a86 Interpreter Semantics

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This document describes the semantics used in the implementation of the a86 interpreter.

1 Syntax

We provide a definition of the abstract syntax of the a86 language.

```
r64
                                                  rdx
               rax
                           rbx
                                      rcx
                rbp
                           rsp
                                      rsi
                                                  rdi
                                                  r11
                r8
                           r9
                                      r10
                           r13
                                      r14
                                                  r15
  \mathbf{r}
               r64
                           eax
  \mathbf{f}
         ::=
               CF
                           ZF
                                      SF
                                                OF
  1
                labels
  \mathbf{z}
                integers
 \mathbf{bc}
                integers\ 0\ through\ 63\ inclusive
 \mathbf{off}
               Offset r z
prog
               Program (inst ...)
```

```
inst
          ::=
                   Label
                                loc
                   Ret
                   Call
                                 (\mathbf{loc} \mid \mathbf{reg})
                   Mov
                                 (reg \mid off)
                                                     (reg \mid off \mid int)
                   Add
                                reg
                                                              off
                                                                      int)
                                                     (reg \mid
                   Sub
                                reg
                                                     (reg \mid off)
                                                                      int)
                   Cmp
                                 (reg \mid off)
                                                     (reg \mid off \mid int)
                   Jmp
                                 (loc \mid reg)
                                 (loc \mid reg)
                   Je
                                 (\mathbf{loc} \mid \mathbf{reg})
                   Jne
                   Jl
                                 (\mathbf{loc} \mid \mathbf{reg})
                   Jg
                                 (\mathbf{loc} \mid \mathbf{reg})
                                 (reg \mid off)
                                                     (reg \mid off \mid int)
                   And
                                                     (reg \mid off \mid int)
                   0r
                                 (reg \mid off)
                                                     (reg \mid off \mid int)
                   Xor
                                 (reg \mid off)
                   Sal
                                reg
                                                     \mathbf{bc}
                                                     \mathbf{bc}
                   Sar
                                reg
                   Push
                                (int \mid reg)
                   Pop
                                reg
                   Lea
                                (reg \mid off)
                                                     loc
```

2 Semantics (Informal)

Before formalizing our semantics, we provide an informal specification. These definitions are based on existing documentation and code. A formal semantics is presented in the next section.

2.1 Program

First, we define a program. A program is a list of instructions, where:

- The list is not empty.
- The first instruction is a Label, which will be used as the entry point of the program (i.e., it is where execution begins).
- Two Label instructions in the same program cannot use the same label name.

2.2 Registers

There are sixteen 64-bit registers, each corresponding to one of the names in **r64**. The special register reference **eax** refers to the lower 32 bits of the **rax** register.

2.3 Flags

There are four single-bit registers, called *flags*, each corresponding to one of the names in **f**. The flags are used for arithmetic and comparison operations:

• OF — Overflow Flag

Set when...

- \dots adding two numbers with the same sign bit and the result has a different sign bit.
- $\dots subtracting$ a negative number from a positive number and the result is negative.
- \dots subtracting a positive number from a negative number and the result is positive.
- SF Sign Flag

Set to the value of the sign bit of the result.

- ZF Zero Flag
 Set if the computed result is exactly 0.
- CF Carry Flag
 Set if the (unsigned) arithmetic operation required an extra bit.

We say a flag "is set" if the value 1 is stored in it, or the flag "is clear" or "is unset" if the value 0 is stored in it instead. We may also use these terms as verbs, i.e., "to set" a flag means to store 1 in it and so on. Note, however, that "set" is overloaded as a verb, since it can also be used to indicate storing a specific value (e.g., "the flag is set to the value of <some computation>").

2.4 Memory

All a86 programs also run with some limited amount of register-external memory called the *stack*. The stack starts at the highest available address space and "grows downwards", which means that adding something to the stack *decrements* the pointer to the current position in the stack.

2.5 Instructions

There are 20 supported instructions in a86, which work as follows:

- Label 1 creates a new label named 1 that points to the next instruction.
- Ret pops an address from the stack and jumps to it.
- Call dst pushes the return address onto the stack, then jumps to the address indicated by the label or register in dst.

• Mov dst src moves the contents of/value at src into dst.

NOTE: Either dst or src may be an offset, but not both.

- Add dst src adds src to dst and writes the result to dst.
- Sub dst src subtracts src from dst and writes the result to dst.
- Cmp a1 a2 compares a1 to a2 by subtracting the former from the latter and sets the flags according to the result:

CF is set if an extra bit was needed to complete the computation.

```
ZF is set if a2 - a1 = 0.
```

SF is set if a2 - a1 < 0.

OF is set if either...

- ... a1 is negative, a2 is positive, and a2 a1 < 0.
- ...a1 is positive, a2 is negative, and a2 a1 > 0.
- Jmp dst jumps to the address at dst.
- Je dst jumps to the address at dst if ZF is set.
- Jne dst jumps to the address at dst if ZF is not set.
- Jl dst jumps to the address at dst if SF and OF have different values.
- Jg dst jumps to the address at dst if SF and OF are set to the same value and ZF is unset.
- And dst src computes the bitwise AND (&) of the operands and stores the result in dst.
- Or dst src computes the bitwise OR (I) of the operands and stores the result in dst.
- Xor dst src computes the bitwise XOR (^) of the operands and stores the result in dst.
- Sal dst i arithmetically shifts the bits in dst to the left by i bits and stores the result in dst. The new bits from the right are 0s, and the CF flag is updated to the value of the most-significant bit during each shift.

NOTE: If i is 1, the OF flag is set to 0 if the most-significant bit of the result is the same as the CF flag. Otherwise, it is set to 1.

• Sar dst src arithmetically shifts the bits in dst to the right by i bits and stores the result in dst. The new bits from the left are duplicated from the original most-significant bit, and the CF flag is updated to the value of the least-significant bit during each shift.

NOTE: The OF flag is always cleared.

- Push src decrements the stack pointer and stores the src operands on the top of the stack.
- Pop dst loads the value from the top of the stack into the dst operand and increments the stack pointer.
- Lea dst 1 loads the address of the label 1 and stores it in dst.