

# Assignment Two – Magic Items Sorter

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C. Marcus DiMarco  
C.DiMarco1@Marist.edu

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## 1 Objectives

- Create a program that sorts all Strings within a given text file using selection sort, insertion sort, merge sort and Quicksort.
- Print the number of comparisons for each sort.

## 2 Conditions

- The relation on the set of Strings must be 'less than', such that for any element a which appears before any other element b, a is less than b. For Strings, 'less than' is defined as preceding alphabetically.
- The array must be shuffled before running the sorts, and the sorts must act on the same shuffle of the array.
- The program must only use custom algorithms for the sorts. Any of these sorts which already exist in the language may not be used.

## 3 Overview

In order to enhance readability, each sorting algorithm is a method contained in the class MarcusSort.

### 3.1 selectionSort()

The most straightforward sorting algorithm in this program is selection sort. Selection sort starts at the first index, finds the minimum value in the array and places it at the current index, repeating until the array is sorted.

```
16 // Selection sort will search for minimum unsorted value and
17 // place it at first unsorted index.
18 public void selectionSort(String[] strings) {
19
20     // Start with counter at 0
21     this.resetCounter();
22
23     // First unsorted index through each pass is i
24     for (int i = 0; i < strings.length; i++) {
25         counter++; // Comparison in for loop
26         int minIndex = i;
27         // Find index of smallest value in remainder of array
28         for (int j = i + 1; j < strings.length; j++) {
29             counter += 2; // Comparisons in for loop and if statement
30             if (strings[minIndex].compareToIgnoreCase(strings[j]) > 0) {
31                 minIndex = j;
32             }
33         }
34         counter++; // Comparison for inner loop termination
35
36         // Swap indices minIndex and i
37         String temp = strings[i];
38         strings[i] = strings[minIndex];
39         strings[minIndex] = temp;
40     }
41     counter++; // Comparison for outer loop termination
42
43     // Print completion message with number of comparisons
44     this.printCompletionMessage();
45 }
```

### 3.2 insertionSort()

Improving on selection sort is insertion sort. Insertion sort relies on taking a key value and placing it in sorted order among all of the previous elements. We start at index 1 and compare it to the value at index 0, sorting those two as needed. Note that after each iteration, indices 0-i are sorted.

```

45 // Insertion sort will sort the array from index 0 through i,
46 // 'inserting' each value into its correct position.
47 public void insertionSort(String[] strings) {
48
49     // Start with counter at 0
50     this.resetCounter();
51
52     // Start comparisons at index 1
53     for (int i = 1; i < strings.length; i++) {
54         counter++; // Comparison in outer loop
55         String keyString = strings[i];
56         for (int j = i - 1; j >= 0; j--) {
57             counter += 2; // Inner loop and if statement
58             // Move all items larger than key forward 1 index
59             if (keyString.compareToIgnoreCase(strings[j]) < 0) {
60                 strings[j + 1] = strings[j];
61                 counter++; // Second comparison in if statement
62                 // If index 0 reached, assign it currentString
63                 if (j == 0) {
64                     strings[j] = keyString;
65                 }
66             } else {
67                 // Insert key at index ahead of first smaller element
68                 strings[j + 1] = keyString;
69                 break;
70             }
71         }
72         counter++; // Comparison for inner loop termination
73     }
74     counter++; // Comparison for outer loop termination
75
76     // Print completion message with number of comparisons
77     this.printCompletionMessage();
78 }

```

### 3.3 mergeSort()

In stark contrast to the other sorting algorithms seen so far, merge sort drastically reduces complexity by implementing a "divide and conquer" tactic, which relies on a simple principle: arrays of size 1 are sorted. Therefore, if we divide the target array into two subarrays, and continue to divide those until we reach arrays of size 1, we can recombine the subarrays in sorted order and greatly reduce the amount of work needed. It uses two operations, `sortThenMerge()` and `merge()`.

```

80  /**
81   * mergeSort() is an abstraction for cleaner user interaction. It will reset
82   * the counter, call the recursive merge function, and then print the
83   * completion message with the number of comparisons.
84   * @param strings
85   */
86  public void mergeSort(String[] strings) {
87
88      // Start with counter at 0
89      this.resetCounter();
90
91      // Nest recursive function inside for readability and proper
92      // counter/print calls
93      this.sortThenMerge(strings, 0, strings.length - 1);
94
95      // Print completion message with number of comparisons
96      this.printCompletionMessage();
97  }
98
99  // Merge sort will divide the array recursively until subarrays are
100 // of size 1, then merge the subarrays together in sorted order.
101 private void sortThenMerge(String[] strings, int leftIndex, int rightIndex) {
102
103     counter++; // Increment the if comparison
104     if (leftIndex < rightIndex) {
105         int midpoint = (leftIndex + rightIndex) / 2;
106         sortThenMerge(strings, leftIndex, midpoint);
107         sortThenMerge(strings, midpoint + 1, rightIndex);
108         merge(strings, leftIndex, midpoint, rightIndex);
109     }
110 }
111
112 // Merge will take an array, create two sorted subarrays, and
113 // sort the elements in place.
114 private void merge(String[] strings, int leftIndex, int midpoint,
115                   int rightIndex) {
116
117     // Create two (sorted, due to recursion) subarrays
118     String[] stringsL = new String[midpoint - leftIndex + 1];
119     String[] stringsR = new String[rightIndex - midpoint];
120
121     // Populate the subarrays, incrementing for each comparison
122     for (int i = 0; i < stringsL.length; i++) {
123         counter++;
124         stringsL[i] = strings[leftIndex + i];
125     }

```

```

126     counter++;    // Increment the loop exit comparison
127
128     for (int j = 0; j < stringsR.length; j++) {
129         counter++;
130         stringsR[j] = strings[midpoint + j + 1];
131     }
132     counter++;    // Increment the loop exit comparison
133
134     int i = 0;
135     int j = 0;
136     int k = leftIndex;
137     // For the selected range, sort until one subarray is exhausted
138     while (i < stringsL.length && j < stringsR.length) {
139         counter += 2; // Increment for loop and if statements
140         if (stringsL[i].compareToIgnoreCase(stringsR[j]) < 0) {
141             strings[k++] = stringsL[i++];
142         } else {
143             strings[k++] = stringsR[j++];
144         }
145     }
146     // Append the remaining (sorted) subarray
147     if (i == stringsL.length) {
148         for (; k <= rightIndex; k++) {
149             counter++;
150             strings[k] = stringsR[j++];
151         }
152         counter += 2; // Increment for loop exit and if
153     } else {
154         for (; k <= rightIndex; k++) {
155             counter++;
156             strings[k] = stringsL[i++];
157         }
158         counter += 2; // Increment for loop exit and if
159     }
160 }

```

### 3.4 quickSort()

Building off of the "divide and conquer" principle in merge sort, Quicksort looks to make things even more efficient by beginning the conquering during the dividing instead of waiting until after. Quicksort selects a pivot value and divides the elements in the array around it, with all smaller elements to its left and all larger elements to its right. This means that after each call, any pivot value is in its correct index.

```

162  /**
163   * quickSort() is an abstraction for cleaner user interaction. It will reset
164   * the counter, call the recursive quicksort function, and print the
165   * completion message with the number of comparisons.
166   * @param strings
167   */
168  public void quickSort(String[] strings) {
169
170      // Set counter to 0
171      this.resetCounter();
172
173      // Nest recursive function inside for readability and proper
174      // counter/print calls
175      this.quickSort(strings, 0, strings.length - 1);
176
177      // Print completion message with number of comparisons
178      this.printCompletionMessage();
179  }
180
181  // Quicksort will take an array and recursively divide it around
182  // a pivot value, sorting the elements around the pivots
183  private void quickSort(String[] strings, int leftIndex, int rightIndex) {
184
185      counter++; // Increment for if comparison
186      if (leftIndex < rightIndex) {
187          // Quicksort the elements to either side of the partition
188          int partition = partition(strings, leftIndex, rightIndex);
189          quickSort(strings, leftIndex, partition - 1);
190          quickSort(strings, partition + 1, rightIndex);
191      }
192  }
193
194  // Partitions an array around a pivot value and places all values
195  // smaller than the pivot to its left, then places pivot at its sorted index
196  private int partition(String[] strings, int leftIndex, int rightIndex) {
197
198      String pivotString = strings[rightIndex];
199
200      int sortedIndex = leftIndex;
201
202      for (int i = leftIndex; i <= rightIndex; i++) {
203          counter += 2; // Increment for loop and if statement
204          if (strings[i].compareToIgnoreCase(pivotString) < 0) {
205              // If element at i is smaller than pivot, place in the
206              // left half of the array
207              String tempString = strings[sortedIndex];

```

```

208         strings[sortedIndex++] = strings[i];
209         strings[i] = tempString;
210     }
211 }
212 counter++;    // Increment for loop exit comparison
213
214 // Swap pivot string and value at the pivot's sorted index
215 strings[rightIndex] = strings[sortedIndex];
216 strings[sortedIndex] = pivotString;
217 return sortedIndex;
218 }

```

### 3.5 Other helper functions

Below are the remaining functions present in MarcusSort().

```

220 // Shuffle routine based on the Knuth or Fisher–Yates, but not
221 // Rosanna, shuffle
222 public void notRosannaShuffle(String[] strings) {
223
224     for (int i = strings.length - 1; i >= 0; i--) {
225         // Get random index in the remaining length
226         int randomIndex = (int) Math.floor(Math.random() * i);
227         if (randomIndex == i) {
228             continue;
229         } else {
230             // If the random index is not i, swap their values
231             String tempString = strings[i];
232             strings[i] = strings[randomIndex];
233             strings[randomIndex] = tempString;
234         }
235     }
236 }
237
238 // Getter for counter
239 public int getCounter() {
240     return counter;
241 }
242
243 // Controlling setter functionality to only reset to 0
244 public void resetCounter() {
245     counter = 0;
246 }
247

```

```

248 // Message to print upon completing sort which includes counter
249 public void printCompletionMessage() {
250     System.out.println("Sort_complete!_Number_of_comparisons:_\"
251         + counter);
252 }

```

### 3.6 Assignment2

With the methods we have just defined, we are ready to create our sorter. We'll need some imported libraries: namely, java.io.File to import a file, java.io.FileNotFoundException to account for errors finding the input file, and java.util.Scanner to read the file.

```

1  /**
2   * A program which reads a constant number of Strings
3   * from a file and sorts them using various methods,
4   * comparing the efficiency of the sorts by
5   * printing the number of comparisons for each.
6   */
7  import java.io.File;
8  import java.io.FileNotFoundException;
9  import java.util.Scanner;
10
11 public class Assignment2 {
12     public static void main(String[] args) {
13         final int NUM_OF_ITEMS = 666; // Length of file as constant
14         String[] magicItems = new String[NUM_OF_ITEMS]; // Array of file strings
15         MarcusSort sorter = new MarcusSort(); // Instance of MarcusSort
16
17         // Try/catch block for file import and reading
18         try {
19             File file = new File("magicitems.txt");
20             Scanner read = new Scanner(file);
21             for (int i = 0; i < NUM_OF_ITEMS; i++) {
22                 magicItems[i] = read.nextLine();
23             }
24             read.close();
25         } catch (FileNotFoundException e) {
26             System.out.println("Whoops!_Couldn't_find_magicitems.txt");
27             e.printStackTrace();
28         }
29
30         // Shuffle array before beginning sorts
31         sorter.notRosannaShuffle(magicItems);
32

```



```

33 // Create second array to store the shuffle
34 String[] magicItemsShuffled = new String[NUM_OF_ITEMS];
35 System.arraycopy(magicItems, 0, magicItemsShuffled, 0, NUM_OF_ITEMS);
36
37 // Sort using selection sort, print comparisons, reset shuffle
38 sorter.selectionSort(magicItems);
39 System.arraycopy(magicItemsShuffled, 0, magicItems, 0, NUM_OF_ITEMS);
40
41 // Sort using insertion sort, print comparisons, reset shuffle
42 sorter.insertionSort(magicItems);
43 System.arraycopy(magicItemsShuffled, 0, magicItems, 0, NUM_OF_ITEMS);
44
45 // Sort using merge sort, print comparisons, reset shuffle
46 sorter.mergeSort(magicItems);
47 System.arraycopy(magicItemsShuffled, 0, magicItems, 0, NUM_OF_ITEMS);
48
49 // Sort using quicksort, print comparisons
50 sorter.quickSort(magicItems);
51 }
52 }

```

## 4 Results

Due to the pseudorandom nature of `notRosannaShuffle()`, exact results will vary from execution to execution. The below table shows the results obtained during 5 trials, along with the average number of comparisons and the expected order of growth for each sort.

Sort Performance (in number of comparisons)				
	selectionSort()	insertionSort()	mergeSort()	quickSort()
Trial 1	444,223	334,807	22,004	15,454
Trial 2	444,223	344,659	22,011	15,834
Trial 3	444,223	338,621	22,025	17,330
Trial 4	444,223	336,871	22,051	15,246
Trial 5	444,223	337,681	22,033	14,963
Average	444,223	338,527.8	22,024.8	15,765.4
$O(g(n))$	$O(n^2)$	$O(n^2)$	$O(n\log(n))$	$O(n\log(n))$

The first two algorithms belong to  $O(n^2)$  due to their nested loops; ignoring constant factors and lower order terms, they iterate over the array during each iteration of the sort, and the sort iterates for each term in the array. The second pair of algorithms expect to run in  $O(n\log(n))$  time due to their trademark "linearithmic" approach, "divide and conquer". The division is logarithmic, and the conquering/merging is linear.