

**CS 4124**  
**Solutions to Homework Assignment 1**  
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[25] 1. Textbook Problem 5 in B.1 on Page 520.

Prove Lemma A.4: For all alphabets  $\Sigma$  and all languages  $L \subseteq \Sigma^*$ , the equivalence relation  $\equiv_L$  is right-invariant.

The lemma and the definition of the equivalence relation are on Page 505.

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*Proof.* Let  $\Sigma$  be an alphabet and  $L \subseteq \Sigma^*$  and  $\equiv_L$  be an equivalence relation as defined by (A.3). Let  $x, y, z, w \in \Sigma^*$  and  $x \equiv_L y$  then we have  $xz, yz \in L$  as  $zw \in \Sigma^*$  we have  $xzw, yzw \in L$  which implies  $xz \equiv_L yz$ .  $\square$

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[25] 2. Textbook Problem 3(c) in B.2 on Page 522.

Design an FA  $M$ , with alphabet  $\Sigma = \{0, 1\}$ , that recognizes (c) the set of all strings that contain the string 011, in that order, but not necessarily consecutively.

Be certain you understand the set (language) of  $M$  before you start designing. Constructing some examples of strings both in  $L(M)$  and not in  $L(M)$  can be helpful.

You may use an algebraic specification or a transition diagram (labeled directed graph) specification to present your design for  $M$ . Be certain to explain why your design works.

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[25] 3. Consider the OA  $M_4$  in Figure 3.4 on Page 60. Give a complete and careful algebraic specification

$$M_4 = (Q, \Sigma, \delta, q_0, F)$$

for  $M_4$ .

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[25] 4. Textbook Problem 7(c) in B.2 on Pages 522 and 523.

Use a (fooling set)-plus-(Continuation Lemma) argument to prove that the following language is *not* Regular:

$$L_5 = \{a^i b^j c^k \mid i = j \text{ OR } j = k\}.$$

Your argument will be a proof by contradiction.

The Continuation Lemma is Lemma 3.2 on Page 57, while the fooling set kind of argument is first utilized on Pages 76 and 77.

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