuration

5 Conclusion

We conclude...