

Lecture 7: Dynamic Programming I

Overview

- Memoization & subproblems; bottom up
- Fibonacci
- Shortest paths

Dynamic Programming (DP)

- *Powerful algorithmic design technique*
- *Large class of seemingly exponential problems have a polynomial solution (“only”) via DP*
- *Particularly for optimization problems (min / max)*

DP ~ «careful brute force»

DP ~ subproblems + «re-use»

Fibonacci Numbers

Problem

$$F_1 = F_2 = 1;$$
$$F_n = F_{n-1} + F_{n-2}$$

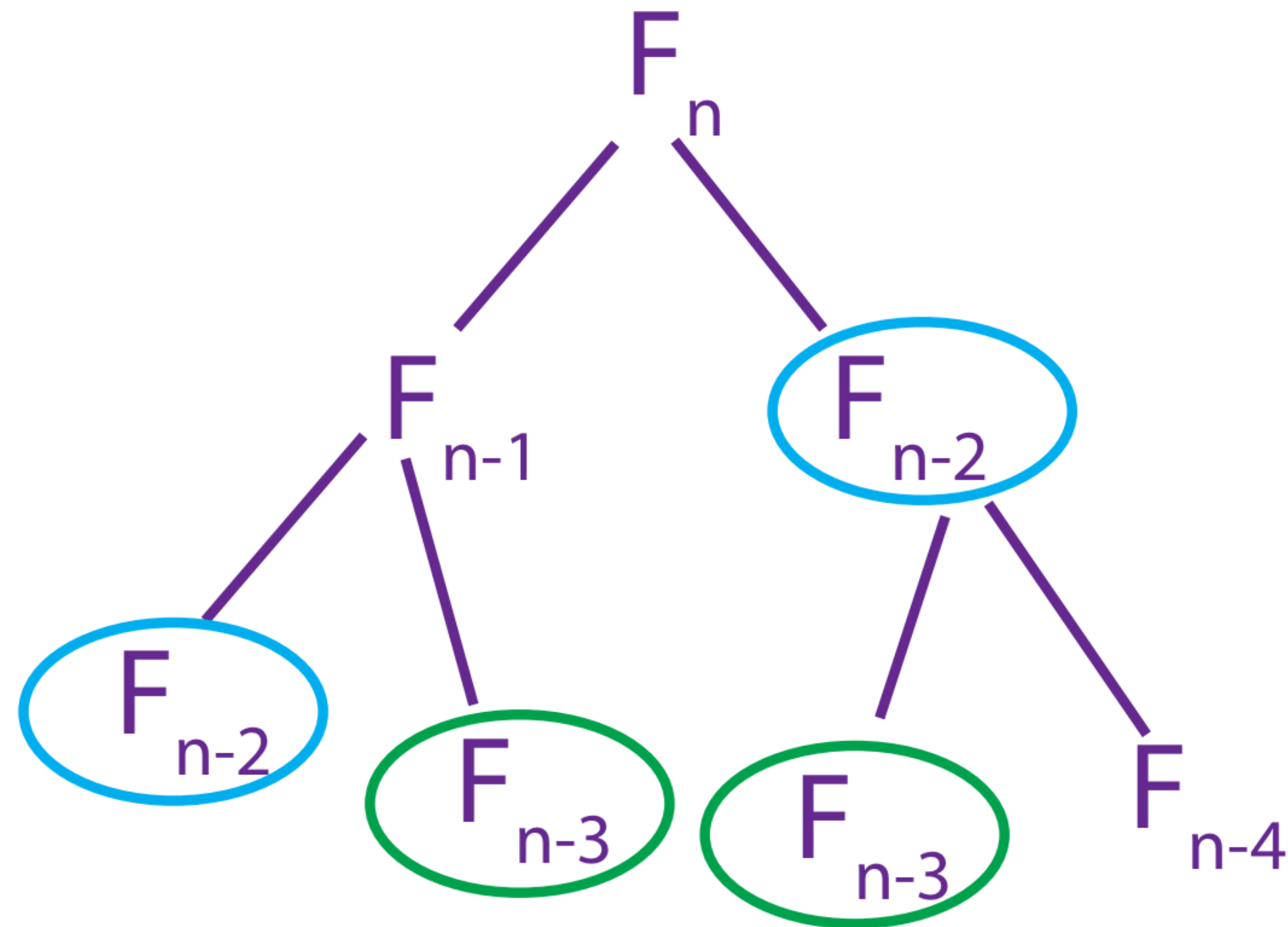
Goal: compute F_n

Solution: Naive Algorithm:

```
fib(n):  
  if  $n \leq 2$ :  $f = 1$   
  else:  $f = \text{fib}(n - 1) + \text{fib}(n - 2)$   
  return  $f$ 
```

$$\begin{aligned} T(n) &= T(n-1) + T(n-2) + O(1) \\ &\geq F_n \approx \phi^n \\ &\geq 2T(n-2) + O(1) \geq 2^{n/2} \end{aligned}$$

Fibonacci Numbers Memoized Algorithm



Solution: Memoized Algorithm:

memo = { }

fib(n):

```
if n in memo: return memo[n]
```

else: if $n \leq 2$: $f = 1$

```
else: f = fib(n - 1) + fib(n - 2)
```

$$\text{memo}[n] = f$$

```
return f
```

Fibonacci Numbers Memoized Algorithm

- fib(k) only recurses first time called, $\forall k$
- only ***n*** nonmemoized calls:
 $k = n, n-1, \dots, 1$
- memoized calls free ($\Theta(1)$ time)
- $\Theta(1)$ time per call (ignoring recursion)

Solution: Memoized Algorithm:

```
memo = {}
```

```
fib(n):
```

```
    if n in memo: return memo[n]
```

```
    else: if  $n \leq 2$  :  $f = 1$ 
```

```
           else:  $f = \text{fib}(n - 1) + \text{fib}(n - 2)$ 
```

```
           memo[n] = f
```

```
    return f
```

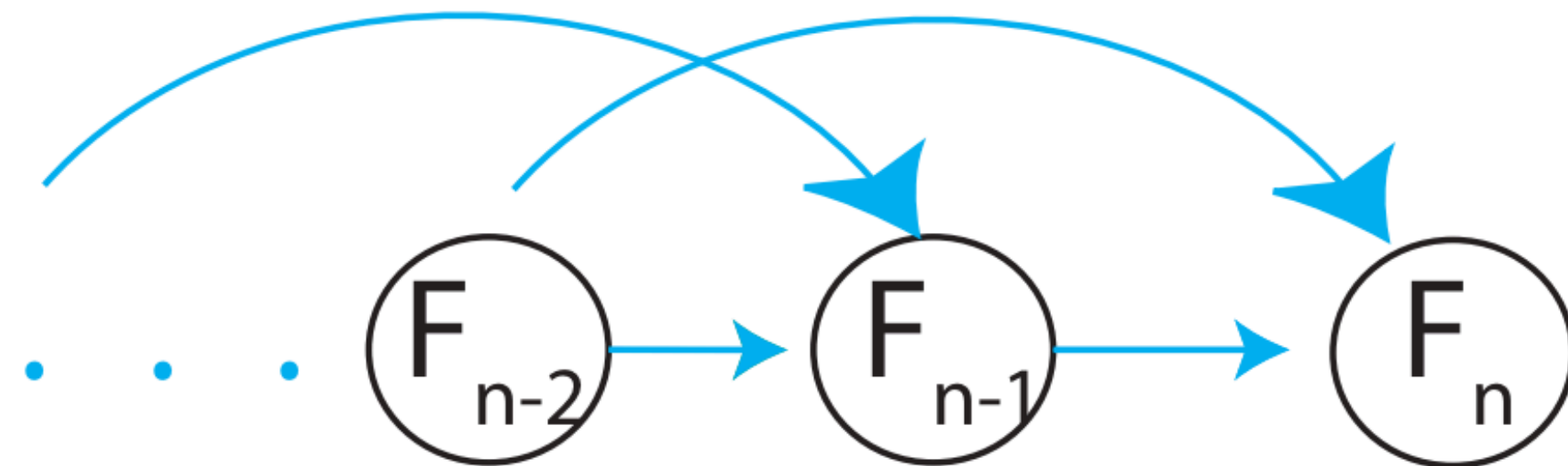
Memoized DP Algorithm

DP ~ recursion + memoization

- memoize (remember) & re-use solutions to subproblems that help solve problem
 - in Fibonacci, subproblems are F_1, F_2, \dots, F
- $\text{time} = \# \text{ of subproblems} \cdot \text{time/subproblem}$
- Fibonacci: # of subproblems is n , and time/subproblem is $\Theta(1) = \Theta(n)$

Fibonacci Numbers Bottom-up DP Algorithm

- exactly the same computation as memoized DP (recursion “unrolled”)
- in general: topological sort of subproblem dependency DAG



- practically faster: no recursion
- can save space: just remember last 2 fibs

Solution: Memoized Algorithm:

```
fib = {}  
for k in [1, 2, . . . , n]:  
    if k ≤ 2: f = 1  
    else: f = fib[k - 1] + fib[k - 2]  
    fib[k] = f  
return fib[n]
```