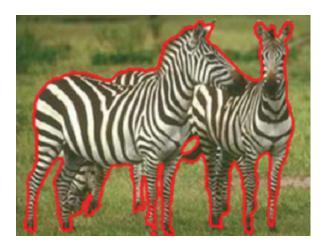
### Introduction to Spectral Graph Theory and Graph Clustering

Chengming Jiang

ECS 231 Spring 2016 University of California, Davis

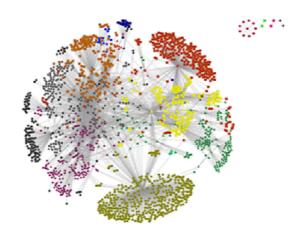
## **Motivation**

Image partitioning in computer vision



## **Motivation**

Community detection in network analysis

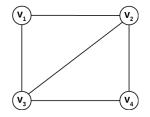


#### **Outline**

- I. Graph and graph Laplacian
  - Graph
  - Weighted graph
  - Graph Laplacian
- II. Graph clustering
  - Graph clustering
  - Normalized cut
  - Spectral clustering

An (undirected) **graph** is a pair G = (V, E), where

- $ightharpoonup \{v_i\}$  is a set of vertices;
- ightharpoonup E is a subset of  $V \times V$ .



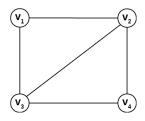
- ▶ An edge is a pair  $\{u, v\}$  with  $u \neq v$  (no self-loop);
- ▶ There is at most one edge from u to v (simple graph).

▶ For every vertex  $v \in V$ , the **degree** d(v) of v is the number of edges adjacent to v:

$$d(v) = |\{u \in V | \{u, v\} \in E\}|.$$

▶ Let  $d_i = d(v_i)$ , the **degree matrix** 

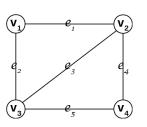
$$D = D(G) = \mathsf{diag}(d_1, \ldots, d_n).$$



$$D = \left(\begin{array}{cccc} 2 & 0 & 0 & 0 \\ 0 & 3 & 0 & 0 \\ 0 & 0 & 3 & 0 \\ 0 & 0 & 0 & 2 \end{array}\right).$$

▶ Given a graph G = (V, E), with |V| = n and |E| = m, the incidence matrix  $\tilde{D}(G)$  of G is an  $n \times m$  matrix with

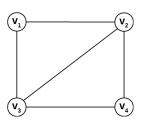
$$\tilde{d}_{ij} = \begin{cases} 1, & \text{if } \exists k \text{ s.t. } e_j = \{v_i, v_k\} \\ 0, & \text{otherwise} \end{cases}$$



$$\tilde{D}(G) = \begin{array}{ccccc} v_1 & e_2 & e_3 & e_4 & e_5 \\ v_1 & 1 & 0 & 0 & 0 \\ v_2 & 1 & 0 & 1 & 1 & 0 \\ v_3 & 0 & 1 & 1 & 0 & 1 \\ v_4 & 0 & 0 & 0 & 1 & 1 \end{array}$$

▶ Given a graph G = (V, E), with |V| = n and |E| = m, the adjacency matrix A(G) of G is a symmetric  $n \times n$  matrix with

$$a_{ij} = \begin{cases} 1, & \text{if } \{v_i, v_j\} \in E \\ 0, & \text{otherwise} \end{cases}.$$



$$A(G) = \left(\begin{array}{cccc} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{array}\right)$$

A weighted graph is a pair G = (V, W) where

- $V = \{v_i\}$  is a set of vertices and |V| = n;
- $W \in \mathbb{R}^{n \times n}$  is called *weight matrix* with

$$w_{ij} = \begin{cases} w_{ji} \ge 0 & \text{if } i \ne j \\ 0 & \text{if } i = j \end{cases}$$

The underlying graph of G is  $\widehat{G} = (V, E)$  with

$$E = \{\{v_i, v_j\} | w_{ij} > 0\}.$$

- ▶ If  $w_{ij} \in \{0,1\}$ , W = A, the adjacency matrix of  $\widehat{G}$ .
- ▶ Since  $w_{ii} = 0$ , there is no self-loops in  $\widehat{G}$ .

▶ For every vertex  $v_i \in V$ , the **degree**  $d(v_i)$  of  $v_i$  is the sum of the weights of the edges adjacent to  $v_i$ :

$$d(v_i) = \sum_{j=1}^n w_{ij}.$$

▶ Let  $d_i = d(v_i)$ , the **degree matrix** 

$$D = D(G) = \operatorname{diag}(d_1, \dots, d_n).$$

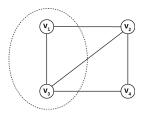
Remark:

Let d = diag(D) and denote  $\mathbf{1} = (1, \dots, 1)^T$ , then  $d = W\mathbf{1}$ .

▶ Given a subset of vertices  $A \subseteq V$ , we define the **volume**  $\operatorname{vol}(A)$  by

$$vol(A) = \sum_{v_i \in A} d(v_i) = \sum_{v_i \in A} \sum_{j=1}^{n} w_{ij}.$$

- ▶ If vol(A) = 0, all the vertices in A are isolated.
- Example:



If 
$$A=\{v_1,v_3\}$$
, then 
$$\operatorname{vol}(A)=d(v_1)+d(v_3)$$
 
$$=(w_{12}+w_{13})+ (w_{31}+w_{32}+w_{34})$$

▶ Given two subsets of vertices  $A, B \subseteq V$ , we define the **links** links(A, B) by

$$\operatorname{links}(A,B) = \sum_{v_i \in A, v_j \in B} w_{ij}.$$

- ► A and B are not necessarily distinct;
- ▶ Since W is symmetric, links(A, B) = links(B, A);
- $ightharpoonup \operatorname{vol}(A) = \operatorname{links}(A, V).$

▶ The quantity cut(A) is defined by

$$\operatorname{cut}(A) = \operatorname{links}(A, V - A).$$

▶ The quantity assoc(A) is defined by

$$\operatorname{assoc}(A) = \operatorname{links}(A, A).$$

- ightharpoonup cut(A) measures how many links escape from A;
- ▶ assoc(A) measures how many links stay within A;
- $ightharpoonup \operatorname{cut}(A) + \operatorname{assoc}(A) = \operatorname{vol}(A).$

# 1.3 Graph Laplacian

Given a weighted graph G=(V,W), the (graph) Laplacian L of G is defined by

$$L = D - W$$
.

where D is the degree matrix of G.

#### Remark

 $D = \mathsf{diag}(W \cdot \mathbf{1}).$ 

## I.3 Graph Laplacian

#### Properties of Laplacian

- 1.  $x^T L x = \frac{1}{2} \sum_{i,j=1}^n w_{ij} (x_i x_j)^2$  for  $\forall x \in \mathbb{R}^n$ .
- 2.  $L \geq 0$  if  $w_{ij} \geq 0$  for all i, j;
- 3.  $L \cdot 1 = 0$ ;
- 4. If the underlying graph of G is connected, then

$$0 = \lambda_1 < \lambda_2 \le \lambda_3 \le \ldots \le \lambda_n;$$

5. If the underlying graph of G is connected, then the dimension of the nullspace of L is 1.

# I.3 Graph Laplacian

#### **Proofs**

**Property 1.** Since L = D - W, we have

$$x^{T}Lx = x^{T}Dx - x^{T}Wx$$

$$= \sum_{i=1}^{n} d_{i}x_{i}^{2} - \sum_{i,j=1}^{n} w_{ij}x_{i}x_{j}$$

$$= \frac{1}{2} \left(\sum_{i}^{n} d_{i}x_{i}^{2} - 2\sum_{i,j=1}^{n} w_{ij}x_{i}x_{j} + \sum_{j=1}^{n} d_{j}x_{j}^{2}\right)$$

$$= \frac{1}{2} \left(\sum_{i,j=1}^{n} w_{ij}x_{i}^{2} - 2\sum_{i,j=1}^{n} w_{ij}x_{i}x_{j} + \sum_{i,j=1}^{n} w_{ij}x_{j}^{2}\right)$$

$$= \frac{1}{2} \sum_{i=1}^{n} w_{ij}(x_{i} - x_{j})^{2}.$$

# 1.3 Graph Laplacian

#### Property 2.

- ▶ Since  $L^T = D W^T = D W = L$ , L is symmetric.
- ▶ Since  $x^T L x = \frac{1}{2} \sum_{i,j=1}^n w_{ij} (x_i x_j)^2$  and  $w_{ij} \ge 0$  for all i, j, we have  $x^T L x \ge 0$ .

### Property 3.

$$L \cdot 1 = (D - W)1 = D1 - W1 = d - d = 0.$$

Properties 4 and 5: skip for now, see §2.2 of [Gallier'14].

### **Outline**

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  - Normalized cut
  - Spectral clustering

# **II.1 Graph clustering**

k-way partitioning: given a weighted graph G=(V,W), find a partition  $A_1,A_2,\ldots,A_k$  of V, such that

- $A_1 \cup A_2 \cup \ldots \cup A_k = V;$
- $A_1 \cap A_2 \cap \ldots \cap A_k = \varnothing;$
- ▶ for any i and j, the edges between  $(A_i, A_j)$  have low weight and the edges within  $A_i$  have high weight.

If k = 2, it is a two-way partitioning.

# **II.1 Graph clustering**

► Recall: (two-way) cut:

$$\mathsf{cut}(A) = \mathsf{links}(A, \bar{A}) = \sum_{v_i \in A, v_i \in \bar{A}} w_{i,i}$$

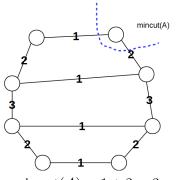
where 
$$\bar{A} = V - A$$
.

# **II.1 Graph clustering problems**

The mincut is defined by

$$\min \mathsf{cut}(A) = \min \sum_{v_i \in A, v_i \in \bar{A}} w_{ij}$$

In practice, the mincut easily yields unbalanced partitions.



$$\min \mathsf{cut}(A) = 1 + 2 = 3;$$

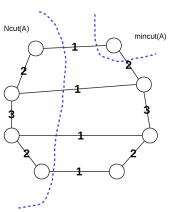
The **normalized cut**<sup>1</sup> is defined by

$$\operatorname{Ncut}(A) = \frac{\operatorname{cut}(A)}{\operatorname{vol}(A)} + \frac{\operatorname{cut}(\bar{A})}{\operatorname{vol}(\bar{A})}.$$

<sup>&</sup>lt;sup>1</sup>Jianbo Shi and Jitendra Malik, 2000

#### **Minimal Ncut:**

## $\min \mathsf{Ncut}(A)$ .



$$\min \mathsf{Ncut}(A) = \tfrac{4}{3+6+6+3} + \tfrac{4}{3+6+6+3} = \tfrac{4}{9}.$$

1. Let  $x = (x_1, \dots, x_n)$  be the *indicator vector*, such that

$$x_i = \left\{ \begin{array}{ll} 1 & \text{if } v_i \in A \\ -1 & \text{if } v_i \in \bar{A} \end{array} \right..$$

2. Then

$$\begin{array}{l} (\mathbf{1} + x)^T D(\mathbf{1} + x) = 4 \sum_{v_i \in A} d_i = 4 \cdot \text{vol}(A); \\ (\mathbf{1} + x)^T W(\mathbf{1} + x) = 4 \sum_{v_i \in A, v_j \in A} w_{ij} = 4 \cdot \operatorname{assoc}(A). \\ (\mathbf{1} + x)^T L(\mathbf{1} + x) = 4 \cdot (\operatorname{vol}(A) - \operatorname{assoc}(A)) = 4 \cdot \operatorname{cut}(A); \\ \text{and} \\ (\mathbf{1} - x)^T D(\mathbf{1} - x) = 4 \sum_{v_i \in \bar{A}} d_i = 4 \cdot \operatorname{vol}(\bar{A}); \\ (\mathbf{1} - x)^T W(\mathbf{1} - x) = 4 \sum_{v_i \in \bar{A}, v_j \in \bar{A}} w_{ij} = 4 \cdot \operatorname{assoc}(\bar{A}). \\ (\mathbf{1} - x)^T L(\mathbf{1} - x) = 4 \cdot (\operatorname{vol}(\bar{A}) - \operatorname{assoc}(\bar{A})) = 4 \cdot \operatorname{cut}(\bar{A}). \\ \text{(Please verify it after class.)} \end{array}$$

3. Ncut(A) can now be written as

$$\begin{aligned} \mathsf{Ncut}(A) &= \frac{1}{4} \left( \frac{(\mathbf{1} + x)^T L (\mathbf{1} + x)}{k \mathbf{1}^T D \mathbf{1}} + \frac{(\mathbf{1} - x)^T L (\mathbf{1} - x)}{(1 - k) \mathbf{1}^T D \mathbf{1}} \right) \\ &= \frac{1}{4} \cdot \frac{((\mathbf{1} + x) - b(\mathbf{1} - x))^T L ((\mathbf{1} + x) - b(\mathbf{1} - x))}{b \mathbf{1}^T D \mathbf{1}}. \end{aligned}$$

where k = vol(A)/vol(V), b = k/(1-k) and  $\text{vol}(V) = \mathbf{1}^T D \mathbf{1}$ .

4. Let y = (1 + x) - b(1 - x), we have

$$Ncut(A) = \frac{1}{4} \cdot \frac{y^T L y}{b \mathbf{1}^T D \mathbf{1}}$$

where

$$y_i = \left\{ \begin{array}{ll} 2 & \text{if } v_i \in A \\ -2b & \text{if } v_i \in \bar{A} \end{array} \right..$$

5. Since 
$$b = k/(1-k) = \operatorname{vol}(A)/\operatorname{vol}(\bar{A})$$
, we have

$$\begin{split} \frac{1}{4}y^TDy &= \sum_{v_i \in A} d_i + b^2 \sum_{v_i \in \bar{A}} d_i = \operatorname{vol}(A) + b^2 \operatorname{vol}(\bar{A}) \\ &= b(\operatorname{vol}(\bar{A}) + \operatorname{vol}(A)) = b \cdot \mathbf{1}^T D \mathbf{1}. \end{split}$$

In addition,

$$y^T D \mathbf{1} = y^T \mathbf{d} = 2 \cdot \sum_{v_i \in A} d_i - 2b \cdot \sum_{v_i \in \bar{A}} d_i$$
$$= 2 \cdot \text{vol}(A) - 2b \cdot \text{vol}(\bar{A}) = 0$$

6. The minimal normalized cut is to solve the following binary optimization:

$$y = \arg\min_{y} \quad \frac{y^{T}Ly}{y^{T}Dy}$$

$$s.t. \quad y(i) \in \{2, -2b\}$$

$$y^{T}D\mathbf{1} = 0$$

$$(1)$$

7. Relaxation

$$y = \arg\min_{y} \quad \frac{y^{T}Ly}{y^{T}Dy}$$

$$s.t. \quad \mathbf{y} \in \mathbb{R}^{n}$$

$$y^{T}D\mathbf{1} = 0$$
(2)

#### Variational principle

- ▶ Let  $A, B \in \mathbb{R}^{n \times n}$ ,  $A^T = A$ ,  $B^T = B > 0$  and  $\lambda_1 \le \lambda_2 \le \dots \lambda_n$  be the eigenvalues of  $Au = \lambda Bu$  with corresponding eigenvectors  $u_1, u_2, \dots, u_n$ ,
- ▶ then

$$\min_{x} \frac{x^{T} A x}{x^{T} B x} = \lambda_{1} , \quad \arg \min_{x} \frac{x^{T} A x}{x^{T} B x} = u_{1}$$

and

$$\min_{x^TBu_1=0}\frac{x^TAx}{x^TBx}=\lambda_2 \text{ , } \quad \arg\min_{x^TBu_1=0}\frac{x^TAx}{x^TBx}=u_2.$$

More general form exists.

- For the matrix pair (L, D), it is known that  $(\lambda_1, y_1) = (0, 1)$ .
- ▶ By the variational principle, the relaxed minimal Ncut (2) is equivalent to finding the second smallest eigenpair  $(\lambda_2, y_2)$  of

$$Ly = \lambda Dy \tag{3}$$

- ▶ L is extremely sparse and D is diagonal;
- ▶ Precision requirement for eigenvectors is low, say  $\mathcal{O}(10^{-3})$ .

Image segmentation: original graph



Image segmentation: heatmap of eigenvectors

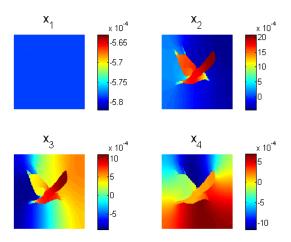


Image segmentation: result of min Ncut



#### Ncut remaining issues

- ➤ Once the indicator vector is computed, how to search the splitting point that the resulting partition has the minimal Ncut(A) value?
- ► How to use the extreme eigenvectors to do the *k*-way partitioning?

The above two problems are addressed in spectral clustering algorithm.

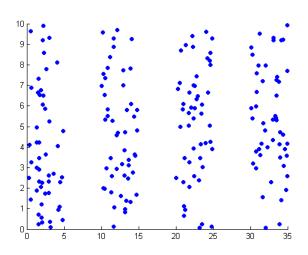
Spectral clustering algorithm [Ng et al, 2002]

Given a weighted graph G = (V, W),

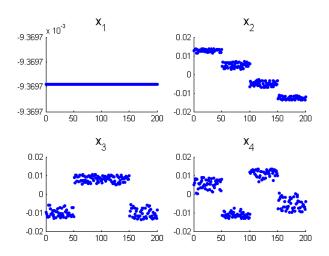
- 1. compute the normalized Laplacian  $L_{\rm n}=D^{-\frac{1}{2}}(D-W)D^{-\frac{1}{2}};$
- 2. find k eigenvectors  $X = [x_1, \ldots, x_k]$  corresponding to the k smallest eigenvalues of  $L_n$ ;
- 3. form  $Y \in \mathbb{R}^{n \times k}$  by normalizing each row of X such that  $Y(i,:) = X(i,:)/\|X(i,:)\|$ ;
- 4. treat each Y(i,:) as a point, cluster them into k clusters via K-means with label  $c_i = \{1, ..., k\}$ .

The label  $c_i$  indicates the cluster that  $v_i$  belongs to.

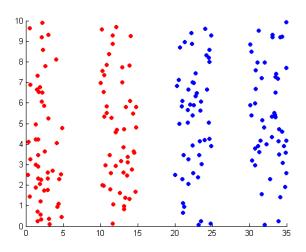
## Synthetic example: original data



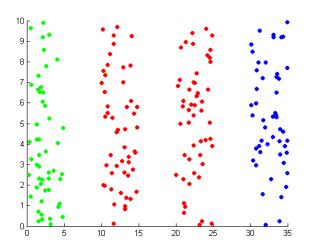
### Synthetic example: computed eigenvectors



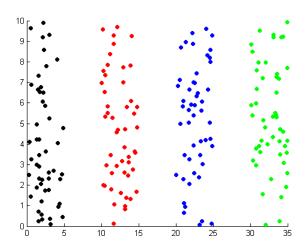
Synthetic example: 2-way clustering



### Synthetic example: 3-way clustering



Synthetic example: 4-way clustering



#### References

- 1. Jean Gallier, Notes on elementary spectral graph theory applications to graph clustering using normalized cuts, 2013.
- 2. Jianbo Shi and Jitendra Malik, *Normalized cuts and image segmentation*, 2000.
- 3. Andrew Y Ng, Michael I. Jordan and Yair Weiss, *On spectral clustering: Analysis and an algorithm*, 2001