

①  $\phi = \frac{1+\sqrt{5}}{2}$   $j=2, m=2$  where  $\text{cost}(A) = \phi \cdot \text{opt}(G)$

let 2 machines have weights where  $\frac{w_i}{s_i} = 1$  :

$w_1 = \phi, s_1 = \phi$  ;  $w_2 = 1, s_2 = 1$  where  $\text{PoA} = \frac{\text{cost}(A)}{\text{opt}(G)} = \phi$

Proof

NE exists at  $A(1)=2, A(2)=1$

where:  $\frac{w_1}{s_2} = \frac{\phi}{1} \leq \frac{w_1+w_2}{s_1} = \frac{1+\phi}{\phi} = \phi$

and  $\frac{w_2}{s_1} = \frac{1}{\phi} \leq \frac{w_1+w_2}{s_2} = \frac{1+\phi}{1}$

where  $\text{cost}(A) = \max\left\{\frac{w_2}{s_1}, \frac{w_1}{s_2}\right\} = \phi$

this implies  $\text{opt}(G) = 1$

$$\begin{aligned} \text{OPT}(G) &= \min\left\{\frac{\sum w}{s_1}, \frac{\sum w}{s_2}, \max\left\{\frac{w_1}{s_2}, \frac{w_2}{s_1}\right\}, \max\left\{\frac{w_1}{s_1}, \frac{w_2}{s_2}\right\}\right\} \\ &= \max\left\{\frac{w_1}{s_1}, \frac{w_2}{s_2}\right\} = \max\{1, 1\} \\ &= 1 \end{aligned}$$

which holds  $\text{PoA} = \frac{\text{cost}(A)}{\text{opt}(G)} = \frac{\phi}{1} = \phi$

$$M = M_{\text{norm}}$$

$$(2) \quad u_i = a_j x_{ij} \quad \text{for } a_j \in \mathbb{Z}^+ \text{ \& } x_{ij} \in \mathbb{R}^+$$

Claim:

$$p_{ij} = \frac{\sum_{i=1}^n e_i}{\sum_{j=1}^m a_j} = \frac{\sum e_i}{\sum v_i} = \frac{\text{endowment}}{\text{availability}}$$

Since  $u_i$  is the same for all traders, then:

$$x_{ij} = a_j \cdot \frac{e_i}{\sum e_i} \quad \text{where trader } i \text{ gets a fraction of the total endowment.}$$

→ assuming traders spend all their money and sell all goods;

$$\sum_j x_{ij} \cdot p_{ij} = \sum_j \frac{a_j}{\sum_j a_j} \cdot e_i = e_i \cdot \frac{1}{\sum_j a_j} \cdot \sum_j a_j = e_i$$

$$\text{and} \quad \sum_i x_{ij} = \sum_i a_j \cdot \frac{e_i}{\sum_i e_i} = a_j \cdot \frac{\sum_i e_i}{\sum_i e_i} = a_j$$

which gives us a market equilibrium



③ <sup>(sellers)</sup>  $N$ -bidders with cost bid:  $c_i$ . The asking bid. is  $\hat{c}_i$   
 let  $p_i$  be the payment for bidder  $i$  where each bidder  
 sells only one item at a min bid. Assume a subset  $S$  of lowest bidders  
 Our modified vcs mechanism is:

$$S^* = \underset{S}{\operatorname{argmin}} \left( \sum_{i=1}^k \hat{c}_i \right)$$

where  $S$  is the set of the lowest bidders and  $|S|=k$ .  
 where there are  $k$  identical items.

The payment for bidder  $i$  given  $k-1$  bids in  $S^*$  is:

$$p_i = \min_{\{S', i \notin S'\}} \left( \sum_{j \in S'} \hat{c}_j \right) - \left[ \left( \sum_{j \in S^*} \hat{c}_j \right) - \hat{c}_i \right]$$

This shows that  $p_i$  is equal to the  $c_{k+1}$  where  
 bidder  $i' \in S^*$  such that

$$p_i = \begin{cases} c_{k+1}, & \text{if } i \in S^* \\ 0, & \text{otherwise} \end{cases} \text{ where } |S|=k, i \in N$$