

CS 4124  
Solutions to Homework Assignment 5  
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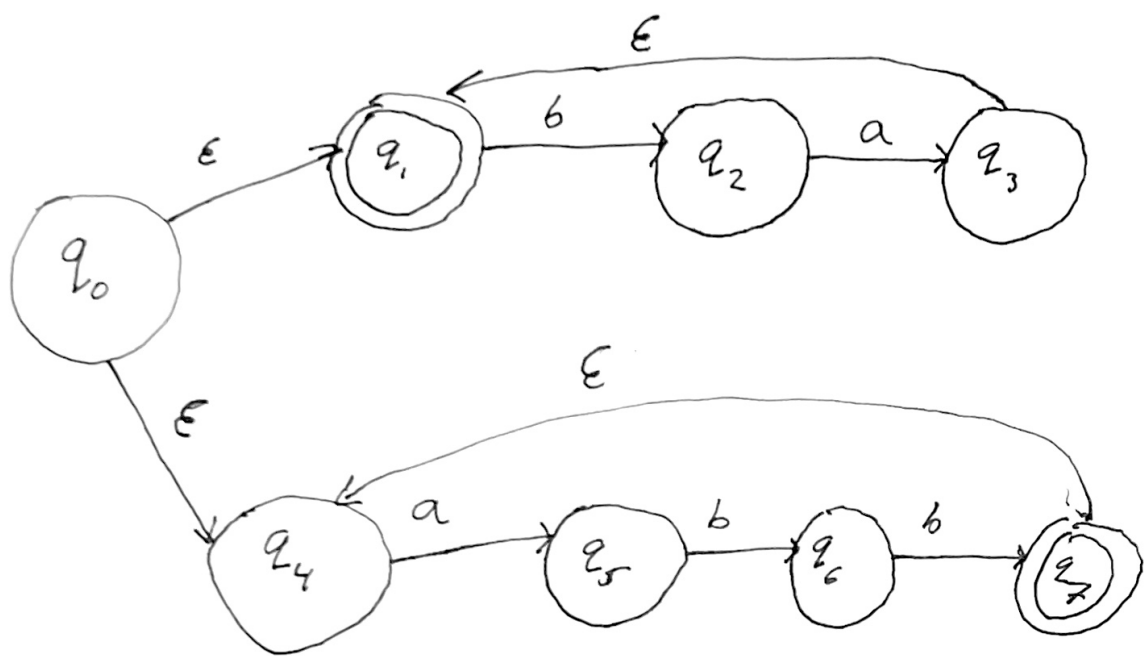
[50] 1. Let  $\mathcal{R}_1$  be the (simplified) regular expression  $(ba)^* + (abb)(abb)^*$ .

**Construct an  $\varepsilon$ -NFA  $N_1$  that accepts the language denoted by  $\mathcal{R}_1$ . You should employ the construction given in class or in the textbook for inspiration, but you do not have to follow the construction precisely. Use reason to construct your  $N_1$  and justify your reasoning. Give  $N_1$  as a labeled directed graph or state diagram. Please draw it neatly!**

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This was constructed by first creating two  $\varepsilon$ -NFA the first one located at the top recognizes  $(ba)^*$  this is seen as it accepts the empty string and if it does receive a string in order for it to accept it must take the path  $q_1 \rightarrow q_2 \rightarrow q_3 \rightarrow q_1$  which can only be traversed by the string  $ba\varepsilon$ . The epsilon ensures that  $(ba)^*$  will be recognized.

The second  $\varepsilon$ -NFA is located at the bottom and recognizes  $(abb)(abb)^*$  this is seen as it first must traverse  $q_4 \rightarrow q_5 \rightarrow q_6 \rightarrow q_7$  to reach an accepting state. This can only be done by the string  $abb$  after it is in state  $q_7$  there is an epsilon transition back to  $q_4$  which ensures that  $(abb)(abb)^*$  will also be recognized. Both these  $\varepsilon$ -NFA are combined with the initial state  $q_0$  and a  $\varepsilon$  transition to their initial states.



[50] 2. Let  $N_2$  be the  $\varepsilon$ -NFA in Figure ??.

A. Compute the  $\varepsilon$ -reachability set  $E(q)$  of each state  $q$  of  $N_2$ .

$$E(q_0) = \{q_0, q_1, q_2\}$$

$$E(q_1) = \{q_1\}$$

$$E(q_2) = \{q_2\}$$

B.

$bb$	$q_0 \xrightarrow{b} q_1 \xrightarrow{b} q_2$
$cb$	$q_0 \xrightarrow{\varepsilon} q_1 \xrightarrow{c} q_1 \xrightarrow{b} q_2$
$\varepsilon$	$q_0 \xrightarrow{\varepsilon} q_2$
$c$	$q_0 \xrightarrow{c} q_2$
$b$	$q_0 \xrightarrow{\varepsilon} q_1 \xrightarrow{b} q_2$
$ab$	$q_0 \xrightarrow{\varepsilon} q_1 \xrightarrow{a} q_0 \xrightarrow{\varepsilon} q_1 \xrightarrow{b} q_2$
$cc$	$q_0 \xrightarrow{\varepsilon} q_1 \xrightarrow{c} q_0 \xrightarrow{c} q_2$
$ac$	$q_0 \xrightarrow{\varepsilon} q_1 \xrightarrow{a} q_0 \xrightarrow{c} q_2$
$ca$	$q_0 \xrightarrow{\varepsilon} q_1 \xrightarrow{c} q_1 \xrightarrow{a} q_0 \xrightarrow{\varepsilon} q_2$
$ba$	$q_0 \xrightarrow{b} q_1 \xrightarrow{a} q_0 \xrightarrow{\varepsilon} q_2$
$a$	$q_0 \xrightarrow{\varepsilon} q_1 \xrightarrow{a} q_0 \xrightarrow{\varepsilon} q_2$

C. Use the power set construction to obtain a DFA  $M_2$  equivalent to  $N_2$ . Give  $M_2$  as a labeled directed graph or state diagram. Please draw it neatly!

