

Chapter 1

Library c04_poly

Capítulo 4 - Polymorphism and Higher-Order Functions (Poly)

From *LF* Require Export *c03_lists*.

Listas polimórficas

```
Inductive list (X : Type) : Type :=  
  | nil : list X  
  | cons : X → list X → list X.
```

Função Repetir

```
Fixpoint repeat (X : Type) (x : X) (count : nat) : list X :=  
  match count with  
  | 0 ⇒ nil X  
  | S count' ⇒ cons X x (repeat X x count')  
end.
```

Example *test_repeat1* : repeat nat 4 2 = cons nat 4 (cons nat 4 (nil nat)).

Example *test_repeat2* : repeat bool false 1 = cons bool false (nil bool).

Exercise: 2 stars (mumble_grumble)

```
Inductive mumble : Type :=  
  | a : mumble  
  | b : mumble → nat → mumble  
  | c : mumble.
```

```
Inductive grumble (X : Type) : Type :=  
  | d : mumble → grumble X  
  | e : X → grumble X.
```

Check d (b a 5).

Check d mumble (b a 5).

Check d bool (b a 5).

Check e bool true.

Check e mumble $(b \ c \ 0)$.

Check e bool $(b \ c \ 0)$.

Check c .

Função Repetir

```
Fixpoint repeat' X x count : list X :=
  match count with
  | 0 => nil X
  | S count' => cons X x (repeat' X x count')
  end.
```

Example $test_repeat'1 : repeat' \text{ nat } 4 \ 2 = cons \text{ nat } 4 (cons \text{ nat } 4 (nil \text{ nat}))$.

Example $test_repeat'2 : repeat' \text{ bool } false \ 1 = cons \text{ bool } false (nil \text{ bool})$.

```
Fixpoint repeat'' X x count : list X :=
  match count with
  | 0 => nil _
  | S count' => cons _ x (repeat'' _ x count')
  end.
```

Example $test_repeat''1 : repeat'' \text{ nat } 4 \ 2 = cons \text{ nat } 4 (cons \text{ nat } 4 (nil \text{ nat}))$.

Example $test_repeat''2 : repeat'' \text{ bool } false \ 1 = cons \text{ bool } false (nil \text{ bool})$.

Definições

Definition $list123 := cons \text{ nat } 1 (cons \text{ nat } 2 (cons \text{ nat } 3 (nil \text{ nat})))$.

Definition $list123' := cons _ 1 (cons _ 2 (cons _ 3 (nil _)))$.

Argumentos

Definition $list123'' := cons \ 1 (cons \ 2 (cons \ 3 \ nil))$.

Função Repetir

```
Fixpoint repeat''' {X : Type} (x : X) (count : nat) : list X :=
  match count with
  | 0 => nil
  | S count' => cons x (repeat''' x count')
  end.
```

Listas polimórficas

```
Inductive list' {X : Type} : Type :=
  | nil' : list'
  | cons' : X -> list' -> list'.
```

Função Concatenar

```

Fixpoint app {X : Type} (l1 l2 : list X) : (list X) :=
  match l1 with
  | nil => l2
  | cons h t => cons h (app t l2)
  end.

```

Função Reverter

```

Fixpoint rev {X : Type} (l : list X) : list X :=
  match l with
  | nil => nil
  | cons h t => app (rev t) (cons h nil)
  end.

```

Example *test_rev1* : *rev (cons 1 (cons 2 nil)) = (cons 2 (cons 1 nil))*.

Example *test_rev2*: *rev (cons true nil) = cons true nil*.

Função Comprimento

```

Fixpoint length {X : Type} (l : list X) : nat :=
  match l with
  | nil => 0
  | cons _ l' => S (length l')
  end.

```

Example *test_length1*: *length (cons 1 (cons 2 (cons 3 nil))) = 3*.

Notações

Notation "*x :: y*" := (*cons x y*) (at level 60, right associativity).

Notation "*[]*" := *nil*.

Notation "*[x ; .. ; y]*" := (*cons x .. (cons y []) ..*).

Notation "*x ++ y*" := (*app x y*) (at level 60, right associativity).

Definition *list123'''* := [*1; 2; 3*].

Exercise: 2 stars, optional (poly-exercises)

Theorem *app_nil_r* : $\forall (X : \text{Type}), \forall l : \text{list } X,$
 $l ++ [] = l$.

Theorem *app_assoc* : $\forall A (l m n : \text{list } A),$
 $l ++ m ++ n = (l ++ m) ++ n$.

Lemma *app_length* : $\forall (X : \text{Type}) (l1 l2 : \text{list } X),$
 $\text{length } (l1 ++ l2) = \text{length } l1 + \text{length } l2$.

Exercise: 2 stars, optional (more_poly-exercises)

Theorem *rev_app_distr* : $\forall X (l1 l2 : \text{list } X),$
 $\text{rev } (l1 ++ l2) = \text{rev } l2 ++ \text{rev } l1$.

Theorem *rev_involutive* : $\forall X : \text{Type}, \forall l : \text{list } X,$
rev (*rev* *l*) = *l*.

Pares polimórficos

Inductive *prod* *X Y* :=
| *pair* : *X* \rightarrow *Y* \rightarrow *prod X Y*.

Notações

Notation "(*x* , *y*)" := (*pair x y*).

Notation "*X* * *Y*" := (*prod X Y*) : *type_scope*.

Projeção *x*

Definition *fst* {*X Y* : *Type*} (*p* : *X* \times *Y*) : *X* :=
 match p with
 | (*x*, *y*) \Rightarrow *x*
 end.

Projeção *y*

Definition *snd* {*X Y* : *Type*} (*p* : *X* \times *Y*) : *Y* :=
 match p with
 | (*x*, *y*) \Rightarrow *y*
 end.

Função Combinar

Fixpoint *combine* {*X Y* : *Type*} (*lx* : *list X*) (*ly* : *list Y*) : *list (X* \times *Y)* :=
 match lx, ly with
 | [], - \Rightarrow []
 | -, [] \Rightarrow []
 | *x* :: *tx*, *y* :: *ty* \Rightarrow (*x*, *y*) :: (*combine tx ty*)
 end.

Exercise: 1 star, optional (*combine_checks*)

Check @*combine*.

Exercise: 2 stars, recommended (*split*)

Função Dividir

Fixpoint *split* {*X Y* : *Type*} (*l* : *list (X* \times *Y)*) : (*list X*) \times (*list Y*) :=
 match l with
 | [] \Rightarrow ([], [])
 | (*x*, *y*) :: *t* \Rightarrow **let** *rest* := *split t* **in** ((*x* :: *fst rest*), (*y* :: *snd rest*))
 end.

Example *test_split*: *split* [(1,*false*);(2,*false*)] = ([1;2],[*false*;false]).

Option

Module *OptionPlayground*.

```

Inductive option (X : Type) : Type :=
| Some (x : X)
| None.

```

End OptionPlayground.

Erro

```

Fixpoint xth_error {X : Type} (l : list X) (x : nat) : option X := match l with | [] =>
None | a :: l' => if eqb x 0 then Some a else xth_error l' (pred x) end.

```

Example test_xth_error1 : xth_error 4;5;6;7 0 = Some 4. Proof. simpl. reflexivity. Qed.

Example test_xth_error2 : xth_error [1];[2] 1 = Some 2. Proof. simpl. reflexivity. Qed.

Example test_xth_error3 : xth_error true 2 = None. Proof. simpl. reflexivity. Qed.

```

Definition hd_error {X : Type} (l : list X) : option X := match l with | [] => None | h
:: t => Some h end.

```

Check @hd_error.

Example test_hd_error1 : hd_error 1;2 = Some 1. Proof. simpl. reflexivity. Qed.

Example test_hd_error2 : hd_error [1];[2] = Some 1. Proof. simpl. reflexivity. Qed.

Funções

```

Definition doit3times {X:Type} (f:X→X) (n:X) : X :=
f (f (f n)).

```

Check @doit3times.

Example test_doit3times': doit3times negb true = false.

```

Fixpoint filter {X : Type} (test : X → bool) (l : list X) : (list X) :=
match l with
| [] => []
| h :: t => if test h then h :: (filter test t)
            else filter test t
end.

```

Example test_filter1: filter evenb [1;2;3;4] = [2;4].

Definition length_is_1 {X : Type} (l : list X) : bool := eqb (length l) 1.

Example test_filter2: filter length_is_1 [[1; 2]; [3]; [4]; [5;6;7]; []; [8]] = [[3]; [4]; [8]].

```

Definition countoddmembers' (l:list nat) : nat :=
length (filter oddb l).

```

Example test_countoddmembers'1: countoddmembers' [1;0;3;1;4;5] = 4.

Example test_countoddmembers'2: countoddmembers' [0;2;4] = 0.

Example test_countoddmembers'3: countoddmembers' nil = 0.

Example test_anon_fun': doit3times (fun n => n × n) 2 = 256.

Example test_filter2': filter (fun l => eqb (length l) 1) [[1; 2]; [3]; [4]; [5;6;7]; []; [8]] = [[3]; [4]; [8]].

Exercise: 2 stars (filter_even_gt7)

Definition *filter_even_gt7* ($l : \text{list nat}$) : $\text{list nat} := \text{filter } (\text{fun } n \Rightarrow \text{andb } (\text{evenb } n) (\text{negb } (\text{leb } n 6))) l$.

Example *test_filter_even_gt7_1* : *filter_even_gt7* [1;2;6;9;10;3;12;8] = [10;12;8].

Example *test_filter_even_gt7_2* : *filter_even_gt7* [5;2;6;19;129] = [].

Exercise: 3 stars (partition)

Definition *partition* { $X : \text{Type}$ } ($\text{test} : X \rightarrow \text{bool}$) ($l : \text{list } X$) : $\text{list } X \times \text{list } X := (\text{filter } \text{test } l, \text{filter } (\text{fun } n \Rightarrow \text{negb } (\text{test } n)) l)$.

Example *test_partition1* : *partition oddb* [1;2;3;4;5] = ([1;3;5], [2;4]).

Example *test_partition2* : *partition* ($\text{fun } x \Rightarrow \text{false}$) [5;9;0] = ([], [5;9;0]).

Fixpoint *map* { $X \ Y : \text{Type}$ } ($f : X \rightarrow Y$) ($l : \text{list } X$) : ($\text{list } Y$) :=
 match l with
 | [] \Rightarrow []
 | $h :: t \Rightarrow (f \ h) :: (\text{map } f \ t)$
 end.

Example *test_map1* : *map* ($\text{fun } x \Rightarrow \text{plus } 3 \ x$) [2;0;2] = [5;3;5].

Example *test_map2* : *map oddb* [2;1;2;5] = [false;true;false;true].

Example *test_map3* : *map* ($\text{fun } n \Rightarrow [\text{evenb } n; \text{oddb } n]$) [2;1;2;5] = [[true;false];[false;true];[true;false];[false;true]].

Exercise: 3 stars (map_rev)

Theorem *map_list_eq_map_app* :

$\forall (X \ Y : \text{Type})$
 ($f : X \rightarrow Y$)
 ($l1 \ l2 : \text{list } X$),
 $\text{map } f (l1 ++ l2) = \text{map } f \ l1 ++ \text{map } f \ l2$.

Theorem *map_rev* : $\forall (X \ Y : \text{Type}) (f : X \rightarrow Y) (l : \text{list } X),$
 $\text{map } f (\text{rev } l) = \text{rev } (\text{map } f \ l)$.

Exercise: 2 stars, recommended (flat_map)

Fixpoint *flat_map* { $X \ Y : \text{Type}$ } ($f : X \rightarrow \text{list } Y$) ($l : \text{list } X$)
 : ($\text{list } Y$) :=

 match l with
 | nil \Rightarrow nil
 | $h :: t \Rightarrow f \ h ++ \text{flat_map } f \ t$
 end.

Example *test_flat_map1* : *flat_map* ($\text{fun } n \Rightarrow [n;n;n]$) [1;5;4] = [1; 1; 1; 5; 5; 5; 4; 4; 4].

Exercise: 2 stars, optional (implicit_args)

Fixpoint *filter'* ($X : \text{Type}$) ($\text{test} : X \rightarrow \text{bool}$) ($l : \text{list } X$)
 : ($\text{list } X$) :=

```

match l with
| [] ⇒ []
| h :: t ⇒ if test h then h :: (filter' - test t)
            else filter' - test t
end.

Fixpoint map' (X Y : Type) (f : X → Y) (l : list X)
           : (list Y) :=
  match l with
  | nil ⇒ nil
  | h :: t ⇒ f h :: map' _ _ f t
  end.

Fixpoint fold {X Y : Type} (f : X → Y → Y) (l : list X) (b : Y)
           : Y :=
  match l with
  | nil ⇒ b
  | h :: t ⇒ f h (fold f t b)
  end.

Example fold_example1 : fold mult [1;2;3;4] 1 = 24.
Example fold_example2 : fold andb [true;true;false;true] true = false.
Example fold_example3 : fold app [[1];[];[2;3];[4]] [] = [1;2;3;4].
  Exercise: 1 star, advanced (fold_types_different)

Definition constfun {X : Type} (x : X) : nat → X :=
  fun (k : nat) ⇒ x.

Definition ftrue := constfun true.

Example constfun_example1 : ftrue 0 = true.
Example constfun_example2 : (constfun 5) 99 = 5.

Check plus.

Definition plus3 := plus 3.

Check plus3.

Example test_plus3 : plus3 4 = 7.
Example test_plus3' : doit3times plus3 0 = 9.
Example test_plus3'' : doit3times (plus 3) 0 = 9.
  Exercise: 2 stars (fold_length)

Definition fold_length {X : Type} (l : list X) : nat :=
  fold (fun _ n ⇒ S n) l 0.

Example test_fold_length1 : fold_length [4;7;0] = 3.

```

Prova da corretude de *fold_length*.

Theorem *fold_length_correct* : $\forall X (l : \text{list } X),$
fold_length *l* = *length* *l*.

Exercise: 3 stars (*fold_map*)

Definition *fold_map* {*X Y* : **Type**} (*f* : *X* \rightarrow *Y*) (*l* : *list X*) : *list Y* :=
fold (*fun next acc* \Rightarrow *f next* :: *acc*) *l* [].

Theorem *fold_map_correct* : $\forall (X Y : \text{Type}) (f : X \rightarrow Y) (l : \text{list } X),$
fold_map *f* *l* = *map* *f* *l*.

Exercise: 2 stars, advanced (*currying*)

Definition *prod_curry* {*X Y Z* : **Type**}
(*f* : *X* \times *Y* \rightarrow *Z*) (*x* : *X*) (*y* : *Y*) : *Z* := *f* (*x*, *y*).

Definition *prod_uncurry* {*X Y Z* : **Type**}
(*f* : *X* \rightarrow *Y* \rightarrow *Z*) (*p* : *X* \times *Y*) : *Z* := *f* (*fst* *p*) (*snd* *p*).

Example *test_map1'* : *map* (*plus* 3) [2;0;2] = [5;3;5].

Check @*prod_curry*.

Check @*prod_uncurry*.

Theorem *uncurry_curry* : $\forall (X Y Z : \text{Type}) (f : X \rightarrow Y \rightarrow Z) x y,$
prod_curry (*prod_uncurry* *f*) *x* *y* = *f* *x* *y*.

Theorem *curry_uncurry* : $\forall (X Y Z : \text{Type}) (f : (X \times Y) \rightarrow Z) (p : X \times Y),$
prod_uncurry (*prod_curry* *f*) *p* = *f* *p*.

Exercise: 4 stars, advanced (*church-numerals*)

Module *Church*.

Definition *cnat* := $\forall X : \text{Type}, (X \rightarrow X) \rightarrow X \rightarrow X.$

Definition *zero* : *cnat* :=
fun (*X* : **Type**) (*f* : *X* \rightarrow *X*) (*x* : *X*) \Rightarrow *x*.

Definition *one* : *cnat* :=
fun (*X* : **Type**) (*f* : *X* \rightarrow *X*) (*x* : *X*) \Rightarrow *f* *x*.

Definition *two* : *cnat* :=
fun (*X* : **Type**) (*f* : *X* \rightarrow *X*) (*x* : *X*) \Rightarrow *f* (*f* *x*).

Definition *three* : *cnat* := @*doit3times*.

Exercise: 1 star, advanced (*church_succ*)

Definition *succ* (*n* : *cnat*) : *cnat* :=
fun (*X* : **Type**) (*f* : *X* \rightarrow *X*) (*x* : *X*) \Rightarrow *f* (*n* *X* *f* *x*).

Example *succ_1* : *succ* *zero* = *one*.

Example *succ_2* : *succ* *one* = *two*.

Example *succ_3* : *succ two = three*.

Example *succ_4* : *succ (succ two) = succ three*.

Exercise: 1 star, advanced (*church_plus*)

Definition *plus* (*n m* : *cnat*) : *cnat* :=

fun (*X* : Type) (*f* : *X* → *X*) (*x* : *X*) ⇒ *n X f* (*m X f x*).

Example *plus_1* : *plus zero one = one*.

Example *plus_2* : *plus two three = plus three two*.

Example *plus_3* : *plus (plus two two) three = plus one (plus three three)*.

Exercise: 2 stars, advanced (*church_mult*)

Definition *mult* (*n m* : *cnat*) : *cnat* :=

fun (*X* : Type) (*f* : *X* → *X*) (*x* : *X*) ⇒ *n X* (*m X f*) *x*.

Example *mult_1* : *mult one one = one*.

Example *mult_2* : *mult zero (plus three three) = zero*.

Example *mult_3* : *mult two three = plus three three*.

Erro

Definition *exp* (*n m* : *cnat*) : *cnat* := fun (*X* : Type) (*f* : *X* -> *X*) (*x* : *X*) => (*m* (*X* -> *X*) (fun *y* => (fun *z* => (*n X y z*))) *f*) *x*.

Check (*exp three two*) (@list bool) (map negb).

Example *exp_1* : *exp two two = plus two two*. Proof. reflexivity. Qed.

Example *exp_2* : *exp three two = plus (mult two (mult two two)) one*. Proof. reflexivity. Qed.

Example *exp_3* : *exp three zero = one*. Proof. reflexivity. Qed.

End *Church*.