

Summary of “Typed Arithmetic Expressions”

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1 INTRODUCTION

The chapter “Typed Arithmetic Expressions” from the Book “Types and Programming languages”[1] is about proofs about a typed programming language. The programming language from the chapter “Untyped Arithmetic Expressions” is reused and extended with a type system that differentiates between natural numbers (Nat) and boolean values (Bool).

Nat is defined as 0, $\text{succ } t_1$, and $\text{pred } t_1$ if t_1 is a Nat. Bool is defined as *true*, *false*, and $\text{iszero } t_1$ if t_1 is a Nat. Additionally, the construct *if* t_1 *then* t_2 *else* t_3 is defined to return a value of the type of t_3 if t_1 is a Bool.

If a term is typeable, it has at most one type and there is always just one derivation tree for it.

The chapter “Safety = Progress + Preservation” is about proving that a term cannot reach a “stuck state” where the evaluation rules don’t tell us, what to do next. This can be done using the progress and preservation theorems. Both Theorems together prove that a well-typed term cannot reach a stuck state.

Progress Theorem: For the case $t = \text{succ } t_1$ $t_1 : \text{Nat}$ there either is a $t'1$ or t_1 is a value. If t_1 is a nv t has to be too. If $t_1 \rightarrow t'1$ by E-SUCC $\text{succ } t_1 \rightarrow \text{succ } t'1$.

REFERENCES

- [1] Benjamin C. Pierce. 2002. *Types and Programming Languages* (1st ed.). The MIT Press.