

# Lecture: Dynamic Programming III

# Overview

- Subproblems for strings
- edit distance (& longest common subseq)
- knapsack

# 5 Easy Steps to Dynamic Programming:

- (1) **define subproblems** | *count # subproblems*
- (2) **guess (part of solution)** | *count # choices*
- (3) **relate subproblems solutions** | *compute (time/subproblem)*
- (4) **recurse or memoize or build DP table bottom up** | *time =  
(time/subproblem) \* #subproblems*
- (5) **solve original problem** | a subproblem or a combination of subproblems

# Examples:

## 1) subproblems

# subproblems

## 2) guess

# choices

## 3) recurrence

time / subproblem

## 4) topological order

total time

## 5) original subproblem

extra time

# Fibonacci

$F_k$  for  $1 \leq k \leq n$

$n$

nothing to guess

1

$F_k = F_{k-1} + F_{k-2}$

$O(1)$  to compute

for  $k = 1, \dots, n$

$O(n)$

$F_n$

$O(1)$

# Shortest Paths

$\delta_k(s, v)$  for  $v \in V$ ,  $0 \leq k < |V|$

$V^2$

edge into  $V$  (if any)

$\text{indegree}(v) + 1$

$\delta_k(s, v) = \min\{\delta_{k-1}(s, u) + w(u, v)\}$

$O(1 + \text{indegree}(v))$

for  $k=0, 1, \dots, |V| - 1$

$O(VE)$

$\delta_{|V|-1}(s, v)$  for  $v \in V$

$O(V)$

# Subproblems for string and sequences:

(1) **suffixes**  $x[i:]$  |  $O(|x|)$

(2) **prefixes**  $x[:j]$  |  $O(|x|)$

(3) **substrings**  $x[i:j]$  |  $O(|x|^2)$

# Edit distance

## **Description:**

Given two strings  $x$  &  $y$ , what is the cheapest possible sequence of character edits (insert  $c$ , delete  $c$ , replace  $c \rightarrow c'$ ) to transform  $x$  into  $y$ ?

- cost of edit depends only on characters alphabet
- for example more common typo  $a \rightarrow s$ , then less cost
- cost of sequence = sum of costs of edits

# Edit distance

## Description:

Given two strings  $x$  &  $y$ , what is the cheapest possible sequence of character edits (insert  $c$ , delete  $c$ , replace  $c \rightarrow c'$ ) to transform  $x$  into  $y$ ?

- If *insert* & *delete* cost 1, replace costs 0, minimum ***edit distance*** equivalent to finding ***longest common subsequence***. Note that a subsequence is sequential but not necessarily contiguous.
- **HIEROGLYPHOLOGY** vs. **MICHAELANGELO**  $\Rightarrow$  **HELLO**

# Edit distance

## **Subproblems for multiple strings/sequences :**

- combine suffix/prefix/substring subproblems
- multiply state spaces



# Examples:

## 1) subproblems

# subproblems

## 2) guess

# choices

## 3) recurrence

time / subproblem

# Edit distance

$c(i, j) = \text{edit-distance}(x[i:], y[j:])$  for  $0 \leq i < |x|$ ,  $0 \leq j < |y|$

$\Theta(|x| \cdot |y|)$  subproblems

**guess whether, to turn x into y (3 choices):**

- $x[i]$  deleted
- $y[j]$  inserted
- $x[i]$  replaced by  $y[j]$

**$c(i, j) = \text{maximum of:}$**

- $\text{cost}(\text{delete } x[i]) + c(i + 1, j)$ , if  $i < |x|$ ,
- $\text{cost}(\text{insert } y[j]) + c(i, j + 1)$ , if  $j < |y|$ ,
- $\text{cost}(\text{replace } x[i] \rightarrow y[j]) + c(i + 1, j + 1)$ , if  $i < |x|$  &  $j < |y|$

# Examples:

4) topological order  
total time

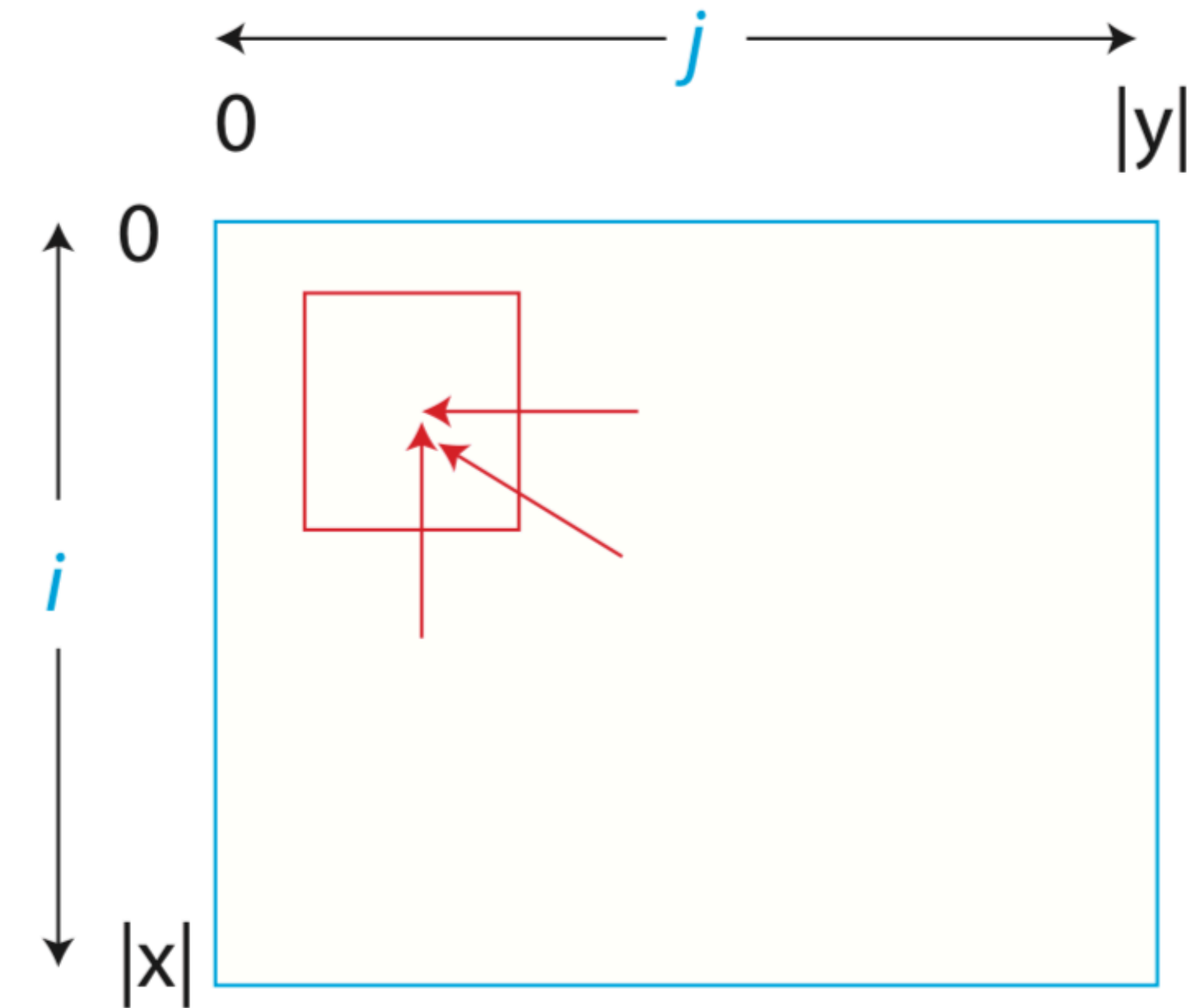
# Edit distance

DAG in 2D table:

- bottom-up OR right to left
- only need to keep last 2 rows/columns  
 $\Rightarrow$  linear space
- total time =  $\Theta(|x| \cdot |y|)$

5) original subproblem

$c(0, 0)$



# Knapsack

## Description:

Knapsack of size  $S$  you want to pack

- item  $i$  has integer size  $s_i$  & real value  $v_i$
- goal: choose subset of items of maximum total value subject to total size  $\leq S$

# Examples:

## 1) subproblems

# subproblems

## 2) guess

# choices

## 3) recurrence

time / subproblem

# Knapsack

value for *suffix i*:

$\Theta(\text{number of items})$  subproblems

guessing = whether to include item *i* or not

$\Rightarrow$  # choices = 2

$DP[i] = \max(DP[i + 1], v_i + DP[i + 1])$  , if  $s_i \leq S$

- not enough information. Don't know how much space is left!

# Examples:

## 1) subproblem

# subproblems

## 2) guess

# choices

## 3) recurrence

time / subproblem

## (4) topological order

total time

# Knapsack

value for *suffix i* given knapsack of size *X*:

$O(nS)$

guessing = whether to include item *i* or not

$\Rightarrow$  # choices = 2

$DP[i, X] = \max(DP[i + 1, X], v_i + DP[i + 1, X - s_i])$  , if  $s_i \leq X$

$DP[n, X] = 0$

for *i* in  $n, \dots, 0$ : for *X* in  $0, \dots, S$

total time =  $O(nS)$

# Examples:

1) subproblems

# Knapsack

DP [0, S]

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# Questions