



ECE 250 Algorithms and Data Structures

# Lab 1: Linked Lists

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# Outline

- List and Node Classes
- The Constructor
- Simple Function Definitions
- Push Front Definition
- Pop Front Definition
- Stepping through Linked Lists
- The Destructor
- Assignment and Copy Constructor
- Summary

# List and Node Classes

- A linked list is a container where each object is stored in separate *node*
- Each node must, in addition to storing an object, must also include a reference/pointer to the next node
  - In C++, this is the address of the next node

# List and Node Classes

- We must dynamically create the nodes in a linked list
- Thus, because **new** returns a pointer, the logical manner in which to track a linked lists is through a pointer
- A **Node** class must store the data object and a reference to the next node (also a pointer)

# The Node Class

- The node must store an **object** and a **pointer**:
  - In this case, the object is an **integer**

```
class Node {  
    private:  
        int    element;  
        Node *next_node;  
    public:  
        Node( int = 0, Node * = 0 );  
  
        int retrieve() const;  
        Node *next() const;  
};
```

- In Project 1, the type is given by the template:  
**Object element;**

# The Node Class

- The constructor assigns the two member variables based on the arguments

```
Node::Node( int e, Node *n ):element( e ), next_node( n ) {  
    // empty constructor  
}
```

- The default values for both arguments are both **0**

# The Node Class

- The two member functions are accessors which simply return the **element** and the **next\_node** member variables, respectively

```
int Node::retrieve() const {  
    return element;  
}
```

```
Node *Node::next() const {  
    return next_node;  
}
```

# Project 1 – Creating your VS Project

- Create a VS project name P1( type is c++, General)
- Add source code provided for P1, in course web page.
  - Start by solving part a) – Single List
  - Files needed in the VS project :

- Ece250.h
- Exception.h
- Tester.h



From  
Common  
Source Files

- Single\_list.h
- Single\_node.h
- Single\_node\_tester.h
- Single\_list\_tester.h
- Single\_list\_driver.cpp



# Project 1 – Creating your VS Project

- Compile your VS Project as given. It should compile without errors.
- Complete class `Single_node.h`
- Questions about C++ syntax so far?
  - Template classes
  - Friend classes
  - Constructor , initializing values

# The List Class

- Because each node in a linked list refers to the next, the linked list class needs to store the address of the first node
  - This requires one member variable: a pointer to a node

```
class List {  
    private:  
        Node *list_head;  
        // ...  
};
```

# The List Class

- This is the class definition:

```
class List {  
    private:  
        Node *list_head;  
    public:  
        List();  
  
        // Accessors  
        bool empty() const;  
        int front() const;  
        Node * head() const;  
        int count( int ) const;  
  
        // Mutators  
        void push_front( int );  
        int pop_front();  
        int erase( int );  
};
```

# The List Class

- This is the class in your project 1 – very similar but not the same:

```
template <typename Object>
class Single_list {
private:
    Single_node<Object> *list_head;
    Single_node<Object> *list_tail;
    int node_count;
```

# Linked Lists in Memory

- To begin, let us look at the internal representation of a linked list
- Suppose we want a linked list to store the values

**42**   **95**   **70**   **81**

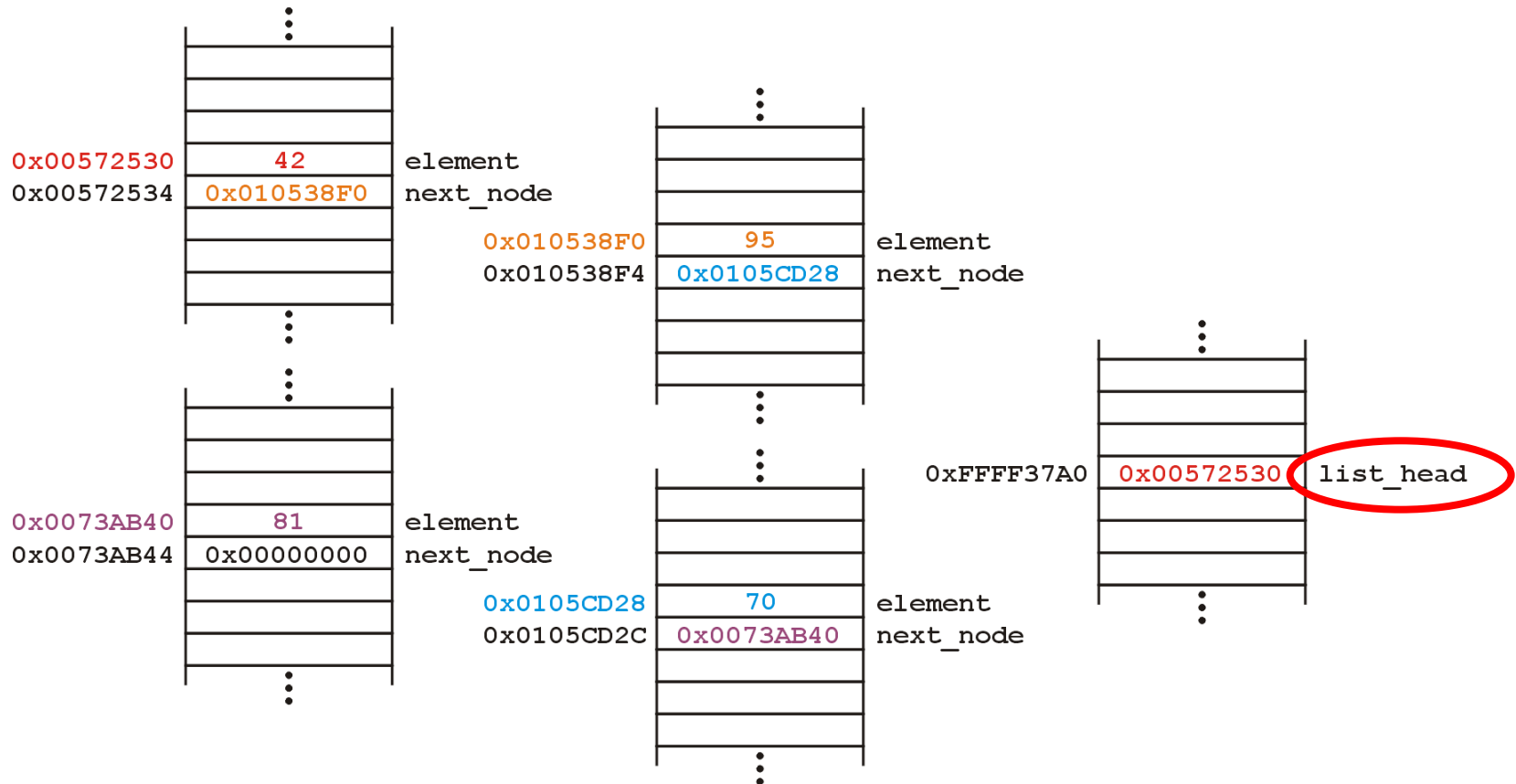
in this order

- This could be the result of:

```
int main() {  
    List ls;  
    ls.push_front( 81 );  
    ls.push_front( 70 );  
    ls.push_front( 95 );  
    ls.push_front( 42 );  
    return 0;  
}
```

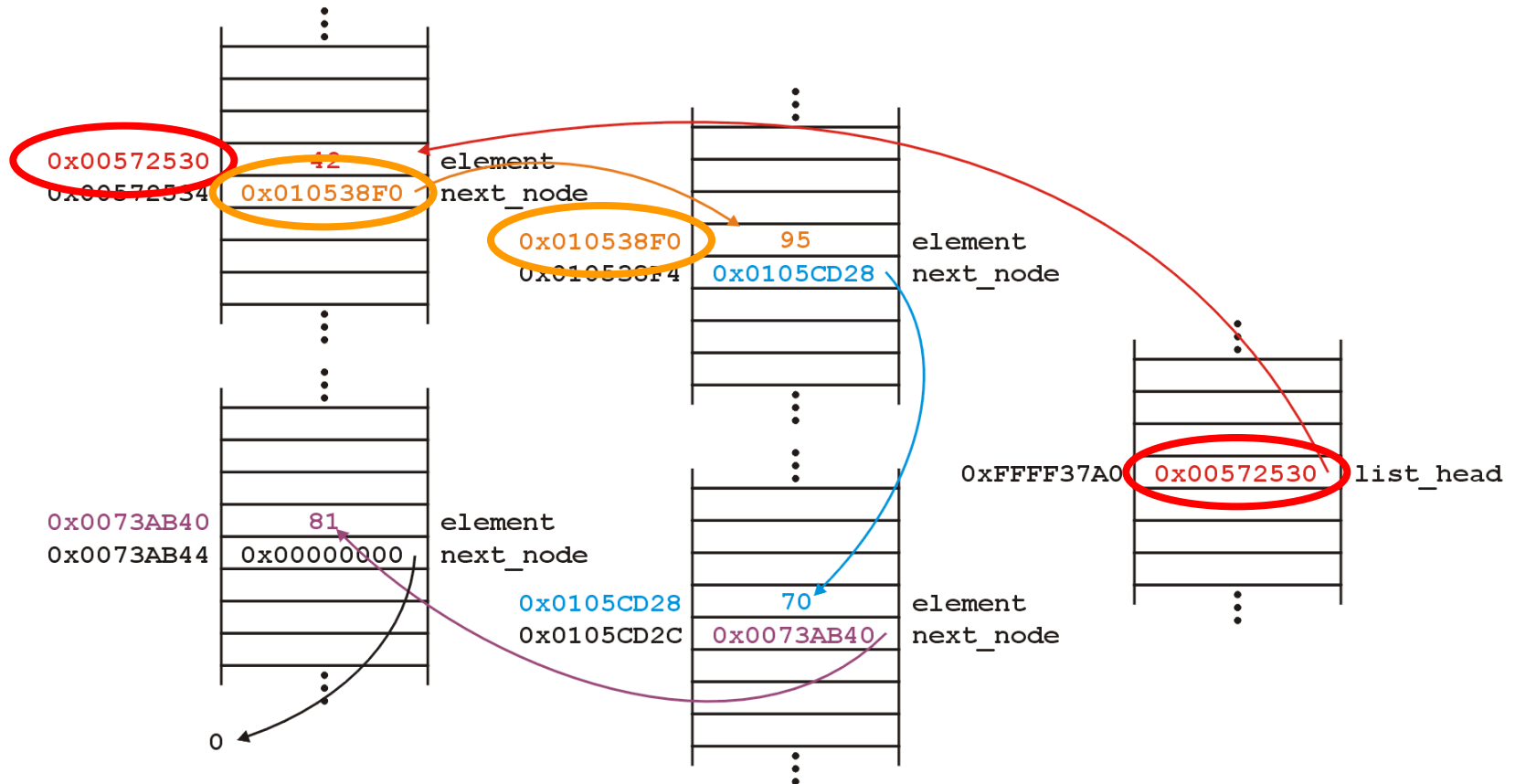
# Linked Lists in Memory

- These nodes could be stored anywhere in memory:
  - For example,



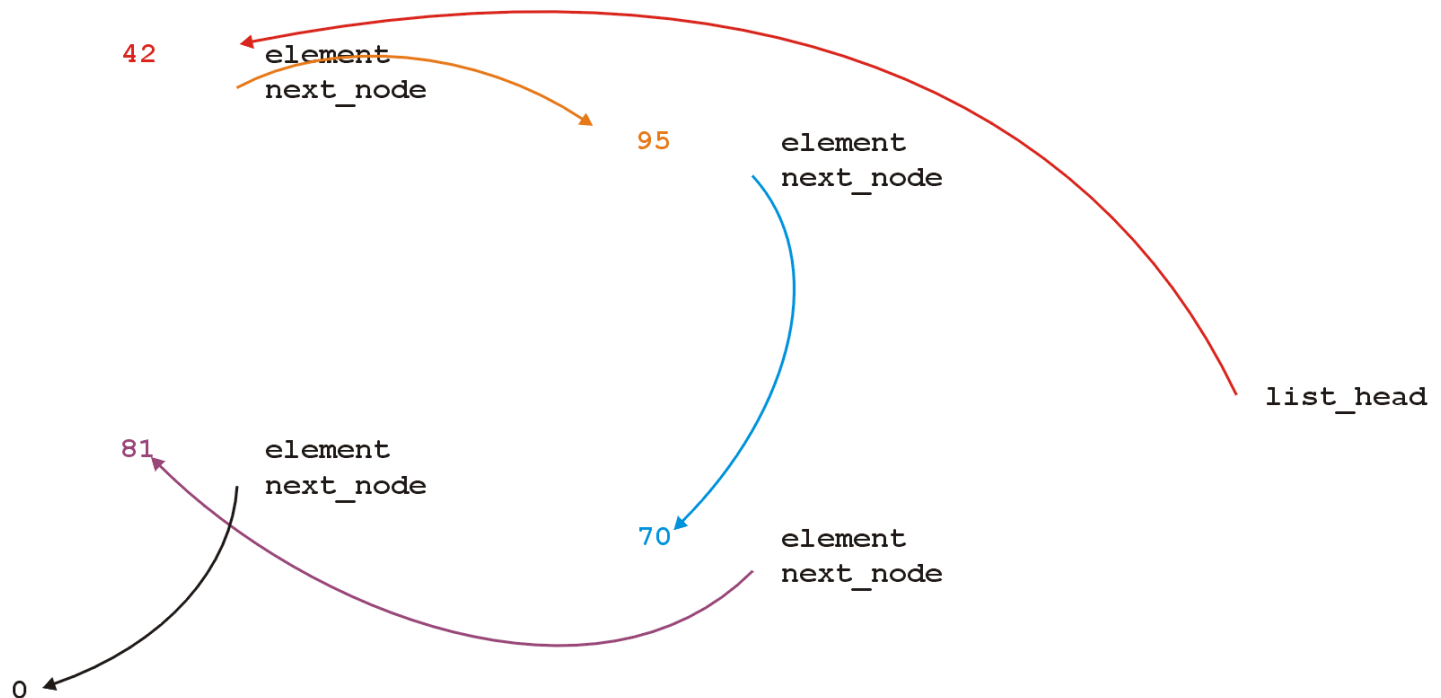
# Linked Lists in Memory

- Notice how each `next_node` address points to the next node in the list?



# Linked Lists in Memory

- Because the addresses are arbitrary, we can remove that information:





# Linked Lists in Memory

- We will clean up the representation as follows:



- We do not specify the addresses because they are arbitrary and:
  - The contents of the circle is the stored object
  - The **next\_node** pointer is represented by an arrow

# Operations on Linked Lists

- First, we want to create a linked list
- We also want to be able to:
  - Insert `void push_front( int );`
  - Access `int front() const;`  
`Node *head() const;`
  - Remove `int pop_front();`the front object stored in the linked list
- Operations on the entire linked list may include:
  - Membership `int count( int ) const;`
  - Erase an object `int erase( int );`
- Additionally, we may wish to check the state: is the linked list empty?  
`bool empty() const;`

# The Constructor

- The constructor initializes the linked list
- We do not count how many objects are in this list, thus:
  - We must rely on the last pointer in the linked list to point to a special value
  - In C++, we use `0` (the zero address)

# The Constructor

- Thus, in the constructor, we assign `0` to `list_head`:

```
List::List():list_head( 0 ) {  
    // empty constructor  
}
```

- We will always ensure that when a linked list is empty, the list head is assigned `0`

# Simple Function Definitions

- Starting with the easier member functions:

```
bool List::empty() const {  
    if ( list_head == 0 ) {  
        return true;  
    } else {  
        return false;  
    }  
}
```

or better yet:

```
bool List::empty() const {  
    return ( list_head == 0 );  
}
```

# Simple Function Definitions

- The member function `Node *head() const` is easy enough to implement:

```
Node *List::head() const {  
    return list_head;  
}
```

- This will always work: if the list is empty, it will return `0`

# Simple Function Definitions

- To get the first element in the linked list, we must access the node to which the `list_head` is pointing
- Because we have a pointer, we must use the `->` operator to call the member function:

```
int List::front() const {  
    return head()->retrieve();  
}
```

# Simple Function Definitions

- The member function `int front() const` requires some additional consideration, however:
  - What if the list is empty?
- If we tried to access a member function of a pointer set to `0`, we would access restricted memory
- The operating system would terminate the running program



# Simple Function Definitions

- Instead, we can use an exception handling mechanism where we throw an exception
- We define a class

```
class underflow {  
    // empty  
};
```

and then we *throw* an instance of this class:

```
throw underflow();
```

# Simple Function Definitions

- Thus, the full function is

```
int List::front() const {  
    if ( empty() ) {  
        throw underflow();  
    }  
  
    return head()->retrieve();  
}
```

# Simple Function Definitions

- Why is `empty()` better than

```
int List::front() const {  
    if ( list_head == 0 ) {  
        throw underflow();  
    }  
  
    return head()->retrieve();  
}
```

?

- Two benefits:
  - more readable
  - if the implementation changes we do nothing

# Push Front Definition

- Next, let us add an element to the list
- If it is empty, we start with:

`list_head`  $\longrightarrow$  0

and, if we try to add 81, we should end up with:

`list_head`  $\longrightarrow$  (81)  $\longrightarrow$  0

# Push Front Definition

- To visualize what we must do:
  - We must create a new node which stores:
    - The value **81**, and
    - The address (pointer) **0**
  - We must then assign its address to **list\_head**
- We can do this as follows:

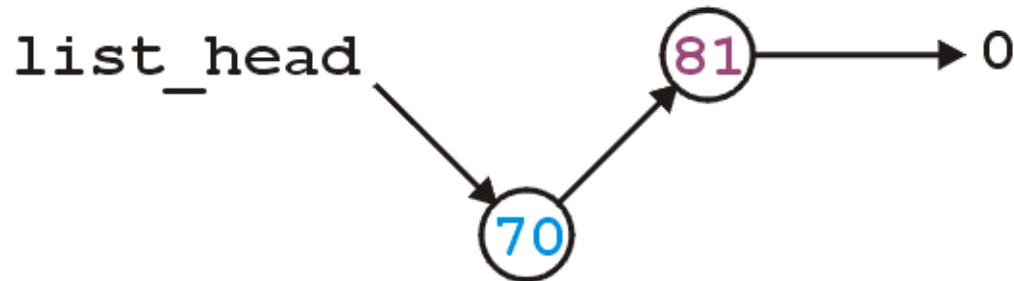
```
list_head = new Node( 81, 0 );
```
- We could also use the default value...

# Push Front Definition

- Suppose however, we already have a non-empty list



- Adding 70, we want:



# Push Front Definition

- To achieve this, we must we must create a new node which stores:
  - The value **70**, and
  - The address of the current list head
    - We must then assign its address to **list\_head**
- We can do this as follows:

```
list_head = new Node( 70, list_head );
```

# Push Front Definition

- Thus, our implementation could be:

```
void List::push_front( int n ) {  
    if ( empty() ) {  
        list_head = new Node( n, 0 );  
    } else {  
        list_head = new Node( n, list_head );  
    }  
}
```



# Push Front Definition

- We could, however, note that when the list is empty, `list_head == 0`, thus we could shorten this to:

```
void List::push_front( int n ) {  
    list_head = new Node( n, list_head );  
}
```

# Push Front Definition

- Are we allowed to do this?

```
void List::push_front( int n ) {  
    list_head = new Node( n, list_head );  
}
```

- Yes: the right-hand side of an assignment statement is evaluated first
- The original value of `list_head` is accessed first before the function call is made

# Pop Front Definition

- Removing from the linked list may appear easier:
  - We simply assign the list head to the next pointer of the first node
- Graphically, given:



we want:



# Pop Front Definition

- Easy enough:

```
int List::pop_front() {  
    list_head = list_head->next();  
}
```

- Unfortunately, we have some problems:
  - We want to return the item stored
  - The list may be empty
  - We still have the memory allocated for the node containing 70

# Pop Front Definition

- We could try:

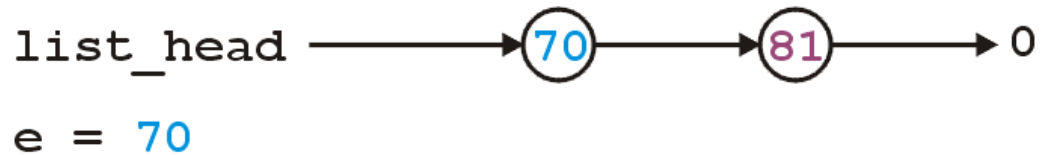
```
int List::pop_front() {  
    if ( empty() ) {  
        throw underflow();  
    }  
  
    int e = front();  
    delete list_head;  
    list_head = list_head -> next();  
    return e;  
}
```

- Let's examine this...

# Pop Front Definition

```
int List::pop_front() {  
    if ( empty() ) {  
        throw underflow();  
    }  
}
```

```
int e = front();
```



```
e = 70
```

```
delete list_head;
```

```
list_head = list_head->next();
```

```
return e;
```

```
}
```

# Pop Front Definition

```
int List::pop_front() {  
    if ( empty() ) {  
        throw underflow();  
    }
```

```
    int e = front();
```

```
    delete list_head;
```

```
    list_head
```

```
    e = 70
```

```
    list_head = list_head->next();
```

```
    return e;
```

```
}
```



# Pop Front Definition

```
int List::pop_front() {  
    if ( empty() ) {  
        throw underflow();  
    }
```

```
    int e = front();
```

```
    delete list_head;
```

```
    list_head = list_head->next();
```

```
    return e;
```

```
}
```

list\_head

e = 70





# Pop Front Definition

- The problem is, we are accessing a node which we have just deleted
- Unfortunately, this will work more than 99% of the time:
  - the running program (process) may still own the memory
- Once in a while it will fail ...
  - ... and it will be almost impossible to debug

# Pop Front Definition

- The correct implementation assigns a temporary pointer to point to the node being deleted:

```
int List::pop_front() {  
    if ( empty() ) {  
        throw underflow();  
    }  
  
    int e = front();  
    Node *ptr = list_head;  
    list_head = list_head->next();  
    delete ptr;  
    return e;  
}
```

# Pop Front Definition

```
int List::pop_front() {  
    if ( empty() ) {  
        throw underflow();  
    }  
}
```

```
int e = front();
```

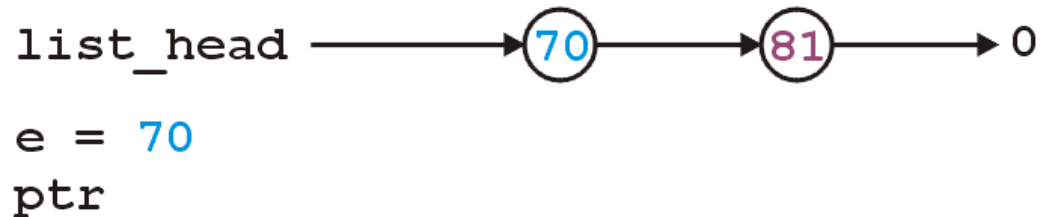
```
Node *ptr = list_head;
```

```
list_head = list_head->next();
```

```
delete ptr;
```

```
return e;
```

```
}
```



# Pop Front Definition

```
int List::pop_front() {  
    if ( empty() ) {  
        throw underflow();  
    }
```

```
    int e = front();
```

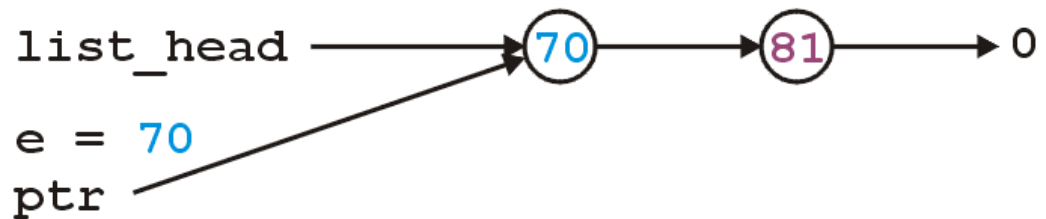
```
    Node *ptr = list_head;
```

```
    list_head = list_head->next();
```

```
    delete ptr;
```

```
    return e;
```

```
}
```



# Pop Front Definition

```
int List::pop_front() {  
    if ( empty() ) {  
        throw underflow();  
    }
```

```
    int e = front();
```

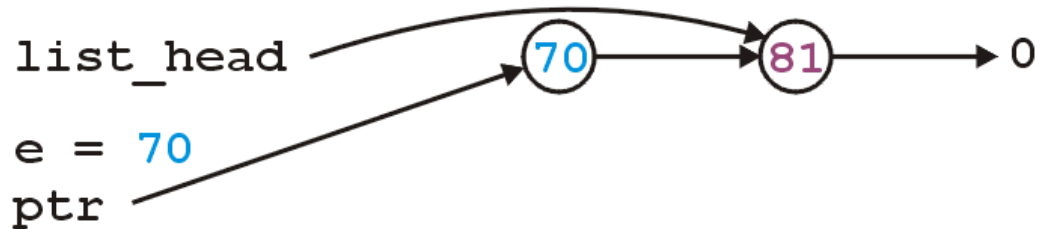
```
    Node *ptr = list_head;
```

```
    list_head = list_head->next();
```

```
    delete ptr;
```

```
    return e;
```

```
}
```



# Pop Front Definition

```
int List::pop_front() {  
    if ( empty() ) {  
        throw underflow();  
    }  
}
```

```
int e = front();
```

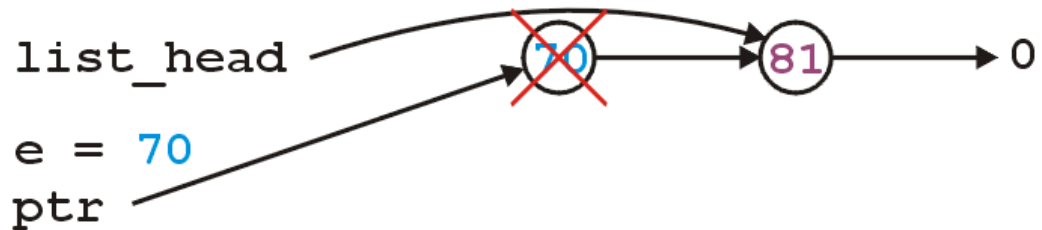
```
Node *ptr = list_head;
```

```
list_head = list_head->next();
```

```
delete ptr;
```

```
return e;
```

```
}
```



## Project 1: Linked List – Accessors and push\_front

- Notice that the constructor is given to you
- Complete accessors in class `Single_list.h` (except `count`)

```
int size() const;
bool empty() const;
Object front() const;
Object back() const;
Single_node<Object> *head() const;
Single_node<Object> *tail() const;
```

- Complete the mutator `push_front`
- Test these functions using commands listed in `Single_list_tester.h`.

# Testing

- Test file with functions implemented so far, ie.

```
new
push_front 3
push_front 7
front 7
back 3
head
  retrieve 7
  next
  retrieve 3
  next0
exit
size 2
empty 0
Exit
```

- Try command “cout” – to list the content of your list
- Notice file int.in. If your code passes this test you have 50 % of the total mark.



# Stepping through Linked Lists

- The next step is to look at member functions which potentially require us to step through the entire list:

```
int count( int ) const;
```

```
int erase( int );
```

- For count, we will return 1 if the object was found and 0 otherwise
- For erase, we return 1 if the object was found and 0 otherwise and we will delete the node.

# Stepping through Linked Lists

- The process of stepping through a linked list can be thought of as being analogous to a for loop:
  - We initialize a temporary pointer with the list head
  - We continue iterating until the pointer equals 0
  - With each step, we set the pointer to point to the next object

# Stepping through Linked Lists

- Thus, we have:

```
for ( Node *ptr = head(); ptr != 0; ptr = ptr->next() ) {  
    // Do something  
    // - use ptr->fn() to call member functions  
    // - use ptr->var to assign/access member variables  
}
```

# Stepping through Linked Lists

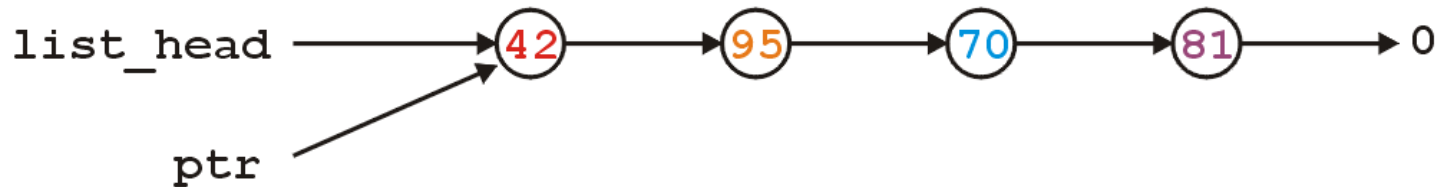
- Analogously:

```
for ( Node *ptr = head(); ptr != 0; ptr = ptr->next() ) {  
    // do something  
    // use ptr->fn() to call member functions  
    // use ptr->var to assign/access member variables  
}
```

```
for ( int i = 0; i < N; ++i ) {  
    // do something  
}
```

# Stepping through Linked Lists

- With the initialization and first iteration of the loop, we have:

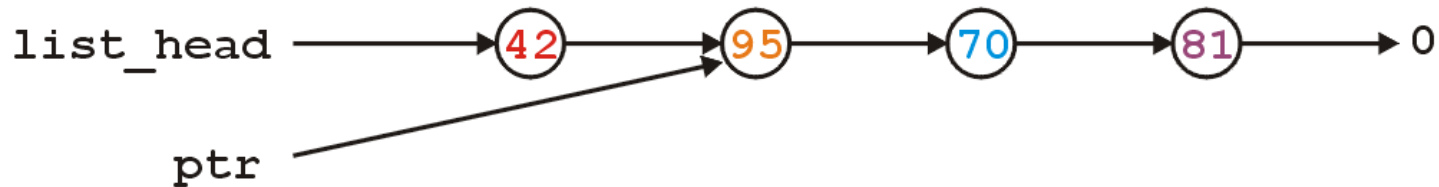


- `ptr != 0` and thus we evaluate the body of the loop and then set `ptr` to the next node: `ptr = ptr->next()`

```
for ( Node *ptr = head(); ptr != 0; ptr = ptr->next() ) {  
    // Do something  
}
```

# Stepping through Linked Lists

- `ptr != 0` and thus we evaluate the loop and increment the pointer

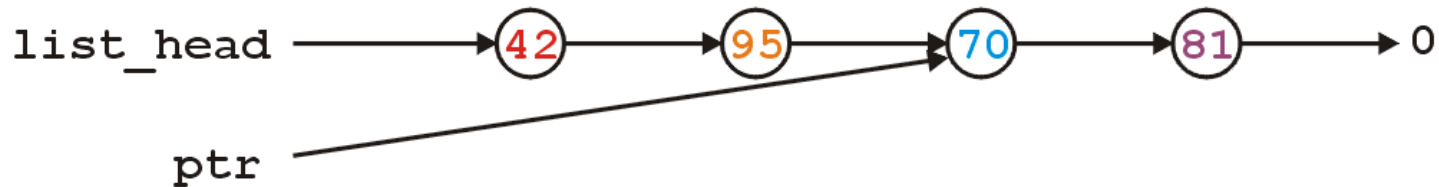


- In the loop, we can access 95 by using `ptr->retrieve()`

```
for ( Node *ptr = head(); ptr != 0; ptr = ptr->next() ) {  
    // Do something  
}
```

# Stepping through Linked Lists

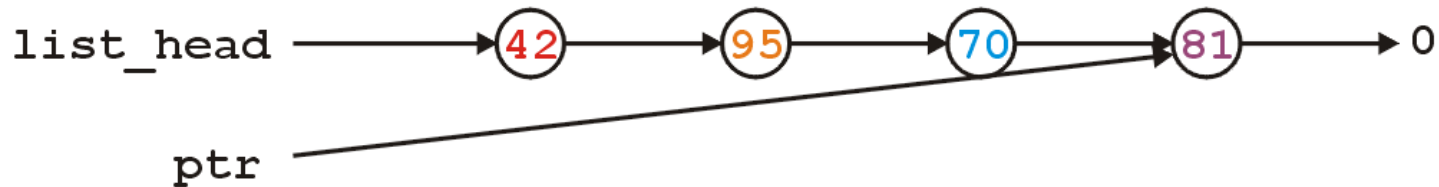
- `ptr != 0` and thus we evaluate the loop and increment the pointer



```
for ( Node *ptr = head(); ptr != 0; ptr = ptr->next() ) {  
    // Do something  
}
```

# Stepping through Linked Lists

- `ptr != 0` and thus we evaluate the loop and increment the pointer

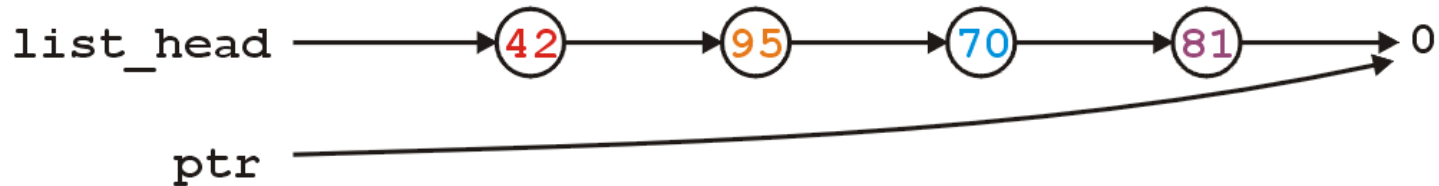


```
for ( Node *ptr = head(); ptr != 0; ptr = ptr->next() ) {  
    // Do something  
}
```



# Stepping through Linked Lists

- Here, we check and find `ptr != 0` is false, and thus we exit the loop



- Because the variable `ptr` was declared inside the loop, we can no longer access it

```
for ( Node *ptr = head(); ptr != 0; ptr = ptr->next() ) {  
    // Do something  
}
```

# Stepping through Linked Lists

- To implement `bool count(int) const;`, we check if the argument matches the retrieved element with each step
  - If we find the argument, we return `true`
  - If we finish the loop, this indicates we did not find the argument: return `false`

# Stepping through Linked Lists

- The implementation:

```
int List::count( int n ) const {  
    for ( Node *ptr = head(); ptr != 0; ptr = ptr->next() ) {  
        if ( ptr->retrieve() == n ) {  
            return 1;  
        }  
    }  
  
    return 0;  
}
```

# Stepping through Linked Lists

- To erase an arbitrary element, i.e., to implement `int erase( int )`, we must update the previous node
- For example, given



if we delete 70, we want to end up with



# Stepping through Linked Lists

- Notice that the `erase` function must modify the member variables of the node prior to the node being removed
- Thus, it must have access to the private instance variable `next_node`
- We could supply the member function

```
void set_next( Node * );
```

however, this would be globally accessible

# Stepping through Linked Lists

- In Java/C#, you could define the `Node` class to be an *inner class*
  - this is an elegant OO solution
- In C++, you explicitly break encapsulation by declaring the class `List` to be a *friend* of the class `Node`:

```
class Node {  
Node * next() const;  
    // ... declaration ...  
    friend class List;  
};
```

# Stepping through Linked Lists

- Now, inside `erase` (a member function of `List`), you can modify all the member variables of any instance of the `Node` class
  - Challenge for `erase`: you will need an additional temporary pointer when removing a node—once you update the linked list, you must delete the removed node

# The Destructor

- We dynamically allocated memory each time we added a new Node into this list
- Suppose we delete a list before it is empty
  - By default, this does not clean up any of those remaining nodes
- Thus, we need to define a destructor:

```
class List {  
    private:  
        Node *list_head;  
    public:  
        List();  
        ~List();  
        // ...etc...  
};
```



# The Destructor

- The destructor has to delete any memory which was allocated with new but not yet deallocated
- This is straight-forward enough:

```
while ( !empty() ) {  
    pop_front();  
}
```
- The overhead of pop\_front is negligible when compared to calling delete

# Project 1: count, erase and destructor

- Complete count
- **Test**
- Complete pop\_front
- Complete erase, and destructor
- **Test**

# Assignment and Copy Constructor

- We now have the functionality of the linked list class
- However, there are other functional operations required for linked lists to interact with C++
- Recall two common C++ operations:
  - Assignment
  - Passing by value (the *copy* constructor)

# Assignment and Copy Constructor

- Suppose you have two objects, be they linked lists or otherwise, for example:

```
List lst1;
```

```
List lst2;
```

```
lst1.push_front( 35 );
```

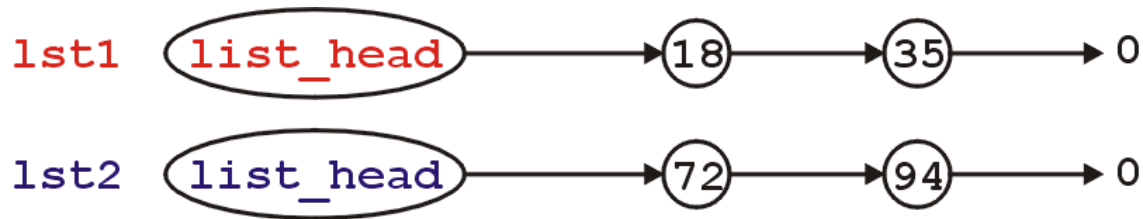
```
lst1.push_front( 18 );
```

```
lst2.push_front( 94 );
```

```
lst2.push_front( 72 );
```

# Assignment and Copy Constructor

- This is the current state:



- If we assign

```
lst2 = lst1;
```

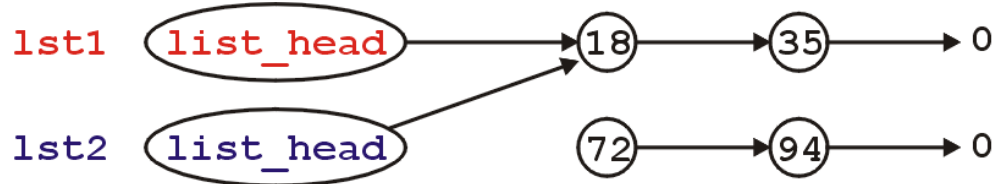
by default, all that occurs is that the member variables from `lst1` are copied to the member variables of `lst2`

# Assignment and Copy Constructor

- Because the only member variable of this class is `list_head`, the value it is storing (the address of the first node) is copied over
- It is equivalent to writing:

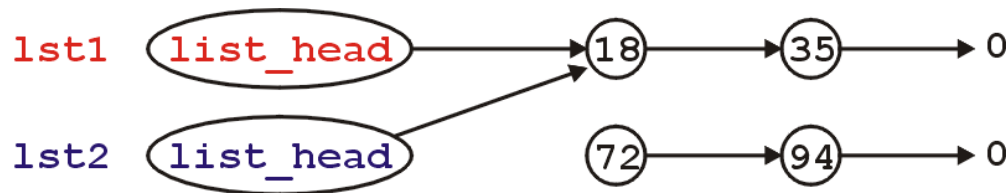
```
lst2.list_head = lst1.list_head;
```

- Graphically:

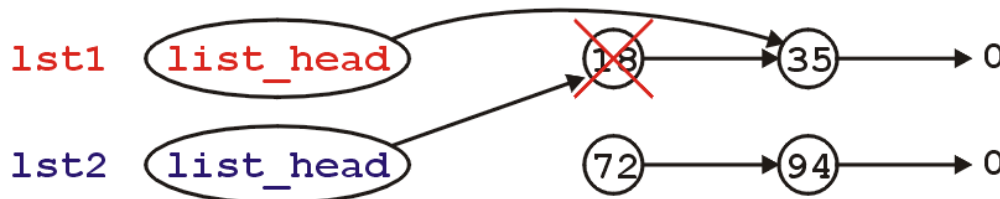


# Assignment and Copy Constructor

- What's wrong with this picture?
  - We no longer have links to either of the nodes storing 72 or 94
  - Also, suppose we call the member function  
`lst1.pop_front();`
  - This only affects the member variable from the object `lst1`
  - Thus, we go from



to



# Assignment and Copy Constructor

- Now, the second list `lst2` is pointing to memory which has been deallocated...
- What is the behavior is we call

```
lst2.pop_front();
```

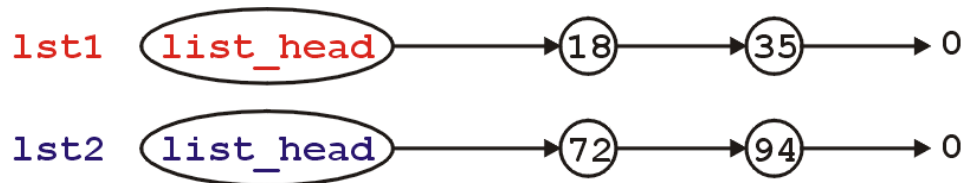
?

- The behaviour is undefined, however, soon this will probably lead to an access violation

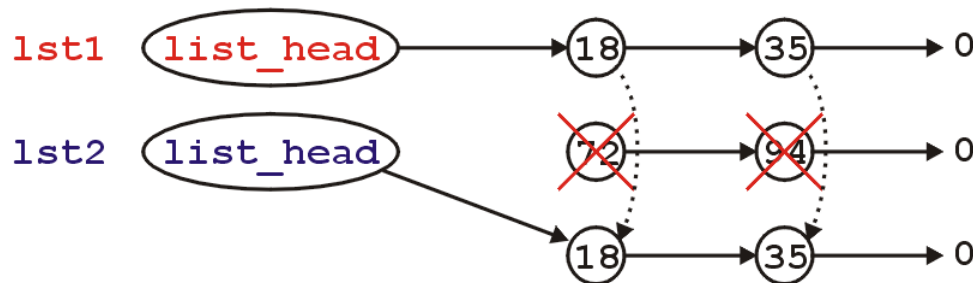


# Assignment and Copy Constructor

- We need to properly implement assignment
- The logical solution would be to:
  - Delete all the nodes in **lst2**
  - Copy all of the nodes in **lst1** to **lst2**
  - Thus, we go from



to



# Assignment and Copy Constructor

- To define assignment, we must overload the function named **operator=**
  - This is how you indicate to the compiler that you are overloading the **=** operator
- Its signature is:  
`List &operator=( List const & );`
- To avoid copies, the argument and return value are passed by reference  
`List &operator=( const List &rhs );`
- Because it is a reference, that means that you must refer use, for example, `rhs.head()` Or `rhs.push_front()`

# Assignment and Copy Constructor

- Because we are overwriting the current object, we must:

```
List &operator=( List const &rhs ) {  
    // clean up this object  
    // make a complete copy of rhs  
  
    return *this;  
}
```

- Here, the name of the argument `rhs` refers to the right-hand side of the assignment

# Assignment and Copy Constructor

- To clean up the current linked list, we need to remove all nodes in this linked list:

```
List &operator=( List const &rhs ) {  
    while ( !empty() ) {  
        pop_front();  
    }  
    // Make a complete copy of rhs  
  
    return *this;  
}
```

# Assignment and Copy Constructor

- However, we have a problem: consider

```
List a;
```

```
a.push_front( 3 );
```

```
a = a;
```

```
List &operator= ( const List &rhs ) {  
    while ( !empty() ) {  
        pop_front();  
    }  
    // Both rhs and *this refer to the same  
    // linked list which is now empty  
    // Make a complete copy of rhs  
  
    return *this;  
}
```

# Assignment and Copy Constructor

- To avoid this problem, we must always make the following check:

```
List &operator= ( const List &rhs ) {  
    if ( this == &rhs ) {  
        return *this;           // No copy is necessary  
    }  
    while ( !empty() ) {  
        pop_front();  
    }  
    // Make a complete copy of rhs  
  
    return *this;  
}
```

# Assignment and Copy Constructor

- To make a copy of the right hand side is a little more difficult...
- We must step through the second linked list and create a new node for each object in the `rhs` list

# Assignment and Copy Constructor

```
List &List::operator= ( const List &rhs ) {
    while ( !empty() ) {
        pop_front();
    }

    if ( rhs.empty() ) {
        return * this;
    }

    list_head = new Node( rhs.front() ); // copy the front element

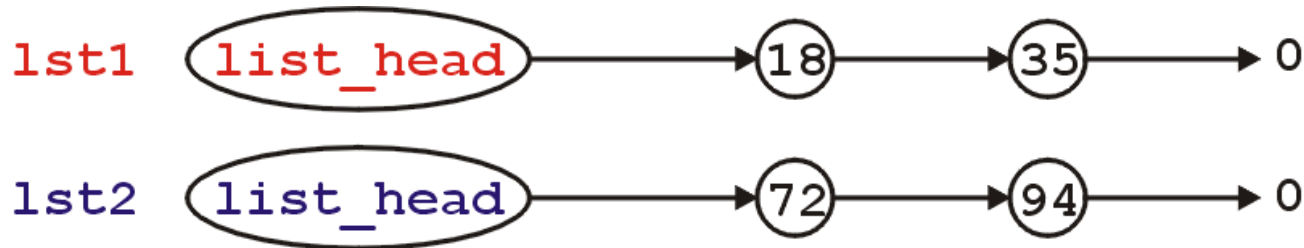
    for ( Node * lhptr = head(), * rhptr = rhs.head();
          rhptr -> next() != 0;
          rhptr = rhptr -> next(), lhptr = lhptr -> next()
    ) {
        lhptr -> next_node = new Node( rhptr -> next() -> retrieve() );
    }

    return * this;
}
```



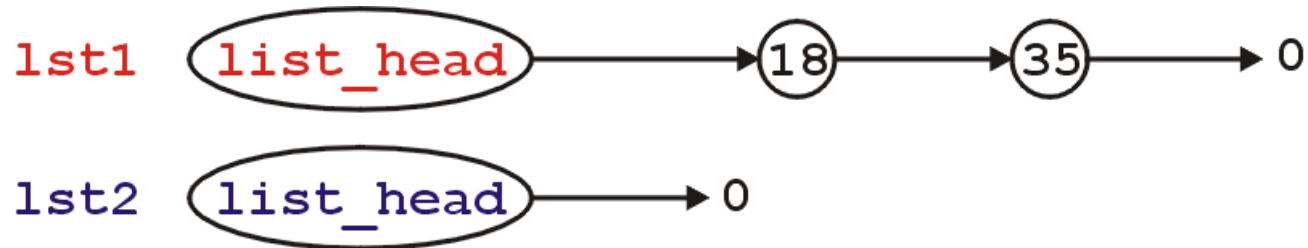
# Assignment and Copy Constructor

- Visually, we are doing the following:
- In assigning `lst2 = lst1;`



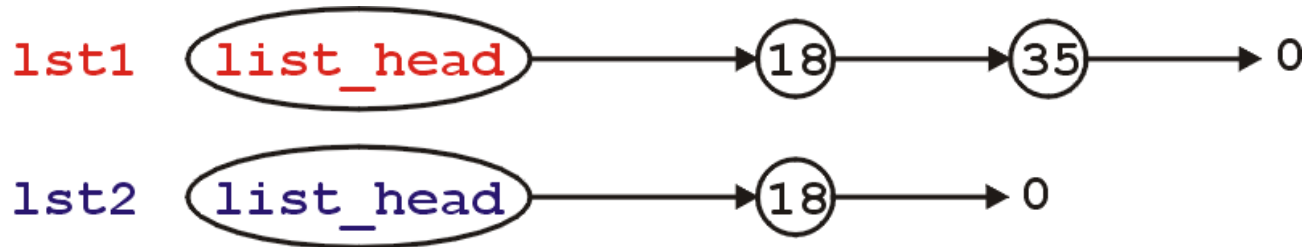
# Assignment and Copy Constructor

- Empty the second list:



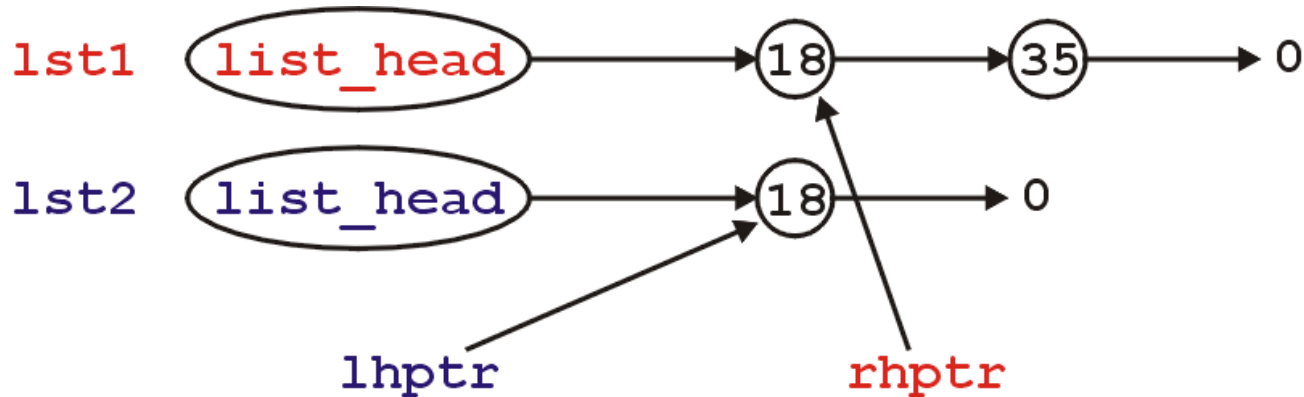
# Assignment and Copy Constructor

- If the 1<sup>st</sup> list is not empty, push the front element of the 1<sup>st</sup> list onto the 2<sup>nd</sup> list:



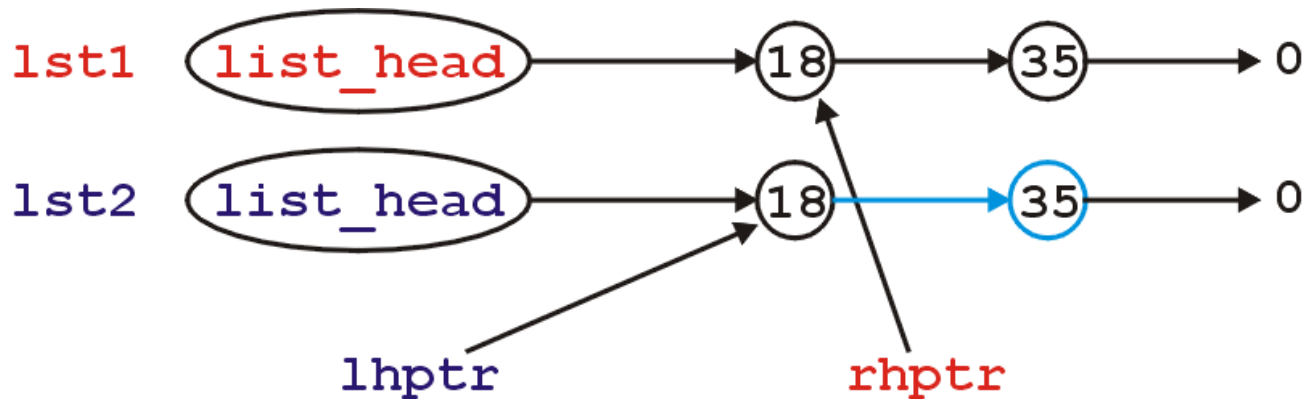
# Assignment and Copy Constructor

- Get two pointers pointing to the first nodes in each of the lists



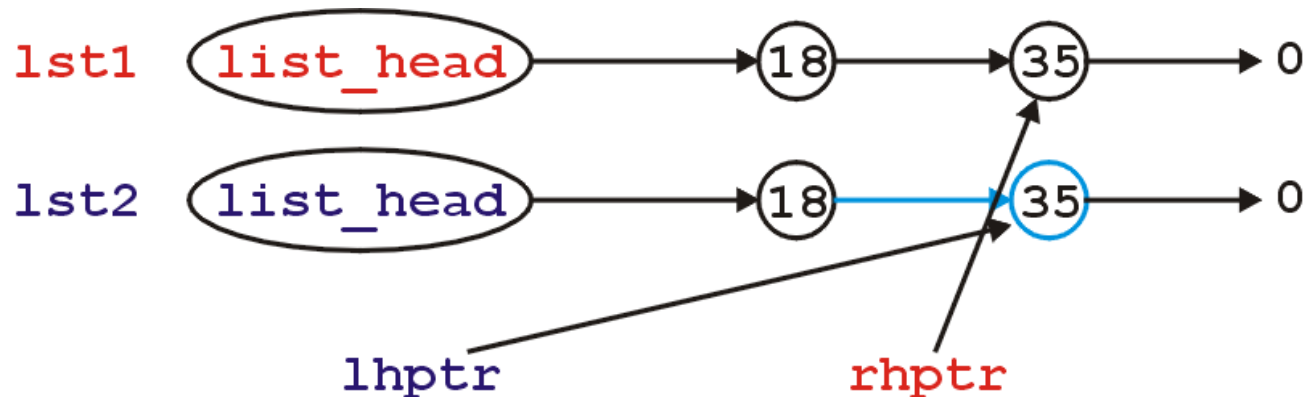
# Assignment and Copy Constructor

- While there is a node following the node pointed to by **rhptr**, add that node to the 2<sup>nd</sup> list:



# Assignment and Copy Constructor

- Then, increment each of the pointers and continue repeating this procedure until there are no more nodes to add



# Assignment and Copy Constructor

- Because in Project 1, you have a tail pointer, you can actually do this a lot more easily
- Like the member function, step through the first linked list, however, just call `push_back( ptr -> retrieve() )` with node in the list

# Assignment and Copy Constructor

- Why do we return `*this`?
- Suppose you have the assignment:  
`lst3 = lst2 = lst1;`
- The assignment operator is right associative: it first evaluates `lst2 = lst1;`
- After this, it must assign `lst3 = ?;`
- `?` is whatever is returned by the assignment operation of `lst2 = lst1;`



# Assignment and Copy Constructor

- We have discussed assignment
- We must now look at a similar problem:
  - Making a copy of an already-existing object (Ex1 and Ex2 below)
  - Passing or returning by value
- These are the same problem: an object is being constructed and initialized with the value of another.

Ex1.

```
List ls1;  
ls1.push_front(57);  
  
List ls2( ls1 );
```

Ex2.

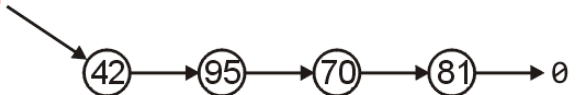
```
List ls1;  
ls1.push_front(57);  
  
List ls2 = ls1;
```

# Assignment and Copy Constructor

- Consider this example where a linked list is passed to a function:
  - Four values are stored in the list

```
List my_marks;  
my_marks.push_front( 81 );  
my_marks.push_front( 70 );  
my_marks.push_front( 95 );  
my_marks.push_front( 42 );
```

```
cout << average( my_marks );
```



```
graph LR; head[my_marks.list_head] --> node42((42)); node42 --> node95((95)); node95 --> node70((70)); node70 --> node81((81)); node81 --> null[0];
```

```
int average( List ls ) {  
  
    int n = 0, sum = 0;  
  
    while ( !ls.empty() ) {  
        sum += ls.pop_front();  
        ++n;  
    }  
  
    return sum/n;  
}
```

# Assignment and Copy Constructor

- By default, pass by value copies values
  - This includes all member variables are copied down including pointers which are just addresses

```
List my_marks;  
my_marks.push_front( 81 );  
my_marks.push_front( 70 );  
my_marks.push_front( 95 );  
my_marks.push_front( 42 );
```

```
cout << average( my_marks );
```

```
int average( List ls ) {
```

```
    int n = 0, sum = 0;
```

```
    while ( !ls.empty() ) {  
        sum += ls.pop_front();  
        ++n;  
    }
```

```
    return sum/n;
```

```
}
```

my\_marks.list\_head



Passed by value

- all member variables passed by value

# Assignment and Copy Constructor

- Now we have two objects which store the same address in the member variable `list_head`

```
List my_marks;
my_marks.push_front( 81 );
my_marks.push_front( 70 );
my_marks.push_front( 95 );
my_marks.push_front( 42 );
```

```
cout << average( my_marks );
```

```
int average( List ls ) {
```

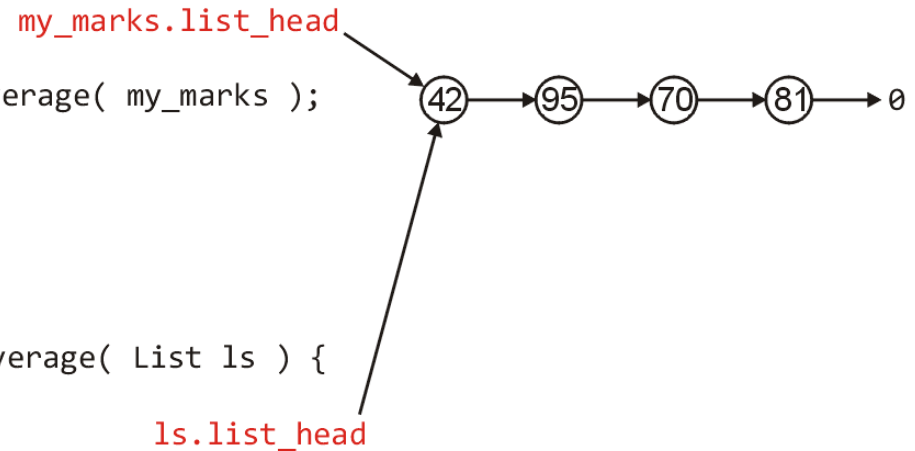
```
    ls.list_head
```

```
    int n = 0, sum = 0;
```

```
    while ( !ls.empty() ) {
        sum += ls.pop_front();
        ++n;
    }
```

```
    return sum/n;
```

```
}
```

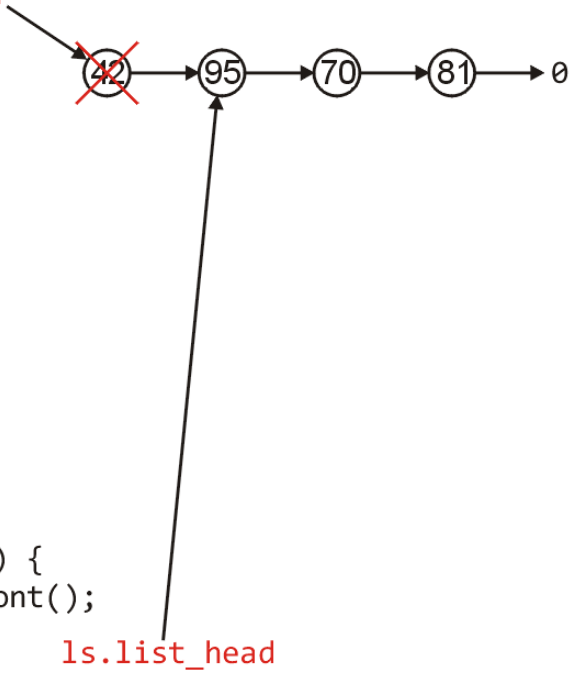


# Assignment and Copy Constructor

- After one iteration of the loop, the pop\_front function is called which deletes the first node

```
List my_marks;
my_marks.push_front( 81 );
my_marks.push_front( 70 );
my_marks.push_front( 95 );
my_marks.push_front( 42 );
```

```
cout << average( my_marks );
```



The diagram shows a linked list with four nodes containing the values 42, 95, 70, and 81. The first node (42) is marked with a large red 'X' and is being pointed to by a red arrow labeled 'my\_marks.list\_head'. The second node (95) is being pointed to by a red arrow labeled 'ls.list\_head'. The list ends with a null pointer (0).

```
int average( List ls ) {

    int n = 0, sum = 0;

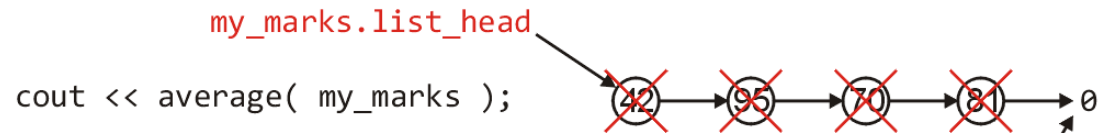
    while ( !ls.empty() ) {
        sum += ls.pop_front();
        ++n;
    }

    return sum/n;
}
```

# Assignment and Copy Constructor

- At the end of the function, all nodes have been deleted

```
List my_marks;  
my_marks.push_front( 81 );  
my_marks.push_front( 70 );  
my_marks.push_front( 95 );  
my_marks.push_front( 42 );
```



```
cout << average( my_marks );
```

```
int average( List ls ) {
```

```
    int n = 0, sum = 0;
```

```
    while ( !ls.empty() ) {  
        sum += ls.pop_front();  
        ++n;  
    }
```

```
    return sum/n;
```

```
}
```


`ls.list_head` points to the first node of the linked list, which is the same as `my_marks.list_head`.

# Assignment and Copy Constructor

- If, in the original function we now attempt to access a node, it no longer exists

```
List my_marks;  
my_marks.push_front( 81 );  
my_marks.push_front( 70 );  
my_marks.push_front( 95 );  
my_marks.push_front( 42 );
```

```
my_marks.list_head  
cout << average( my_marks );  
cout << my_marks.front();
```



```
graph LR; 42((42)) --> 95((95)); 95 --> 70((70)); 70 --> 81((81)); 81 --> 0((0));
```

# Assignment and Copy Constructor


- We can, however, define a copy constructor

- This function is called whenever a copy of the object is being made

```
List my_marks;  
my_marks.push_front( 81 );  
my_marks.push_front( 70 );  
my_marks.push_front( 95 );  
my_marks.push_front( 42 );
```

`my_marks.list_head` →

`cout << average( my_marks );`



```
graph LR; head((42)) --> node1((95)); node1 --> node2((70)); node2 --> node3((81)); node3 --> null((0));
```

If a copy constructor exists, it is called  
`List::List( List const & );`

```
int average( List ls ) {
```

```
    int n = 0, sum = 0;
```

```
    while ( !ls.empty() ) {  
        sum += ls.pop_front();  
        ++n;  
    }
```

```
    return sum/n;
```

```
}
```



# Assignment and Copy Constructor

- In this case, the copy constructor should make a copy of all four nodes

```
List my_marks;  
my_marks.push_front( 81 );  
my_marks.push_front( 70 );  
my_marks.push_front( 95 );  
my_marks.push_front( 42 );
```

```
cout << average( my_marks );
```

The diagram shows a linked list with four nodes containing the values 42, 95, 70, and 81. The nodes are connected by arrows pointing from left to right. The last node points to a null terminator represented by the symbol  $\emptyset$ . A red label 'my\_marks.list\_head' has an arrow pointing to the first node (42).

```
int average( List ls ) {  
    ls.list_head  
  
    int n = 0, sum = 0;  
  
    while ( !ls.empty() ) {  
        sum += ls.pop_front();  
        ++n;  
    }  
  
    return sum/n;  
}
```

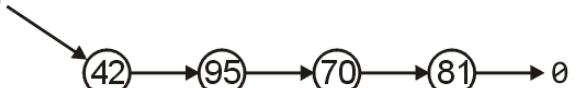
The diagram shows a linked list with four nodes containing the values 42, 95, 70, and 81. The nodes are connected by arrows pointing from left to right. The last node points to a null terminator represented by the symbol  $\emptyset$ . A red label 'ls.list\_head' has an arrow pointing to the first node (42).

# Assignment and Copy Constructor

- Any changes made to this copy do not affect the original linked list

```
List my_marks;
my_marks.push_front( 81 );
my_marks.push_front( 70 );
my_marks.push_front( 95 );
my_marks.push_front( 42 );
```

```
cout << average( my_marks );
```



```
graph LR
    head((my_marks.list_head)) --> n1((42))
    n1 --> n2((95))
    n2 --> n3((70))
    n3 --> n4((81))
    n4 --> null((0))
```

```
int average( List ls ) {
```



```
int n = 0, sum = 0;
```

```
while ( !ls.empty() ) {
    sum += ls.pop_front();
    ++n;
}
```

```
return sum/n;
```

```
}
```

# Assignment and Copy Constructor

- Thus, we require a *copy constructor*
- This is a method called when an object is being copied, either as a result of being passed-by-value to a function or being returned-by-value from a function
- Like a constructor, you should assume the variables are not initialized
- Unlike a constructor, you are not passed arguments to the constructor, rather, you are passed a reference to the object which should be copied

# Assignment and Copy Constructor

- The signature of a copy constructor has a constant reference to the same class as an argument:

```
List( const List & );
```

- The implementation of the copy constructor is almost identical to that of the **operator =**, except you are starting with a fresh object

# Assignment and Copy Constructor

- One simple implementation of a copy constructor is:

```
List::List( const List & lst ):list_head(0) {  
    *this = lst;  
}
```

- Here, all we do is:
  - initialize the array (set `list_head` to 0), and
  - let `operator =` do all the work

# Summary

- First, things to remember:
  - In C++, a member function and a member variable cannot have the same name
  - Common mistakes:

```
ptr->retrieve;  
list->head;
```
- Today, we looked at an example of a linked list implementation.

# Summary

- The Linked List example included
  - An overview of the list and node classes
  - Operations on linked lists
    - The default constructor
    - Getting started by coding simple functions
    - Implementing push front and pop front functions
    - Functions that require stepping through linked lists
    - The required behavior of the list destructor
    - The default and required behavior of operator=
    - The need for a copy constructor

# Summary

- The complete class is:

```
class List {
    private:
        Node * list_head;
    public:
        List();
        List( const List & );
        ~List();

        List & operator = ( const List & );

        bool empty() const;
        int front() const;
        Node * head() const;
        int count( int ) const;

        void push_front( int );
        int pop_front();
        int remove( int );
};
```



# Project 1: assignment operator and copy constructor

- Complete = operator
- Complete copy constructor
- **Test**

# Usage Notes

- These slides are made publicly available on the web for anyone to use
- If you choose to use them, or a part thereof, for a course at another institution, I ask only three things:
  - that you inform me that you are using the slides,
  - that you acknowledge my work, and
  - that you alert me of any mistakes which I made or changes which you make, and allow me the option of incorporating such changes (with an acknowledgment) in my set of slides

Sincerely,

Douglas Wilhelm Harder, MMath

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