

# A Decidability Problem on Planar Protein Folding

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## Abstract

We look into the decidability of whether a hinged configuration locks.

## 1 Introduction

We look into the decidability of continuity on planar configuration space using regular, unitary hexagonal polygons. These polygons can also represent unit disk configurations [3]

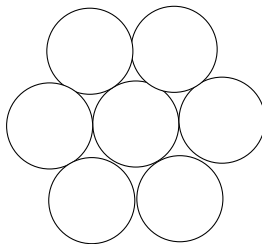


Figure 1: A locked 7 ball configuration

**Motivation** Protein folding, graphite, crystalline structures in metallurgy.

**Outline** Section 2 covers the necessary mathematical concepts to understanding the problem. Section 3 explains the problem, Section 4 covers the results and findings about the problem. Section 5, the conclusion, offers final remarks on the problem.

## 2 Background

Here we review some of the necessary mathematics behind the problem.

### 2.1 SAT Problems

*Problem 2.1* (Satisfiability Problem). Let  $\{x_i\}_{i=1}^n$  be boolean variables, and  $t_i \in \{x_i\}_{i=1}^n \cup \{\bar{x}_i\}_{i=1}^n$ . A *clause* is said to be a disjunction of distinct terms:

$$t_1 \vee \cdots \vee t_{j_k} = C_k$$

Then the *satisfiability problem* is the decidability of a conjunction of a set of clauses, i.e.:

$$\bigwedge_{i=1}^m C_i$$

[7]

### 2.1.1 3-SAT Problems

A 3-SAT problem is a SAT problem with all clauses having only three boolean variables.

## 2.2 Linkages

**Definition 2.1** (Linkage). A collection of fixed-length 1D segments joined at their endpoints to form a graph.

**Definition 2.2** (Graph). An ordered pair  $G = (V, E)$  comprising a set  $V$  of vertices or nodes together with a set  $E$  of edges or lines

**Definition 2.3** (Cycle). A closed walk with no repetitions of vertices or edges allowed, other than the repetition of the starting and ending vertex

**Definition 2.4** (Configuration). A specification of the location of all the link endpoints, link orientations and joint angles.[5]

## 2.3 Dissections

*Problem 2.2* (Polygonal Dissection). Given two polygons of equal area,  $P_1$  and  $P_2$ , partition  $P_1$  into smaller pieces,  $\{P_{1,i}\}_{i=1}^n$ , rearrange the pieces to form  $P_2$ . /citefrederickson1997dissections

## 2.4 Circle Packing

**Definition 2.5** (Circle Packing).  $P$  of a planar graph  $G$  is a set of circles with disjoint interiors  $\{C_v\}_{v \in G}$  such that two circles are tangent if and only if the corresponding vertices form an edge. [1]

**Theorem 2.1** (Circle Packing Theorem). *For every connected simple planar graph  $G$  there is a circle packing in the plane whose intersection graph is (isomorphic to)  $G$ . [8]*

### 2.4.1 Circle Packings and Polygonal Linkages

Given a circle of radius  $r$ , we establish the isomorphism to a hexagon by circumscribing the vertices of the regular hexagon.

### 2.4.2 Hinged Polygons

**Definition 2.6** (Polygonal Chain). A polygonal chain  $P = (v_0, v_1, \dots, v_{n-1})$  is a sequence of consecutively joined segments (or edges)  $e_i = v_i v_{i+1}$  of fixed lengths  $l_i = |e_i|$ , in a plane. [2]

Hinged polygons have been researched for decades and related to linkage problems [2, 4].

Consider the locked configuration of figure 4. We can configure the hexagons to be locked by placing hing points as follows: To prove that it is a locked configuration:

(i)

(ii)

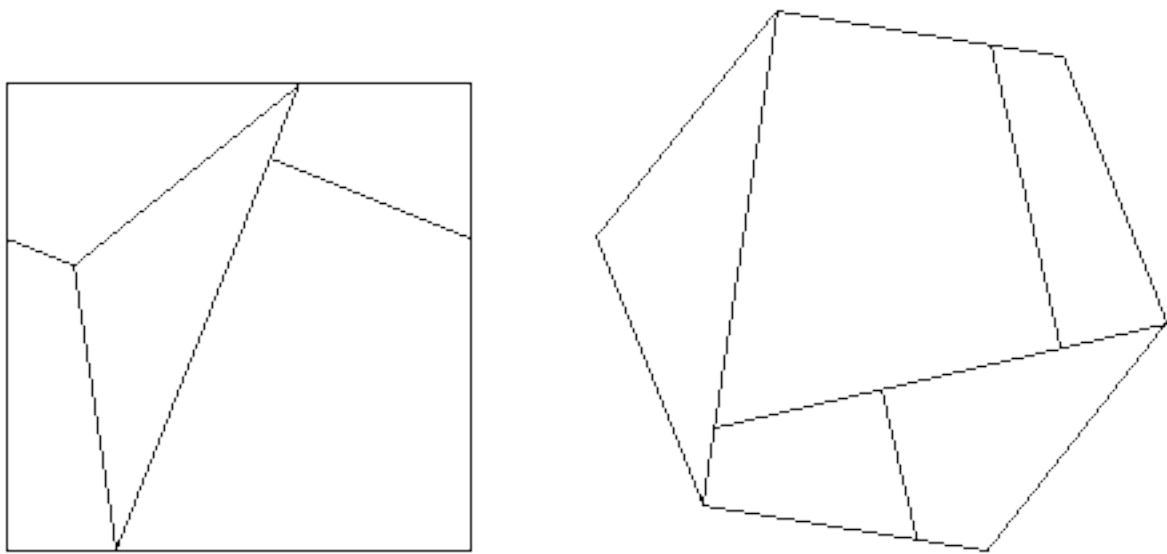


Figure 2: An axample of two polygons of equal area that can be rearranged into the other by the given partition.[6]

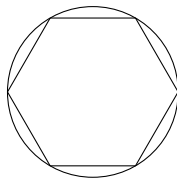


Figure 3: A circumscribed hexagon

- (iii)
- (iv)
- (v)
- (vi)
- (vii)
- (viii)
- (ix)
- (x)

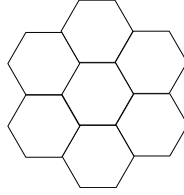


Figure 4: A locked 7 hexagonal configuration. (needs to modify picture by placing red points for hing points.)

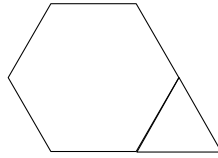


Figure 5: This is the shape that resides in boundary of the lattice.

### 2.4.3 Hinged Hexagons

**The Shapes** Figure 5 is a locking shape: Figure 5 shall reside in the boundary of a lattice and have a hinge point at one vertex where the locking shape and boundary meet.

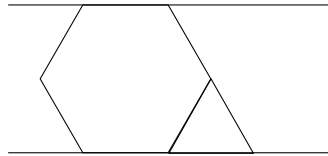


Figure 6: A locking shape in the lattice boundary's channel

**Junctions** We define junctions to be the point three hexagons meet in a hexagonal lattice, e.g. Figure 7.

### Central Scaling

**Junctions in Conjunctive Normal Form** Explain the configurations we're interested in.

### 2.4.4 Configurations and Locked Configurations

## 3 Problem

### 3.1 Problem Statement

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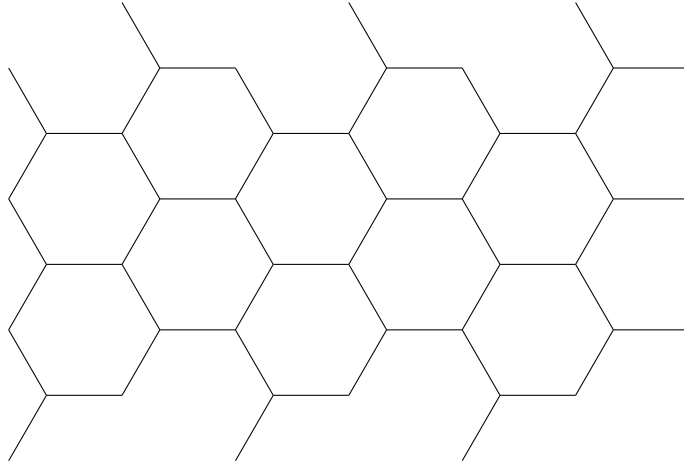


Figure 7: A portion of a hexagonal lattice.

### 3.2 Decidability of Problem

test

### 3.3 Hexagonal Locked Configuration

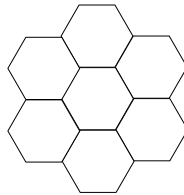


Figure 8: 7 hexagonal configuration

## 4 Conclusion

We conclude..

## References

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- [2] T. Biedl, E. Demaine, M. Demaine, S. Lazard, A. Lubiw, J. O'Rourke, M. Overmars, S. Robbins, I. Streinu, G. Toussaint, and S. Whitesides. Locked and unlocked polygonal chains in 3d, 1999.
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