DyeVC: An Approach for Monitoring and Visualizing Distributed Repositories

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**Abstract:** Software development using distributed version control systems has become more frequent recently. Such systems bring more flexibility, but also greater complexity to manage and monitor the multiple existing repositories as well as their myriad of branches. In this paper, we propose DyeVC, an approach to assist developers and repository administrators in identifying dependencies among clones of distributed repositories. It allows understanding what is going on around one’s clone and depicting the relationship between existing clones. DyeVC was evaluated over open source projects, showing how they could benefit from having such kind of tool in place. We also ran an observational study and a benchmark over DyeVC, and the results were promising: it was considered easy to use and fast for most repository history exploration operations, while providing the expected answers.

**Keywords:** Distributed Version Control, Monitoring, Visualization, Awareness

# Introduction

Version Control Systems (VCS) date back to the 70s, when SCCS emerged [1]. Their primary purpose is to keep software development under control [2]. Along these almost 40 years, VCSs have evolved from a centralized repository with local access (e.g., SCCS and RCS [3]) to a client-server architecture (e.g., CVS [4] and Subversion [5]). More recently, distributed VCSs (DVCS) arose (e.g., Git [6] and Mercurial [7]) allowing clones of the entire repository in different locations. According to a survey conducted by the Eclipse community [8], Git and Github combined usage increased from 6.8% to 42.9% between 2010 and 2014 (a growth greater than 500%). During this same period, Subversion and CVS combined usage decreased from 71% to 34.4%. This clearly shows momentum and a strong tendency in the adoption of DVCSs in the open source community.

Besides these changes from local to client-server and then to a distributed architecture, the concurrency control policy adopted by VCSs also changed from lock-based (pessimistic) to branch-based (optimistic). According to Walrad and Strom [9], creating branches in VCSs is essential to software development because it enables concurrent development, allowing the maintenance of different versions of a system, the customization to different platforms/customers, among other features. DVCSs include better support for working with branches [10], turning the branch creation into a recurring pattern, no matter if this creation is explicitly done by executing a “branch” command or implicitly when a repository is cloned.

However, distributed software development, especially from the geographical perspective [11], brings a set of risk factors, and Configuration Management (CM) is affected by them. The increasing growth of development teams and their distribution along distant locations, together with the proliferation of branches, introduce additional complexity for perceiving actions performed in parallel by different developers. According to Perry et al. [12], concurrent development increases the number of defects in software. Besides, Silva et al. [13] say that branches are frequently used for promoting isolation among developers, postponing the perception of conflicts that result from changes made by co-workers. These conflicts are noticed only after pulling changes in the context of DVCSs. Moreover, Brun et al. [14] show that, even using modern DVCSs, conflicts during merges are frequent, persistent, and appear not only as overlapping textual edits (i.e., physical conflicts) but also as subsequent build (i.e., syntactic conflicts) and test failures (i.e., semantic conflicts).

By enabling repository clones, DVCSs expand the branching possibilities discussed by Appleton et al. [15], allowing several repositories to coexist with fragments of the project history. This may lead to complex topologies where changes can be sent to or received from any clone. This scenario generates traffic similar to that of peer-to-peer applications. In practice, projects impose some restrictions over this topology freedom. However, it can be still much more complex than the traditional client-server topology found in centralized VCS.

With this diversity of topologies, managing the evolution of a complex system becomes a tough task, making it difficult to find answers to the following questions: (Q1) Which clones were created from a repository? (Q2) What are the communication paths among different clones? (Q3) Which changes are under work in parallel (in different clones or different branches) and which of them are available to be incorporated into others’ clones? Most of the existing works, such as Palantir [16], FASTDash [17], Lighthouse [13], CollabVS [18], Safe-Commit [19], Crystal [14], and WeCode [20], deal with question Q3, giving to the developers awareness of concurrent changes. However, they do not provide an overview of the topology of repositories, indicating which commits belong to which clones. This overview is essential to understand the distributed evolution of the project.

In order to answer the aforementioned questions, we propose DyeVC[[1]](#footnote-1), a novel monitoring and visualization approach for DVCS that gathers information about different repositories and presents them visually to the user. DyeVC allows developers to perceive how their repository evolved over time and how this evolution compares to the evolution of other repositories in the project. DyeVC’s main goal is two-fold: increasing the developers’ knowledge of what is going on around their repository and the repositories of their teammates, and enabling repository administrators to visualize the relationship between existing clones.

This paper extends a previous conference paper [21] by including a more thorough discussion about our approach, including how DyeVC discovers the topology and a formal definition of the process underneath DyeVC. We also added an automatic collapsing feature, allowing DyeVC to deal with larger repositories. The performance study was expanded to present an evaluation of the automatic collapsing feature. Finally, we included a deeper comparison of DyeVC with its related work. This paper is organized as follows: Section 2 shows a motivational example. Section 3 presents the DyeVC approach. Section 4 presents the technologies used in our prototype implementation. Section 5 describes the evaluation of DyeVC. Section 6 discusses related work and Section 7 concludes the paper and presents some future work.

# Motivational Example

Figure 1 shows a scenario with some developers, each one owning a clone of the repository originally created at Xavier Institute. Xavier Institute acts like a central repository, where code developed by all teams is integrated, tested, and released to production. There is a team working at Xavier Institute, led by Professor Xavier, and a remote developer (Storm) that periodically receives updates from the Institute. Outside the Institute, Wolverine leads a remote team located in a different site, which is constantly synchronized with the Institute. Solid lines in Figure 1 indicate data being pushed, whereas dotted lines indicate data being pulled. Thus, for example, Rogue can both pull updates from Gambit and push updates to him, and Beast can pull updates from Rogue, but cannot push updates to her.

Each one of the developers has a complete copy of the repository. Luckily, this scenario has a CM Plan in action, otherwise each one would be able to send and receive updates to and from any other, leading to a total of different possibilities of communication (where n is the number of developers in the topology). In practice, however, this limit is not reached: while interaction amongst some developers is frequent, it may happen that others have no idea about the existence of some coworkers. It occurs with Mystique and Nightcrawler, for example, where there is no direct communication.

As an example, from a developer’s point of view, like Beast, how can he know at a given moment if there are commits in Rogue, in Gambit, or in Nightcrawler clones that were not pulled yet? Alternatively, would be the case that there are local commits pending to be pushed to Gambit? Beast could certainly periodically pull changes from his peers, checking if there were updates available, but this would be a manual procedure, prone to be forgotten. It would be more practical if Beast could have an up to date knowledge of his peers, warning him about any local or remote updates that had not been synchronized yet. On the other hand, from an administrator’s point of view, how can she know which are the existing clones of a project and how they relate among each other? How can she know if there are pending commits to be sent from a staging repository to a production one?

# DyeVC Approach

Aiming at supporting both developers and repository administrators in understanding the interaction of repository clones, the main features of DyeVC include: (1) a mechanism to gather information from a set of clones and (2) a set of extensible views with different levels of detail, which let DyeVC users visualize this information. We detail in the following sub-section how DyeVC gathers information from DVCSs. Next, we discuss how this information is presented using different levels of detail. Finally, we show what happens behind the scenes, discussing the algorithm involved in the data synchronization process.

## Information Gathering

DyeVC continuously gathers information from interrelated clones, starting from clones registered by the user. As shown in Figure 2, for each registered clone *rep*, DyeVC transparently creates a local clone *rep’* in the user’s home folder to fetch data from all of the peers that *rep* communicates with. Data is gathered by DyeVC instances running at each user machine and is stored in a central document database. In this way, information from one DyeVC instance is made available to every other instance in the topology.

DyeVC gathers information from registered clones in the user’s machine and also from their peers, which are clones that communicate with them. Since there is a communication path between a registered clone and its peers (either to push or pull data), we are able to analyze the commits that exist in these peers. This allows us to present a broader topology visualization that contains not only registered clones, but also those that have a push or pull relationship with them. DyeVC finds out related clones by looking at the remote repositories registered in the DVCS configuration. More details on how data is gathered are explained in section 3.3.

Figure 3 shows how DyeVC discovers the topology from the nodes where it is running and the registered clones. Blue nodes represent registered clones, where DyeVC is running, yellow nodes represent known clones located at nodes where DyeVC is not running, dashed nodes and dashed lines represent clones and communication paths that are not known yet. Suppose a scenario where the existing clones and interdependencies are shown in Figure 3.a. After installing DyeVC and registering clone 3, DyeVC finds out that this clone communicates with clones 1, 2, and 4 (either by pushing to or pulling from them), as shown in Figure 3.b. Later on, clone 4 is registered and clone 5 is included as a known clone in the topology (Figure 3.c). Clone 6 is the next to be registered, allowing DyeVC to discover that clone 7 also exists, as well as the communication between clone 6 and clone 1, which was already a known clone (Figure 3.d). Assuming that no other clones are registered, the known topology is shown in Figure 3.e. Notice that, although only clones 3, 4, and 6 were registered, DyeVC is also aware of the existence of clones 1, 2, 5, and 7. Only clone 8 will not be known, as well as some communication paths between clones that were not registered (1-2, 1-5, 1-8, and 7-8).

DyeVC finds out related clones by looking at the remote repositories, which are registered in Git’s *config* file of each clone. Figure 4 shows an example of this configuration, taken from a local clone of the *DyeVC* project, where there is a remote named *origin*, which is located at *github.com/gems-uff/dyevc*.

Data stored at the central database follows the metamodel presented in Figure 5. A *Project* groups repository clones of the same system. Clones are stored as *RepositoryInfo* and are identified by an id and a meaningful clone name provided by the user. A *RepositoryInfo* has a list of clones to which it pushes data and a list of clones from which it pulls data. These lists are represented respectively by the self-associations *pushesTo* and *pullsFrom*. Finally, a *RepositoryInfo* stores the hostname where it resides (e.g., a server name or localhost) and its path (be it an operating system path or a URL).

Branches are part of a *RepositoryInfo*. A *Branch* has a name and a boolean attribute *isTracked*, which is true if the branch tracks a remote branch. A *RepositoryInfo* may have one or many branches (it must have at least one branch, which is the main one). A *Branch* has two associations with *CommitInfo*: through the first association, a *Branch* knows which commit is its head and, conversely, a commit knows which branches point to it as a head. The second association represents which commits are reachable from a given branch and, conversely, the branches from which the commit is reachable.

The finer grain of information is the *CommitInfo*, which represents each commit in the topology. A commit is identified by a hash code and it refers to its parents (except for the first commit in the repository, which does not have any parent). As each commit may not exist in all clones of the topology, we store the list of clones where each commit can be found (*foundIn* association). We also store the committer, the commit message, and whether the commits belong to tracked branches or to non-tracked branches.

## Information Visualization

DyeVC presents information in four different levels of detail: Level 1 shows high-level notifications about registered repositories; Level 2 shows the whole topology of a given project. Level 3 zooms into the branches of the repository, showing the status of each tracked branch. Lastly, Level 4 zooms into the commits of the repository, showing a visual log with information about each commit. The following sections discuss these levels.

### Level 1: Notifications

In Level 1, our approach periodically monitors registered repositories and presents notifications whenever a change is detected in any known peer. The period between subsequent runs is configurable, and notifications are presented in the system notification area, in a non-obtrusive way. Figure 6 shows an example of this kind of notification, where DyeVC detected changes in two different repositories. The notification shows the repository id, the clone name, and the project (system) name. Clicking on the balloon opens DyeVC main screen.

### Level 2: Topology

Aiming at helping answering questions Q1 and Q2, we present a topology view showing all repositories for a given project (Figure 7), where each node represents a known clone. A blue computer represents the current user clone and black computers represent other clones where DyeVC is running. Servers represent central repositories, that do not pull from nor push to any other clone, or clones where DyeVC is not running. Both kinds of nodes use the same representation because, once DyeVC is not running at a given clone, we cannot infer the *pushesTo* and *pullsFrom* lists, which will thus be empty as in a server.

Each edge in the graph represents a relationship between two repositories. Continuous edges mean that the source clone pushes to the destination clone, whereas dashed edges mean that the destination clone pulls from the source clone. The edge labels show two numbers separated by a dash. The first and second numbers represent how many commits in tracked and non-tracked branches of the source clone are missing in the destination clone, respectively. The edge colors are used to represent the synchronization status: green edges mean that both clones are synchronized (i.e., the destination clone has all commits in the source clone), whereas red edges mean that the pair is not synchronized and indicates the direction that is missing commits. For example, it is possible to observe in Figure 7 that the current user clone (blue computer) is hosted at *cmcdell* and is named dyevc. This clone pulls from *gems-uff/dyevc*, which is located at github.com, and there are four tracked commits ready to be pulled (i.e., commits that exist in the remote repository and do not exist locally). It also pushes to the same peer, having five tracked commits ready to be pushed.

### Level3: Tracked branches

Aiming at answering question Q3, DyeVC’s main screen (see Figure 8) shows Level 3 information, allowing one to depict the status of each tracked branch of registered repositories regarding their peers. This information is complemented with that of Level 4, shown in the next section.

The status evaluation considers the existing commits in each repository individually. Due to the nature of DVCS, old data is almost never deleted and commits are cumulative. Thus, if commit *N*is created over commit *N – 1*, the existence of commit *N* in a given repository implies that commit *N – 1* also exists in the repository. In this way, by using set theory it is possible to subtract the set of commits in the local repository from the set of commits in its peers, resulting in the set of commits not pulled yet. In this case, local repository will be *behind* its peers (arrow down in Figure 8). Conversely, subtracting the sets in the inverse order will result in the set of commits not pushed yet, meaning that local repository is ahead its peers (arrow up). When both sets are empty, local repository is synchronized (green checkmark in Figure 8) and when both sets have elements, it is both ahead and behind its peer (arrow up and down in Figure 8).

| **Repository** | **Wolverine** | **Gambit** | **Rogue** | **Nightcrawler** | **Beast** |
| --- | --- | --- | --- | --- | --- |
| Commits | 10; 11 | 10; 11 | 10; 12 | 10; 11; 13 | 10 |

Table 1. Existing Commits in Each Repository

Let us assume that each commit is represented by an integer number to illustrate how our approach works. At a giving moment, the local repositories of each developer have the commits shown in Table 1. Considering the synchronization paths presented in the right-hand-side of Figure 1, the perception of each developer regarding to their known peers is shown in Table 2. Notice that the perceptions are not symmetric. For instance, as Gambit does not pull updates from Nightcrawler, there is no sense in giving him information regarding Nightcrawler. Furthermore, it is uncommon to have a scenario where pushes are performed from a developer to another (such as the one between Beast and Gambit). Generally, what happens is that a developer pulls from another (for example, between Gambit and Nightcrawler). This avoids inadvertent inclusion of commits inside others’ clones. Although infrequent, this scenario helps in understanding the need to have awareness about who are the peers in a project and what are their interdependencies.

| **Repository** | **Wolverine** | **Gambit** | **Rogue** | **Nightcrawler** | **Beast** |
| --- | --- | --- | --- | --- | --- |
| Wolverine | - | - | - | - | - |
| Gambit | check_32 | - | - | - | - |
| Rogue | - | aheadbehind_ylw_32 | - | - | - |
| Nightcrawler | - | check_32 | aheadbehind_ylw_32 | - | - |
| Beast | - | behind_ylw_32 | behind_ylw_32 | behind_ylw_32 | - |

Table 2. Status of Each Repository Based on Known Remote Repositories

### Level 4: Commits

Level 4 complements information of Level 3 to provide an answer to Question Q3. Differently from the usual repository version graph, it presents a combined version graph of the whole topology (Figure 9). Each vertex in the graph represents either a known commit in the topology, which is named after its hash’s five initial characters (e.g., the node labeled *2e10a*in Figure 9), or a collapsed node, representing several commits blended together. We implement two ways of collapsing nodes in order to provide a better understanding over huge amounts of data: manual and automatic. Manually collapsed nodes are named after the number of contained nodes, such as the white node containing 118 commits and the green node containing 24 commits in Figure 9). Automatically collapsed nodes have ellipses before and after the number of contained nodes in their names (if the first collapse of Figure 9 were automatic, its name would be “…118…”). Automatic collapse are detailed in Section 3.2.5.

Thicker borders denote that the commit is a branch’s head (e.g., commit *ea6a4*). Commits are drawn according to their precedence order. Thus, if a commit *N*is created over a commit *N – 1*, then commit *N* will be located to the right of commit *N – 1*. For each commit, DyeVC presents the information described in Figure 5 (gathered from the central database), along with information that is read in real time from the repository metadata, such as branches that point to that commit and affected files (added, edited, and deleted).

This visualization contains all commits of all clones in an integrated graph. Each commit is painted according to its existence in the local repository and in the peers’ repositories. Ordinary commits that exist locally and in all peers are painted in white. Green commits are ready to be pushed, as they exist locally but do not exist in peers of the push list. Yellow commits need attention because they exist in at least one peer in the pull list, but do not exist locally, meaning that they may be pulled. Red commits do not exist locally and are not available to be pulled, as they exist only in clones that are not peers. Finally, gray commits belong to non-tracked branches, so they can neither be pushed nor pulled. Heads in these branches are not identified with thicker borders.

This visualization can easily have thousands of nodes, one for each commit in the topology. Nevertheless, despite the high number of nodes, users are generally interested in the most recent commits. As we show the commits following a chronological order, from left to right, most recent commits will be at the right part of the visualization. DyeVC positions the graph so that these commits are shown when opening the visualization.

### Automatic Collapse

We identified two common node structures that can be automatically collapsed: sequential and parallel. The former contains a sequence of commits of the same type, where each of them has degree two, i.e., nodes with just one ancestor and one successor. This kind of structure can be collapsed because it does not represent any additional information besides the fact that some sequential work was performed. Figure 10 shows examples of sequences of commits, highlighted in red, which could be collapsed, producing the graph shown in Figure 11 (still in red). On the other hand, the later contains one fork node and one merge node, with at most one (regular or collapsed) node in each branch, between the fork and the merge nodes. Figure 11 shows examples highlighted in yellow of this parallel structure. The result of the collapse is shown in Figure 12.

In order to benefit from both sequential and parallel collapse strategies together, we implemented an iterative algorithm that work in phases. Each iteration is executed in linear time complexity. The first phase collapses sequential structures and the second phase collapse parallel structures. These phases can be repeated, as the collapse of parallel structures may lead to new sequential structures. For instance, after applying parallel collapses over the graph shown in Figure 11, a new sequential structure is formed, as illustrated in Figure 12. The iteration would lead to a new collapse, and so on. As previously discussed, collapses are performed just for commits of the same type (same color, discussed in section 3.2.4), reducing the size of the graph without compromising the quality of the information shown in the graph.

## Behind the Scenes

The process underneath DyeVC can be formally defined using Set Theory. We can define a project p as a tuple , where R is the set of all cloned repositories of p monitored by DyeVC, C is the set of all commits of p, is the set of commits of p in the DyeVC database, and is the set of named branches of p. Each repository is a tuple , where is the set of repositories that is allowed to push to, is the set of repositories that is allowed to pull from, is the set of commits in in the previous execution of DyeVC, and is the set of commits in in the current execution of DyeVC (Figure 13.a). Each commit has a set of parent commits . Commits are organized in a directed acyclic graph (Figure 13.b), where the first commit of the project has no parent (e.g., commit A in Figure 13.b), revision commits have only one parent (e.g., commit B in Figure 13.b), and merge commits have two or more parents (e.g., commit I in Figure 13.b). All reachable commits from form its history, including itself and the transitive closure over its parents (e.g., {A, B, E, F, H, I, J} is the history of commit J in Figure 13.b). The history of is formally defined as:

Each branch is a tuple , where name is the name of and is the tip (i.e., head) of . Consequently, contains all reachable commits of .

With this foundation established, we can now formalize the process of updating commits in the topology. For a local repository being monitored by DyeVC, the rare situations where a commit is deleted can be formally defined as:

Each locally deleted commit should be removed from if no other repository still contains this commit. Conversely, the new commits in since the previous monitoring cycle can be formally defined as:

Each locally added commit that is not already in the database () should be inserted in . This verification is necessary because some of the locally added commits might have already been inserted into the database by another instance of DyeVC.

Moreover, we can formalize the identification of repositories that contain a specific commit and the repositories that are ahead or behind of a given repository. This information is necessary for building some of our visualizations. We formally define the repositories that contain a commit as:

We formally define from which repositories is ahead or behind as:

Finally, we can also formalize the commits that are ahead or behind two specific repositories and the branches that a commit belongs to. This is also necessary for some of our visualizations. Considering two repositories , we formally define the commits ahead or behind regarding as:

We formally define the branches that a commit belongs to as:

# Implementation

We implemented our approach as a Java application launched via Java Web Start Technology. It currently monitors Git repositories, as it is the most used DVCS nowadays [8]. The source code and the link to download the tool via Java Web Start can be found at https://github.com/gems-uff/dyevc. The tool gathers information from repositories using JGit library[[2]](#footnote-2), which allows using our approach without having a Git client installed.

Gathered information is stored in a central document database running MongoDB. We hosted our database on a free MongoDB instance provided by MongoLab. We did not use MongoDB proprietary API, which would demand opening specific ports to connect to MongoDB. Instead, we opted to use MongoLab’s RESTful (*Representational State Transfer*) API. RESTful APIs [22] have the advantage of being available using standard HTTP and HTTPS protocols. In this way, our approach can be used in environments protected with firewalls without major problems. In order to use this RESTful API, we implemented a *MongoLab Provider*, which translates the application methods into RESTful commands and vice-versa. It also serializes/deserializes the application objects to/from JSON (*JavaScript Object Notation*) representations to be used through the RESTful commands.

We present the gathered information as a series of graphs by using the JUNG (*Java Universal Network/Graph*) library[[3]](#footnote-3), from which DyeVC inherits the ability to extend existing layouts and filters. All graphs present similar behavior, allowing the window to be zoomed in or out, whether the user wants to see details of a particular area or an overview of the entire graph. By changing the window mode from *transforming* to *picking*, it is possible to select a group of nodes and collapse them into one node, or simply drag them into new positions to have a better understanding of parts with too many crossing lines.

# Evaluation

In order to evaluate our approach we performed two experiments and an observational study. First, we conducted a *post-hoc* analysis over the JQuery project[[4]](#footnote-4), an open-source project, aiming at checking if DyeVC can help answering questions Q1-Q3. Next, we conducted an observational study involving four participants that used DyeVC. This study also used the JQuery project. Finally, we run DyeVC over some open-source projects of different sizes and from different sources, aiming at evaluating the scalability of our approach.

## Post-hoc Study

We conducted a *post-hoc* analysis using a real open source project to demonstrate that our approach can help answering questions Q1-Q3. The selected project, JQuery, began in 2006 and had 6,222 commits by the time of the evaluation. We reconstructed the repository history, simulating the actions that occurred in the past. We do not replicate the repository history here, due to its size, but it is publicly available at Github. Automatically generated comments helped us to depict specific flows. For example, the comment “*Merge branch 'master' of https://github.com/scottjehl/ jquery into scottjehl-master*” tells us that there was a user named “*scottjehl*” and that the merge operation was done at a branch called “*scottjehl-master*”. Although one might perform a merge manually and insert a different text in the comment, this did not compromise our analysis because we had a focus on depicting some of the merge situations, and not all of them.

Due to the operating mode of Git, some details are missing, but these details do not compromise our analysis. The first one is the moment when a clone arises or deceases. This information does not exist anywhere in the repository. We inferred the creation of clones by looking at the commit messages (a commit by developer X led to the creation of a clone named X). Clones created at a given time stayed alive for the rest of the analysis.

The second missing detail is that, although we had the commit dates and times in the repository history, these dates and times were not guaranteed to be correct. This occurs because DVCSs do not have a central clock. Each commit is registered with the local time at the machine where the clone is located, which could lead to commits in the history with a predecessor in the future, depending on when and where each commit was performed. This missing detail is not important, because the order of commits is not depicted using their times, but using the pointers that Git maintains from a commit to its parents, as discussed in section 3.1. We can use these dates, but not as an authoritative information.

We chose a moment in time when three developers were involved, performing commits and merging changes in the repository. We created three clones for these developers, named after their user names: *jeresig*, *adam*, and *aakosh*. Figure 14 shows the topology view on Sep 24 2010, when *aakosh* had 121 commits pending to be pushed to the central repository (hereafter called *central-repo*). Figure 15 shows part of *aakosh’s* commit history and how DyeVC represents commits pending to be pushed (green nodes).

Later on, *aakoch* pushed his commits to *central-repo***.** In the meantime, both *adam* and *jeresig* committed some changes. Before they pushed their work to *central-repo*, *adam’s* last commit was on Jun 21 2010 and *jeresig’s* on Sep 27 2010.At this moment, we registered them to be monitored by DyeVC. Figure 16 shows the topology view after this registration on Sep 27 2010. Here, we can see that *aakoch* was synchronized with *central-repo*, whereas *adam* and *jeresig* had pending actions.

At this point, we can revisit questions Q1 and Q2:

Q1: *Which clones were created from a repository?* DyeVC’s topology view (Figure 16) shows all the clones where it is running, and also discovers other clones connected to them, even if it is not running there.

Q2: *What are the communication paths among different clones?* DyeVC’s topology view (Figure 16) shows the dependencies between the peers in the topology, as well as the number of commits ahead or behind in each of these clones.

*Adam* had 121 commits to pull from *central-repo*, what is corroborated by the details of his tracked branches (master branch in Figure 17). He also had a non-tracked commit pending to be pushed. Non-tracked commits are not shown in the tracked branches view, but we can see them in gray in the commit history views. Figure 9 shows the collapsed commit history for *jeresig*, where we can see adam’s non tracked commit with hash *a2bd8*.

The repository history leads us to think that *jeresig* is a core developer of this project, because he performed most of the merges to the master branch. Looking at Figure 16, we see that he had 26 commits pending to be pushed to *central-repo*. These 26 commits can be seen at *aakoch’s* commit history (Figure 18) as red commits, since they could not be pulled by *aakoch* until *jeresig* has pushed them to *central-repo*. There was also a commit in central-repo pending to be pulled by *jeresig*. If we look back at Figure 9 we see that the only yellow commit is *a0887*, made by *aakoch*. This tells us that *jeresig* pulled changes from *central-repo* just before *aakoch* pushed commit *a0887*. If we look at Figure 19, we see that all the pending commits (those that were pending to be pushed and pulled) are related to the same branch (master). This tells us that, if *jeresig* wanted to push these commits to *central-repo*, he would have to perform a pull operation before.

This analysis helps us revisit and answer Q3:

Q3: Which changes are under work in parallel (in different clones or different branches) and which of them are available to be incorporated into others’ clones? New commits in tracked branches of peers can be easily found by looking at Level 3 information (tracked branches, shown in Figure 17 and Figure 19). This view shows to which branch these commits are related and how many new commits exist. If we want to look at each commit individually, we can look at Level 4 information (commit history, shown in Figure 15 and Figure 18) and notice the yellow nodes. Additionally, Level 4 information can be used to find new commits in repositories that are not peers (red nodes), or new commits in non-tracked branches (gray nodes).

## Observational Study

We conducted an observational study over the same project used in the post hoc analysis (JQuery) to assess the capability of the visualizations provided by DyeVC in supporting developers and repository administrators. The study was conducted with four volunteers, which had previous experience with DVCS. They were graduate students from the Software Engineering research area at Universidade Federal Fluminense (UFF). Four sessions were conducted, each of them with one subject. The study was divided in two phases (without and with DyeVC), each one with two scenarios, where the subject had to answer questions related to usual work with DVCS. In Scenario 1, the subject played the developer role, working in a clone named *aakoch*. In Scenario 2, the subject played the repository administrator role. The following questions were posed: Q1.1 What is the status of your clone, compared to the central repository? Q1.2 Who else is working in the JQuery project (other clones)? Q1.3 Which files were modified in commit *5d454*? Q2.1 What are the existing clones for JQuery project? Q2.2 Which clones are synchronized with the central repository? Q2.3 How many commits in tracked branches are pending to be sent to the central repository? Q2.4 Is there any commit in non-tracked branches? Where?

In Phase 1 (without DyeVC), DyeVC was not in place and the subject answered the questions using any desired DVCS client among the ones available in the computer used in the experiment: *gitk***,** *Tortoise Git***,** *Git Bash*, and *SourceTree*. Participants were allowed to access the Internet and search any other procedure or tool that could help in answering the questions. After that, the subject watched a 10-minute video presenting DyeVC and started Phase 2 (with DyeVC), which consisted in answering the same questions with the help of DyeVC. The possible answers in Phase 2 were either “keep the answer of Phase 1”, meaning that using DyeVC did not change the subject perception, or a different answer, otherwise.

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| **Subject** | **Scenario 1** | | **Scenario 2** | |
| **Phase 1** | **Phase 2** | **Phase 1** | **Phase 2** |
| P1 | 14 | 5 | - | 6 |
| P2 | 13 | 6 | - | 5 |
| P3 | 3 | 2 | - | 4 |
| P4 | 10 | 2 | - | 10 |

Table 3. Time spent (in minutes) to answer each question

Table 3 presents the time spent by each subject to answer each question in both scenarios and both phases. It is possible to notice, by looking at Table 3, that all subjects took less time to complete Scenario 1 (developer role) in Phase 2 (with DyeVC). For Scenario 2 (admin role), times for Phase 1 (without DyeVC) are not shown because none of the subjects managed to answer the questions without using DyeVC.

The overall results of this study were positive. In Phase 1 (without DyeVC), each subject used a different tool and followed a different procedure to find answers regarding DVCS usage. Subjects were able to answer correctly questions Q1.1 and Q1.3 whether using DyeVC or not. In addition, further questions were answered correctly only by using DyeVC.

The subjects also answered an exit questionnaire. All subjects found easy to interact with DyeVC, to identify related repositories, and to use the operations available. They consensually elected the topology visualization as the most helpful visualization in DyeVC. In addition, by using the Product Reaction Cards[[5]](#footnote-5), 3 out of the 4 subjects stated that DyeVC is helpful and easy to use.

## Performance Study

In order to evaluate the scalability of our approach, we measured the time spent to perform the most common DyeVC operations. We used projects of different sizes and hosted in different Git servers. Table 4 shows the monitored projects (name and hosting service), the repository metrics – number of commits, disk usage, and number of files – and the time spent to run some background and foreground operations in DyeVC. All measurements were taken in the same period of the day and from the same machine, a Core Duo CPU at 2.53 GHz, with 4GB RAM running Windows 8.1 Professional 64 bits, connected to the internet at 35 Mbit/s.

| **Repository** | **Hosting** | **Repository metrics** | | | **Foreground operations** | | | **Background operations times (s)** | | | |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
| **# commits** | **Size (MB)** | **# files** | **Commit History** | | **Topology** | **Insert 1st** | **Insert 2nd** | **Check**  **Branches** | **Update**  **Topology** |
| **Time (s)** | **Memory Usage[[6]](#footnote-6)** | **Time (s)** |
| DyeVC | github.com | 187 | 1.0 | 539 | 3.5 | 15 | 2.7 | 12.4 | 16.1 | 1.7 | 4.4 |
| SAPOS | github.com | 702 | 7.0 | 685 | 5.6 | 19 | 3.2 | 20.8 | 22.6 | 1.8 | 5.2 |
| jgit | eclipse.org | 2,979 | 10.0 | 1,595 | 18.4 | 512 | 3.4 | 42.4 | 46.0 | 5.9 | 6.8 |
| egit | eclipse.org | 3,775 | 27.0 | 1,478 | 21.3 | 559 | 3.7 | 49.6 | 46.6 | 4.2 | 7.3 |
| jquery | github.com | 5,518 | 20.0 | 253 | 65.0 | 1,121 | 4.1 | 40.0 | 37.4 | 1.4 | 9.4 |
| Tortoise Git | code.google.com | 6,166 | 85.0 | 3,220 | 68.0 | 492 | 4.2 | 39.0 | 36.0 | 1.6 | 9.6 |
| Gitextensions | github.com | 6,417 | 448.0 | 1,549 | 73.0 | 1,529 | 17.0 | 155.8 | 129.0 | 1.6 | 10.6 |
| Drupal | drupal.org | 23,922 | 84.4 | 9,290 | - | - | 18.0 | 102.0 | 95.0 | 2.0 | 18.0 |
| ExpressoLivre | gitorious.org | 25,822 | 141.0 | 20,729 | - | - | 18.2 | 110.0 | 102.0 | 2.1 | 19.3 |
| Git | github.com | 35,260 | 98.0 | 2,656 | - | - | 19.4 | 196.0 | 158.6 | 3.4 | 40.0 |

Table 4. Scalability results of DyeVC for repositories with different sizes

We measured the main operations of our approach: “Insert 1st”, invoked when the user includes the first repository of a given system to be monitored; “Insert 2nd”, invoked when the user includes a repository to be monitored in a system that already have registered repositories; “Commit History”, invoked when the user requests to see the commit history of a given repository; “Topology”, invoked when the user wants to see the topology of repositories of a given system; “Check Branches”, invoked periodically to check all the monitored repositories, searching for ahead or behind commits; and “Update Topology”, invoked periodically to update the topology information in the central database. This last operation updates the existing repositories, their peers, and the existing commits, marking in which repositories each commit is found.

It is possible to notice that the “Commit History” operation has no values for the last three repositories. This occurs because, as the number of commits increases, more memory is used to calculate the commit history graph. The current algorithm has an *O(x2)* space complexity (being *x* the number of commits). Our experiment computer was configured with a 2 GB maximum Java Heap Size, which let us analyze repositories with up to 6K commits. This limitation occurs mainly because of JUNG.

Table 5 shows the correlation between each repository size metric and the DyeVC operations execution time, according to the Spearman’s rank correlation coefficient [23]. This correlation coefficient measures the monotonic relation between two variables and ranges from *-1* to *1*. Values of *1* or *-1* mean that each variable is a perfect (increasing or decreasing) monotone function of the other. A value of *0* means that there is no correlation between the variables.

|  |  |  |  |
| --- | --- | --- | --- |
| **Operation** | **# commits** | **Size** | **# files** |
| Insert 1st | 0.85 | 0.83 | 0.76 |
| Insert 2nd | 0.85 | 0.83 | 0.76 |
| Check Branches | 0.07 | -0.05 | 0.72 |
| Update Topology | 1.00 | 0.88 | 0.52 |
| CommitHistory | 1.00 | 0.96 | -0.04 |
| Topology | 1.00 | 0.88 | 0.52 |

Table 5. Spearman’s rank correlation coefficient between repository size metrics and DyeVC operations time

Looking at Table 5, it is possible to notice that, except for the “Check Branches” operation, all other operation times are strongly correlated to the number of commits and repository size. This is due to the nature of these operations, which update or show information about all commits in the repository. On the other hand, except for the “Commit History” operation, all other operation times correlate with the number of files. This is also expected due to the nature of “Commit History” operation, which does not dig into the changed files.

### Automatic Collapse Evaluation

In addition, we studied the impact of the automatic collapse algorithm in the “Commit History” operation performance. These experiments were performed at a later time in comparison with the results obtained in previous section. Consequently, the repository metrics are slightly different. The repository size, number of commits, and number of files are generally higher, as shown in Table 6.

The design of the experiment was as follows. First, the “Commit History” operation is performed without using the automatic collapse. Afterwards, sequential and parallel collapse strategies described in section 3.2.5 were used in order to simplify the structure of the commit graph, collapsing the corresponding node structures. The execution of the sequential strategy is the first stage, and the parallel strategy is the second stage of each iteration of the automatic collapse algorithm. Moreover, after each stage, running time and memory consumption were measured. The experiment was executed in a Core i7 CPU at 2.00 GHz, with 16GB of RAM running Windows 7 64 bits.

|  |  |  |  |
| --- | --- | --- | --- |
| **Repository** | **Characteristics** | | |
| **Size (MB)** | **# files** | **# commits** |
| DyeVC | 3.2 | 745 | 228 |
| SAPOS | 18.8 | 668 | 1,245 |
| jgit | 39.3 | 1,902 | 4,741 |
| egit | 63.6 | 1,779 | 4,983 |
| jquery | 29.2 | 296 | 7,291 |
| Gitextensions | 94.9 | 1,710 | 8,146 |
| Tortoise Git | 168 | 3,518 | 8,442 |
| Drupal | 176 | 10,285 | 38,047 |
| ExpressoLivre | 366 | 21,592 | 27,079 |
| Git | 104 | 3,026 | 46,794 |

Table 6. Characterization of the repositories used in the evaluation of the automatic collapse algorithm

We evaluated the capability of the automatic collapse algorithm to reduce the number of nodes in the commit graph without compromising the quality of the information shown in the graph (collapses are performed just for commits of the same type, i.e., same color discussed in section 3.2.4).

Table 7 shows the reduction achieved by two iterations of the algorithm. The number of iterations was set to two due to empirical observation that there was almost no reduction after a second iteration. With two iterations, the algorithm is capable of reducing the number of nodes by an average of 73% compared to the original graph. In some cases, such as Drupal or ExpressoLivre, which are repositories that we could not analyze before, the nodes reduction surpassed 90%, allowing us to visualize their commit history graph after the automatic collapse process.

|  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- |
| **Repository** | **Before Collapse** | **Iteration 1** | | | **Iteration 2** | | |
| **1st stage** | **2nd stage** | **Reduction (%)** | **1st stage** | **2nd stage** | **Reduction (%)** |
| DyeVC | 228 | 73 | 47 | 67.98 | 32 | 32 | 85.96 |
| SAPOS | 1,245 | 456 | 404 | 63.37 | 378 | 375 | 69.88 |
| jgit | 4,741 | 3,015 | 2,751 | 36.41 | 2,635 | 2,635 | 44.42 |
| egit | 4,983 | 3,007 | 2,564 | 39.65 | 2,347 | 2,329 | 53.26 |
| jquery | 7,291 | 867 | 709 | 88.11 | 609 | 603 | 91.73 |
| Gitextensions | 8,146 | 4,083 | 3,833 | 49.88 | 3,702 | 3,684 | 54.78 |
| Tortoise Git | 8,442 | 1,466 | 945 | 82.63 | 497 | 482 | 94.29 |
| Drupal | 38,047 | 903 | 697 | 97.63 | 563 | 557 | 98.54 |
| ExpressoLivre | 27,079 | 3,008 | 2,792 | 88.89 | 2,669 | 2,669 | 90.14 |
| Git | 46,794 | 24,459 | 24,216 | 47.73 | 24,094 | 24,094 | 48.51 |

Table 7. Reduction of the number of nodes by the automatic collapse algorithm

Furthermore, we analyze the running time and memory consumption of the “Commit History” operation. In particular, data collected for repositories that were visualized before and after collapse are represented in Figure 20 and Figure 21 using boxplots. The figures show that, the more collapse stages we execute, the less time is needed to represent the commit history, and the less memory is consumed for this purpose. This can be explained by the fact that the automatic collapse algorithm is linear and very fast comparing to the subsequent visualization process, and speeds up the presentation of the commit graph. Using this method, significantly lower running times and memory consumption values are obtained, compared with values before the automatic collapse. It was possible to visualize repositories with tens of thousands of nodes (Drupal and ExpressoLivre), which could not be represented before, without applying collapse process.

Git was the only repository that was not represented visually, even after the automatic collapse. The main contributing factors for this fact are: its high number of nodes, its low nodes reduction rates, and its intrinsic complexity. Additional automatic collapse heuristics seem necessary to accommodate such factors.

## Threats to Validity

While we have taken care to minimize threats to the validity of the experiment, some factors can influence the results. The usage of a *post-hoc* analysis to evaluate a real project may not reflect the exact sequence of events that occurred, although the outcome did not change. For example, when we say that *aakosh*, at some moment, had 121 commits pending to be pushed to the central repository, these commits could have been pushed at once or by a series of smaller pushes. Moreover, only one project was selected to perform the analysis, what imposes limitations from a generalization standpoint. Furthermore, we used an open source project to perform the *post-hoc* analysis, but the *modus operandi* of peers may be different in academic or industrial contexts.

In the observation study, the selection of subjects was done by asking for volunteers from students in the same research group of the author. This was necessary due to time and people restrictions. Therefore, this group might not be representative and can be biased. Moreover, there were few subjects in this study. Thus, the results may have been influenced by the size and by specific characteristics of the group. Furthermore, subjects performed tasks involving DyeVC right after knowing the approach, giving no time to subjects to assimilate the tool. Results may have been influenced by this lack of time to mature the necessary knowledge to use the approach efficiently. In addition, subjects could have answered questions in Phase 2 faster than in Phase 1 due to their learning regarding the scenario.

Finally, there is a risk regarding the instrumentation used to measure the response times during the performance evaluation. As we used a database stored over the Internet, connectivity issues and network instability may have affected the response times.

# Related Work

According to Diehl [24], software visualization can be separated into three aspects: structure, behavior, and evolution. DyeVC relates primarily with the evolution aspect, more specifically with studies that aim at improving the awareness of developers that work with distributed software development. A recent work by Steinmacher et al. [25] presents a systematic review of awareness studies, which we used to perform a forward and backward snowballing. The approaches obtained after the snowballing were divided into four groups. The first group (“*Commit notification*”) includes approaches that notify commit activities. The second group (“*Awareness of concurrent changes*”) comprises approaches that not only give the developer awareness of concurrent changes, but also inform them about conflicts. The third group (“*Repository visualization*”) includes approaches that visualize repository information. Finally, the fourth group (“*DVCS clients*”) contains commercial and open source DVCS clients.

The first group contains tools such as SVNNotifier[[7]](#footnote-7), SCMNotifier[[8]](#footnote-8), Commit Monitor[[9]](#footnote-9), SVN Radar[[10]](#footnote-10), Hg Commit Monitor[[11]](#footnote-11) and Elvin [26]. The primary focus of these approaches is on increasing the developer’s perception of concurrent work by showing notifications whenever other developers perform actions. The approaches in this group do not identify related repositories and do not provide information in different levels of details, such as status, branches, and commits. DyeVC provides these different levels of details, as shown in Section 3.2.

The second group comprises approaches that give the developer awareness of concurrent changes, sometimes informing them if conflicts are likely to occur. This group includes tools such as Palantir [16], CollabVS [18], Crystal [14], Lighthouse [13], FASTDash [17], and WeCode [20]. Among these, only Crystal and FASTDash work with DVCSs. Crystal detects physical, syntactic, and semantic conflicts in Mercurial and Git repositories (provided that the user informs the compiling and testing commands), but does not precisely deal with repositories that pull updates from more than one peer. FASTDash does not detect conflicts directly, as the previous cited studies, but provides awareness of potential conflicts, such as two programmers editing the same region of the same source file in repositories stored in Microsoft Team Foundation Server. Although DyeVC primary focus is not to detect conflicts, it can be combined with such approaches to allow conflicts and metrics analysis over DVCS.

The third group includes approaches that visualize repository information. Each approach has a different visualization focus, such as program structures [27], classes [28], lines [29], authors [30], and branch history [31][[12]](#footnote-12), [[13]](#footnote-13). The latter have the same focus of DyeVC’s Commit History visualization, but dealing only with the local repository, not showing, for example, where a given commit can be found in related repositories.

Finally, the fourth group includes commercial / open source DVCS clients, which allows one to execute operations on repositories / clones (push, pull, checkout, commit, etc.) and also visualizing the repository history, i.e., the commits, along with their attributes (comment, date, affected files, committer, etc.). For example, some Git clients include *gitk*[[14]](#footnote-14)**,***TortoiseGit*[[15]](#footnote-15), *EGit for Eclipse*[[16]](#footnote-16)**,** and *SourceTree*[[17]](#footnote-17). The data about commits shown by these tools varies, but generally involves the committer name, message, date, affected files, and a visual representation of the history. These tools, though, have no knowledge regarding peers. For this reason, they do not present commits from other clones and do not include information about where each commit can be found. It is worth noticing that we could not find any similar work showing the dependencies among several clones of a DVCS.

Table 8 compares DyeVC with each group used to classify related work presented in this section, according to the following features: notifications (Which types of notification the approach supports?); CVCS (Does the approach support CVCS?); DVCS (Does the approach support DVCS?); related repositories (Does the approach identifies related repositories?); levels of detail (Does the approach present information in different levels of detail?); multiple peers (Does the approach support repositories with multiple peers, i.e., multiple pull / push destinations?); commits in peer nodes (Does the approach detects commits in peer nodes, i.e., nodes that have a direct communication path among each other?); commits in non-peer nodes (Does the approach detect commits in non-peer nodes, i.e., nodes that do not have a direct communication path among each other?); multiple branches (Does the approach support multiple branches in DVCS?); topology (Does the approach supply any topology visualization that shows dependencies among repositories?); and, finally, commit History (Does the approach allow visualizing only a partial commit history, showing only local commits, or does it allow visualizing a full commit history, including commits in other repositories that were not synchronized yet, or that are in non-tracked branches?).

|  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- |
| **Feature** | **Commit notification** | **Awareness of concurrent changes** | **Repository visualization** | **DVCS clients** | **DyeVC** |
| Notifications | new commits | Conflicts | No | No | Status change against peers |
| CVCS | Yes | Yes | Yes | No | No |
| DVCS | Some[[18]](#footnote-18) | Some[[19]](#footnote-19) | Some[[20]](#footnote-20) | Yes | Yes |
| Related repositories | No | No | No | No | Yes |
| Levels of detail | No | No | No | No | Yes |
| Multiple peers | No | No | No | No | Yes |
| Commits in peer nodes | No | Some[[21]](#footnote-21) | Some[[22]](#footnote-22) | No | Yes |
| Commits in non-peer nodes | No | No | No | No | Yes |
| Multiple branches | No | No | No | Yes | Yes |
| Topology | No | No | No | No | Yes |
| Commit history | No | No | Some[[23]](#footnote-23) / Partial[[24]](#footnote-24) | Partial24 | Full |

Table 8.Comparing DyeVC features with related work

All in all, among related work, *Crystal* is the most similar to DyeVC, and deserves a deeper comparison. Both approaches work with DVCSs (besides Git, *Crystal* also supports Mercurial) and use working copies to perform analyses, but there are major differences between them. *Crystal*’s goal is to identify conflicts among pairs of repositories, whereas *DyeVC’s* goal is to provide awareness regarding the existing peers and their synchronization, in different levels. To identify repositories, *Crystal* demands the user to point out all repositories they want to compare, whereas *DyeVC* requires that some of the repositories be registered and it automatically looks at configuration files to discover all the repositories that one pushes to or pulls from. The repository comparison in *Crystal* is from one repository against all the other together, whereas *DyeVC* analyzes each repository against each other, providing a pairwise view and a combined view of the history. Finally, the allowed actions in *Crystal* include the ability to *push*, *pull*, *compile*, and *test* a repository, whereas *DyeVC* allows one to visualize branches status, topology, and history. In this way, we see potential to have both tools working together to better provide awareness and safety when working with DVCS.

# Conclusion

In this paper we presented DyeVC, an approach that identifies the status of a repository in contrast with its peers, which are dynamically found in an unobtrusive way. We have evaluated DyeVC on a real project, showing that it can be used to answer questions that arise when working with DVCSs. The observational study results were promising: DyeVC was considered easy to use and fast for most repository history exploration operations, while providing the expected answers. This provides initial evidence that DyeVC could effectively help developers and repository administrators by saving time and by supporting answering questions regarding DVCS usage that could not be answered before. We have also evaluated DyeVC’s performance over repositories of different sizes, and we found out that the time and space complexity of the approach are directly related to the number of commits in the repository, especially in the view levels with finer granularity.

A number of future researches arise from this work. Different visualizations can be developed to show the commit history, compacting it, for example, by automatically collapsing contiguous nodes representing commits with the same level of accessibility. DyeVC could gather additional metadata, for example, to create a visualization showing conflicts that would happen when merging two or more branches. This data could also be used to mine information in the repositories, revealing usage patterns or presenting metrics. Finally, some optimization should be done to allow DyeVC work with larger repositories.

# DECLARATIONS

## List of Abbreviations

* Application Programming Interface (API)
* Configuration Management (CM)
* Central Processing Unit (CPU)
* Centralized Version Control Systems (CVCS)
* Distributed Version Control Systems (DVCS)
* Hypertext Transfer Protocol (HTTP)
* HTTP Secure (HTTPS)
* Java Universal Network/Graph (JUNG)
* JavaScript Object Notation (JSON)
* Megabyte (MB)
* Random Access Memory (RAM)
* Representational State Transfer (RESTful)
* Unified Modeling Language (UML)
* Version Control Systems (VCS)

## Availability of Data and Materials

The source code and the link to download DyeVC via Java Web Start can be found at https://github.com/gems-uff/dyevc. All projects used in the experiment are available in their respective repositories, described in Table 4.

## Competing Interests

The authors declare that they have no competing interests.

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## Author’s Contributions

Cristiano Cesario contributed in the design and implementation of DyeVC and in the design of the automatic collapse algorithm, and was responsible for running the post-hoc, observational, and performance evaluations. Rubén Interian contributed in the design and implementation of the automatic collapse algorithm, and was responsible for running the automatic collapse evaluation. Leonardo Murta contributed in the design of DyeVC, the automatic collapse algorithm, and the experiments, and was responsible for the DyeVC formalization. All three authors contributed in writing the paper.

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Figure 1. Development scenario involving some developers

Figure 2. How DyeVC gathers information

Figure 3. DyeVC discovering the topology

Figure 4. Remote repository configuration in Git’s *config* file

Figure 5. Metamodel used to store DyeVC data

Figure 6. DyeVC showing notifications in notification area

Figure 7. Topology view for a given project

Figure 8. DyeVC main screen

Figure 9. Collapsed commit history

Figure 10. Sequential structures before automatic collapse

Figure 11. Sequential structures after automatic collapse and Parallel structures before automatic collapse

Figure 12. Parallel structures after automatic collapse

Figure 13. UML class diagram representing the DyeVC formalization (a) and a directed acyclic graph of commits (b)

Figure 14. First monitored repository in Topology view (Sep 24 2010)

Figure 15. aakoch’s commit history showing commits pending to be pushed

Figure 16. Three monitored repositories in Topology view (Sep 27 2010)

Figure 17. Adam’s tracked branches

Figure 18. Aakoch’s commit history

Figure 19. Jeresig’s tracked branches

Figure 20. Boxplot showing the running time of “Commit History” operation depending on the number of executed collapse stages

Figure 21. Boxplot showing the memory consumption of “Commit History” operation depending on the number of executed collapse stages

1. Dye is commonly used in cells to observe the cell division process. As an analogy, DyeVC allows developers to observe how a Version Control repository evolved over time. [↑](#footnote-ref-1)
2. http://www.eclipse.org/jgit/ [↑](#footnote-ref-2)
3. http://jung.sourceforge.net/ [↑](#footnote-ref-3)
4. https://github.com/jquery/jquery [↑](#footnote-ref-4)
5. Developed by and © 2002 Microsoft Corporation. All rights reserved. [↑](#footnote-ref-5)
6. Memory usage was measured in MB during the execution of “Commit History” operation. [↑](#footnote-ref-6)
7. http://svnnotifier.tigris.org/ (2012) [↑](#footnote-ref-7)
8. https://github.com/pocorall/scm-notifier (2012) [↑](#footnote-ref-8)
9. http://tools.tortoisesvn.net/CommitMonitor.html (2013) [↑](#footnote-ref-9)
10. http://code.google.com/p/svnradar/ (2011) [↑](#footnote-ref-10)
11. http://www.fsmpi.uni-bayreuth.de/~dun3/hg-commit-monitor (2009) [↑](#footnote-ref-11)
12. Visugit: https://github.com/hozumi/visugit [↑](#footnote-ref-12)
13. GitHub’s Network Graph: https://github.com/blog/39-say-hello-to-the-network-graph-visualizer [↑](#footnote-ref-13)
14. http://git-scm.com/docs/gitk [↑](#footnote-ref-14)
15. https://code.google.com/p/tortoisegit/ [↑](#footnote-ref-15)
16. http://eclipse.org/egit/ [↑](#footnote-ref-16)
17. http://www.sourcetreeapp.com/ [↑](#footnote-ref-17)
18. Exceptions are SCM Notifier and Hg Commit Monitor [↑](#footnote-ref-18)
19. Exception is Crystal [↑](#footnote-ref-19)
20. Exceptions are VisGi, Visugit, and GitHub's Network Graph [↑](#footnote-ref-20)
21. Exception is Lighthouse [↑](#footnote-ref-21)
22. Exception is GitHub's Network Graph [↑](#footnote-ref-22)
23. Visugit and GitHub's Network Graph [↑](#footnote-ref-23)
24. Approaches allow visualizing only local commits. Commits in other repositories that were not synchronized yet, or that are in non-tracked branches, are not shown. [↑](#footnote-ref-24)