

Outlining a Demand Driven Execution Model for CnC

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Motivation

- Remove unneeded computation

Outline

- Step inverse
- Demand driven model
- Future work

Preliminary: Step Inverse

- A step definition represents a map from step tags to output collection tags
- Remark: this map is **one-to-one** if we do not output multiple items into a single collection
- Then we can find the inverse, the map from an item collection tag to the step tag that outputs it.

Abridged example

```
[ int data: i ];  
  
// Init: Set data[1] and prescribe process[1].  
  
( process: x )  
    <- [ data : x ]  
    -> [ data : x + 1 ],  
        ( process: x + 1 ) $when(x + 1 < 5);  
  
// Final: get item data[5]
```

Inverse of 'process'

```
( process: x )  
  <- [ data : x ]  
  -> [ data : x + 1 ],  
      ( process: x + 1 )  
      $when(x + 1 < 5);
```

```
{  
  'data': [{x: t1 - 1}],  
  'process': [{x:  
    Piecewise(  
      (t1 - 1, t1 < 5),  
      (nan, True))  
    }]  
}
```

Step inverse uses

- Value of item is wrong?
 - Step inverse to blame the step that put that item
- Deadlock?
 - Use event graph to highlight deadlocked items
 - Run step inverse to find the responsible step

Deadlock blame example

Init: Put data[0], prescribe process[1].

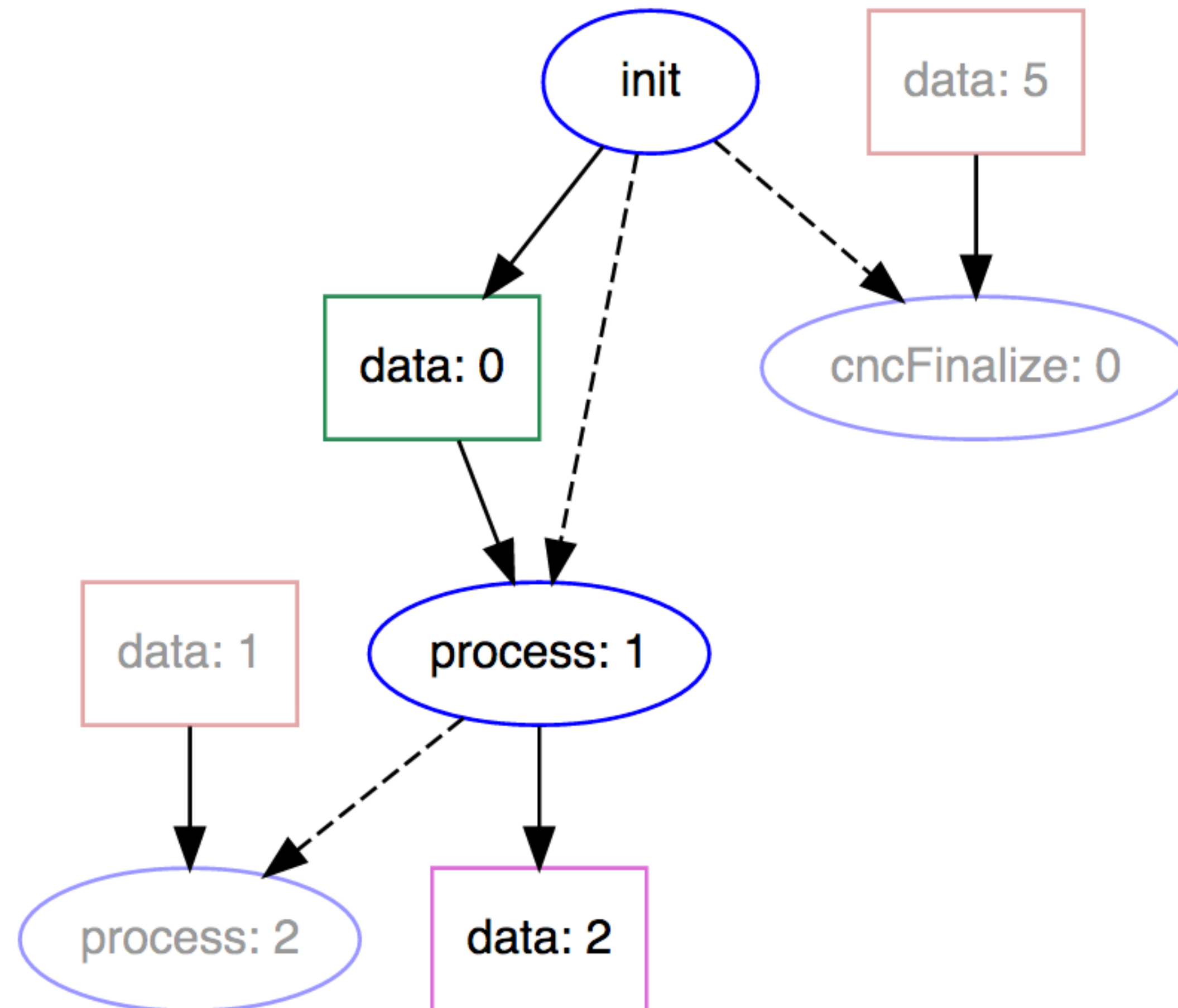
(process: x)

<- [data : x - 1]

-> [data : x + 1],

(process: x + 1) \$when(x + 1 < 5);

Deadlocked graph



Auto-blame output

Performing automatic blame on potentially deadlocked
items from event log: ['data@1']

{

 'data@1': {

 'process': {

 'x': 0

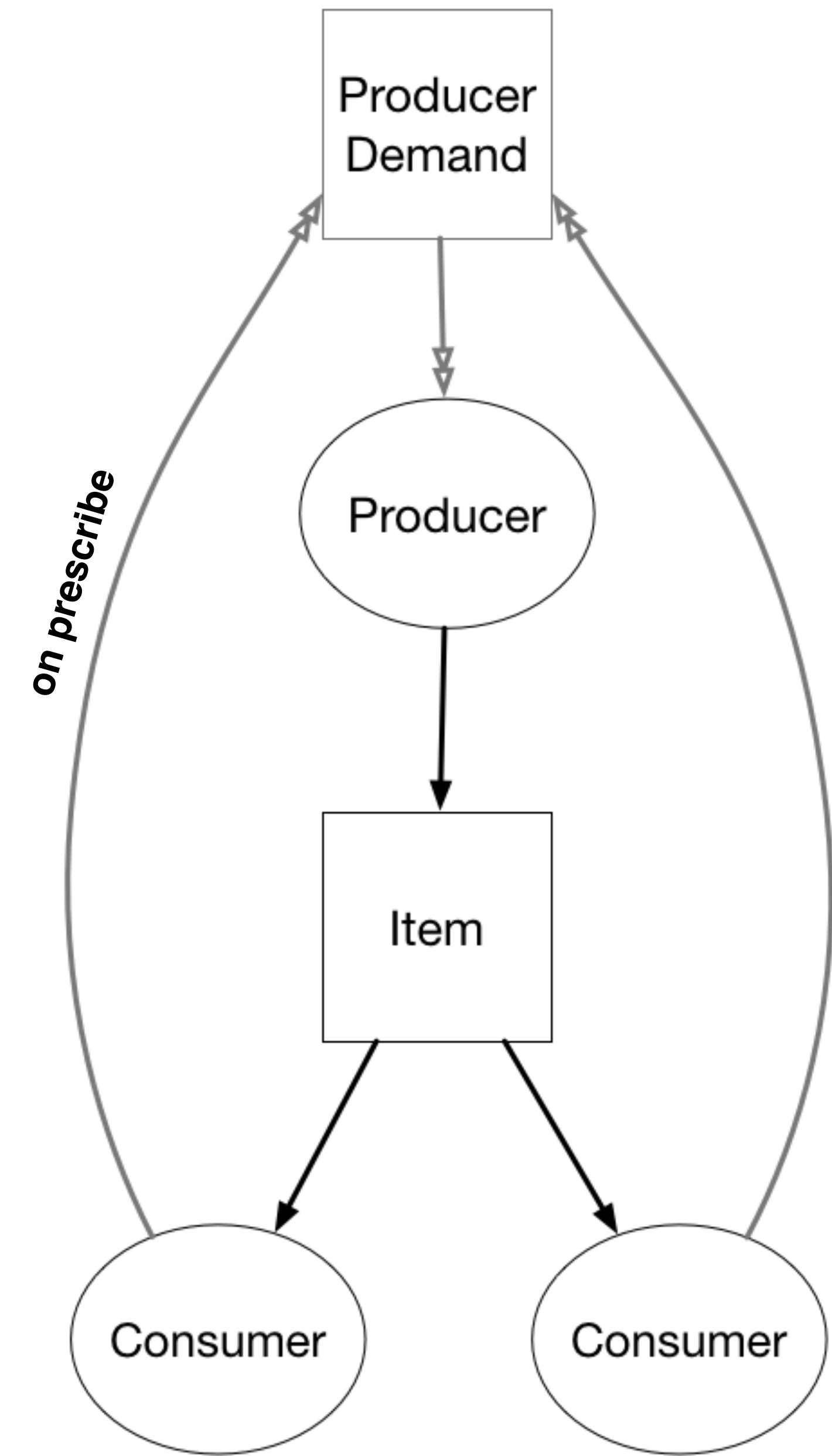
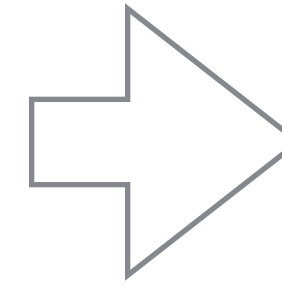
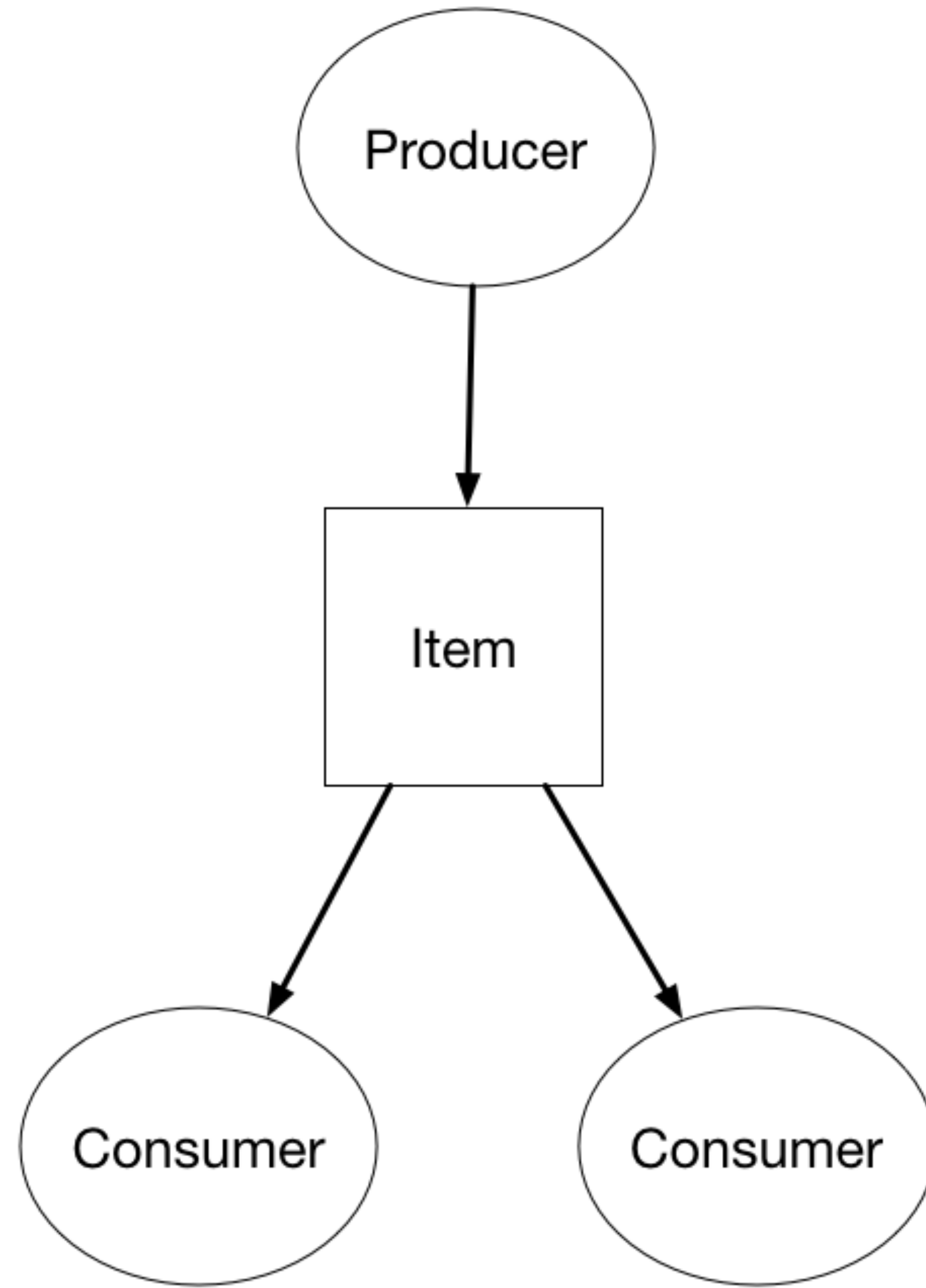
 }}}

Demand driven rules

- Definition. Only run step when we need its output.
- Condition. Collection producer is unambiguous.

Implementation

- Shadow collections to track demand



Example

```
[ int X: i ];
```

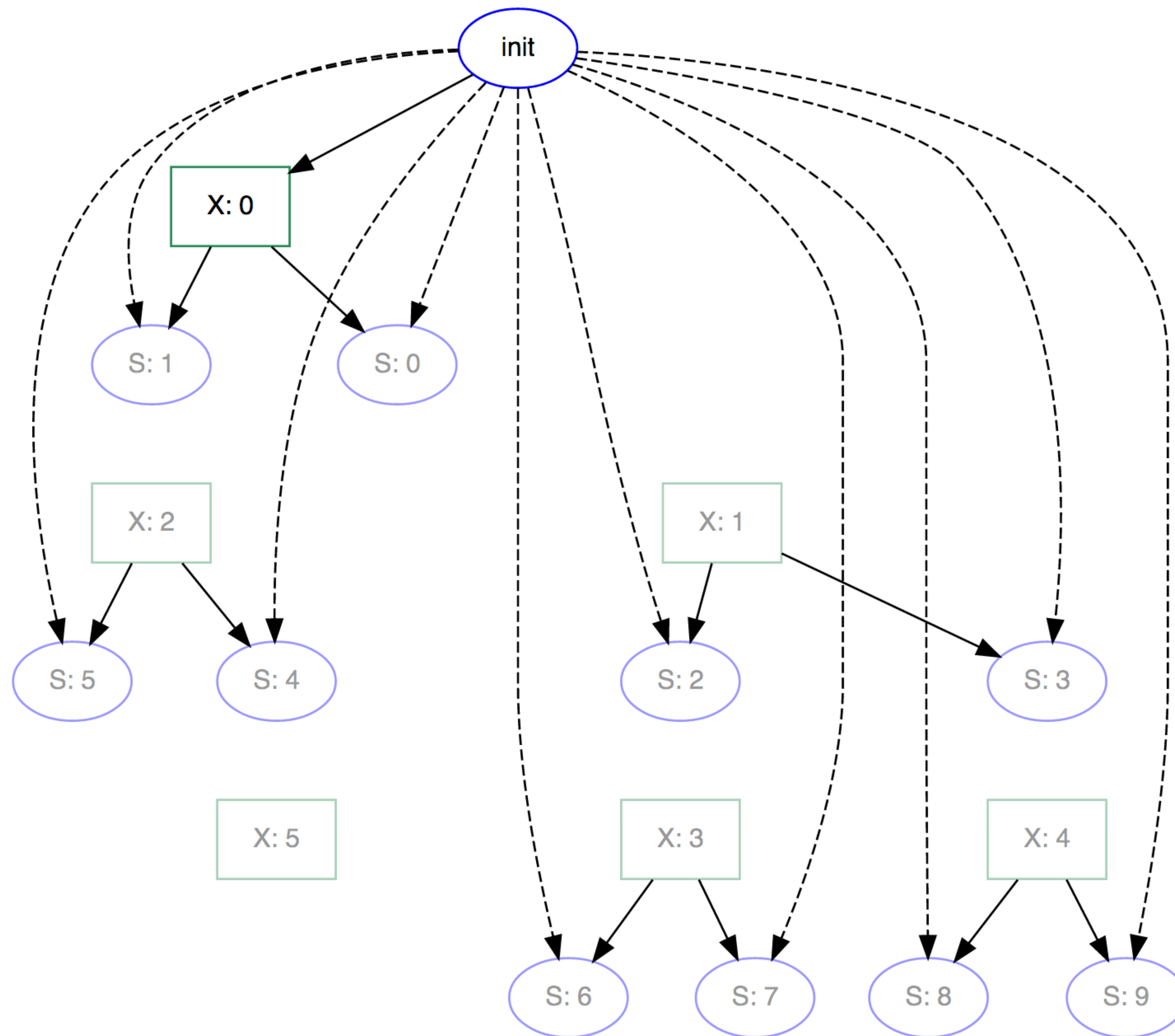
```
( $initialize: () ) -> ( S: $range(0, 10) ), [ X: 0 ];
```

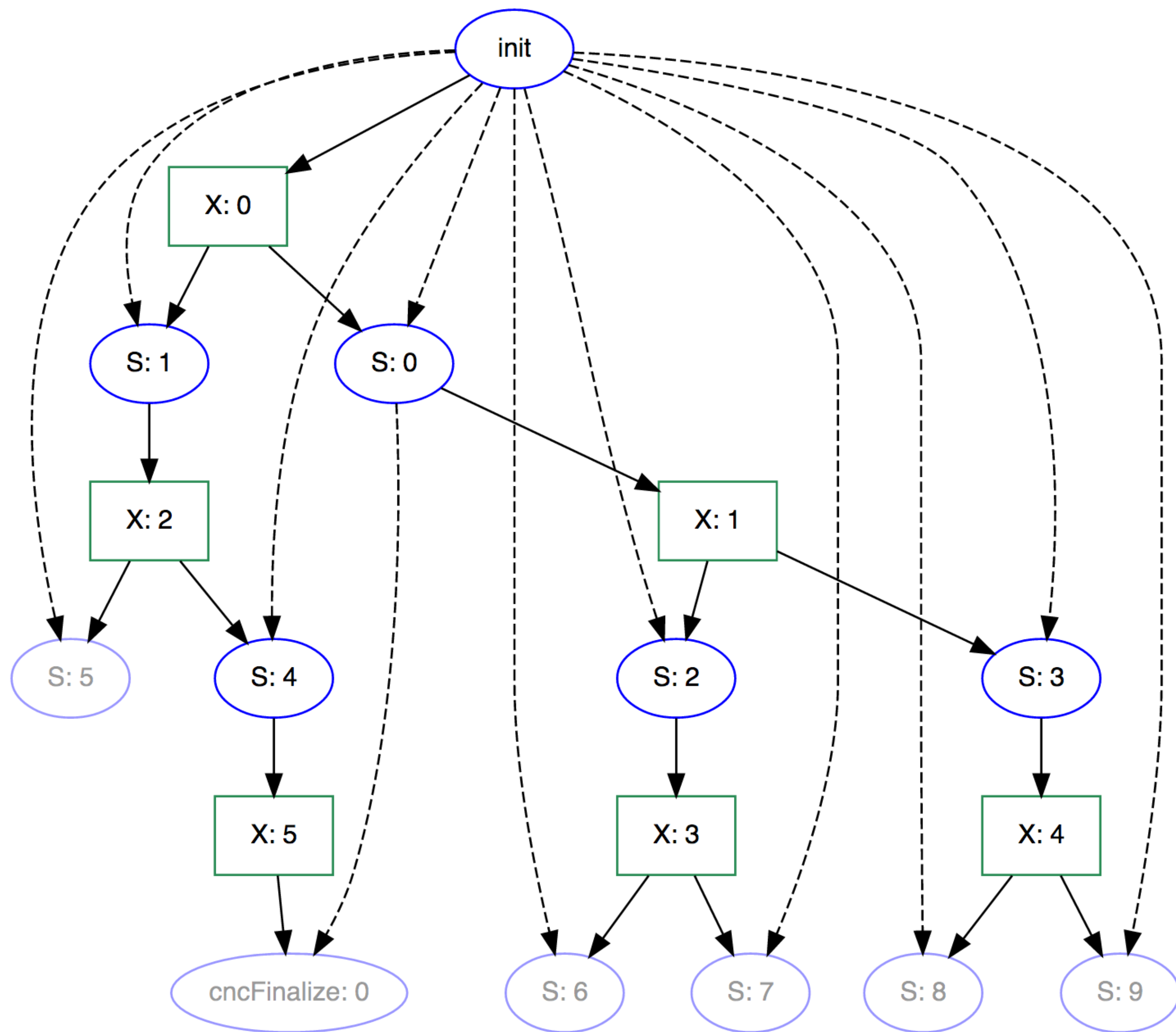
```
( S: i )
```

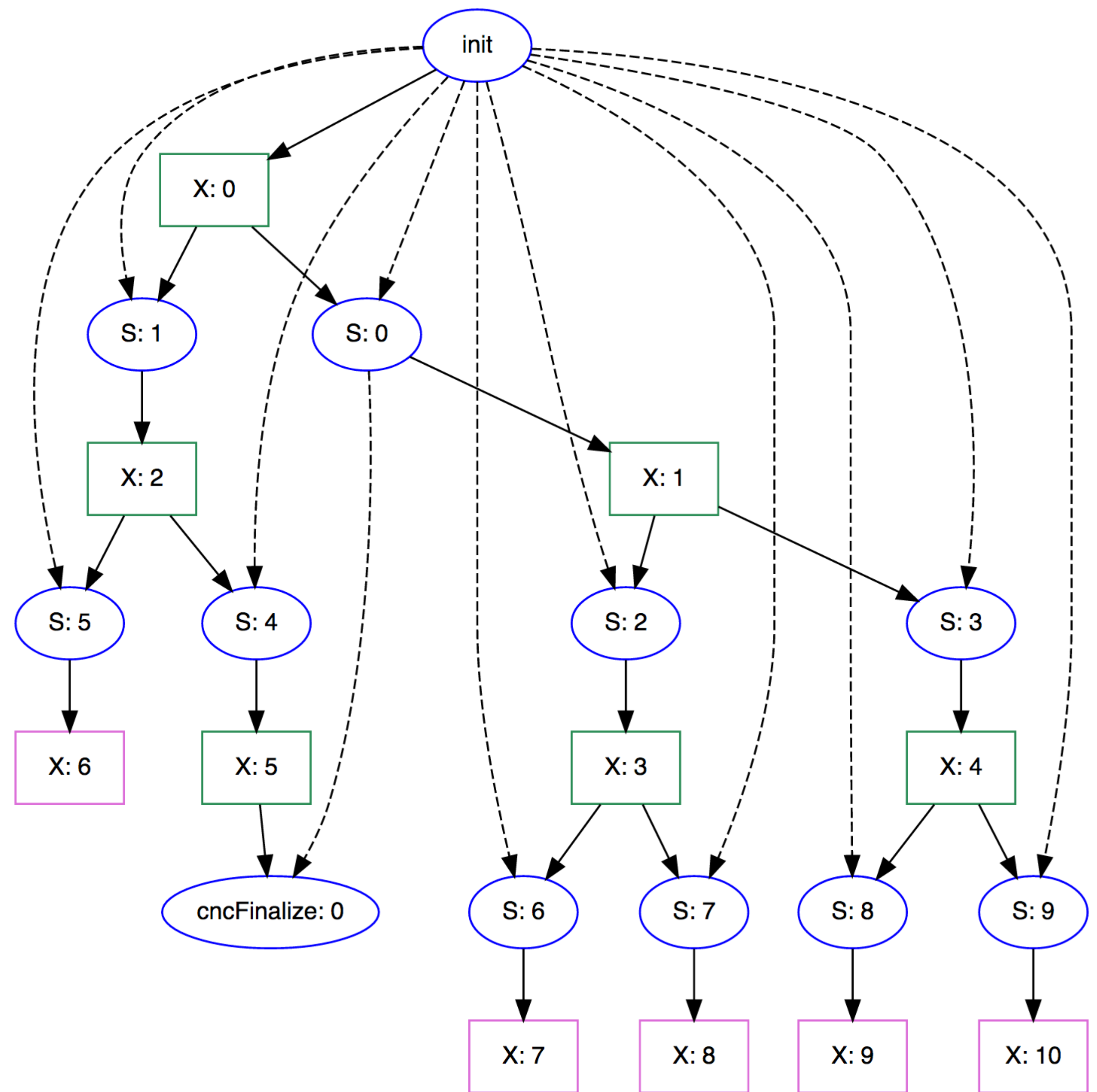
```
  <- [ X: i / 2 ]
```

```
  -> [ X: i + 1 ];
```

```
( $finalize: () ) <- [ X: 5 ];
```







Example

```
[ int X: i ];
```

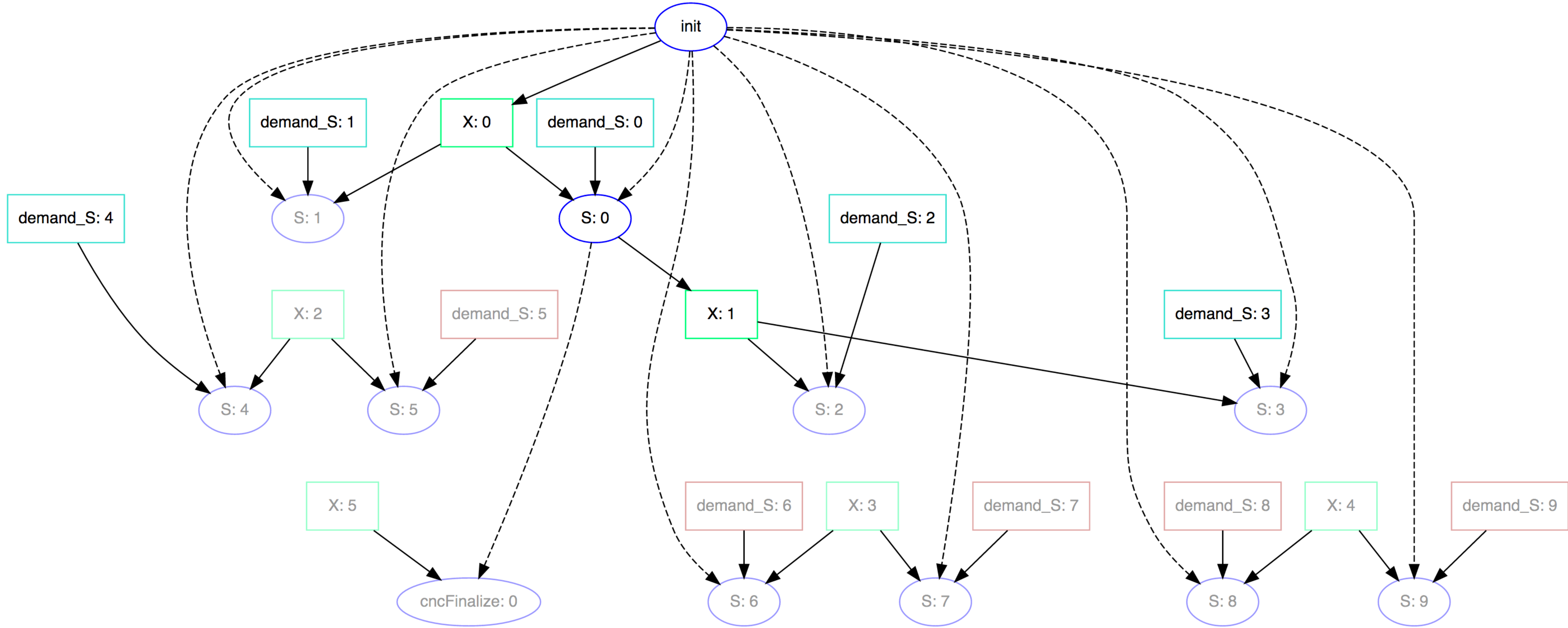
```
( $initialize: () ) -> ( S: $range(0, 10) ), [ X: 0 ];
```

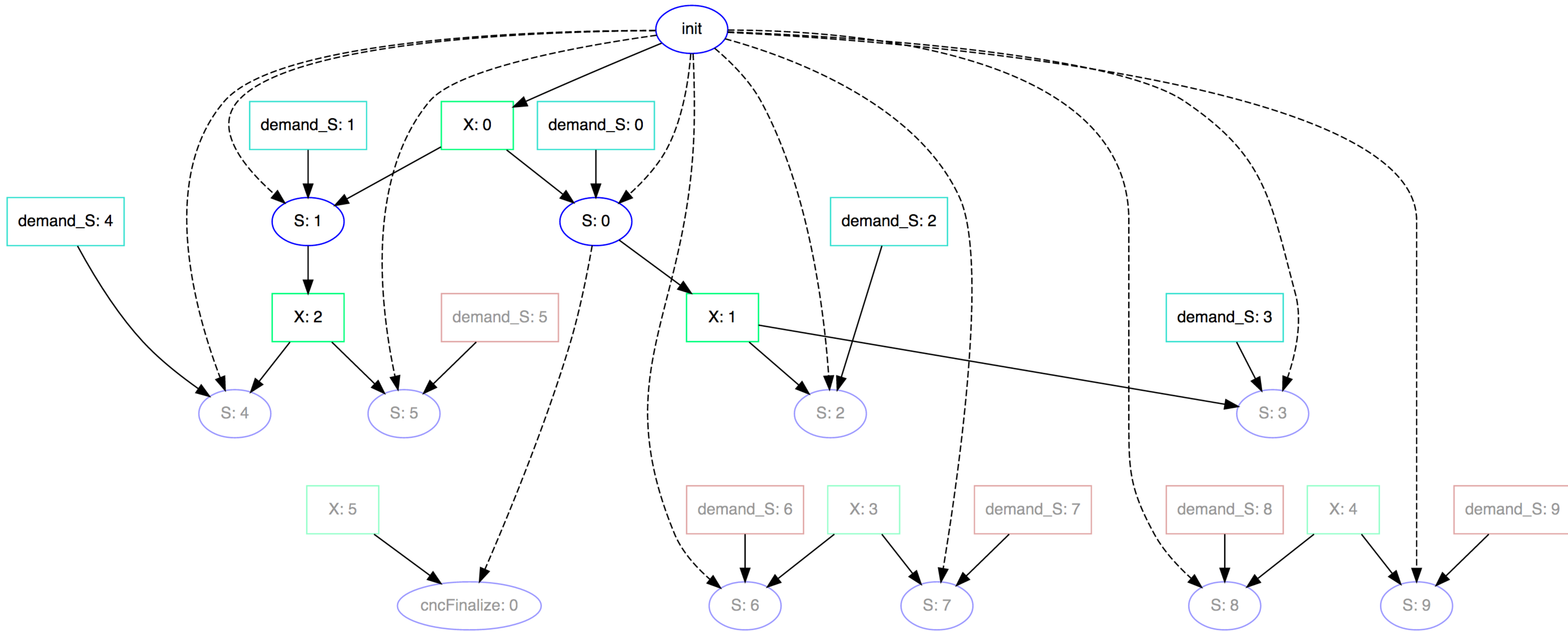
```
( S: i )
```

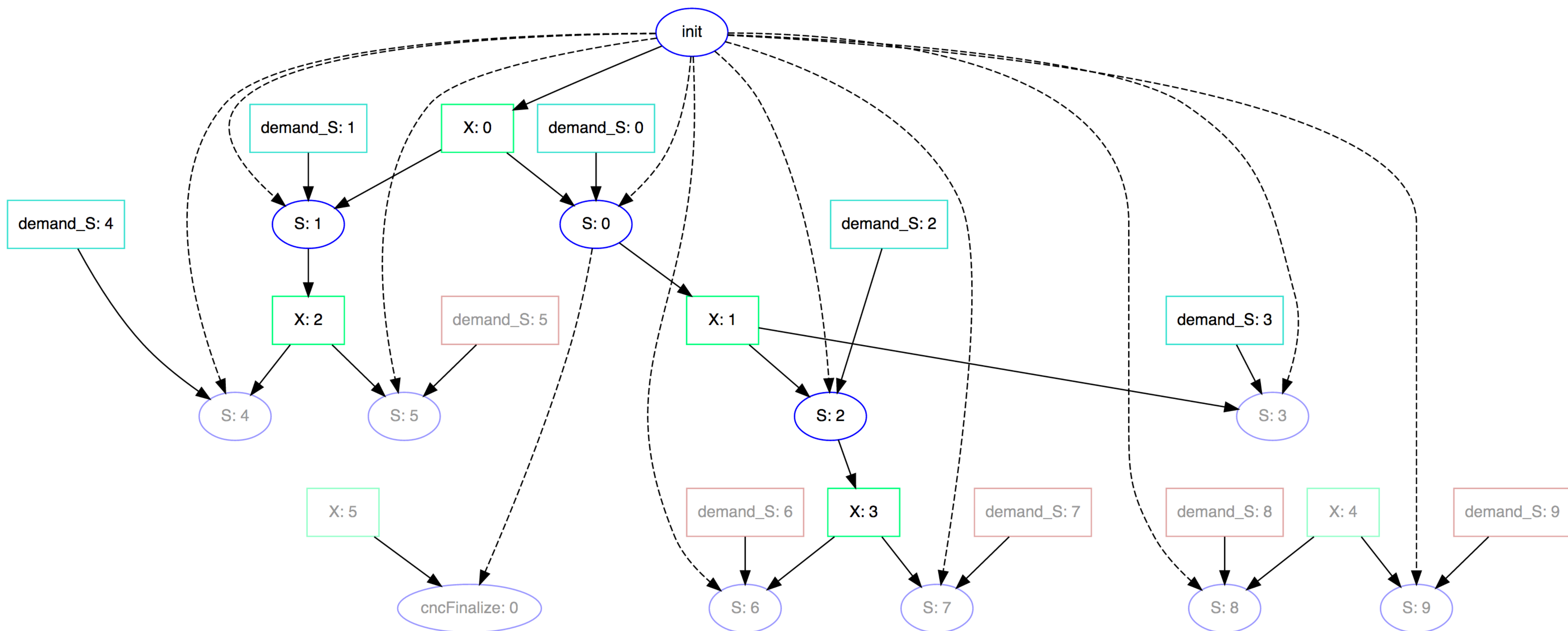
```
  <- [ X: i / 2 ]
```

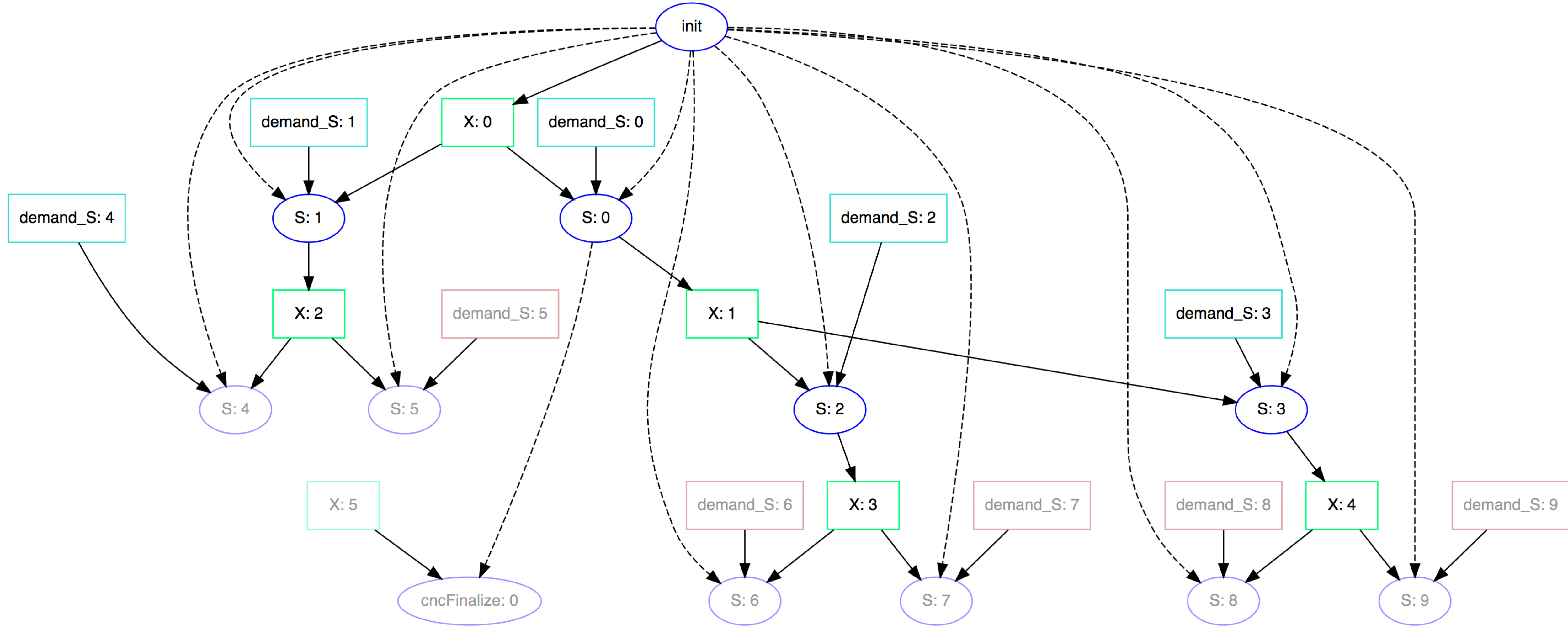
```
  -> [ X: i + 1 ];
```

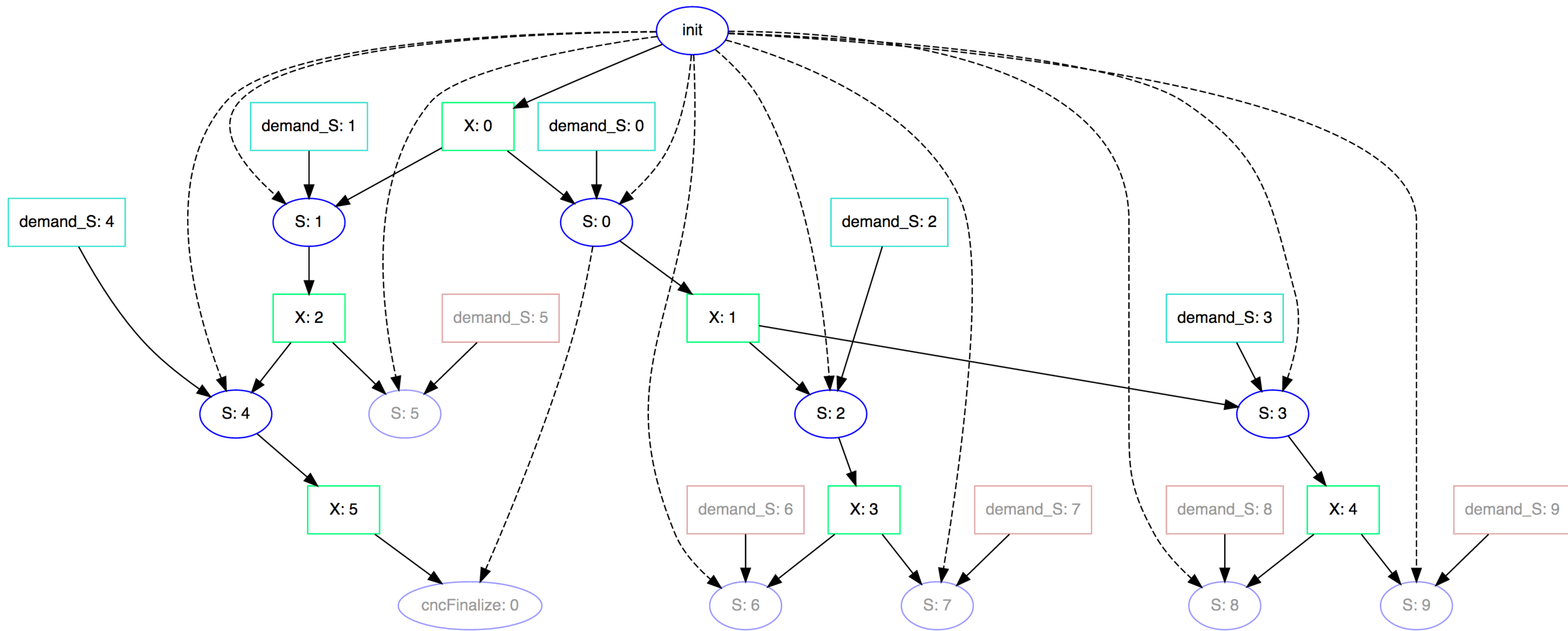
```
( $finalize: () ) <- [ X: 5 ];
```

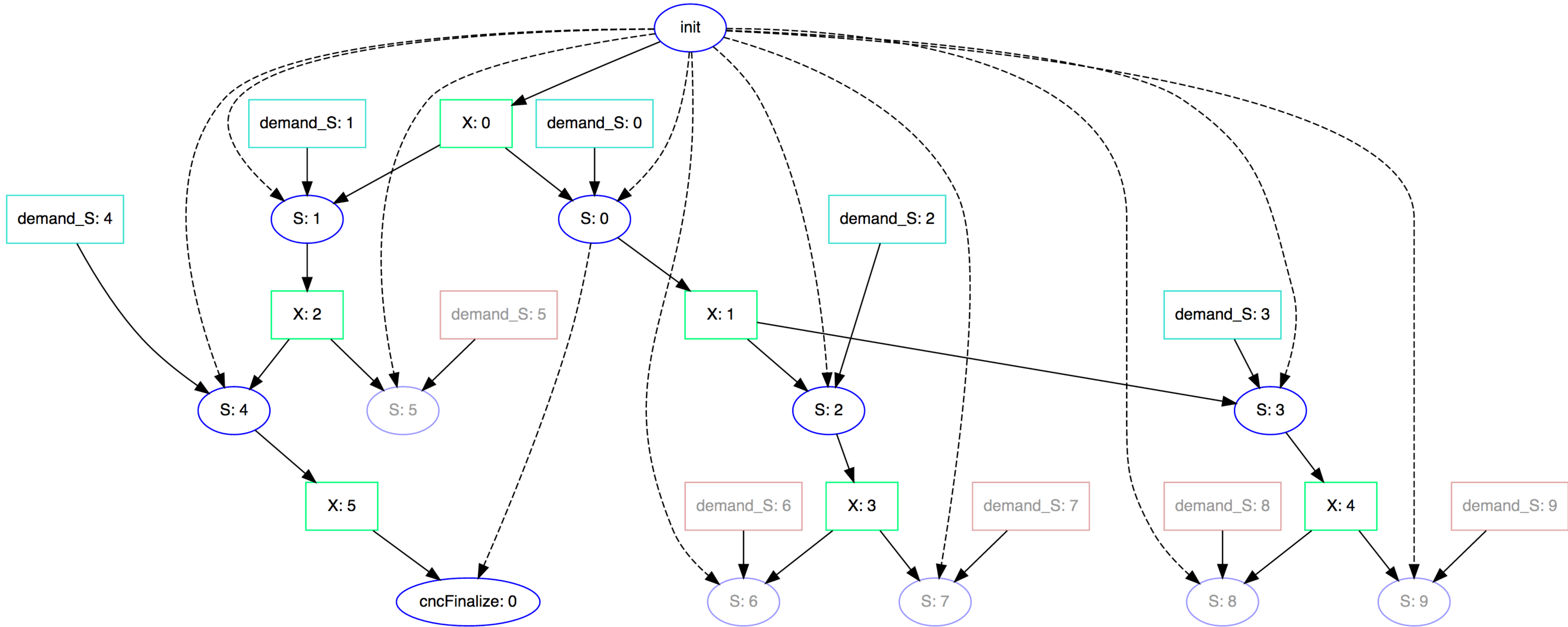










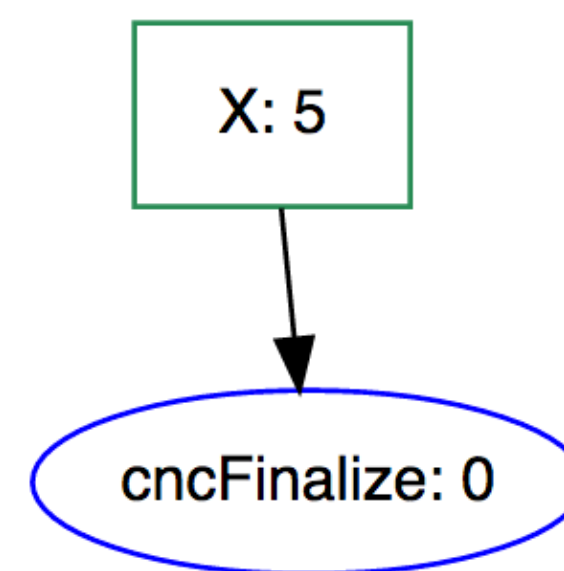


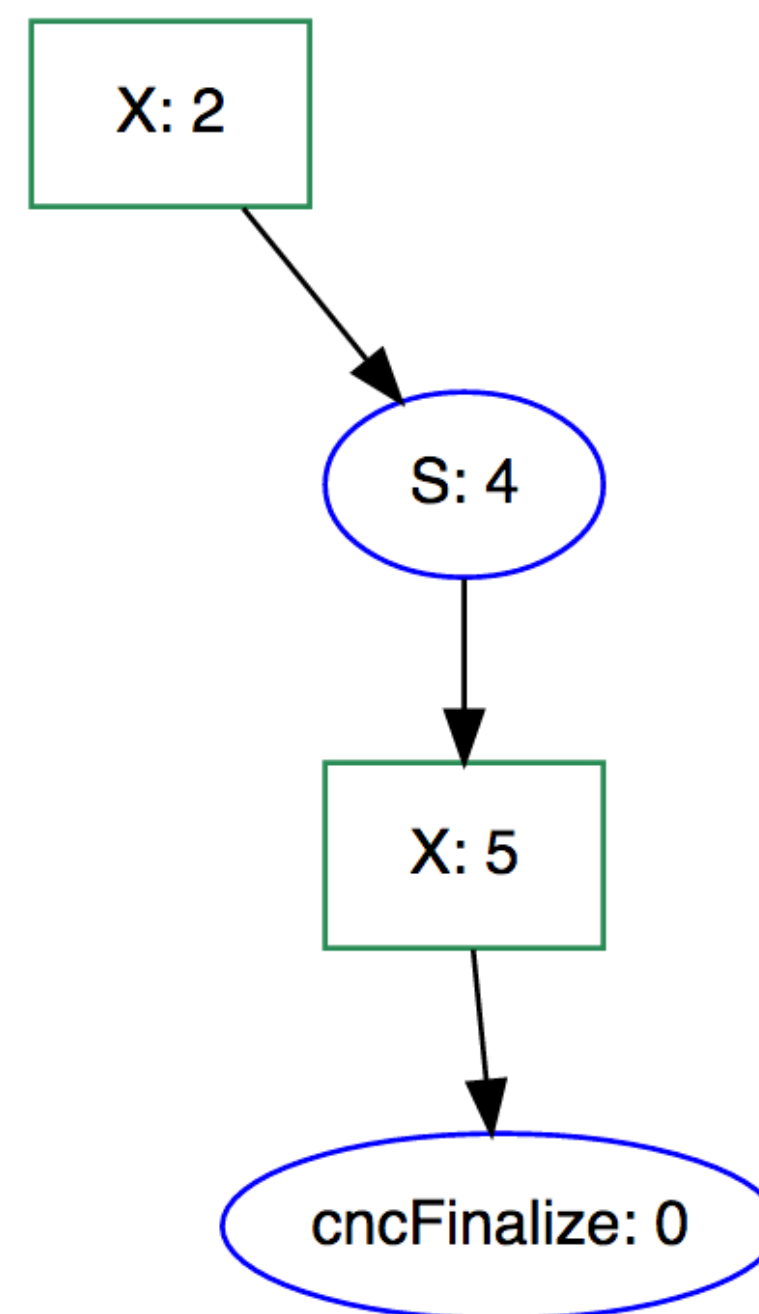
Future directions

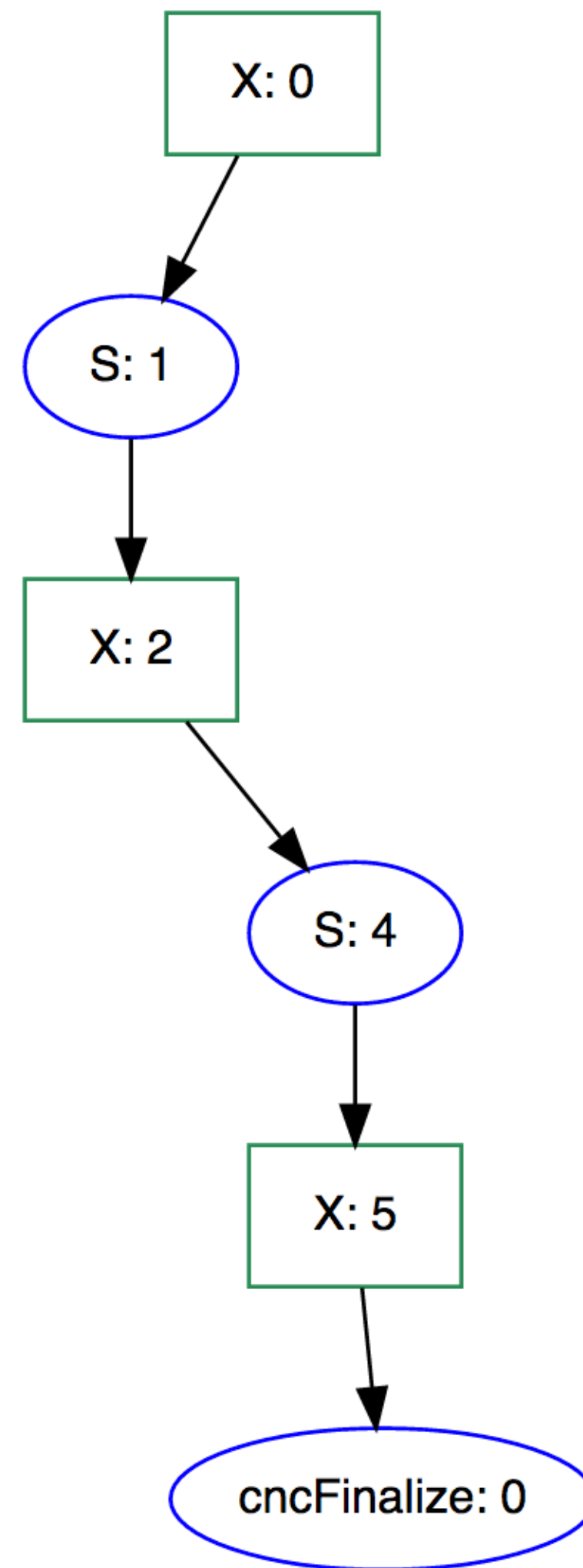
- Statically compute entire execution for some programs
- Combine with speculative execution

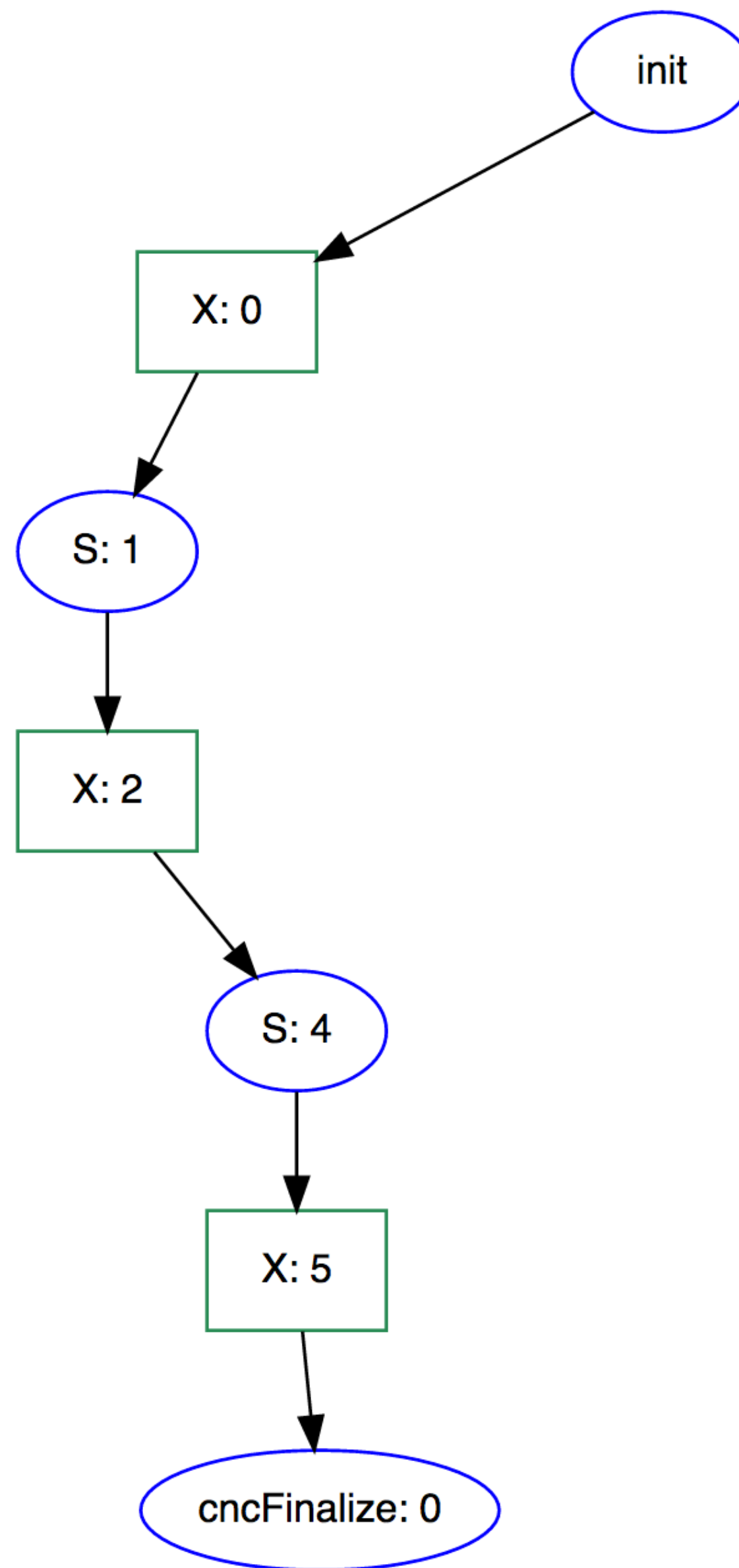
Future directions

- If all collections unambiguous,
 - Derive demand from finalize step back to inputs
 - Simulate entire execution graph before any execution
 - Steps and items visited are necessary to reach finalize
 - Eliminates prescribes, become pure dataflow (I/O) model









Demand with speculation

- Speculative execution: run steps whose inputs are available, but the program has not requested
- Used if program has both unambiguous and ambiguous collections
- Priority scheme
 - 1. Demanded steps; their output is necessary
 - 2. Regular prescribed steps; user thinks their output is necessary
 - 3. Speculative steps; only if we have extra resources

Wrap up

- Step inverses
- Demand driven execution
- Future possible work