

CALIFORNIA STATE UNIVERSITY, NORTHRIDGE

PROTEIN FOLDING: PLANAR CONFIGURATION SPACES OF DISC
ARRANGEMENTS AND HINGED POLYGONS

A thesis submitted in partial fulfillment of the requirements for the degree of
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by

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ABSTRACT

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Chapter 1

Realizability Problems for Weighted Trees

In this chapter our goal is to prove Theorem ?? : It is NP-Hard to decide whether a given tree with positive vertex weights is the contact graph of a disk arrangements with specified radii. This chapter's approach to proving Theorem ?? introduces an ordered weighted tree T , perturbed ordered weight tree T_ε , the Hausdorff distance, and then prove a lemma which shows that hexagons can be approximated by an ordered disk contact graph corresponding to the weighted tree T_ε .

1.1 Hausdorff Distance

Let A and B be sets in the plane. The *directed Hausdorff distance* is:

$$d(A, B) = \sup_{a \in A} \inf_{b \in B} \|a - b\| \quad (1.1)$$

$d(A, B)$ finds the furthest point $a \in A$ from any point in B . *Hausdorff distance* is

$$D(A, B) = \max \{d(A, B), d(B, A)\} \quad (1.2)$$

In Figure 1.1, we have two sets X and Y and illustrate $d(X, Y)$ and $d(Y, X)$. From this, it is possible to calculate the Hausdorff distance between X and Y .

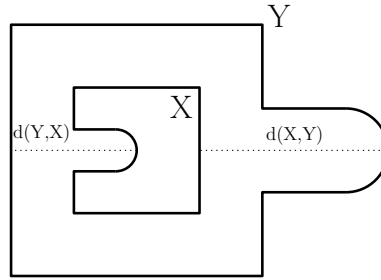


Figure 1.1: An illustrative example of $d(X, Y)$ and $d(Y, X)$ where X is the inner curve, and Y is the outer curve.

ε -approximation The weighted graph, G , is an ε -approximation of a polygon P if the Hausdorff distance between every realization of G as a contact graph of disks and a congruent copy of P is at most epsilon. A weighted graph G is said to be a $O(f(x))$ -approximation of a polygon P if there is a positive constant M such that for all sufficiently large values of x the Hausdorff distance between every realization such realization of G as a contact graph of disks and a congruent copy of P is at $M \cdot |f(x)|$. A weighted graph G is said to be a *stable* if it has the property that for every two such realizations of G , the distance between the centers of the corresponding disks is at most ε after a suitable rigid transformation.

Suppose we have a unit disk U and we have a grid overlayed on the disk with side length δ . Let $S_1(\delta)$ be the union of grid squares completely in the interior of U . Let $S_2(\delta)$ be the union of squares that intersect U . The Hausdorff distance of U and $S_1(\delta)$ is at most $H(S_1(\delta), U) \leq \sqrt{2}\delta$. Similarly, the Hausdorff distance of U and $S_2(\delta)$ is at most $H(S_2(\delta), U) \leq \sqrt{2}\delta$. For any $\varepsilon > 0$, choose a δ such that $\sqrt{2}\delta \leq \varepsilon$. Then the Hausdorff distance between U and $S_1(\delta)$, U and $S_2(\delta)$ is:

$$\begin{aligned} H(S_1(\delta), U) &\leq \sqrt{2}\delta \\ H(S_2(\delta), U) &\leq \sqrt{2}\delta \end{aligned}$$

Similiarily, the Hausdorff distance of U and $S_2(\delta)$ is at most $H(S_2(\delta), U) = \sqrt{2}\delta$.

Lemma 1. *For every $\varepsilon > 0$ and $x > 0$, there exists an ordered weighted tree T_ε and regular hexagon h of side length x such that:*

1. *Every realization r of T_ε as an ordered disk contact graph where the radii of the disks equal the vertex weights, approximates the hexagon in the sense that:*

$$H(r(T_\varepsilon), h) = \varepsilon$$

2. *The number of nodes in T_ε and the weights are polynomial in ε and x .*

Lemma 1 is rather restrictive; it begs the question of whether it can be generalized. Some ways this lemma could be generalized are: (1) if all weights are equal, (2) order does not matter, and (3) if we relax the regular hexagon to any arbitrary polygon.

1.2 Weighted Trees T_k

In this section we describe a particular family of unit weight trees and corresponding contact graphs disk arrangements called *snowflakes*. Note that we regard snowflakes with unit weight as a weight of r . For $i \in \mathbb{N}$, the construction of the snowflake tree, T_i , is as follows:

- Let v_0 be a vertex that has six paths attached to it: p_1, p_2, \dots, p_6 . Each path has i vertices.
- In botany, the stalk that attaches to a stem of a plant is called a *petiole*; petioles usually have leaves attached to their ends. Our snowflake will incorporate these concepts by incorporating petioles and leaves. Petioles will be paths that extend off the arms. *Leaves* will be paths of one vertex that extend off the petioles and arms. We will now attach paths (petioles) onto every other path p_1, p_3 , and p_5 :
 - Every third vertex on that path has two petioles attached, one petiole on each side of p_k ; otherwise the vertex on p_k will have two leaves, one on either side.
 - The number of vertices that lie on a petiole attached to the j^{th} vertex of p_k is $i - j$.
 - The first vertex of the $i - j$ vertices has one petiole attached; the remaining $i - j - 1$ vertices contain two paths. Each of these paths contain only one vertex. These paths are called *leaves*.
- Attach leaves to the vertices on paths p_2, p_4 , and p_6 .

An drawaing example of this snowflake description is shown in Figure 1.2

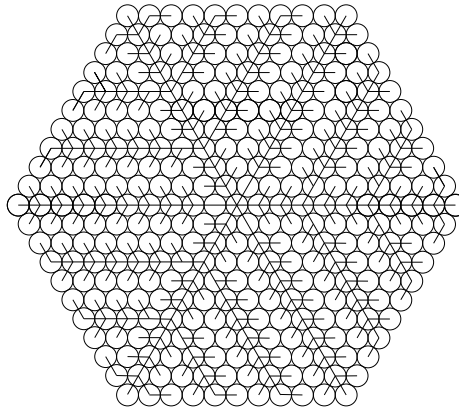


Figure 1.2: A contact graph that resembles the shape of concentric hexagons.

A *perfectly weighted snowflake tree* is a snowflake tree with all vertices having weight r . A *perturbed snowflake tree* is a snowflake tree with all vertices having weight of 1 with the exception of v_0 ; in a perturbed snowflake tree, v_0 will have a weight of $r + \zeta(\varepsilon)$. The value of $\zeta(\varepsilon)$ will be determined later on in the proof of the Lemma 1. For our analysis, all realizations of any snowflake, perfect or perturbed, shall have the disk corresponding to v_0 is centered at the origin. We can assume that a neighbor of v_0 is on the positive x-axis.

Perfectly Weighted Snowflake Tree. Consider the graph of the triangular lattice with unit distant edges:

$$\begin{aligned} V &= \left\{ a \cdot (1, 0) + b \cdot \left(\frac{1}{2}, \frac{\sqrt{3}}{2} \right) : a, b \in \mathbb{Z} \right\} \\ E &= \{ \{u, v\} : \|u - v\| = 1 \text{ and } u, v \in V \} \end{aligned}$$

The graph, $G = (V, E)$ is said to be the *unit distance graph* of the triangular lattice. We can show that no two distinct edges cross. First suppose that there were two distinct edges that crossed, $\{u_1, v_1\}$ and $\{u_2, v_2\}$. By strict triangle inequality of the sides and diagonals of the quadrilateral (u_1, u_2, v_1, v_2) , we have the following result:

$$|u_1, u_2| + |u_2, v_1| + |v_1, v_2| + |u_1, v_2| < 4 < 2(|u_1, u_2| + |u_1, v_2|) = 4$$

which is a contradiction.

The perfectly weighted snowflake tree that is a subgraph over the *unit distance graph*, $G = (V, E)$, of the triangular lattice. For the remainder of the thesis, a *snowflake*, S_i is a realization of a weighted tree T_i . To show this, for any S_i , fix $v_0 = 0 \cdot (1, 0) + 0 \cdot \left(\frac{1}{2}, \frac{\sqrt{3}}{2} \right) = (0, 0) \in V$ at origin. Next consider the six paths attached from origin. Fix each consecutive path $\frac{\pi}{3}$ radians away from the next such that the following points like on the corresponding paths: $(1, 0) \in p_1, \left(\frac{1}{2}, \frac{\sqrt{3}}{2} \right) \in p_2, \left(-\frac{1}{2}, \frac{\sqrt{3}}{2} \right) \in p_3, (-1, 0) \in p_4, \left(-\frac{1}{2}, -\frac{\sqrt{3}}{2} \right) \in p_5, \left(\frac{1}{2}, -\frac{\sqrt{3}}{2} \right) \in p_6$. For S_i , there are i vertices on each path.

We define the six paths from origin as follows:

$$\begin{aligned} p_1 &= \{ a \cdot (1, 0) \mid a = 1, 2, \dots, i \} \\ p_2 &= \left\{ a \cdot \left(\frac{1}{2}, \frac{\sqrt{3}}{2} \right) \mid a = 1, 2, \dots, i \right\} \\ p_3 &= \left\{ a \cdot \left(-\frac{1}{2}, \frac{\sqrt{3}}{2} \right) \mid a = 1, 2, \dots, i \right\} \\ p_4 &= \{ a \cdot (-1, 0) \mid a = 1, 2, \dots, i \} \\ p_5 &= \left\{ a \cdot \left(-\frac{1}{2}, -\frac{\sqrt{3}}{2} \right) \mid a = 1, 2, \dots, i \right\} \\ p_6 &= \left\{ a \cdot \left(\frac{1}{2}, -\frac{\sqrt{3}}{2} \right) \mid a = 1, 2, \dots, i \right\} \end{aligned}$$

For S_i there exists i vertices on each path. We shall denote the i^{th} vertex on the j^{th} path as $v_{j,i}$. For each path defined above, the paths are defined as a set of vectors, $\vec{v} = a \cdot \vec{p}$ for some $a \in \mathbb{N}$ and $\vec{p} \in \mathbb{R}^2$. By setting $a = 1, 2, \dots, i$, we obtain points that are contained in V . For $j = 1, 3, 5$ and $\ell = 3b \leq i$ where $b \in \mathbb{N}$, there exists two paths attached to each vertex $v_{j,\ell}$. For S_i , each petiole attached to the ℓ^{th} vertex of p_j , there are $i - \ell$ vertices. For each vertex v on a petiole, which is not in the paths p_1, p_3 , or p_5 , there are two *leaves* on either side of the vertex; each leaf is a vertex that has an edge with v . The exceptions to the two leaves rule is on the first and last vertices of the petiole off of p_1, p_3 , or p_5 . In these exception, attach one leaf to the side of the vertex that is closest to center vertex v_0 .

The triangular lattice is symmetric under rotation about v_0 by $\frac{\pi}{3}$ radians. For each vertex $v_{1,l}$ and $l = 3b \leq i$ where $b \in \mathbb{N}$, we place two petioles from it; the first petiole $\frac{\pi}{3}$ above p_1 at $v_{1,l}$ and $\frac{-\pi}{3}$ below p_1 at $v_{1,l}$ and call these petioles $p_{1,l}^+$ and $p_{1,l}^-$ respectively. With respect to $v_{1,l}$, one unit along $p_{1,l}^+$ is a point on the triangular lattice and similarly so on $p_{1,l}^-$. Without loss of generality, for each vertex v of the petiole which are not in p_1 has two associated leaf nodes v^+ and v^- ; v^+ is placed $\frac{\pi}{3}$ and one unit above v and v^- is placed $\frac{-\pi}{3}$ and one unit below v . Thus all leaf nodes are in the triangular lattice. This shows that each of the $i - k$ vertices on $p_{1,l}^-$, $p_{1,l}^+$, and leaves are in V . By rotating all of the paths along p_1 by $\frac{2\pi}{3}$ and $\frac{4\pi}{3}$, we obtain the paths p_3 and p_5 respectively, completing the construction.

In Figure 1.2, we have a set of unit radii disks arranged in a manner that outlines the perfectly weighted snowflake description above.

1.2.1 Perturbed Weighted Trees T_ε

In a perfect snowflake there are contacts in the disk arrangement that do not reflect the the corresponding edge set of the tree. By modifying the weight of the center disk corresponding to v_0 to differ from all others weights in a snowflake, the disks around the central disk will have additional area of realization and removes the issue of having unintended constacts that do not reflect a give tree's edge relations. This claim is shown as a result of Lemma 2 later.

Given $\varepsilon > 0$ and $x > 0$, we define T_ε as follows: the tree T_i with weight ε for every vertex except v_0 ; v_0 has a weight $\varepsilon + \zeta(\varepsilon)$ for some $\zeta(\varepsilon) > 0$ that is specified later. For T_i , let

$$i = \left\lceil \frac{x}{\varepsilon} \right\rceil.$$

A perturbed weighted tree T_ε can be realized as a disk touching graph (a disk arrangement). A perturbed snowflake realization has some distinct qualities from perfect snowflake realizations. The angular relationships between adjacent vertices may vary.

Modification of S_1 . We will show for any $\varepsilon > 0$ and arbitrary position of vertices, the placement of vertices is close to canonical position. In order to show this, we show the components of a perturbed snowflake in arbitrary position are close to canonical position. The argument comprises of three parts: (1) Showing that the perturbation of S_1 is small, (2) show that the displacement along the arms for all S_i for $i \geq 1$ is small, and (3) show that the displacement along the petioles is small.

Given a instance of a perturbed snowflake with v_0 having weight $\varepsilon + \zeta(\varepsilon)$ where $\varepsilon > 0$, vertices neighboring v_0 each have a range of placement on the plane when realized as a disk arrangement. Figure ?? shows a realization of S_1 and illustrates one such example of possible gaps, $\zeta(\varepsilon)$, that could be created between adjacent disks of S_1 in a perfect snowflake.

Displacement on S_1 is small. Note that (1) the adjacent disks in a perfect snowflake may or may not be adjacent in a given perturbed snowflake of S_1 and (2) $S_1 \subseteq S_i$ for any $i \in \mathbb{N}$. Given a snowflake in arbitrary position with n unit segments per arm, the arms of the snowflake has a maximal length of n , end to end, if in canonical position; otherwise, the arm will have an end to end length less than n . Figure 1.3 shows an arm of a tree in arbitrary position corresponds to a compression and shift of vertices. The arm realized in arbitrary position in Figure 1.3 is analogous to a tree realized in arbitrary position where vertices are in a different position than canonical. Our goal is to show that for any $\varepsilon > 0$ and $x > 0$, the position of the center v of any disk in the disk arrangement in any arbitrary realization corresponding to T_ε has a small displacement, i.e. v is found in an open ball referenced at the canonical position of v , v_c :

$$v \in b_{\chi(\varepsilon, x)}(v_c).$$

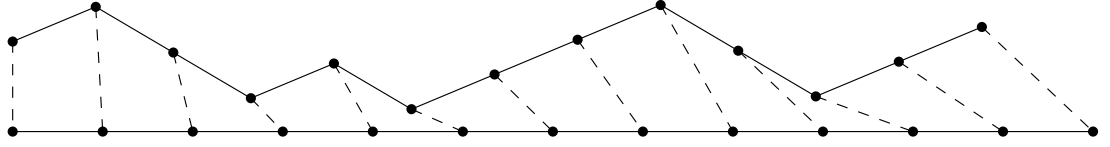


Figure 1.3: The polyline at the bottom represents a snowflake arm in canonical position. The polyline above represents a snowflake arm in non-canonical position.

Suppose we are given an ordered weighted tree T_ϵ such that the corresponding disk arrangement is a perturbed S_1 . In a perfect snowflake of S_1 the six disks around the central disk kiss each other. The angle formed from the center of the central disk to the centers of any two adjacent disks is $\frac{\pi}{3}$. The side lengths of the equilateral triangle formed by the centers of three adjacent disks, one of which is the central disk, is 2ϵ . For a perturbed S_1 the central disk is weighted $\epsilon + \zeta(\epsilon)$. This can yield a change of angular displacement $\frac{\pi}{3}$ to $\frac{\pi}{3} \pm 2\chi$. To find the bounds of how large or small χ can be, we show the trigonometric relation of the half angle of the triangle corresponding to three adjacent disks (See Figure 1.4).

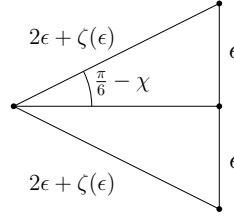


Figure 1.4: This figure depicts a triangle corresponding to the center of the central disk and two adjacent disks.

$$\begin{aligned}
 \sin\left(\frac{\pi}{6} - \chi\right) &= \frac{\epsilon}{2\epsilon + \zeta(\epsilon)} \\
 \frac{1}{2} \cos \chi = \sin \frac{\pi}{6} \cos \chi &= \frac{\epsilon}{2\epsilon + \zeta(\epsilon)} + \cos \frac{\pi}{6} \sin \chi = \frac{\epsilon}{2\epsilon + \zeta(\epsilon)} + \frac{\sqrt{3}}{2} \sin \chi \\
 &\iff \\
 \frac{\epsilon}{2\epsilon + \zeta(\epsilon)} + \frac{\sqrt{3}}{2} \left(\chi - \frac{\chi^3}{6}\right) &\leq \frac{\epsilon}{2\epsilon + \zeta(\epsilon)} + \frac{\sqrt{3}}{2} \sin \chi = \frac{1}{2} \cos \chi \\
 &\iff \\
 \frac{\sqrt{3}}{2} \left(\chi - \frac{\chi^3}{6}\right) &\leq \frac{1}{2} - \frac{\epsilon}{2\epsilon + \zeta(\epsilon)} \quad \text{if } \chi < 1 \\
 \frac{5\sqrt{3}}{12} \chi &\leq \frac{2\epsilon + \zeta(\epsilon) - 2\epsilon}{2(2 + \zeta(\epsilon))} \\
 \chi &\leq \frac{6}{5\sqrt{3}} \frac{\zeta(\epsilon)}{2 + \zeta(\epsilon)}
 \end{aligned}$$

The coordinates of the centers of the disks of the construction are close to canonical position following the angular displacement argument above. That is if u is a center of a disk in a disk arrangement corresponding to T_ϵ , u lies in a ball $b_{\chi(\epsilon)}(u_c)$ where u_c is the canonical position of the u .

For any $\epsilon > 0$, the bounds for angular displacement formed at the center of the central disk and two adjacent disks is:

$$\frac{\pi}{3} - \frac{12\zeta(\epsilon)}{5\sqrt{3}} \leq \frac{\pi}{3} - 2\chi = 2\chi_{\min} \leq 2\chi \leq 2\chi_{\max} = \frac{\pi}{3} + 2\chi \leq \frac{\pi}{3} + \frac{12\zeta(\epsilon)}{5\sqrt{3}}$$

Displacement on the arms is small. Without loss of generality, consider a path of vertices along a distinct arm, petiole, or leaf (v_1, v_2, \dots, v_n) . In the case that there are two paths attached to each vertex, four different angles about the vertex are formed. For the i^{th} vertex, the counter clockwise order of the angles are $(\alpha_{i+1}, \gamma_i, \delta_i, \beta_{i+1})$ (See Figure 1.6 for reference). We want to establish an angular relationship between two consecutive vertices in a path.

Lemma 2. *Let (a, b, c, d) be a polygonal path in the plane such that the unit disks centered at a, b, c , and d are interior-disjoint. Then the sum of the two angles on each side at the two interior vertices is greater than π . Then the sum of the two angles on each side at the two interior vertices is greater than π .*

Proof. Without loss of generality, consider the two angles on the left side at the two interior vertices, $\angle abc$ and $\angle bcd$. We have $|ab| = |bc| = |cd| = 2$, since we have a disk arrangement of unit disks. If (a, b, c, d) is a rhombus, then $|ad| = 2$ and $\angle abc + \angle bcd = \pi$. Hence $|ad| > 2$ implies $\angle abc + \angle bcd > \pi$. \square

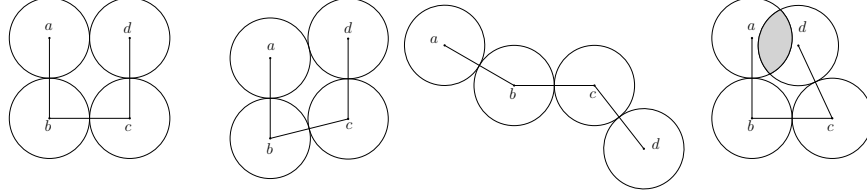


Figure 1.5: (a) Four disks centered along a polygonal path (a, b, c, d) in various drawings.

We can now say that for any two consecutive vertices (v_i, v_{i+1}) along a path, the following angular relationship from Lemma 2:

$$\pi < \alpha_i + \gamma_i$$

$$\pi < \beta_i + \delta_i$$

This result shows that a perturbed snowflake removes the issue of having unintended constacts that do not reflect a give tree's edge relations that a perfect snowflake has.

In Figure 1.6(a), we have a parallelogram with four unit disks centered, one on each of the vertices of the parallelogram. If $\alpha_i + \gamma_{i+1} < \pi$, One of the disks will overlap with another. For any arbitrary realization of a perturbed snowflake, any two consecutive angles formed between two adjacent vertices along a path must satisfy the following constraint

$$\alpha_i + \gamma_{i+1} \geq \pi.$$

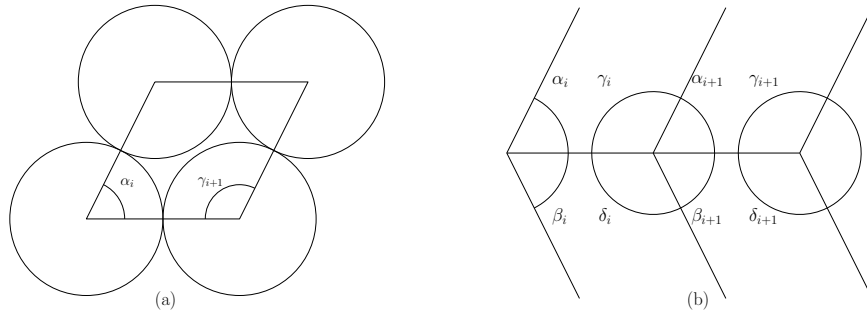


Figure 1.6: (a) Four disks of the snowflake shown where the top two disks can either be leaves off petioles or off the arms. (b) An arm depicted at the i^{th} and $(i+1)^{\text{st}}$ vertex.

To show that the angular displacement along the arm is small, we extend the angular argument on the perturbed S_1 and by induction, show that it is small for all i . Denote the angles on the concave side of the i^{th} vertex as α_i and β_i and the convex side of the $(i+1)^{\text{st}}$ vertex as γ_i and δ_i respectively (see Figure 1.6(b) for

reference). For any vertex, the sum of angles about the vertex is 2π , e.g.:

$$\gamma_i + \delta_i + \alpha_{i+1} + \beta_{i+1} = 2\pi$$

Suppose we numbered the disks about the central disk 1 through 6. Without loss of generality, the angles α_0 and β_0 correspond to the angles formed between the central angle, disks i and $i+1$ and disks $i+1$ and $i+2$ respectively, for $i = 1, 2, 3$. The bounds for α_0 and β_0 are the same as 2χ in the earlier argument, i.e.:

$$\begin{aligned} \frac{\pi}{3} - \frac{12\zeta(\varepsilon)}{5\sqrt{3}} &\leq \alpha_0 \leq \frac{\pi}{3} + \frac{12\zeta(\varepsilon)}{5\sqrt{3}} \\ \frac{\pi}{3} - \frac{12\zeta(\varepsilon)}{5\sqrt{3}} &\leq \beta_0 \leq \frac{\pi}{3} + \frac{12\zeta(\varepsilon)}{5\sqrt{3}} \end{aligned}$$

We know that $\alpha_0 + \beta_0 \leq \frac{2\pi}{3} + \frac{24\zeta(\varepsilon)}{5\sqrt{3}}$. We also know that in canonical position:

Together, we have the following result:

$$\begin{aligned} 2\pi &= \alpha_0 + \gamma_0 + \beta_0 + \delta_0 \\ 2\pi &= \alpha_0 + \gamma_0 + (2\pi - \alpha_1 - \beta_1) \\ \alpha_1 + \beta_1 &= \alpha_0 + \gamma_0 \\ &\leq \frac{2\pi}{3} + \frac{24\zeta(\varepsilon)}{5\sqrt{3}} \end{aligned}$$

And so the error bounds on $\frac{\pi}{3} + 2\chi$ hold in general for α_i and β_i for all i .

$$\alpha_i + \beta_i \leq \frac{2\pi}{3} + \frac{24\zeta(\varepsilon)}{5\sqrt{3}}$$

Displacement on the petioles is small. Note that the petioles have the same geometric structure as the arms; the exception is the number of leaves on each side of the petioles. Since we've shown that the geometric shape in arbitrary position is already close to canonical position for any $\varepsilon > 0$, the same argument applies here for the petioles.

We have shown the displacements of all components of the perturbed snowflake are small for any $\varepsilon > 0$. This shows that the structure has stability in preserving any information encoded with it.

Proof of Lemma 1

Proof. For any $\varepsilon > 0$, we construct an ordered weighted tree T_ε and regular hexagon of side length x .

Let the corresponding disk arrangement for T_ε be S_j where

$$j = \frac{x}{2\varepsilon + 1}.$$

The Hausdorff distance between the regular hexagon that is the convex hull for the centers of the disks in the disk arrangement and the union of the disks themselves is

$$H(r(T_\varepsilon), h) = \left(\frac{2}{\sqrt{3}} - 1 \right) \zeta(\varepsilon)$$

The number of disk in the contact graph corresponding is a polynomial $d(x, \varepsilon) \leq \frac{4x}{\varepsilon}$.

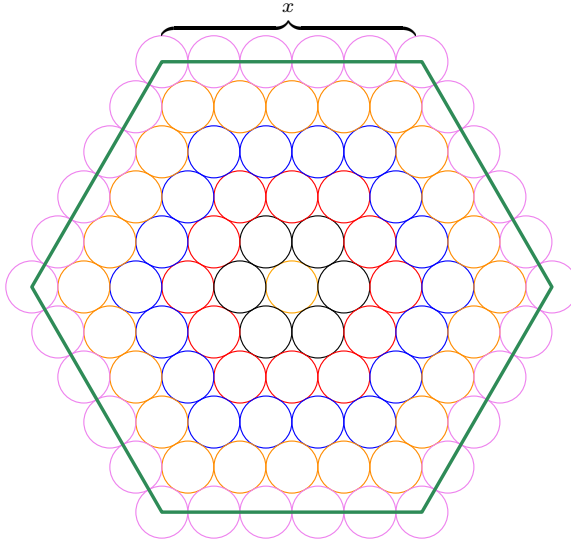


Figure 1.7: A regular hexagon of side length x as the convex hull of the centers of a disk arrangement in canonical position. □

Proof of Theorem ??

Proof. Given an instance of a P3SAT boolean formula, we can use the snowflake reduction of the modified auxiliary construction. For any center of a disk in any realization the displacement of the center is in an open ball $b_{\zeta(\epsilon)}(c)$ where c is the position of the center in canonical position. Using Lemma 1, we can approximate any hexagon with a tree T_ϵ . In the modified auxiliary construction in Chapter 3, we had four types of hexagons with different side lengths and the skinny rhombus. We can scale the weights (radii) of the corresponding ordered weighted disk contact graph corresponding to T_ϵ to the rigid frame, obstacle, flag, and half sized hexagons in a modified auxiliary construction accordingly. The rhombus can be approximated by a chain of obstacle hexagons.

By approximating the polygons in the modified auxiliary construction with the snowflake, we show that Theorem ?? is a corollary by applying Lemma 1. □

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