

Lecture 1

Median Finding

Overview

- Divide & Conquer Overview

- Master Theorem

- Median Finding Trivial Solution



- Hoar Algorithm

- Median Finding (Blum, Floyd, Pratt, Rivest and Tarjan)

Divide & Conquer

Given a problem of size n divide it into subproblems of size $\frac{n}{b}$, $a \geq 1$, $b > 1$. Solve each subproblem recursively. Combine solutions of subproblems to get overall solution.

$$T(n) = aT\left(\frac{n}{b}\right) + [\text{work for merge}]$$

Master Theorem Definition

Master Theorem Examples

Median Finding Problem Definition

Median Finding Trivial Solution

Median Finding $O(n)$ in average case by Toni Hoar

Select(*S*, *i*)

1. Pick $x \in S$ // choose a pivot element. We can do it even randomly
2. $B = \{y \in S \mid y < x\}$ // Set of values less than x
3. $C = \{y \in S \mid y > x\}$ // Set of values greater than x
4. $D = \{y \in S \mid y == x\}$ // Set of values equal to x
5. Compute $k = \text{rank}(x)$ // We know how many elements less than x , so we know the rank!
6. if $k = i$:
7. return x
8. else if $k > i$:
9. return *Select*(B , i)
10. else if $k < i$:
11. return *Select*(C , $i - k$)

Median Finding $O(n)$ in average case by Toni Hoar

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Example

Target array: $l = [9, 1, 0, 2, 3, 4, 6, 8, 7, 10, 5]$

$\text{len}(l) == 11 \rightarrow$ we find the 6th smallest element.

Choose pivot randomly: $l[3] = 2$.

Divide the array according to the pivot:

$[1, 0, 2], [9, 3, 4, 6, 8, 7, 10, 5]$

We find the 3rd ($6 - \text{rank}(x)$) smallest element in the right array:

$[9, 3, 4, 6, 8, 7, 10, 5]$

Choose pivot randomly: $l[2] = 6$

Divide the array according to the pivot:

$[3, 4, 5, 6], [9, 7, 10]$

We find the 3rd smallest element in the left array:

$[3, 4, 5, 6]$

Choose pivot randomly: $l[1] = 4$

Divide the array according to the pivot:

$[3, 4], [5, 6]$

We find the ($3 - 2$) smallest element in the right array:

$[5, 6]$

Base variant where we just choose the min or max.

We need min return 5

$i = 6$
 $k = 3$

$k = 4$
 $i = 3$
 $k > i$

$i = 3, k = 2$

$\text{Select}(S, i)$

1. Pick $x \in S$ randomly $| O(1)$
2. $B = \{y \in S \mid y < x\}$
3. $C = \{y \in S \mid y > x\}$
4. $D = \{y \in S \mid y == x\}$
5. Compute $k = \text{rank}(x)$ $O(h)$
6. if $k = i$:
7. return x
8. else if $k > i$:
9. return $\text{Select}(B, i)$
10. else if $k < i$:
11. return $\text{Select}(C, i - k)$

$$T(h) = O(h) + T(\quad)$$

$$T(h) = h + \frac{h}{2} + \frac{h}{4} + \frac{h}{8} + \dots = 2h$$

Median Finding $O(n)$ in average case

Time Analysis

Worst case:

$$T(n) = n + (n-1) + (n-2) + \dots \sim n^2 = O(n^2)$$

1 $[1, 2, 3, \dots, n]$

2 $[2, 3, \dots, n]$

3 $[3, \dots, n]$

Median Finding $O(n)$ deterministic

Select(S, i)

1. Pick $x \in S$ // Need to pick it cleverly!
 - 1.1 Arrange S into columns of size 5 (In $\lceil n/5 \rceil$ cols)
 - 1.2 Sort each column (bigger elements on top) (linear time)
 - 1.3 Find "median of medians" as x

Median Finding $O(n)$ deterministic time analysis

Select(S, i)

1. Pick $x \in S$ // Need to pick it cleverly!
 - 1.1 Arrange S into columns of size 5 (In $\lfloor n/5 \rfloor$ cols)
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2. $B = \{y \in S \mid y < x\}$
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5. Compute $k = \text{rank}(x)$
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$O(h)$

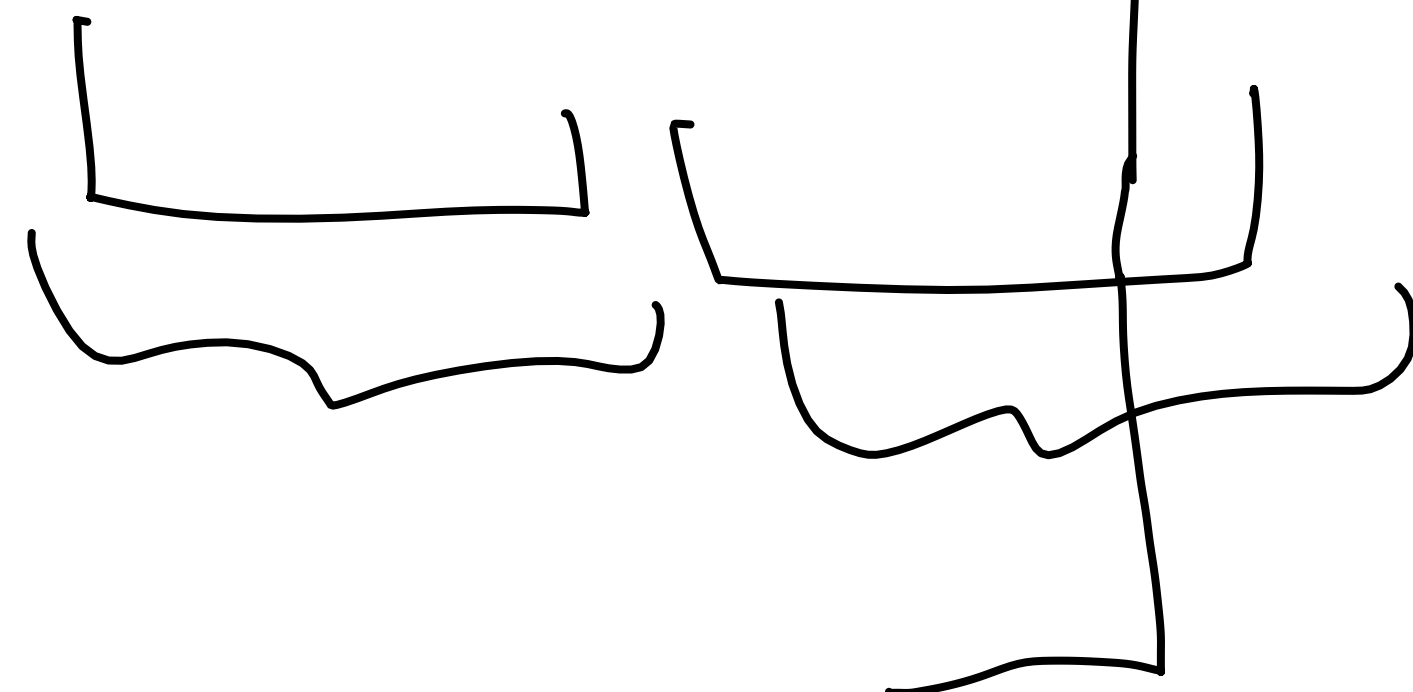
$O(h)$

$O(h)$

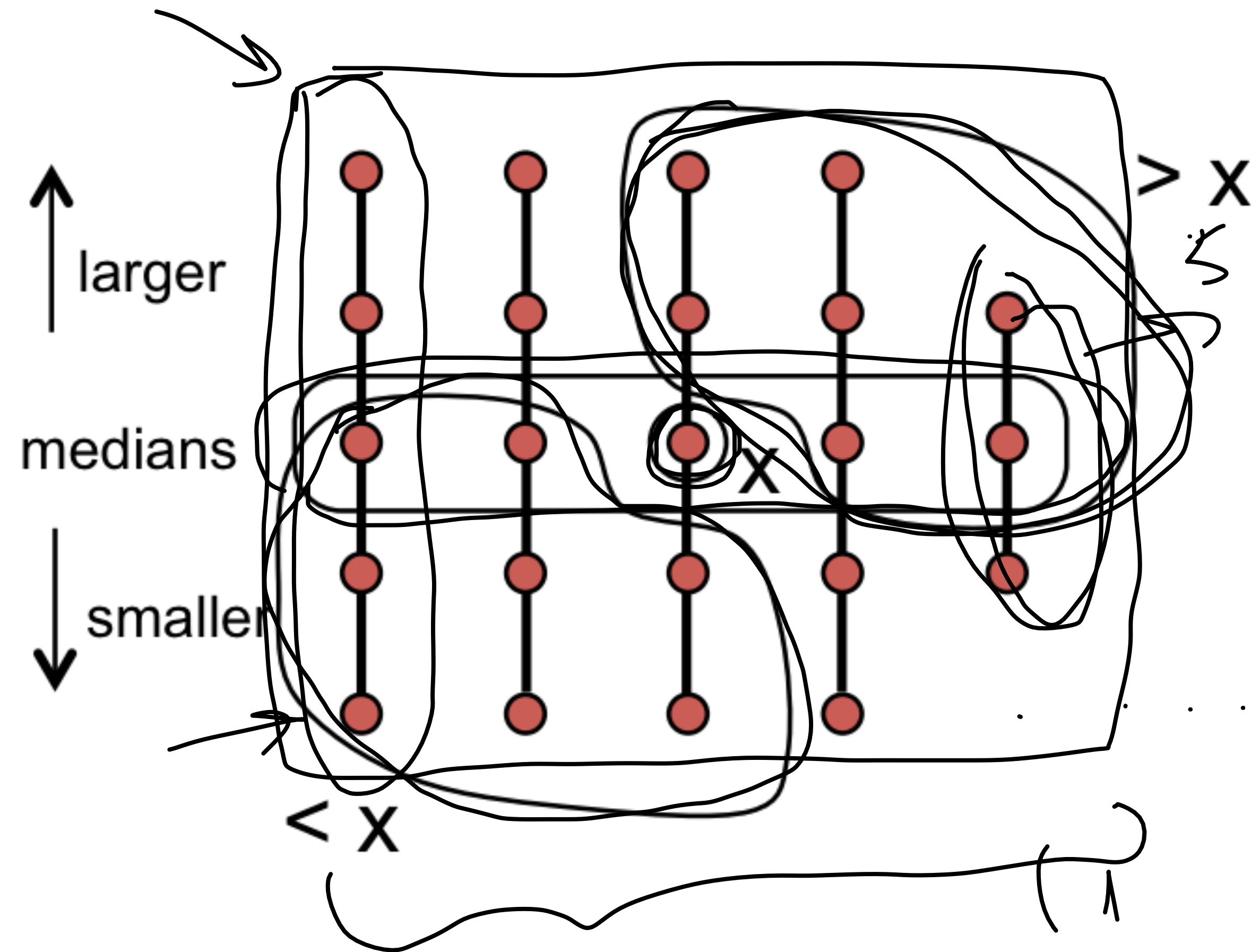
$2 \cdot h$

$h=5$

$n \log n$



Median Finding $O(n)$ deterministic time analysis



How many el-s $> x$

$$\left\lceil 3 \left(\frac{n}{10} \right) \right\rceil > x$$

$$\left\lceil 3 \left(\frac{n}{10} \right) \right\rceil$$

$$\left(\frac{7n}{10} \right)$$

Median Finding $O(n)$ deterministic time analysis

$$T(n) \leq T(\lceil n/5 \rceil) + T\left(\frac{7}{10}n + 6\right) + O(n).$$

find
pivot ± 3

choose
a value of x