

CS539-Spring 2018, Final-Exam  
 2nd May, 2018.  
 Closed Books. No Calculators  
 2:00pm - 4:00pm

Prove All Results. Total points: 75+75 (Take-Home).  
 Sub-questions that must be solved at home are distinctly marked.

1. (a) Find a correlated equilibrium (that is not a mixed NE if possible) in the following hot-potato routing game:

Table 1: Routing Game

Payoff	HotPotato	Planned
HotPotato	-5,-5	-2,-6
Planned	-6,-2	-1,-1

- (b) Also compute mixed Nash solutions and compare.

(15+15 pts)

2. (a) Suppose in non-atomic routing in networks, the delay function on edges is of the form  $af_e^2 + b_e$  where  $f_e$  is the flow on edge  $e$ . Show an example network where the price of anarchy is exactly

$$\frac{1}{1 - \frac{2}{3\sqrt{3}}}$$

- (b) **Take-home:** Show that  $PoA \leq \frac{1}{1 - \frac{2}{3\sqrt{3}}}$ .

(10+20 pts)

*Hint: Use Pigou's Example*

3. (a) Recall the local connection game defined by a set of vertices and graph  $G$  where edges are to be added to build connections, suppose the cost to user  $u$  (i.e. user at node  $u$ ) is defined to be  $C_u = \alpha n_u + \beta \sum_v dist(u, v)$  where  $n_u$  is the number of edges incident to  $u$ , that  $u$  purchases, and  $dist(u, v)$  is the shortest distance from  $u$  to  $v$ .  $\alpha, \beta \geq 0$ .

Can a path that connects all nodes be a Nash equilibrium solution. If so, then derive a condition on  $\alpha$  for a path to occur as a Nash equilibrium, otherwise prove that a path can never be a Nash equilibrium.

- (b) **Take-Home:** Suppose the cost function is defined to be  $\alpha n_u + \sum_v \sqrt{\text{dist}(u,v)}$ . Characterize NE solutions as a function of  $\alpha$ .

(15+15 pts)

4. (a) In a cloud computing environment,  $k$  machines are required to complete the task of finding a solution to a physics problem. Machine owners bid to provide machines for the job where the  $i$ th machine owner has valuation  $v_i$  for the machine. Show an incentive compatible method to determine the winners.
- (b) **Take-Home:** Suppose you need to consider only the cost of connecting the machines. The machines are represent by a vertex set of a graph and edges between vertices represent communication between machines. Each edge is controlled by a player who declares a cost for using the edge. The goal is to determine a connected subgraph to communicate amongst all nodes (hint: Minimum Spanning Tree is one such subgraph). Determine an incentive compatible mechanism for this problem.

(10+20 pts)

5. (a) In the Fisher linear market equilibrium model, let  $M$  be a market with  $n$  traders and  $m$  goods, with each trader  $i$  having endowment  $m_i$  and  $u_{ij} = v_{ij}x_{ij} = r_j s_i x_{ij}$  for good  $j$  where  $r_j, s_i \in \mathbb{Z}^+$ , and  $x_{ij} \in \mathbb{R}^+$  i.e. Show that the prices of 2 goods  $j$  and  $j'$  acquired by a trader is in the ratio  $r_j/r_{j'}$ , i.e.  $p_j/p_{j'} = r_j/r_{j'}$ . The available quantity of good  $j$  is  $a_j$ .
- (b) **Take-home:** Use this to determine the price of the goods and then determine an allocation at Market equilibrium.

(10+20 pts)