Matching Theory Notes

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Lemma 1.0.1. For any graph, G, $\alpha(G) + \tau(G) = |V(G)|$.

Proof. Let G be an arbitrary graph and let M be an arbitrary point cover where $|M| = \tau(G)$. Then V(G) - M is an independent set of points. This is true because if V(G) - M was not an independent set of points then there are at least two points $u, v \in V(G) - M$ where u and v are adjacent. If they are adjacent the line incident with u and v is not covered by M which is a contradiction. As V(G) - M is an independent set of points we get the inequality

$$\alpha(G) \ge |V(G) - M| = |V(G)| - \tau(G)$$

Now assume that L is an arbitrary independent set of points where $|L| = \alpha(G)$. Then V(G) - L is a point cover of G. If V(G) - L is not a point cover of G then there exists a line l such that l is not covered by V(G) - L. However this would imply that l is incident with two points in L this is a contradiction on the fact L is independent points hence V(G) - L is a point cover of G. As V(G) - L is a point cover of G we get the inequality

$$\tau(G) \le |V(G) - L| = |V(G)| - \alpha(G)$$

Combining both inequalities we get $\alpha(G) \ge |V(G)| - \tau(G)$ and $\tau(G) \le |V(G)| - \alpha(G)$. This implies that

$$\alpha(G) + \tau(G) = |V(G)|$$

. \square

Lemma 1.0.2. For any graph G with no isolated points, $v(G) + \rho(G) = |V(G)|$.

Proof. Let G be an arbitrary graph with no isolated points and let C be a line cover of G where $|C| = \rho(G)$. Let $\langle C \rangle$ be the graph formed from lines the set of lines C and the set of points V(C). We have that $\langle C \rangle$ is a union of stars. This is because if $\langle C \rangle$ was not a union of stars. Then there would be two points in the graph that that are adjacent two one or more point in $\langle C \rangle$. If they are adjacent to a single point then removing either of the lines incident would create a smaller minimal line cover. If they are adjacent to two or more points then removing any of the incident lines would create a smaller minimal line cover. Now if we let n be the number of components in $\langle C \rangle$ we get that $n = |V(G)| - \rho(G)$ which arises because in each star there is one more point than there are lines. If we take one line from each star we get a matching hence

$$v(G) \ge |V(G)| - \rho(G)$$

Now let M be an arbitrary matching of G where |M| = v(G), and U be the set of points that are not covered by M. We get that U is an independent set of points because if U was not an independent set of points then there would be two points $u, v \in U$ such that u and v are adjacent. If u and v are adjacent then there is a line incident with u and v that is not covered by M which is a contradiction. As U is an independent set of points and G has no isolated points we get that |U| = |V(G)| - 2v(G). Let S be a line covering of U we get that $M \cup S$ is a line covering. Then we get the inequality

$$\rho(G) \le |M \cup S| = v(G) + |V(G)| - 2v(G) = |V(G)| - v(G)$$

Combining both inequalities we get

$$v(G) + \rho(G) = |V(G)|$$

Exercise 1.0.3.

- 1. A minimal line cover is minimum if and only if it contains a maximum matching.
- 2. A maximal matching is maximum if and only if it is contained in a minimum line cover.

Proof. (1)

Let L be a line cover where $|L| = \rho(G)$. As L is a minimum line cover we have that it forms the lines from a union of stars on G. We have that the number of stars is v(G) which comes from selecting one matching from each star. Hence L contains a maximum matching.

Let L be a minimal line cover that contains a maximum matching. Then as L is a union of stars we get a maximum matching by selecting a line from each star. Hence $|L| = |V(G)| - v(G) = \rho(G)$ hence it is a minimum line cover.

Proof. (2) Let M be a maximum matching. Let L be an line cover obtained by taking an arbitrary line covering of the edges incident with V(G) - V(M) and M. We have that $|L| = |V(G)| - 2v(G) + v(G) = \rho(G)$ hence L contains a minimum line cover.

Let M be a maximal matching that is contained in a minimum line cover L. Then we have that the matchings in M are formed from individual lines of the stars formed by L. Hence $|M| = V(G) - \rho(G) = v(G)$ hence M is a maximum matching. \square

Exercise 1.0.4. For any graph G, $v(G) \leq \tau(G)$

Proof. Let G be an arbitrary graph and let M be a maximum matching in G. As M is a maximum matching we get that it is contained in a minimum line cover L. As L is a minimum line cover we have that it is the union of stars where each matching in M comes from a single line in each star. We get that the number of stars is $\tau(G)$ hence $v(G) \leq \tau(G)$

Theorem 1.1.1. König's Minimax Theorem: If G is bipartite, then $v(G) = \tau(G)$

Proof. Remove lines from G as long as possible while keeping $\tau(G)$ the same. Denote the resulting minimal graph by G'. Hence $\tau(G') = \tau(G)$, but for any line $l \in E(G)$ we have $\tau(G'-l) < \tau(G)$. Suppose that G' does not consist of independent lines. Then there are two lines $l_1, l_2 \in E(G')$ that are both incident with a point $a \in V(G')$. As G' is minimal we have that there exists a point cover of G' - x we denote this point cover by S_x where $|S_x| = \tau(G') - 1$. We also have the point cover S_y which is covering G' - y with $|S_x| = |S_y|$. Form the induced subgraph G'' of G', $G'' = G'[\{a\} \cup S_x \cup S_y \cap \overline{S_x \cap S_y}]$ and let $t = |S_x \cap S_y|$. Then we have $|V(G'')| = 2(\tau(G') - 1 - t) + 1$. As G'' is a subgraph of a bipartite graph we have that it is bipartite. Hence we have a set T which is a point cover of G'' where $|T| = \tau(G') - 1 - t$. We claim that the set of points $T' = T \cup (S_x \cap S_y)$ covers G'. Let z be a line in G' if $z \neq x$ or $z \neq y$ then we have that z is covered by S_x and S_y hence it is covered by $S_x \cap S_y$ or it connects $S_x - S_y$ to $S_y - S_x$. If z = x or z = y then we have it is covered by T. So $\tau(G') \leq |T'| = |T \cup (S_x \cap S_y)| = |T| + |S_x \cap S_y| \leq \tau(G') - 1 - t + t = \tau(G') - 1$ a contradiction. Hence G' is an independent set of lines. Hence

$$\tau(G) = \tau(G') = v(G') \le v(G)$$

Theorem 1.1.2. In any (possibly infinite) bipartite graph there exists a matching M and a point cover P such that every line in M contains exactly one point in P and every point in P is contained in exactly one line of M.

Proof. Let G be an arbitrary bipartite graph and let L be a minimum point cover of G. Create the new graph G' by removing lines from G as long as $\tau(G') = \tau(G)$. We have that G' consists of an independent set of lines from each of the points in the point cover $\tau(G')$ select one line. As G' consists of an independent set of lines we have a maximum matching by König's Minimax Theorem which satisfies the implication.

If X is any set of points in a graph G, we denote the set of all points which are adjacent to at least one point of X as $\Gamma(X)$.

Theorem 1.1.3. (P. Hall's THeorem). Let G = (A, B) be a bipartite graph. Then G has a matching of A into B if and only if $|\Gamma(X)| \ge |X|$ for all $X \subseteq A$.

Proof. Let G = (A, B) be a bipartite graph and suppose that there is a matching A to B and suppose that for some subgraph $X \subseteq A$ we have $|\Gamma(X)| < |X|$. If X is any subgraph of A and there exists a matching from A to B we have that there exists a matching from X to B as well. As $|\Gamma(X)| < |X|$ we have that there is at least 1 point in X that is not adjacent to any point in B hence it can not be a complete matching.

Suppose $|\Gamma(X)| \geq |X|$ for all $X \subseteq A$. Then using mathematical induction on |A|. If |A| = 0 then we would have a complete matching. If |A| = 1 then by our assumption we have that $|\Gamma(A)| \geq |A|$ hence there is a matching. Now suppose that for all $X \subset A$, $X \neq \emptyset$, $|X| < |\Gamma(X)|$ holds. Let $a \in A, b \in B$ be two adjacent points. Let G' = G - a - b and let X be any subset of A - a. If $X = \emptyset$, then $|X| = |\Gamma_{G'}(X)|$ so assume $X \neq \emptyset$. Since $X \neq A$, $|X| < |\Gamma_{G'}(X)|$ based on the assumption.