

Planar Configuration Spaces of Disk Arrangements and Hinged Polygons

Clinton Bowen

Cal State Northridge

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Motivation: Hinged Dissection

- * *Dudeney Problem*: Can a square and an equilateral triangle of the same area have a common dissection into four pieces?
- * Can finite collection of polygons of equal area has a common hinged dissection?

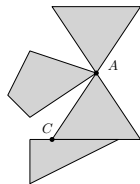
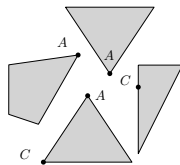


Figure: A collection of polygons with hinges.

Motivation: Hinged Dissection

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Figure: blah

Protein Folding

Protein folding is the process in which a protein chain acquires its 3-dimensional structure.

- * Proteins in an organism fold into a specific geometric pattern (sometimes referred as its *native state*).
- * Geometric patterns can determine a protein's function and behavior.

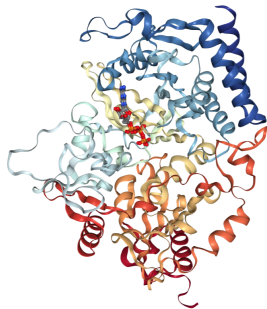


Figure: The structure of rat cytosolic PEPCK variant E89A in complex with oxalic acid and GTP [?].

Logic Engine Realized as Hinged Polygons.pdf

- * Suppose we are given an Boolean formula with m clauses and n variables in 3-CNF form, Φ , we construct the polygonal linkage similarly to the logic engine.

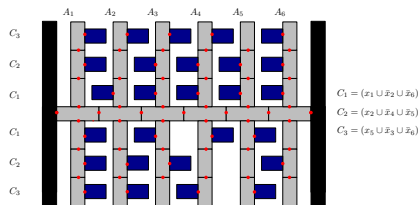


Figure: Caption Text

- * Breu and Kirkpatrick [?] proved that it is NP-hard to decide whether a graph G is the contact graph of unit disks in the plane, i.e., recognizing *coin graphs* is NP-hard; see also [?].

Planar 3SAT

- * Define the *associated graph* $A(\Phi)$ as follows: the vertices correspond to the variables and clauses in Φ . We place an edge in the graph if variable x_i appears in clause C_j .
- * Given a Boolean formula Φ in 3-CNF such that its associated graph is planar, decide whether it is satisfiable is a *3-SAT problem*.

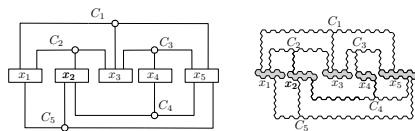


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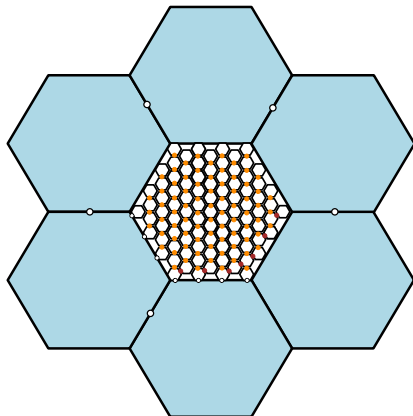


Figure: Caption Text

Transmitter Gadget

- * A *transmitter gadget* is constructed for each edge $\{x_i, C_j\}$ of the graph $A(\Phi)$; it consists of a sequence of junctions and corridors from a variable gadget's junction to a clause junction.

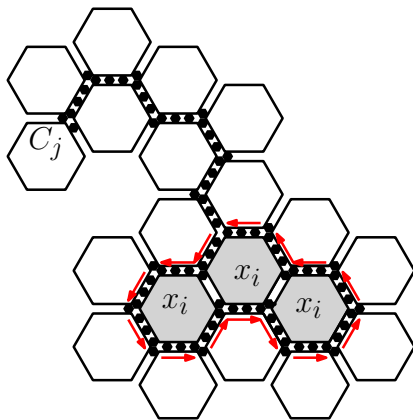
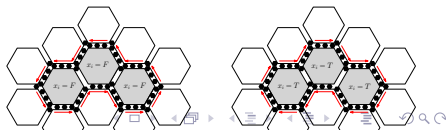


Figure: Caption Text



Variable Gadget

- * Variable x_i corresponds to a cycle in the associated graph $\tilde{A}(\Phi)$, which has been embedded as a cycle in the hexagonal tiling, with corridors and junctions. In each junction along this cycle, attach a small hexagon in the common boundary of the two corridors in the cycle.

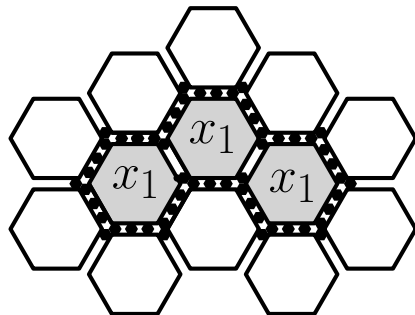


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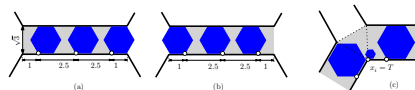


Figure: Caption Text

Clause Junction Gadget

- * The *clause gadget* lies at a junction adjacent to three transmitter gadgets (see Fig. ?? and Section ??). At such a junction, we attach a unit line segment to an arbitrary vertex of the junction, and a small hexagon of side length $\frac{1}{3}$ to the other end of the segment.

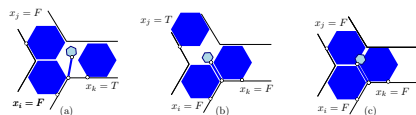
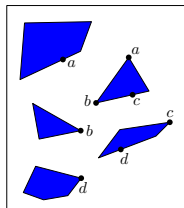


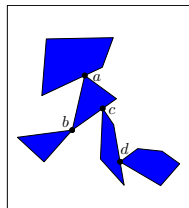
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Motivation: Weighted Trees and Disk Arrangements

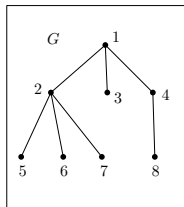
- * Is it strongly NP-hard to decide whether a polygonal linkage whose hinge graph is a *tree* can be realized?
- * Is it NP-Hard to decide whether a given ordered tree with positive vertex weights is the contact graph of a disk arrangements with specified radii?



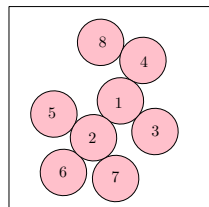
(a)



(b)



(c)



(d)

Figure: Caption Text

Modified Auxiliary Gadget

- * The modified auxiliary gadget channels and junctions in a hexagonal grid enclosed by six frame hexagons.
- * A comment about NAE3SAT....

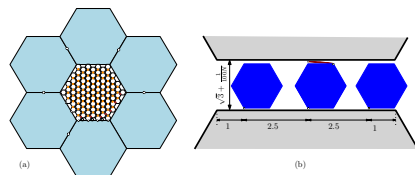


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- * item 1
- * item 2

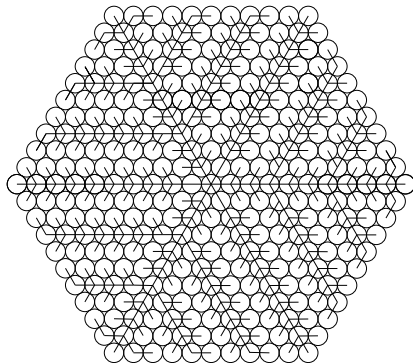


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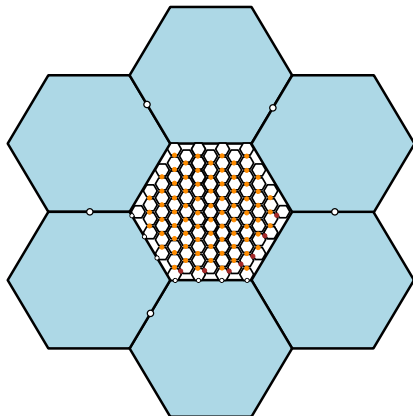


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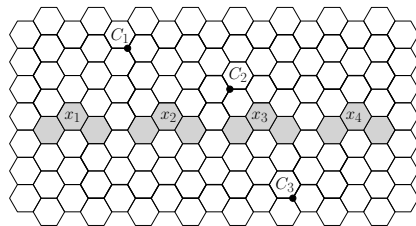


Figure: Caption Text

Problem

It is strongly NP-hard to decide whether a polygonal linkage whose inge graph is a *tree* can be realized.

It is strongly NP-hard to decide whether a polygonal linkage whose inge graph is a *tree* can be realized with fixed orientation.

It is NP-Hard to decide whether a given ordered tree with positive vertex weights is the contact graph of a disk arrangements with pecified radii.