

CS 4124
Solutions to Homework Assignment 4
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[30] 1. Let $\Sigma = \{0, 1\}$, the binary alphabet. Let

$$F = \{h \mid h : \Sigma^* \rightarrow \Sigma^*\},$$

the set of functions from strings to strings.

Prove that F is uncountable.

HINT: Use the fact that Σ^* is countably infinite and employ a proof by diagonalization.

My proof strategy:

I will prove that the set of functions $F' = \{h : h : \Sigma^* \rightarrow \{0, 1\}\}$ is uncountable and as this set is a subset of F that would imply that F is uncountable. I will do this by constructing a bijection between $H : F' \rightarrow \mathcal{P}(\Sigma^*)$ (the power set of Σ^*) in the last homework we proved that no set can have a bijection with its power set and as Σ^* is countable we get that $\mathcal{P}(\Sigma^*)$ is uncountable hence if a bijection exists H exists we would have that F' is uncountable.

Let $H : F' \rightarrow \mathcal{P}(\Sigma^*)$ where $H(f) = \{w \in \Sigma^* : f(w) = 1\}$ (informally this is the set of elements from Σ^* who get mapped to 1). Now to show that this is a bijection I will come up with an inverse $H^{-1} : \mathcal{P}(\Sigma^*) \rightarrow F'$

[30] 2. Let

$$E = \{w \in \{0, 1\}^* \mid \ell(w) \text{ is even}\},$$

the language of binary strings of even length. See Figure 1 for a sample program that decides E , just for L^AT_EX formatting purposes. Let

$$L = \{x \mid \text{program } x \text{ decides } E\}$$

be the language of programs that decide E .

Prove that L is undecidable by using m-reducibility.

[40] 3. Let

$$\text{DOUBLE} = \{ww \mid w \in \{a, b\}^*\}.$$

Let

$$W = \{x \mid \text{program } x \text{ decides DOUBLE}\}$$

be the language of programs that decide DOUBLE.

Program P Input w if $\ell(w)$ is even then accept else reject
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Figure 1: A sample program to decide E , just to show how to format a program in L^AT_EX.

A. Is DOUBLE decidable? Justify your answer.

B. Is W decidable? Prove your answer.
